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1 Overview of the OpenMP API

The collection of compiler directives, library routines, and environment variables that this document describes collectively define the specification of the OpenMP Application Programming Interface (OpenMP API) for parallelism in C, C++, and Fortran programs.

This specification provides a model for parallel programming that is portable across architectures from different vendors. Compilers from numerous vendors support the OpenMP API. More information about the OpenMP API can be found at the following web site

http://www.openmp.org

The directives, library routines, environment variables, and tool support that this document defines allow users to create, to manage, to debug and to analyze parallel programs while permitting portability. The directives extend the C, C++, and Fortran base languages with single program multiple data (SPMD) constructs, tasking constructs, device constructs, worksharing constructs, and synchronization constructs, and they provide support for sharing, mapping, and privatizing data. The functionality to control the runtime environment is provided by library routines and environment variables. Compilers that support the OpenMP API often include command line options to enable or to disable interpretation of some or all OpenMP directives.

1.1 Scope

The OpenMP API covers only user-directed parallelization, wherein the programmer explicitly specifies the actions to be taken by the compiler and runtime system in order to execute the program in parallel. OpenMP-compliant implementations are not required to check for data dependences, data conflicts, race conditions, or deadlocks. Compliant implementations also are not required to check for any code sequences that cause a program to be classified as non-conforming. Application developers are responsible for correctly using the OpenMP API to produce a conforming program. The OpenMP API does not cover compiler-generated automatic parallelization.
1.2 Glossary

1.2.1 Threading Concepts

thread An execution entity with a stack and associated threadprivate memory.

OpenMP thread A thread that is managed by the OpenMP implementation.

thread number A number that the OpenMP implementation assigns to an OpenMP thread. For threads within the same team, zero identifies the primary thread and consecutive numbers identify the other threads of this team.

idle thread An OpenMP thread that is not currently part of any parallel region.

thread-safe routine A routine that performs the intended function even when executed concurrently (by more than one thread).

processor Implementation-defined hardware unit on which one or more OpenMP threads can execute.

device An implementation-defined logical execution engine.

COMMENT: A device could have one or more processors.

host device The device on which the OpenMP program begins execution.

target device A device with respect to which the current device performs an operation, as specified by a device construct or an OpenMP device memory routine.

parent device For a given target region, the device on which the corresponding target construct was encountered.

1.2.2 OpenMP Language Terminology

base language A programming language that serves as the foundation of the OpenMP specification.

COMMENT: See Section 1.7 for a listing of current base languages for the OpenMP API.

base program A program written in a base language.

preprocessed code For C/C++, a sequence of preprocessing tokens that result from the first six phases of translation, as defined by the base language.

program order An ordering of operations performed by the same thread as determined by the execution sequence of operations specified by the base language.
COMMENT: For versions of C and C++ that include base language support for threading, program order corresponds to the sequenced before relation between operations performed by the same thread.

**structured block**
For C/C++, an executable statement, possibly compound, with a single entry at the top and a single exit at the bottom, or an OpenMP construct.

For Fortran, a *strictly structured block*, or a *loosely structured block*.

**structured block sequence**
A *structured block*, or, for C/C++, a sequence of two or more executable statements that together have a single entry at the top and a single exit at the bottom.

**strictly structured block**
A single Fortran **BLOCK** construct, with a single entry at the top and a single exit at the bottom.

**loosely structured block**
A block of executable constructs, where the first executable construct is not a Fortran **BLOCK** construct, with a single entry at the top and a single exit at the bottom, or an OpenMP construct.

COMMENT: In Fortran code, when a *strictly structured block* appears within an OpenMP construct, that OpenMP construct does not usually require a paired **end directive** to define the range of the OpenMP construct, while an OpenMP construct that contains a *loosely structured block* relies on the paired **end directive** to define the range of the OpenMP construct.

**compilation unit**
For C/C++, a translation unit.

For Fortran, a program unit.

**enclosing context**
For C/C++, the innermost scope enclosing an OpenMP directive.

For Fortran, the innermost scoping unit enclosing an OpenMP directive.

**directive**
A base language mechanism to specify OpenMP program behavior.

COMMENT: See Section 2.1 for a description of OpenMP directive syntax in each base language.

**white space**
A non-empty sequence of space and/or horizontal tab characters.

**OpenMP program**
A program that consists of a base program that is annotated with OpenMP directives or that calls OpenMP API runtime library routines.

**conforming program**
An OpenMP program that follows all rules and restrictions of the OpenMP specification.

**implementation code**
Implicit code that is introduced by the OpenMP implementation.

**metadirective**
A directive that conditionally resolves to another directive.
declarative directive  An OpenMP directive that may only be placed in a declarative context and results in one or more declarations only; it is not associated with the immediate execution of any user code or implementation code. For C++, if a declarative directive applies to a function declaration or definition and it is specified with one or more C++ attribute specifiers, the specified attributes must be applied to the function as permitted by the base language. For Fortran, a declarative directive must appear after any USE, IMPORT, and IMPLICIT statements in a declarative context.

executable directive  An OpenMP directive that appears in an executable context and results in implementation code and/or prescribes the manner in which associated user code must execute.

informational directive  An OpenMP directive that is neither declarative nor executable, but otherwise conveys user code properties to the compiler.

utility directive  An OpenMP directive that is neither declarative nor executable, but otherwise facilitates interactions with the compiler and/or supports code readability.

stand-alone directive  An OpenMP executable directive that has no associated user code, but may produce implementation code resulting from clauses in the directive.

construct  An OpenMP executable directive (and for Fortran, the paired end directive, if any) and the associated statement, loop nest or structured block, if any, not including the code in any called routines. That is, the lexical extent of an executable directive.

combined construct  A construct that is a shortcut for specifying one construct immediately nested inside another construct. A combined construct is semantically identical to that of explicitly specifying the first construct containing one instance of the second construct and no other statements.

composite construct  A construct that is composed of two constructs but does not have identical semantics to specifying one of the constructs immediately nested inside the other. A composite construct either adds semantics not included in the constructs from which it is composed or provides an effective nesting of the one construct inside the other that would otherwise be non-conforming.

constituent construct  For a given combined or composite construct, a construct from which it, or any one of its constituent constructs, is composed.

COMMENT: The constituent constructs of a target teams distribute parallel for simd construct are the following constructs: target, teams distribute parallel for simd, teams, distribute parallel for simd, distribute, parallel for simd, parallel, for simd, for, and simd.
leaf construct  For a given combined or composite construct, a constituent construct that is not itself a combined or composite construct.

COMMENT: The leaf constructs of a

target teams distribute parallel for simd construct are the following constructs: target, teams, distribute, parallel, for, and simd.

combined target construct  A combined construct that is composed of a target construct along with another construct.

region  All code encountered during a specific instance of the execution of a given construct, structured block sequence or OpenMP library routine. A region includes any code in called routines as well as any implementation code. The generation of a task at the point where a task generating construct is encountered is a part of the region of the encountering thread. However, an explicit task region that corresponds to a task generating construct is not part of the region of the encountering thread unless it is an included task region. The point where a target or teams directive is encountered is a part of the region of the encountering thread, but the region that corresponds to the target or teams directive is not.

COMMENTS:

A region may also be thought of as the dynamic or runtime extent of a construct or of an OpenMP library routine.

During the execution of an OpenMP program, a construct may give rise to many regions.

active parallel region  A parallel region that is executed by a team consisting of more than one thread.

inactive parallel region  A parallel region that is executed by a team of only one thread.

active target region  A target region that is executed on a device other than the device that encountered the target construct.

inactive target region  A target region that is executed on the same device that encountered the target construct.

sequential part  All code encountered during the execution of an initial task region that is not part of a parallel region corresponding to a parallel construct or a task region corresponding to a task construct.

COMMENTS:

A sequential part is enclosed by an implicit parallel region.
Executable statements in called routines may be in both a sequential part and any number of explicit parallel regions at different points in the program execution.

**primary thread** An OpenMP thread that has thread number 0. A primary thread may be an initial thread or the thread that encounters a parallel construct, creates a team, generates a set of implicit tasks, and then executes one of those tasks as thread number 0.

**parent thread** The thread that encountered the parallel construct and generated a parallel region is the parent thread of each of the threads in the team of that parallel region. The primary thread of a parallel region is the same thread as its parent thread with respect to any resources associated with an OpenMP thread.

**child thread** When a thread encounters a parallel construct, each of the threads in the generated parallel region’s team are child threads of the encountering thread. The target or teams region’s initial thread is not a child thread of the thread that encountered the target or teams construct.

**ancestor thread** For a given thread, its parent thread or one of its parent thread’s ancestor threads.

**descendent thread** For a given thread, one of its child threads or one of its child threads’ descendent threads.

**team** A set of one or more threads participating in the execution of a parallel region.

**league** The set of teams created by a teams construct.

**contention group** An initial thread and its descendent threads.

**implicit parallel region** An inactive parallel region that is not generated from a parallel construct. Implicit parallel regions surround the whole OpenMP program, all target regions, and all teams regions.

**initial thread** The thread that executes an implicit parallel region.

**initial team** The team that comprises an initial thread executing an implicit parallel region.

**nested construct** A construct (lexically) enclosed by another construct.

**closely nested construct** A construct nested inside another construct with no other construct nested between them.
explicit region A *region* that corresponds to either a *construct* of the same name or a library routine call that explicitly appears in the program.

nested region A *region* (dynamically) enclosed by another *region*. That is, a *region* generated from the execution of another *region* or one of its nested *regions*.

COMMENT: Some nestings are *conforming* and some are not. See Section 2.22 for the restrictions on nesting.

closely nested region A *region nested* inside another *region* with no *parallel region nested* between them.

strictly nested region A *region nested* inside another *region* with no other *explicit region nested* between them.

all threads All OpenMP *threads* participating in the *OpenMP program*.

current team All *threads* in the *team* executing the innermost enclosing *parallel region*.

encountering thread For a given *region*, the *thread* that encounters the corresponding *construct*.

all tasks All *tasks* participating in the *OpenMP program*.

current team tasks All *tasks* encountered by the corresponding *team*. The *implicit tasks* constituting the *parallel region* and any *descendent tasks* encountered during the execution of these *implicit tasks* are included in this set of tasks.

generating task For a given *region*, the task for which execution by a *thread* generated the *region*.

binding thread set The set of *threads* that are affected by, or provide the context for, the execution of a *region*.

The *binding thread set* for a given *region* can be *all threads* on a specified set of devices, *all threads* in a *contention group*, *all primary threads* executing an enclosing *teams region*, the *current team*, or the *encountering thread*.

COMMENT: The *binding thread set* for a particular *region* is described in its corresponding subsection of this specification.

binding task set The set of *tasks* that are affected by, or provide the context for, the execution of a *region*.

The *binding task set* for a given *region* can be *all tasks*, the *current team tasks*, *all tasks of the current team that are generated in the region*, the *binding implicit task*, or the *generating task*.

COMMENT: The *binding task set* for a particular *region* (if applicable) is described in its corresponding subsection of this specification.
binding region  The enclosing region that determines the execution context and limits the scope of
the effects of the bound region is called the binding region.

Binding region is not defined for regions for which the binding thread set is all
threads or the encountering thread, nor is it defined for regions for which the binding
task set is all tasks.

orphaned construct  A construct that gives rise to a region for which the binding thread set is the current
team, but is not nested within another construct that gives rise to the binding region.

worksharing construct  A construct that divides the work within its structured block into partitions, each of
which is executed exactly once by one of the threads in the team executing the
construct.

device construct  An OpenMP construct that accepts the device clause.

device routine  A function (for C/C++ and Fortran) or subroutine (for Fortran) that can be executed
on a target device, as part of a target region.

foreign runtime environment  A runtime environment that exists outside the OpenMP runtime with which the
OpenMP implementation may interoperate.

foreign execution context  A context that is instantiated from a foreign runtime environment in order to facilitate
execution on a given device.

foreign task  A unit of work executed in a foreign execution context.

indirect device invocation  An indirect call to the device version of a procedure on a device other than the host
device, through a function pointer (C/C++), a pointer to a member function (C++) or
a procedure pointer (Fortran) that refers to the host version of the procedure.

place  An unordered set of processors on a device.

place list  The ordered list that describes all OpenMP places available to the execution
environment.

place partition  An ordered list that corresponds to a contiguous interval in the OpenMP place list. It
describes the places currently available to the execution environment for a given
parallel region.

place number  A number that uniquely identifies a place in the place list, with zero identifying the
first place in the place list, and each consecutive whole number identifying the next
place in the place list.

thread affinity  A binding of threads to places within the current place partition.

SIMD instruction  A single machine instruction that can operate on multiple data elements.

SIMD lane  A software or hardware mechanism capable of processing one data element from a
SIMD instruction.
SIMD chunk  A set of iterations executed concurrently, each by a SIMD lane, by a single thread by means of SIMD instructions.

memory  A storage resource to store and to retrieve variables accessible by OpenMP threads.

memory space  A representation of storage resources from which memory can be allocated or deallocated. More than one memory space may exist.

memory allocator  An OpenMP object that fulfills requests to allocate and to deallocate memory for program variables from the storage resources of its associated memory space.

handle  An opaque reference that uniquely identifies an abstraction.

1.2.3 Loop Terminology

canonical loop nest  A loop nest that complies with the rules and restrictions defined in Section 2.11.1.

loop-associated directive  An OpenMP executable directive for which the associated user code must be a canonical loop nest.

associated loop  A loop from a canonical loop nest that is controlled by a given loop-associated directive.

loop nest depth  For a canonical loop nest, the maximal number of loops, including the outermost loop, that can be associated with a loop-associated directive.

logical iteration space  For a loop-associated directive, the sequence \(0, \ldots, N - 1\) where \(N\) is the number of iterations of the loops associated with the directive. The logical numbering denotes the sequence in which the iterations would be executed if the set of associated loops were executed sequentially.

logical iteration  An iteration from the associated loops of a loop-associated directive, designated by a logical number from the logical iteration space of the associated loops.

logical iteration vector space  For a loop-associated directive with \(n\) associated nested loops, the set of \(n\)-tuples \((i_1, \ldots, i_n)\). For the \(k^{th}\) associated loop, from outermost to innermost, \(i_k\) is its logical iteration number as if it was the only associated loop.

logical iteration vector  An iteration from the associated nested loops of a loop-associated directive, where \(n\) is the number of associated loops, designated by an \(n\)-tuple from the logical iteration vector space of the associated loops.

lexicographic order  The total order of two logical iteration vectors \(\omega_a = (i_1, \ldots, i_n)\) and \(\omega_b = (j_1, \ldots, j_n)\), denoted by \(\omega_a \leq_{\text{lex}} \omega_b\), where either \(\omega_a = \omega_b\) or \(\exists m \in \{1, \ldots, n\}\) such that \(i_m < j_m\) and \(i_k = j_k\) for all \(k \in \{1, \ldots, m - 1\}\).
The partial order of two logical iteration vectors \( \omega_a = (i_1, \ldots, i_n) \) and \( \omega_b = (j_1, \ldots, j_n) \), denoted by \( \omega_a \leq_{\text{product}} \omega_b \), where \( i_k \leq j_k \) for all \( k \in \{1, \ldots, n\} \).

A construct that is replaced by the loops that result from applying the transformation as defined by its directive to its associated loops.

A loop that is generated by a loop transformation construct and is one of the resulting loops that replace the construct.

A loop that includes at least one SIMD chunk.

For a loop nest, a loop for which a loop bound references the iteration variable of a surrounding loop in the loop nest.

A loop that has no intervening code between it and the body of its surrounding loop. The outermost loop of a loop nest is always perfectly nested.

A loop nest, consisting of loops that may be associated with the same loop-associated directive, that has cross-iteration dependences. An iteration is dependent on one or more lexicographically earlier iterations.

COMMENT: The ordered clause parameter on a worksharing-loop directive identifies the loops associated with the doacross loop nest.

### 1.2.4 Synchronization Terminology

A point in the execution of a program encountered by a team of threads, beyond which no thread in the team may execute until all threads in the team have reached the barrier and all explicit tasks generated by the team have executed to completion. If cancellation has been requested, threads may proceed to the end of the canceled region even if some threads in the team have not reached the barrier.

An action that cancels (that is, aborts) an OpenMP region and causes executing implicit or explicit tasks to proceed to the end of the canceled region.

A point at which implicit and explicit tasks check if cancellation has been requested. If cancellation has been observed, they perform the cancellation.

An operation that a thread performs to enforce consistency between its view and other threads’ view of memory.

The set of devices for which a flush operation may enforce memory consistency.

Properties that determine the manner in which a flush operation enforces memory consistency. These properties are:
• *strong*: flushes a set of variables from the current thread’s temporary view of the memory to the memory;

• *release*: orders memory operations that precede the flush before memory operations performed by a different thread with which it synchronizes;

• *acquire*: orders memory operations that follow the flush after memory operations performed by a different thread that synchronizes with it.

**COMMENT:** Any flush operation has one or more flush properties.

**strong flush** A flush operation that has the *strong flush property*.

**release flush** A flush operation that has the *release flush property*.

**acquire flush** A flush operation that has the *acquire flush property*.

**atomic operation** An operation that is specified by an atomic construct or is implicitly performed by the OpenMP implementation and that atomically accesses and/or modifies a specific storage location.

**atomic read** An atomic operation that is specified by an atomic construct on which the read clause is present.

**atomic write** An atomic operation that is specified by an atomic construct on which the write clause is present.

**atomic update** An atomic operation that is specified by an atomic construct on which the update clause is present.

**atomic captured update** An atomic update operation that is specified by an atomic construct on which the capture clause is present.

**atomic conditional update** An atomic update operation that is specified by an atomic construct on which the compare clause is present.

**read-modify-write** An atomic operation that reads and writes to a given storage location.

**COMMENT:** Any atomic update is a read-modify-write operation.

**sequentially consistent atomic construct** An atomic construct for which the seq_cst clause is specified.

**non-sequentially consistent atomic construct** An atomic construct for which the seq_cst clause is not specified

**sequentially consistent atomic operation** An atomic operation that is specified by a sequentially consistent atomic construct.
1.2.5 Tasking Terminology

**task** A specific instance of executable code and its data environment that the OpenMP implementation can schedule for execution by threads.

**task region** A region consisting of all code encountered during the execution of a task.

COMMENT: A **parallel region** consists of one or more implicit task regions.

**implicit task** A task generated by an implicit parallel region or generated when a parallel construct is encountered during execution.

**binding implicit task** The implicit task of the current thread team assigned to the encountering thread.

**explicit task** A task that is not an implicit task.

**initial task** An implicit task associated with an implicit parallel region.

**current task** For a given thread, the task corresponding to the task region in which it is executing.

**encountering task** For a given region, the current task of the encountering thread.

**child task** A task is a child task of its generating task region. A child task region is not part of its generating task region.

**sibling tasks** Tasks that are child tasks of the same task region.

**descendent task** A task that is the child task of a task region or of one of its descendent task regions.

**task completion** A condition that is satisfied when a thread reaches the end of the executable code that is associated with the task and any allow-completion event that is created for the task has been fulfilled.

COMMENT: Completion of the initial task that is generated when the program begins occurs at program exit.

**task scheduling point** A point during the execution of the current task region at which it can be suspended to be resumed later; or the point of task completion, after which the executing thread may switch to a different task region.

**task switching** The act of a thread switching from the execution of one task to another task.

**tied task** A task that, when its task region is suspended, can be resumed only by the same thread that was executing it before suspension. That is, the task is tied to that thread.

**untied task** A task that, when its task region is suspended, can be resumed by any thread in the team. That is, the task is not tied to any thread.
<table>
<thead>
<tr>
<th>Term</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>undeferred task</strong></td>
<td>A task for which execution is not deferred with respect to its generating task region. That is, its generating task region is suspended until execution of the structured block associated with the undeferred task is completed.</td>
</tr>
<tr>
<td><strong>included task</strong></td>
<td>A task for which execution is sequentially included in the generating task region. That is, an included task is undeferred and executed by the encountering thread.</td>
</tr>
<tr>
<td><strong>merged task</strong></td>
<td>A task for which the data environment, inclusive of ICVs, is the same as that of its generating task region.</td>
</tr>
<tr>
<td><strong>mergeable task</strong></td>
<td>A task that may be a merged task if it is an undeferred task or an included task.</td>
</tr>
<tr>
<td><strong>final task</strong></td>
<td>A task that forces all of its child tasks to become final and included tasks.</td>
</tr>
<tr>
<td><strong>task dependence</strong></td>
<td>An ordering relation between two sibling tasks: the dependent task and a previously generated predecessor task. The task dependence is fulfilled when the predecessor task has completed.</td>
</tr>
<tr>
<td><strong>dependent task</strong></td>
<td>A task that because of a task dependence cannot be executed until its predecessor tasks have completed.</td>
</tr>
<tr>
<td><strong>mutually exclusive tasks</strong></td>
<td>Tasks that may be executed in any order, but not at the same time.</td>
</tr>
<tr>
<td><strong>predecessor task</strong></td>
<td>A task that must complete before its dependent tasks can be executed.</td>
</tr>
<tr>
<td><strong>task synchronization construct</strong></td>
<td>A taskwait, taskgroup, or a barrier construct.</td>
</tr>
<tr>
<td><strong>task generating construct</strong></td>
<td>A construct that generates one or more explicit tasks that are child tasks of the encountering task.</td>
</tr>
<tr>
<td><strong>target task</strong></td>
<td>A mergeable and untied task that is generated by a device construct or a call to a device memory routine and that coordinates activity between the current device and the target device.</td>
</tr>
<tr>
<td><strong>taskgroup set</strong></td>
<td>A set of tasks that are logically grouped by a taskgroup region.</td>
</tr>
</tbody>
</table>
1.2.6 Data Terminology

**variable**
A named data storage block, for which the value can be defined and redefined during the execution of a program.

COMMENT: An array element or structure element is a variable that is part of another variable.

**scalar variable**
For C/C++, a scalar variable, as defined by the base language.
For Fortran, a scalar variable with intrinsic type, as defined by the base language, excluding character type.

**aggregate variable**
A variable, such as an array or structure, composed of other variables.

**array section**
A designated subset of the elements of an array that is specified using a subscript notation that can select more than one element.

**array item**
An array, an array section, or an array element.

**shape-operator**
For C/C++, an array shaping operator that reinterprets a pointer expression as an array with one or more specified dimensions.

**implicit array**
For C/C++, the set of array elements of non-array type $T$ that may be accessed by applying a sequence of [ ] operators to a given pointer that is either a pointer to type $T$ or a pointer to a multidimensional array of elements of type $T$.
For Fortran, the set of array elements for a given array pointer.

COMMENT: For C/C++, the implicit array for pointer p with type $T$ (*p)[10] consists of all accessible elements p[i][j], for all $i$ and $j=0,1,...,9$.

**base pointer**
For C/C++, an lvalue pointer expression that is used by a given lvalue expression or array section to refer indirectly to its storage, where the lvalue expression or array section is part of the implicit array for that lvalue pointer expression.
For Fortran, a data pointer that appears last in the designator for a given variable or array section, where the variable or array section is part of the pointer target for that data pointer.

COMMENT: For the array section (*p0).x0[k1].p1->p2[k2].x1[k3].x2[4][0:n], where identifiers pi have a pointer type declaration and identifiers xi have an array type declaration, the base pointer is: (*p0).x0[k1].p1->p2.

**named pointer**
For C/C++, the base pointer of a given lvalue expression or array section, or the base pointer of one of its named pointers.
For Fortran, the base pointer of a given variable or array section, or the base pointer of one of its named pointers.
COMMENT: For the array section
(*p0).x0[k1].p1->p2[k2].x1[k3].x2[4][0:n], where identifiers p\textsubscript{i} have a
pointer type declaration and identifiers x\textsubscript{i} have an array type declaration,
the named pointers are: p0, (*p0).x0[k1].p1, and (*p0).x0[k1].p1->p2.

containing array  For C/C++, a non-subscripted array (a containing array) that appears in a given
lvalue expression or array section, where the lvalue expression or array section is part
of that containing array.

For Fortran, an array (a containing array) without the POINTER attribute and
without a subscript list that appears in the designator of a given variable or array
section, where the variable or array section is part of that containing array.

COMMENT: For the array section
(*p0).x0[k1].p1->p2[k2].x1[k3].x2[4][0:n], where identifiers p\textsubscript{i} have a
pointer type declaration and identifiers x\textsubscript{i} have an array type declaration,
the containing arrays are: (*p0).x0[k1].p1->p2[k2].x1 and
(*p0).x0[k1].p1->p2[k2].x1[k3].x2.

base array  For C/C++, a containing array of a given lvalue expression or array section that does
not appear in the expression of any of its other containing arrays.

For Fortran, a containing array of a given variable or array section that does not
appear in the designator of any of its other containing arrays.

COMMENT: For the array section
(*p0).x0[k1].p1->p2[k2].x1[k3].x2[4][0:n], where identifiers p\textsubscript{i} have a
pointer type declaration and identifiers x\textsubscript{i} have an array type declaration,
the base array is: (*p0).x0[k1].p1->p2[k2].x1[k3].x2.

named array  For C/C++, a containing array of a given lvalue expression or array section, or a
containing array of one of its named pointers.

For Fortran, a containing array of a given variable or array section, or a containing
array of one of its named pointers.

COMMENT: For the array section
(*p0).x0[k1].p1->p2[k2].x1[k3].x2[4][0:n], where identifiers p\textsubscript{i} have a
pointer type declaration and identifiers x\textsubscript{i} have an array type declaration,
the named arrays are: (*p0).x0, (*p0).x0[k1].p1->p2[k2].x1, and
(*p0).x0[k1].p1->p2[k2].x1[k3].x2.

base expression  The base array of a given array section or array element, if it exists; otherwise, the
base pointer of the array section or array element.

COMMENT: For the array section
(*p0).x0[k1].p1->p2[k2].x1[k3].x2[4][0:n], where identifiers p\textsubscript{i} have a
pointer type declaration and identifiers x\textsubscript{i} have an array type declaration,
the base expression is: (*p0).x0[k1].p1->p2[k2].x1[k3].x2.
More examples for C/C++:

- The base expression for x[i] and for x[i:n] is x, if x is an array or pointer.
- The base expression for x[5][i] and for x[5][i:n] is x, if x is a pointer to an array or x is 2-dimensional array.
- The base expression for y[5][i] and for y[5][i:n] is y[5], if y is an array of pointers or y is a pointer to a pointer.

Examples for Fortran:

- The base expression for x(i) and for x(i:j) is x.

attached pointer A pointer variable in a device data environment to which the effect of a map clause assigns the address of an object, minus some offset, that is created in the device data environment. The pointer is an attached pointer for the remainder of its lifetime in the device data environment.

simply contiguous array section An array section that statically can be determined to have contiguous storage or that, in Fortran, has the CONTIGUOUS attribute.

structure A structure is a variable that contains one or more variables.

For C/C++: Implemented using struct types.
For C++: Implemented using class types.
For Fortran: Implemented using derived types.

string literal For C/C++, a string literal.

For Fortran, a character literal constant.

private variable With respect to a given set of task regions or SIMD lanes that bind to the same parallel region, a variable for which the name provides access to a different block of storage for each task region or SIMD lane.

A variable that is part of another variable (as an array or structure element) cannot be made private independently of other components. If a variable is privatized, its components are also private.

shared variable With respect to a given set of task regions that bind to the same parallel region, a variable for which the name provides access to the same block of storage for each task region.

A variable that is part of another variable (as an array or structure element) cannot be shared independently of the other components, except for static data members of C++ classes.
A \textit{variable} that is replicated, one instance per \textit{thread}, by the OpenMP implementation. Its name then provides access to a different block of storage for each \textit{thread}.

A \textit{variable} that is part of another variable (as an array or structure element) cannot be made \textit{threadprivate} independently of the other components, except for static data members of C++ classes. If a \textit{variable} is made \textit{threadprivate}, its components are also \textit{threadprivate}.

The set of \textit{threadprivate variables} associated with each \textit{thread}.

The \textit{variables} associated with the execution of a given \textit{region}.

The initial \textit{data environment} associated with a device.

An address of an object that may be referenced on a \textit{target device}.

An \textit{implementation defined handle} that refers to a \textit{device address}.

An original \textit{variable} in a \textit{data environment} with a corresponding \textit{variable} in a device \textit{data environment}.

COMMENT: The original and corresponding \textit{variables} may share storage.

An operation that defines how variables of given type are to be mapped or updated with respect to a device \textit{data environment}.

A \textit{mapper} that is defined by a \textit{declare mapper} directive.

The process that determines the final map types of the map operations that result from mapping a variable with a \textit{user-defined mapper}.

A \textit{type} that is valid for a \textit{mapped variable}. If a type is composed from other types (such as the type of an array or structure element) and any of the other types are not mappable then the type is not mappable.

COMMENT: Pointer types are \textit{mappable} but the memory block to which the pointer refers is not \textit{mapped}.

For C, the type must be a complete type.

For C++, the type must be a complete type.

In addition, for class types:

• All member functions accessed in any \textit{target} region must appear in a declare \textit{target} directive.

For Fortran, no restrictions on the type except that for derived types:
• All type-bound procedures accessed in any target region must appear in a
declare target directive.

defined For variables, the property of having a valid value.
For C, for the contents of variables, the property of having a valid value.
For C++, for the contents of variables of POD (plain old data) type, the property of
having a valid value.
For variables of non-POD class type, the property of having been constructed but not
subsequently destructed.
For Fortran, for the contents of variables, the property of having a valid value. For
the allocation or association status of variables, the property of having a valid status.

COMMENT: Programs that rely upon variables that are not defined are
non-conforming programs.

class type For C++, variables declared with one of the class, struct, or union keywords.

static storage duration For C/C++, the lifetime of an object with static storage duration, as defined by the
base language.

For Fortran, the lifetime of a variable with a SAVE attribute, implicit or explicit, a
common block object or a variable declared in a module.

1.2.7 Implementation Terminology

supported active levels of parallelism An implementation-defined maximum number of active parallel regions that may
enclose any region of code in the program.

OpenMP API support Support of at least one active level of parallelism.

nested parallelism support Support of more than one active level of parallelism.

internal control variable A conceptual variable that specifies runtime behavior of a set of threads or tasks in
an OpenMP program.

COMMENT: The acronym ICV is used interchangeably with the term
internal control variable in the remainder of this specification.

OpenMP Additional Definitions document A document that exists outside of the OpenMP specification and defines additional
values that may be used in a conforming program. The OpenMP Additional
Definitions document is available at http://www.openmp.org/.
compliant implementation An implementation of the OpenMP specification that compiles and executes any conforming program as defined by the specification.

COMMENT: A compliant implementation may exhibit unspecified behavior when compiling or executing a non-conforming program.

unspecified behavior A behavior or result that is not specified by the OpenMP specification or not known prior to the compilation or execution of an OpenMP program.

Such unspecified behavior may result from:

- Issues documented by the OpenMP specification as having unspecified behavior.
- A non-conforming program.
- A conforming program exhibiting an implementation-defined behavior.

implementation defined Behavior that must be documented by the implementation, and is allowed to vary among different compliant implementations. An implementation is allowed to define this behavior as unspecified.

COMMENT: All features that have implementation-defined behavior are documented in Appendix A.

deprecated For a construct, clause, or other feature, the property that it is normative in the current specification but is considered obsolescent and will be removed in the future.

1.2.8 Tool Terminology

tool Code that can observe and/or modify the execution of an application.

first-party tool A tool that executes in the address space of the program that it is monitoring.

third-party tool A tool that executes as a separate process from the process that it is monitoring and potentially controlling.

activated tool A first-party tool that successfully completed its initialization.

event A point of interest in the execution of a thread.

native thread A thread defined by an underlying thread implementation.

tool callback A function that a tool provides to an OpenMP implementation to invoke when an associated event occurs.

registering a callback Providing a tool callback to an OpenMP implementation.
dispatching a callback
at an event

Processing a callback when an associated event occurs in a manner consistent with

the return code provided when a first-party tool registered the callback.

thread state
An enumeration type that describes the current OpenMP activity of a thread. A
thread can be in only one state at any time.

wait identifier
A unique opaque handle associated with each data object (for example, a lock) that
the OpenMP runtime uses to enforce mutual exclusion and potentially to cause a
thread to wait actively or passively.

frame
A storage area on a thread’s stack associated with a procedure invocation. A frame
includes space for one or more saved registers and often also includes space for saved
arguments, local variables, and padding for alignment.

canonical frame
address
An address associated with a procedure frame on a call stack that was the value of the
stack pointer immediately prior to calling the procedure for which the frame
represents the invocation.

runtime entry point
A function interface provided by an OpenMP runtime for use by a tool. A runtime
entry point is typically not associated with a global function symbol.

trace record
A data structure in which to store information associated with an occurrence of an
event.

native trace record
A trace record for an OpenMP device that is in a device-specific format.

signal
A software interrupt delivered to a thread.

signal handler
A function called asynchronously when a signal is delivered to a thread.

async signal safe
The guarantee that interruption by signal delivery will not interfere with a set of
operations. An async signal safe runtime entry point is safe to call from a signal
handler.

code block
A contiguous region of memory that contains code of an OpenMP program to be
executed on a device.

OMPT
An interface that helps a first-party tool monitor the execution of an OpenMP
program.

OMPT interface state
A state that indicates the permitted interactions between a first-party tool and the
OpenMP implementation.

OMPT active
An OMPT interface state in which the OpenMP implementation is prepared to accept
runtime calls from a first party tool and will dispatch any registered callbacks and in
which a first-party tool can invoke runtime entry points if not otherwise restricted.
OMPT pending  An *OMPT interface state* in which the OpenMP implementation can only call functions to initialize a *first party tool* and in which a *first-party tool* cannot invoke *runtime entry points*.

OMPT inactive  An *OMPT interface state* in which the OpenMP implementation will not make any callbacks and in which a *first-party tool* cannot invoke *runtime entry points*.

OMPD  An interface that helps a *third-party tool* inspect the OpenMP state of a program that has begun execution.

OMPD library  A dynamically loadable library that implements the *OMPD* interface.

Image file  An executable or shared library.

Address space  A collection of logical, virtual, or physical memory address ranges that contain code, stack, and/or data. Address ranges within an address space need not be contiguous. An address space consists of one or more *segments*.

Segment  A portion of an address space associated with a set of address ranges.

OpenMP architecture  The architecture on which an OpenMP *region* executes.

Tool architecture  The architecture on which an *OMPD* tool executes.

OpenMP process  A collection of one or more *threads* and *address spaces*. A process may contain *threads* and *address spaces* for multiple OpenMP architectures. At least one thread in an OpenMP process is an OpenMP *thread*. A process may be live or a core file.

Address space handle  A *handle* that refers to an *address space* within an OpenMP process.

Thread handle  A *handle* that refers to an OpenMP *thread*.

Parallel handle  A *handle* that refers to an OpenMP parallel *region*.

Task handle  A *handle* that refers to an OpenMP task *region*.

Descendent handle  An output *handle* that is returned from the *OMPD* library in a function that accepts an input *handle*: the output *handle* is a descendent of the input *handle*.

Ancestor handle  An input *handle* that is passed to the *OMPD* library in a function that returns an output *handle*: the input *handle* is an ancestor of the output *handle*. For a given *handle*, the ancestors of the *handle* are also the ancestors of the *handle*’s descendent.

COMMENT: A tool cannot use a *handle* in an *OMPD* call if any ancestor of the *handle* has been released, except for *OMPD* calls that release it.
tool context  An opaque reference provided by a tool to an OMPD library. A tool context uniquely identifies an abstraction.

address space context  A tool context that refers to an address space within a process.

thread context  A tool context that refers to a native thread.

native thread identifier  An identifier for a native thread defined by a thread implementation.

1.3 Execution Model

The OpenMP API uses the fork-join model of parallel execution. Multiple threads of execution perform tasks defined implicitly or explicitly by OpenMP directives. The OpenMP API is intended to support programs that will execute correctly both as parallel programs (multiple threads of execution and a full OpenMP support library) and as sequential programs (directives ignored and a simple OpenMP stubs library). However, a conforming OpenMP program may execute correctly as a parallel program but not as a sequential program, or may produce different results when executed as a parallel program compared to when it is executed as a sequential program. Further, using different numbers of threads may result in different numeric results because of changes in the association of numeric operations. For example, a serial addition reduction may have a different pattern of addition associations than a parallel reduction. These different associations may change the results of floating-point addition.

An OpenMP program begins as a single thread of execution, called an initial thread. An initial thread executes sequentially, as if the code encountered is part of an implicit task region, called an initial task region, that is generated by the implicit parallel region surrounding the whole program.

The thread that executes the implicit parallel region that surrounds the whole program executes on the host device. An implementation may support other devices besides the host device. If supported, these devices are available to the host device for offloading code and data. Each device has its own threads that are distinct from threads that execute on another device. Threads cannot migrate from one device to another device. Each device is identified by a device number. The device number for the host device is the value of the total number of non-host devices, while each non-host device has a unique device number that is greater than or equal to zero and less than the device number for the host device.

When a target construct is encountered, a new target task is generated. The target task region encloses the target region. The target task is complete after the execution of the target region is complete.

When a target task executes, the enclosed target region is executed by an initial thread. The initial thread executes sequentially, as if the target region is part of an initial task region that is generated by an implicit parallel region. The initial thread may execute on the requested target...
device, if it is available and supported. If the target device does not exist or the implementation
does not support it, all target regions associated with that device execute on the host device.

The implementation must ensure that the target region executes as if it were executed in the data
environment of the target device unless an if clause is present and the if clause expression
evaluates to false.

The teams construct creates a league of teams, where each team is an initial team that comprises
an initial thread that executes the teams region. Each initial thread executes sequentially, as if the
code encountered is part of an initial task region that is generated by an implicit parallel region
associated with each team. Whether the initial threads concurrently execute the teams region is
unspecified, and a program that relies on their concurrent execution for the purposes of
synchronization may deadlock.

If a construct creates a data environment, the data environment is created at the time the construct is
encountered. The description of a construct defines whether it creates a data environment.

When any thread encounters a parallel construct, the thread creates a team of itself and zero or
more additional threads and becomes the primary thread of the new team. A set of implicit tasks,
one per thread, is generated. The code for each task is defined by the code inside the parallel
construct. Each task is assigned to a different thread in the team and becomes tied; that is, it is
always executed by the thread to which it is initially assigned. The task region of the task being
executed by the encountering thread is suspended, and each member of the new team executes its
implicit task. An implicit barrier occurs at the end of the parallel region. Only the primary
thread resumes execution beyond the end of the parallel construct, resuming the task region
that was suspended upon encountering the parallel construct. Any number of parallel
constructs can be specified in a single program.

parallel regions may be arbitrarily nested inside each other. If nested parallelism is disabled, or
is not supported by the OpenMP implementation, then the new team that is created by a thread that
encounters a parallel construct inside a parallel region will consist only of the
encountering thread. However, if nested parallelism is supported and enabled, then the new team
can consist of more than one thread. A parallel construct may include a proc_bind clause to
specify the places to use for the threads in the team within the parallel region.

When any team encounters a worksharing construct, the work inside the construct is divided among
the members of the team, and executed cooperatively instead of being executed by every thread. An
implicit barrier occurs at the end of any region that corresponds to a worksharing construct for
which the nowait clause is not specified. Redundant execution of code by every thread in the
team resumes after the end of the worksharing construct.

When any thread encounters a task generating construct, one or more explicit tasks are generated.
Execution of explicitly generated tasks is assigned to one of the threads in the current team, subject
to the thread’s availability to execute work. Thus, execution of the new task could be immediate, or
deferred until later according to task scheduling constraints and thread availability. Threads are
allowed to suspend the current task region at a task scheduling point in order to execute a different
task. If the suspended task region is for a tied task, the initially assigned thread later resumes
execution of the suspended task region. If the suspended task region is for an untied task, then any
thread may resume its execution. Completion of all explicit tasks bound to a given parallel region is
 guaranteed before the primary thread leaves the implicit barrier at the end of the region.
Completion of a subset of all explicit tasks bound to a given parallel region may be specified
through the use of task synchronization constructs. Completion of all explicit tasks bound to the
implicit parallel region is guaranteed by the time the program exits.

When any thread encounters a \texttt{simd} construct, the iterations of the loop associated with the
construct may be executed concurrently using the SIMD lanes that are available to the thread.

When a \texttt{loop} construct is encountered, the iterations of the loop associated with the construct are
executed in the context of its encountering threads, as determined according to its binding region. If
the \texttt{loop} region binds to a \texttt{teams} region, the region is encountered by the set of primary threads
that execute the \texttt{teams} region. If the \texttt{loop} region binds to a \texttt{parallel} region, the region is
encountered by the team of threads that execute the \texttt{parallel} region. Otherwise, the region is
encountered by a single thread.

If the \texttt{loop} region binds to a \texttt{teams} region, the encountering threads may continue execution
after the \texttt{loop} region without waiting for all iterations to complete; the iterations are guaranteed to
complete before the end of the \texttt{teams} region. Otherwise, all iterations must complete before the
encountering threads continue execution after the \texttt{loop} region. All threads that encounter the
\texttt{loop} construct may participate in the execution of the iterations. Only one of these threads may
execute any given iteration.

The cancel construct can alter the previously described flow of execution in an OpenMP region.
The effect of the cancel construct depends on its \texttt{construct-type-clause}. If a task encounters a
cancel construct with a \texttt{taskgroup construct-type-clause}, then the task activates cancellation
and continues execution at the end of its \texttt{task} region, which implies completion of that task. Any
other task in that \texttt{taskgroup} that has begun executing completes execution unless it encounters a
cancellation point construct, in which case it continues execution at the end of its \texttt{task}
region, which implies its completion. Other tasks in that \texttt{taskgroup} region that have not begun
execution are aborted, which implies their completion.

For all other \texttt{construct-type-clause} values, if a thread encounters a cancel construct, it activates
cancellation of the innermost enclosing region of the type specified and the thread continues
execution at the end of that region. Threads check if cancellation has been activated for their region
at cancellation points and, if so, also resume execution at the end of the canceled region.

If cancellation has been activated, regardless of \texttt{construct-type-clause}, threads that are waiting
inside a barrier other than an implicit barrier at the end of the canceled region exit the barrier and
resume execution at the end of the canceled region. This action can occur before the other threads
reach that barrier.

Synchronization constructs and library routines are available in the OpenMP API to coordinate
tasks and data access in \texttt{parallel} regions. In addition, library routines and environment
variables are available to control or to query the runtime environment of OpenMP programs.
The OpenMP specification makes no guarantee that input or output to the same file is synchronous when executed in parallel. In this case, the programmer is responsible for synchronizing input and output processing with the assistance of OpenMP synchronization constructs or library routines.

For the case where each thread accesses a different file, the programmer does not need to synchronize access.

All concurrency semantics defined by the base language with respect to threads of execution apply to OpenMP threads, unless specified otherwise.

### 1.4 Memory Model

#### 1.4.1 Structure of the OpenMP Memory Model

The OpenMP API provides a relaxed-consistency, shared-memory model. All OpenMP threads have access to a place to store and to retrieve variables, called the *memory*. A given storage location in the memory may be associated with one or more devices, such that only threads on associated devices have access to it. In addition, each thread is allowed to have its own *temporary view* of the memory. The temporary view of memory for each thread is not a required part of the OpenMP memory model, but can represent any kind of intervening structure, such as machine registers, cache, or other local storage, between the thread and the memory. The temporary view of memory allows the thread to cache variables and thereby to avoid going to memory for every reference to a variable. Each thread also has access to another type of memory that must not be accessed by other threads, called *threadprivate memory*.

A directive that accepts data-sharing attribute clauses determines two kinds of access to variables used in the directive’s associated structured block: shared and private. Each variable referenced in the structured block has an original variable, which is the variable by the same name that exists in the program immediately outside the construct. Each reference to a shared variable in the structured block becomes a reference to the original variable. For each private variable referenced in the structured block, a new version of the original variable (of the same type and size) is created in memory for each task or SIMD lane that contains code associated with the directive. Creation of the new version does not alter the value of the original variable. However, the impact of attempts to access the original variable from within the region corresponding to the directive is unspecified; see Section 2.21.4.3 for additional details. References to a private variable in the structured block refer to the private version of the original variable for the current task or SIMD lane. The relationship between the value of the original variable and the initial or final value of the private version depends on the exact clause that specifies it. Details of this issue, as well as other issues with privatization, are provided in Section 2.21.

The minimum size at which a memory update may also read and write back adjacent variables that are part of another variable (as array or structure elements) is implementation defined but is no larger than the base language requires.

A single access to a variable may be implemented with multiple load or store instructions and, thus, is not guaranteed to be atomic with respect to other accesses to the same variable. Accesses to
variables smaller than the implementation defined minimum size or to C or C++ bit-fields may be implemented by reading, modifying, and rewriting a larger unit of memory, and may thus interfere with updates of variables or fields in the same unit of memory.

Two memory operations are considered unordered if the order in which they must complete, as seen by their affected threads, is not specified by the memory consistency guarantees listed in Section 1.4.6. If multiple threads write to the same memory unit (defined consistently with the above access considerations) then a data race occurs if the writes are unordered. Similarly, if at least one thread reads from a memory unit and at least one thread writes to that same memory unit then a data race occurs if the read and write are unordered. If a data race occurs then the result of the program is unspecified.

A private variable in a task region that subsequently generates an inner nested parallel region is permitted to be made shared for implicit tasks in the inner parallel region. A private variable in a task region can also be shared by an explicit task region generated during its execution. However, the programmer must use synchronization that ensures that the lifetime of the variable does not end before completion of the explicit task region sharing it. Any other access by one task to the private variables of another task results in unspecified behavior.

A storage location in memory that is associated with a given device has a device address that may be dereferenced by a thread executing on that device, but it may not be generally accessible from other devices. A different device may obtain a device pointer that refers to this device address. The manner in which a program can obtain the referenced device address from a device pointer, outside of mechanisms specified by OpenMP, is implementation defined.

### 1.4.2 Device Data Environments

When an OpenMP program begins, an implicit target data region for each device surrounds the whole program. Each device has a device data environment that is defined by its implicit target data region. Any declare target directives and directives that accept data-mapping attribute clauses determine how an original variable in a data environment is mapped to a corresponding variable in a device data environment.

When an original variable is mapped to a device data environment and a corresponding variable is not present in the device data environment, a new corresponding variable (of the same type and size as the original variable) is created in the device data environment. Conversely, the original variable becomes the new variable’s corresponding variable in the device data environment of the device that performs a mapping operation.

The corresponding variable in the device data environment may share storage with the original variable. Writes to the corresponding variable may alter the value of the original variable. The impact of this possibility on memory consistency is discussed in Section 1.4.6. When a task executes in the context of a device data environment, references to the original variable refer to the corresponding variable in the device data environment. If an original variable is not currently mapped and a corresponding variable does not exist in the device data environment then accesses to
the original variable result in unspecified behavior unless the `unified_shared_memory`
clause is specified on a `requires` directive for the compilation unit.

The relationship between the value of the original variable and the initial or final value of the
 corresponding variable depends on the `map-type`. Details of this issue, as well as other issues with
mapping a variable, are provided in Section 2.21.7.1.

The original variable in a data environment and a corresponding variable in a device data
environment may share storage. Without intervening synchronization data races can occur.

If a variable has a corresponding variable with which it does not share storage, a write to a storage
location designated by the variable causes the value at the corresponding storage location to
become undefined.

### 1.4.3 Memory Management

The host device, and other devices that an implementation may support, have attached storage
resources where program variables are stored. These resources can have different traits. A memory
space in an OpenMP program represents a set of these storage resources. Memory spaces are
defined according to a set of traits, and a single resource may be exposed as multiple memory
spaces with different traits or may be part of multiple memory spaces. In any device, at least one
memory space is guaranteed to exist.

An OpenMP program can use a `memory allocator` to allocate memory in which to store variables.
This memory will be allocated from the storage resources of the `memory space` associated with the
memory allocator. Memory allocators are also used to deallocate previously allocated memory.
When an OpenMP memory allocator is not used to allocate memory, OpenMP does not prescribe
the storage resource for the allocation; the memory for the variables may be allocated in any storage
resource.

### 1.4.4 The Flush Operation

The memory model has relaxed-consistency because a thread’s temporary view of memory is not
required to be consistent with memory at all times. A value written to a variable can remain in the
thread’s temporary view until it is forced to memory at a later time. Likewise, a read from a
variable may retrieve the value from the thread’s temporary view, unless it is forced to read from
memory. OpenMP flush operations are used to enforce consistency between a thread’s temporary
view of memory and memory, or between multiple threads’ view of memory.

A flush operation has an associated `device-set` that constrains the threads with which it enforces
memory consistency. Consistency is only guaranteed to be enforced between the view of memory
of its thread and the view of memory of other threads executing on devices in its device-set. Unless
otherwise stated, the device-set of a flush operation only includes the current device.

If a flush operation is a strong flush, it enforces consistency between a thread’s temporary view and
memory. A strong flush operation is applied to a set of variables called the `flush-set`. A strong flush
restricts reordering of memory operations that an implementation might otherwise do. Implementations must not reorder the code for a memory operation for a given variable, or the code for a flush operation for the variable, with respect to a strong flush operation that refers to the same variable.

If a thread has performed a write to its temporary view of a shared variable since its last strong flush of that variable, then when it executes another strong flush of the variable, the strong flush does not complete until the value of the variable has been written to the variable in memory. If a thread performs multiple writes to the same variable between two strong flushes of that variable, the strong flush ensures that the value of the last write is written to the variable in memory. A strong flush of a variable executed by a thread also causes its temporary view of the variable to be discarded, so that if its next memory operation for that variable is a read, then the thread will read from memory and capture the value in its temporary view. When a thread executes a strong flush, no later memory operation by that thread for a variable involved in that strong flush is allowed to start until the strong flush completes. The completion of a strong flush executed by a thread is defined as the point at which all writes to the flush-set performed by the thread before the strong flush are visible in memory to all other threads, and at which that thread’s temporary view of the flush-set is discarded.

A strong flush operation provides a guarantee of consistency between a thread’s temporary view and memory. Therefore, a strong flush can be used to guarantee that a value written to a variable by one thread may be read by a second thread. To accomplish this, the programmer must ensure that the second thread has not written to the variable since its last strong flush of the variable, and that the following sequence of events are completed in this specific order:

1. The value is written to the variable by the first thread;
2. The variable is flushed, with a strong flush, by the first thread;
3. The variable is flushed, with a strong flush, by the second thread; and
4. The value is read from the variable by the second thread.

If a flush operation is a release flush or acquire flush, it can enforce consistency between the views of memory of two synchronizing threads. A release flush guarantees that any prior operation that writes or reads a shared variable will appear to be completed before any operation that writes or reads the same shared variable and follows an acquire flush with which the release flush synchronizes (see Section 1.4.5 for more details on flush synchronization). A release flush will propagate the values of all shared variables in its temporary view to memory prior to the thread performing any subsequent atomic operation that may establish a synchronization. An acquire flush will discard any value of a shared variable in its temporary view to which the thread has not written since last performing a release flush, and it will load any value of a shared variable propagated by a release flush that synchronizes with it into its temporary view so that it may be subsequently read. Therefore, release and acquire flushes may also be used to guarantee that a value written to a variable by one thread may be read by a second thread. To accomplish this, the programmer must ensure that the second thread has not written to the variable since its last acquire flush, and that the following sequence of events happen in this specific order:
1. The value is written to the variable by the first thread;
2. The first thread performs a release flush;
3. The second thread performs an acquire flush; and
4. The value is read from the variable by the second thread.

\[\text{Note} – \text{OpenMP synchronization operations, described in Section 2.19 and in Section 3.9, are recommended for enforcing this order. Synchronization through variables is possible but is not recommended because the proper timing of flushes is difficult.}\]

The flush properties that define whether a flush operation is a strong flush, a release flush, or an acquire flush are not mutually disjoint. A flush operation may be a strong flush and a release flush; it may be a strong flush and an acquire flush; it may be a release flush and an acquire flush; or it may be all three.

### 1.4.5 Flush Synchronization and Happens Before

OpenMP supports thread synchronization with the use of release flushes and acquire flushes. For any such synchronization, a release flush is the source of the synchronization and an acquire flush is the sink of the synchronization, such that the release flush synchronizes with the acquire flush.

A release flush has one or more associated release sequences that define the set of modifications that may be used to establish a synchronization. A release sequence starts with an atomic operation that follows the release flush and modifies a shared variable and additionally includes any read-modify-write atomic operations that read a value taken from some modification in the release sequence. The following rules determine the atomic operation that starts an associated release sequence.

- If a release flush is performed on entry to an atomic operation, that atomic operation starts its release sequence.
- If a release flush is performed in an implicit \texttt{flush} region, an atomic operation that is provided by the implementation and that modifies an internal synchronization variable starts its release sequence.
- If a release flush is performed by an explicit \texttt{flush} region, any atomic operation that modifies a shared variable and follows the \texttt{flush} region in its thread’s program order starts an associated release sequence.

An acquire flush is associated with one or more prior atomic operations that read a shared variable and that may be used to establish a synchronization. The following rules determine the associated atomic operation that may establish a synchronization.
• If an acquire flush is performed on exit from an atomic operation, that atomic operation is its
  associated atomic operation.

• If an acquire flush is performed in an implicit flush region, an atomic operation that is
  provided by the implementation and that reads an internal synchronization variable is its
  associated atomic operation.

• If an acquire flush is performed by an explicit flush region, any atomic operation that reads a
  shared variable and precedes the flush region in its thread’s program order is an associated
  atomic operation.

A release flush synchronizes with an acquire flush if the following conditions are satisfied:

• An atomic operation associated with the acquire flush reads a value written by a modification
  from a release sequence associated with the release flush; and

• The device on which each flush is performed is in both of their respective device-sets.

An operation \(X\) simply happens before an operation \(Y\) if any of the following conditions are
satisfied:

1. \(X\) and \(Y\) are performed by the same thread, and \(X\) precedes \(Y\) in the thread’s program order;
2. \(X\) synchronizes with \(Y\) according to the flush synchronization conditions explained above or
   according to the base language’s definition of synchronizes with, if such a definition exists; or
3. Another operation, \(Z\), exists such that \(X\) simply happens before \(Z\) and \(Z\) simply happens before \(Y\).

An operation \(X\) happens before an operation \(Y\) if any of the following conditions are satisfied:

1. \(X\) happens before \(Y\) according to the base language’s definition of happens before, if such a
   definition exists; or
2. \(X\) simply happens before \(Y\).

A variable with an initial value is treated as if the value is stored to the variable by an operation that
happens before all operations that access or modify the variable in the program.

1.4.6 OpenMP Memory Consistency

The following rules guarantee an observable completion order for a given pair of memory
operations in race-free programs, as seen by all affected threads. If both memory operations are
strong flushes, the affected threads are all threads on devices in both of their respective device-sets.
If exactly one of the memory operations is a strong flush, the affected threads are all threads on
devices in its device-set. Otherwise, the affected threads are all threads.

• If two operations performed by different threads are sequentially consistent atomic operations or
  they are strong flushes that flush the same variable, then they must be completed as if in some
  sequential order, seen by all affected threads.
If two operations performed by the same thread are sequentially consistent atomic operations or they access, modify, or, with a strong flush, flush the same variable, then they must be completed as if in that thread’s program order, as seen by all affected threads.

If two operations are performed by different threads and one happens before the other, then they must be completed as if in that happens before order, as seen by all affected threads, if:

- both operations access or modify the same variable;
- both operations are strong flushes that flush the same variable; or
- both operations are sequentially consistent atomic operations.

Any two atomic memory operations from different atomic regions must be completed as if in the same order as the strong flushes implied in their respective regions, as seen by all affected threads.

The flush operation can be specified using the flush directive, and is also implied at various locations in an OpenMP program: see Section 2.19.8 for details.

Note – Since flush operations by themselves cannot prevent data races, explicit flush operations are only useful in combination with non-sequentially consistent atomic directives.

OpenMP programs that:

- Do not use non-sequentially consistent atomic directives;

- Do not rely on the accuracy of a false result from omp_test_lock and omp_test_nest_lock; and

- Correctly avoid data races as required in Section 1.4.1,

behave as though operations on shared variables were simply interleaved in an order consistent with the order in which they are performed by each thread. The relaxed consistency model is invisible for such programs, and any explicit flush operations in such programs are redundant.

1.5 Tool Interfaces

The OpenMP API includes two tool interfaces, OMPT and OMPD, to enable development of high-quality, portable, tools that support monitoring, performance, or correctness analysis and debugging of OpenMP programs developed using any implementation of the OpenMP API.

An implementation of the OpenMP API may differ from the abstract execution model described by its specification. The ability of tools that use the OMPT or OMPD interfaces to observe such differences does not constrain implementations of the OpenMP API in any way.
1.5.1 OMPT

The OMPT interface, which is intended for first-party tools, provides the following:

• A mechanism to initialize a first-party tool;
• Routines that enable a tool to determine the capabilities of an OpenMP implementation;
• Routines that enable a tool to examine OpenMP state information associated with a thread;
• Mechanisms that enable a tool to map implementation-level calling contexts back to their source-level representations;
• A callback interface that enables a tool to receive notification of OpenMP events;
• A tracing interface that enables a tool to trace activity on OpenMP target devices; and
• A runtime library routine that an application can use to control a tool.

OpenMP implementations may differ with respect to the thread states that they support, the mutual exclusion implementations that they employ, and the OpenMP events for which tool callbacks are invoked. For some OpenMP events, OpenMP implementations must guarantee that a registered callback will be invoked for each occurrence of the event. For other OpenMP events, OpenMP implementations are permitted to invoke a registered callback for some or no occurrences of the event; for such OpenMP events, however, OpenMP implementations are encouraged to invoke tool callbacks on as many occurrences of the event as is practical. Section 4.2.4 specifies the subset of OMPT callbacks that an OpenMP implementation must support for a minimal implementation of the OMPT interface.

With the exception of the omp_control_tool runtime library routine for tool control, all other routines in the OMPT interface are intended for use only by tools and are not visible to applications. For that reason, a Fortran binding is provided only for omp_control_tool; all other OMPT functionality is described with C syntax only.

1.5.2 OMPD

The OMPD interface is intended for third-party tools, which run as separate processes. An OpenMP implementation must provide an OMPD library that can be dynamically loaded and used by a third-party tool. A third-party tool, such as a debugger, uses the OMPD library to access OpenMP state of a program that has begun execution. OMPD defines the following:

• An interface that an OMPD library exports, which a tool can use to access OpenMP state of a program that has begun execution;
• A callback interface that a tool provides to the OMPD library so that the library can use it to access the OpenMP state of a program that has begun execution; and
• A small number of symbols that must be defined by an OpenMP implementation to help the tool
find the correct OMPD library to use for that OpenMP implementation and to facilitate
notification of events.

Section 5 describes OMPD in detail.

1.6 OpenMP Compliance

The OpenMP API defines constructs that operate in the context of the base language that is
supported by an implementation. If the implementation of the base language does not support a
language construct that appears in this document, a compliant OpenMP implementation is not
required to support it, with the exception that for Fortran, the implementation must allow case
insensitivity for directive and API routines names, and must allow identifiers of more than six
characters. An implementation of the OpenMP API is compliant if and only if it compiles and
executes all other conforming programs, and supports the tool interface, according to the syntax and
semantics laid out in Chapters 1, 2, 3, 4 and 5. Appendices A and B as well as sections designated
as Notes (see Section 1.8) are for information purposes only and are not part of the specification.

All library, intrinsic and built-in routines provided by the base language must be thread-safe in a
compliant implementation. In addition, the implementation of the base language must also be
thread-safe. For example, ALLOCATE and DEALLOCATE statements must be thread-safe in
Fortran. Unsynchronized concurrent use of such routines by different threads must produce correct
results (although not necessarily the same as serial execution results, as in the case of random
number generation routines).

Starting with Fortran 90, variables with explicit initialization have the SAVE attribute implicitly.
This is not the case in Fortran 77. However, a compliant OpenMP Fortran implementation must
give such a variable the SAVE attribute, regardless of the underlying base language version.

Appendix A lists certain aspects of the OpenMP API that are implementation defined. A compliant
implementation must define and document its behavior for each of the items in Appendix A.

1.7 Normative References

• ISO/IEC 9899:1990, Information Technology - Programming Languages - C.
  This OpenMP API specification refers to ISO/IEC 9899:1990 as C90.
• ISO/IEC 9899:1999, Information Technology - Programming Languages - C.
  This OpenMP API specification refers to ISO/IEC 9899:1999 as C99.
• ISO/IEC 9899:2011, Information Technology - Programming Languages - C.
  This OpenMP API specification refers to ISO/IEC 9899:2011 as C11.
  This OpenMP API specification refers to ISO/IEC 9899:2018 as C18.

  This OpenMP API specification refers to ISO/IEC 14882:1998 as C++98.

  This OpenMP API specification refers to ISO/IEC 14882:2011 as C++11.

  This OpenMP API specification refers to ISO/IEC 14882:2014 as C++14.

  This OpenMP API specification refers to ISO/IEC 14882:2017 as C++17.

  This OpenMP API specification refers to ISO/IEC 14882:2020 as C++20.

  This OpenMP API specification refers to ISO/IEC 1539:1980 as Fortran 77.

  This OpenMP API specification refers to ISO/IEC 1539:1991 as Fortran 90.

  This OpenMP API specification refers to ISO/IEC 1539-1:1997 as Fortran 95.

  This OpenMP API specification refers to ISO/IEC 1539-1:2004 as Fortran 2003.

  This OpenMP API specification refers to ISO/IEC 1539-1:2010 as Fortran 2008.

  This OpenMP API specification refers to ISO/IEC 1539-1:2018 as Fortran 2018. While future versions of the OpenMP specification are expected to address the following features, currently their use may result in unspecified behavior.
  - Declared type of a polymorphic allocatable component in structure constructor
  - `SELECT RANK` construct
  - IEEE comparison predicate in intrinsic relational operators
  - Finalization of an allocatable subobject in intrinsic assignment
– Locality of variables in a **DO CONCURRENT** construct
– **IMPORT** statement extensions
– Assumed-rank dummy argument
– Assumed-type dummy argument
– Interoperable procedure enhancements
– **ASYNCHRONOUS** attribute enhancement

Where this OpenMP API specification refers to C, C++ or Fortran, reference is made to the base language supported by the implementation.

1.8 **Organization of this Document**

The remainder of this document is structured as follows:

- Chapter 2 “Directives”
- Chapter 3 “Runtime Library Routines”
- Chapter 4 “OMPT Interface”
- Chapter 5 “OMPD Interface”
- Chapter 6 “Environment Variables”
- Appendix A “OpenMP Implementation-Defined Behaviors”
- Appendix B “Features History”

Some sections of this document only apply to programs written in a certain base language. Text that applies only to programs for which the base language is C or C++ is shown as follows:

\[ \text{C/C++ specific text...} \]

Text that applies only to programs for which the base language is C only is shown as follows:

\[ \text{C specific text...} \]

Text that applies only to programs for which the base language is C++ only is shown as follows:

\[ \text{C++ specific text...} \]
Text that applies only to programs for which the base language is Fortran is shown as follows:

--- Fortran ---

Fortran specific text...

--- Fortran ---

Where an entire page consists of base language specific text, a marker is shown at the top of the page. For Fortran-specific text, the marker is:

--- Fortran (cont.) ---

For C/C++-specific text, the marker is:

--- C/C++ (cont.) ---

Some text is for information only, and is not part of the normative specification. Such text is designated as a note or comment, like this:

Note – Non-normative text...

COMMENT: Non-normative text...
2 Directives

This chapter describes the syntax and behavior of OpenMP directives.

OpenMP directives are specified with the `#pragma` mechanism provided by the C standard.

OpenMP directives are specified with attribute specifiers or the `#pragma` mechanism provided by the C++ standard.

OpenMP directives are specified with stylized comments that are identified by unique sentinels.

Also, a stylized comment form is available for conditional compilation.

If a directive appears in the declarative part of a module then the behavior is as if that directive appears after any references to that module.

Compilers can therefore ignore OpenMP directives and conditionally compiled code if support of the OpenMP API is not provided or enabled. A compliant implementation must provide an option or interface that ensures that underlying support of all OpenMP directives and OpenMP conditional compilation mechanisms is enabled. In the remainder of this document, the phrase OpenMP compilation is used to mean a compilation with these OpenMP features enabled.

This chapter uses `NULL` as a generic term for a null pointer constant, `true` as a generic term for a non-zero integer value and `false` as a generic term for an integer value of zero.

This chapter uses `NULL` as a generic term for the named constant `C_NULL_PTR`, `true` as a generic term for a logical value of `.TRUE.` and `false` as a generic term for a logical value of `.FALSE.`.
Restrictions

The following restrictions apply to OpenMP directives:

- A declarative directive may not be used in place of a substatement in a selection statement, in place of the loop body in an iteration statement, or in place of the statement that follows a label.

- A declarative directive may not be used in place of a substatement in a selection statement or iteration statement, or in place of the statement that follows a label.

- OpenMP directives, except *simd* and declarative directives, may not appear in pure procedures.

- OpenMP directives may not appear in the *WHERE, FORALL* or *DO CONCURRENT* constructs.

2.1 Directive Format

OpenMP directives for C/C++ may be specified with *#pragma* directives as follows:

```
#pragma omp directive-name [[,] clause[ [,] clause] ... ] new-line
```

Where *directive-name* is the name of the directive and, when specified in the syntax of the directive, any directive-level arguments enclosed in parentheses.

Note – In the following example, *deobj(o)* is the *directive-name*:

```
#pragma omp deobj(o) depend(inout: d)
```

Each *#pragma* directive starts with *#pragma omp*. The remainder of the directive follows the conventions of the C and C++ standards for compiler directives. In particular, white space can be used before and after the #, and sometimes white space must be used to separate the words in a directive. Preprocessing tokens following *#pragma omp* are subject to macro replacement.

Some OpenMP directives may be composed of consecutive *#pragma* directives if specified in their syntax.
In C++11 and higher, all OpenMP directives may be specified with C++ attribute specifiers as follows:

```cpp
[[ omp :: directive( directive-name[, clause[,...] ] ) ]]
```

or

```cpp
[[ using omp : directive( directive-name[, clause[,...] ] ) ]]
```

The above two forms are interchangeable for any OpenMP directive. Some OpenMP directives may be composed of consecutive attribute specifiers if specified in their syntax. Any two consecutive attribute specifiers may be reordered or expressed as a single attribute specifier, as permitted by the base language, without changing the behavior of the OpenMP directive.

Some directives may have additional forms that use the attribute syntax.

Multiple attributes on the same statement are allowed. A directive that uses the attribute syntax cannot be applied to the same statement as a directive that uses the pragma syntax. For any directive that has a paired end directive, including those with a begin and end pair, both directives must use either the attribute syntax or the pragma syntax. Attribute directives that apply to the same statement are unordered. An ordering can be imposed with the `sequence` attribute, which is specified as follows:

```cpp
[[ omp :: sequence( [omp::]directive-attr [, [omp::]directive-attr]... ) ]]
```

where `directive-attr` is any attribute in the `omp` namespace, optionally specified with a `omp::` namespace qualifier, which may be another `sequence` attribute.

The application of multiple attributes in a `sequence` attribute is ordered as if each directive had been written as a `#pragma` directive on subsequent lines.

Note – This is an example of the expected transformation:

```cpp
[[ omp::sequence(directive(parallel), directive(for)) ]]
for(...) {}
// becomes
#pragma omp parallel
#pragma omp for
for(...) {}
```
Directives are case-sensitive.

Each of the expressions used in the OpenMP syntax inside of the clauses must be a valid `assignment-expression` of the base language unless otherwise specified.

Directives may not appear in `constexpr` functions or in constant expressions.

OpenMP directives for Fortran are specified as follows:

```
sentinel directive-name [clause[ [, clause]...]]
```

All OpenMP compiler directives must begin with a directive `sentinel`. The format of a sentinel differs between fixed form and free form source files, as described in Section 2.1.1 and Section 2.1.2.

Directives are case insensitive. Directives cannot be embedded within continued statements, and statements cannot be embedded within directives.

Each of the expressions used in the OpenMP syntax inside of the clauses must be a valid `expression` of the base language unless otherwise specified.

In order to simplify the presentation, free form is used for the syntax of OpenMP directives for Fortran in the remainder of this document, except as noted.

A directive may be categorized as one of the following: a metadirective, a declarative directive, an executable directive, an informational directive, or a utility directive.

Only one `directive-name` can be specified per directive (note that this includes combined directives, see Section 2.16). The order in which clauses appear on directives is not significant. Clauses on directives may be repeated as needed, subject to the restrictions listed in the description of each clause or the directives on which they can appear.

Some clauses accept a `list`, an `extended-list`, or a `locator-list`. A `list` consists of a comma-separated collection of one or more `list items`. An `extended-list` consists of a comma-separated collection of one or more `extended list items`. A `locator-list` consists of a comma-separated collection of one or more `locator list items`.

A `list item` is a variable or an array section. An `extended list item` is a `list item` or a function name. A `locator list item` is any lvalue expression including variables, an array section, or a reserved locator.
A *list item* is a variable that is not coindexed, an array section that is not coindexed, a *named constant*, an associate name that may appear in a variable definition context, or a common block name (enclosed in slashes). An *extended list item* is a *list item* or a procedure name. A *locator list item* is a *list item*, or a reserved locator.

A *named constant* as a *list item* can appear only in clauses where it is explicitly allowed.

When a named common block appears in a *list*, it has the same meaning and restrictions as if every explicit member of the common block appeared in the list. An explicit member of a common block is a variable that is named in a `COMMON` statement that specifies the common block name and is declared in the same scoping unit in which the clause appears. Named common blocks do not include the blank common block.

Although variables in common blocks can be accessed by use association or host association, common block names cannot. As a result, a common block name specified in a data-sharing attribute, a data copying, or a data-mapping attribute clause must be declared to be a common block in the same scoping unit in which the clause appears.

If a list item that appears in a directive or clause is an optional dummy argument that is not present, the directive or clause for that list item is ignored.

If the variable referenced inside a construct is an optional dummy argument that is not present, any explicitly determined, implicitly determined, or predetermined data-sharing and data-mapping attribute rules for that variable are ignored. Otherwise, if the variable is an optional dummy argument that is present, it is present inside the construct.

For all base languages, a *list item*, an *extended list item*, or a *locator list item* is subject to the restrictions specified in Section 2.1.5 and in each of the sections that describe clauses and directives for which the *list*, the *extended-list*, or the *locator-list* appears.

Some clauses and directives accept the use of reserved locators as special identifiers that represent system storage not necessarily bound to any base language storage item. Reserved locators may only appear in clauses and directives where they are explicitly allowed and may not otherwise be referenced in the program. The list of reserved locators is:

```
omp_all_memory
```

The reserved locator `omp_all_memory` is a reserved identifier that denotes a list item treated as having storage that corresponds to the storage of all other objects in memory.

Some directives have an associated structured block or a structured block sequence.
A structured block sequence that consists of more than one statement may appear only for executable directives that explicitly allow it. The corresponding compound statement obtained by enclosing the sequence in { and } must be a structured block and the structured block sequence then should be considered to be a structured block with all of its restrictions.

A structured block:

- may contain infinite loops where the point of exit is never reached;
- may halt due to an IEEE exception;
- may contain calls to `exit()`, `_Exit()`, `quick_exit()`, `abort()` or functions with a `_Noreturn` specifier (in C) or a `noreturn` attribute (in C/C++);
- may be an expression statement, iteration statement, selection statement, or try block, provided that the corresponding compound statement obtained by enclosing it in { and } would be a structured block; and
- may contain `STOP` or `ERROR STOP` statements.

**Restrictions**

Restrictions to structured blocks are as follows:

- Entry to a structured block must not be the result of a branch.
- The point of exit cannot be a branch out of the structured block.
- The point of entry to a structured block must not be a call to `setjmp`.
- `longjmp` must not violate the entry/exit criteria.
- `throw` must not violate the entry/exit criteria.
- `co-await`, `co_yield` and `co_return` must not violate the entry/exit criteria.
• When a BLOCK construct appears in a structured block, that BLOCK construct must not contain any ASYNCHRONOUS or VOLATILE statements, nor any specification statements that include the ASYNCHRONOUS or VOLATILE attributes.

Restrictions on explicit OpenMP regions (that arise from executable directives) are as follows:

• If more than one image is executing the program, any image control statement, ERROR STOP statement, FAIL IMAGE statement, collective subroutine call or access to a coindexed object that appears in an explicit OpenMP region will result in unspecified behavior.

2.1.1 Fixed Source Form Directives

The following sentinels are recognized in fixed form source files:

![omp | comp | $omp]

Sentinels must start in column 1 and appear as a single word with no intervening characters. Fortran fixed form line length, white space, continuation, and column rules apply to the directive line. Initial directive lines must have a space or a zero in column 6, and continuation directive lines must have a character other than a space or a zero in column 6.

Comments may appear on the same line as a directive. The exclamation point initiates a comment when it appears after column 6. The comment extends to the end of the source line and is ignored. If the first non-blank character after the directive sentinel of an initial or continuation directive line is an exclamation point, the line is ignored.

Note – In the following example, the three formats for specifying the directive are equivalent (the first line represents the position of the first 9 columns):

c23456789
 !$omp parallel do shared(a,b,c)

 c$omp parallel do
 c$omp+shared(a,b,c)
 c$omp paraleldoshared(a,b,c)
2.1.2 Free Source Form Directives

The following sentinel is recognized in free form source files:

\[
!$omp
\]

The sentinel can appear in any column as long as it is preceded only by white space. It must appear as a single word with no intervening white space. Fortran free form line length, white space, and continuation rules apply to the directive line. Initial directive lines must have a space after the sentinel. Continued directive lines must have an ampersand (\&) as the last non-blank character on the line, prior to any comment placed inside the directive. Continuation directive lines can have an ampersand after the directive sentinel with optional white space before and after the ampersand.

Comments may appear on the same line as a directive. The exclamation point (!) initiates a comment. The comment extends to the end of the source line and is ignored. If the first non-blank character after the directive sentinel is an exclamation point, the line is ignored.

One or more blanks or horizontal tabs are optional to separate adjacent keywords in \textit{directive-names} unless otherwise specified.

\begin{verbatim}
!23456789
  !$omp parallel do &
    !$omp shared(a,b,c)

  !$omp parallel &
    !$omp&do shared(a,b,c)

  !$omp paralleldo shared(a,b,c)
\end{verbatim}

\textbf{Note} – In the following example the three formats for specifying the directive are equivalent (the first line represents the position of the first 9 columns):
2.1.3 Stand-Alone Directives

Summary
Stand-alone directives are executable directives that have no associated user code.

Description
Stand-alone directives do not have any associated executable user code. Instead, they represent executable statements that typically do not have succinct equivalent statements in the base language. Some restrictions limit the placement of a stand-alone directive within a program. A stand-alone directive may be placed only at a point where a base language executable statement is allowed.

Restrictions
Restrictions to stand-alone directives are as follows:

- A stand-alone directive may not be used in place of a substatement in a selection statement, in place of the loop body in an iteration statement, or in place of the statement that follows a label.

2.1.4 Array Shaping

If an expression has a type of pointer to $T$, then a shape-operator can be used to specify the extent of that pointer. In other words, the shape-operator is used to reinterpret, as an n-dimensional array, the region of memory to which that expression points.

Formally, the syntax of the shape-operator is as follows:

```plaintext
shaped-expression := ([s_1] [s_2] \ldots [s_n]) cast-expression
```

The result of applying the shape-operator to an expression is an lvalue expression with an n-dimensional array type with dimensions $s_1 \times s_2 \ldots \times s_n$ and element type $T$.

The precedence of the shape-operator is the same as a type cast.

Each $s_i$ is an integral type expression that must evaluate to a positive integer.
Restrictions
Restrictions to the shape-operator are as follows:

- The type $T$ must be a complete type.
- The shape-operator can appear only in clauses for which it is explicitly allowed.
- The result of a shape-operator must be a named array of a list item.
- The type of the expression upon which a shape-operator is applied must be a pointer type.

- If the type $T$ is a reference to a type $T'$, then the type will be considered to be $T'$ for all purposes of the designated array.

2.1.5 Array Sections
An array section designates a subset of the elements in an array.

To specify an array section in an OpenMP construct, array subscript expressions are extended with the following syntax:

```c
[ lower-bound : length : stride ] or
[ lower-bound : length ] or
[ lower-bound : length ] or
[ lower-bound : stride ] or
[ lower-bound : ] or
[ lower-bound : ] or
[ : length : stride ] or
[ : length ] or
[ : length ] or
[ : stride ]
[ : ]
[ : ]
```
The array section must be a subset of the original array.

Array sections are allowed on multidimensional arrays. Base language array subscript expressions can be used to specify length-one dimensions of multidimensional array sections.

Each of the \textit{lower-bound}, \textit{length}, and \textit{stride} expressions if specified must be an integral type \textit{expression} of the base language. When evaluated they represent a set of integer values as follows:

\[
\{ \text{lower-bound}, \text{lower-bound} + \text{stride}, \text{lower-bound} + 2 \times \text{stride}, \ldots , \text{lower-bound} + ((\text{length} - 1) \times \text{stride}) \}
\]

The \textit{length} must evaluate to a non-negative integer.

The \textit{stride} must evaluate to a positive integer.

When the size of the array dimension is not known, the \textit{length} must be specified explicitly.

When the \textit{stride} is absent it defaults to 1.

When the \textit{length} is absent it defaults to \(
\lceil \lceil (\text{size} - \text{lower-bound}) / \text{stride} \rceil \rceil
\), where \textit{size} is the size of the array dimension.

When the \textit{lower-bound} is absent it defaults to 0.

The precedence of a subscript operator that uses the array section syntax is the same as the precedence of a subscript operator that does not use the array section syntax.

\begin{itemize}
\item \texttt{a[0:6]}
\item \texttt{a[0:6:1]}
\item \texttt{a[1:10]}
\item \texttt{a[1:]}
\item \texttt{a[:10:2]}
\item \texttt{b[10][::]}
\item \texttt{b[10][::0]}
\item \texttt{c[42][0:6][::]}
\item \texttt{c[42][0:6:2][::]}
\item \texttt{c[1:10][42][0:6]}
\item \texttt{S.c[:100]}
\item \texttt{p->y[:10]}
\item \texttt{this->a[:N]}
\item \texttt{(p+10)[::N]}
\end{itemize}
Assume \( \mathbf{a} \) is declared to be a 1-dimensional array with dimension size 11. The first two examples are equivalent, and the third and fourth examples are equivalent. The fifth example specifies a stride of 2 and therefore is not contiguous.

Assume \( \mathbf{b} \) is declared to be a pointer to a 2-dimensional array with dimension sizes 10 and 10. The sixth example refers to all elements of the 2-dimensional array given by \( \mathbf{b}[10] \). The seventh example is a zero-length array section.

Assume \( \mathbf{c} \) is declared to be a 3-dimensional array with dimension sizes 50, 50, and 50. The eighth example is contiguous, while the ninth and tenth examples are not contiguous.

The final four examples show array sections that are formed from more general base expressions.

The following are examples that are non-conforming array sections:

\[
\begin{align*}
\mathbf{s}[10].x \\
\mathbf{p}[10]\rightarrow y \\
*(\mathbf{xp}[10])
\end{align*}
\]

For all three examples, a base language operator is applied in an undefined manner to an array section. The only operator that may be applied to an array section is a subscript operator for which the array section appears as the postfix expression.

Restrictions

Restrictions to array sections are as follows:

- An array section can appear only in clauses for which it is explicitly allowed.
- A *stride* expression may not be specified unless otherwise stated.
- An element of an array section with a non-zero size must have a complete type.
- The base expression of an array section must have an array or pointer type.
- If a consecutive sequence of array subscript expressions appears in an array section, and the first subscript expression in the sequence uses the extended array section syntax defined in this section, then only the last subscript expression in the sequence may select array elements that have a pointer type.
• If the type of the base expression of an array section is a reference to a type $T$, then the type will be considered to be $T$ for all purposes of the array section.

• An array section cannot be used in an overloaded $[\ ]$ operator.

• If a stride expression is specified, it must be positive.

• The upper bound for the last dimension of an assumed-size dummy array must be specified.

• If a list item is an array section with vector subscripts, the first array element must be the lowest in the array element order of the array section.

• If a list item is an array section, the last part-ref of the list item must have a section subscript list.

2.1.6 Iterators

Iterators are identifiers that expand to multiple values in the clause on which they appear.

The syntax of the iterator modifier is as follows:

```
iterator (iterators-definition)
```

where `iterators-definition` is one of the following:

```
iterator-specifier [, iterators-definition ]
```

where `iterator-specifier` is one of the following:

```
[ iterator-type ] identifier = range-specification
```

where:

• `identifier` is a base language identifier.

• `iterator-type` is a type name.

• `iterator-type` is a type specifier.

• `range-specification` is of the form `begin : end [ : step ]`, where `begin` and `end` are expressions for which their types can be converted to `iterator-type` and `step` is an integral expression.
In an *iterator-specifier*, if the *iterator-type* is not specified then that iterator is of **int** type.

In an *iterator-specifier*, if the *iterator-type* is not specified then that iterator has default integer type.

In a *range-specification*, if the *step* is not specified its value is implicitly defined to be 1.

An iterator only exists in the context of the clause in which it appears. An iterator also hides all accessible symbols with the same name in the context of the clause.

The use of a variable in an expression that appears in the *range-specification* causes an implicit reference to the variable in all enclosing constructs.

The values of the iterator are the set of values $i_0, \ldots, i_{N-1}$ where:

- $i_0 = (\text{iterator-type}) \begin{equation*} \text{begin}, \end{equation*}$
- $i_j = (\text{iterator-type}) (i_{j-1} + \text{step})$, where $j \geq 1$ and
  - if $\text{step} > 0$,
    - $i_0 < (\text{iterator-type}) \begin{equation*} \text{end}, \end{equation*}$
    - $i_{N-1} < (\text{iterator-type}) \begin{equation*} \text{end}, \end{equation*}$ and
    - $(\text{iterator-type}) (i_{N-1} + \text{step}) \geq (\text{iterator-type}) \begin{equation*} \text{end}$;
  - if $\text{step} < 0$,
    - $i_0 > (\text{iterator-type}) \begin{equation*} \text{end}, \end{equation*}$
    - $i_{N-1} > (\text{iterator-type}) \begin{equation*} \text{end}, \end{equation*}$ and
    - $(\text{iterator-type}) (i_{N-1} + \text{step}) \leq (\text{iterator-type}) \begin{equation*} \text{end}$.

The values of the iterator are the set of values $i_1, \ldots, i_N$ where:

- $i_1 = \text{begin}$,
- $i_j = i_{j-1} + \text{step}$, where $j \geq 2$ and
  - if $\text{step} > 0$,
    - $i_1 \leq \text{end}$,
    - $i_N \leq \text{end}$, and
\begin{verbatim}
\text{– } i_N + step > end;
\end{verbatim}

- if \( step < 0 \),
  \begin{itemize}
  \item \( - i_1 \geq end \),
  \item \( - i_N \geq end \), and
  \item \( - i_N + step \leq end \).
\end{itemize}

The set of values will be empty if no possible value complies with the conditions above.

For those clauses that contain expressions that contain iterator identifiers, the effect is as if the list item is instantiated within the clause for each value of the iterator in the set defined above, substituting each occurrence of the iterator identifier in the expression with the iterator value. If the set of values of the iterator is empty then the effect is as if the clause was not specified.

The behavior is unspecified if \( i_j + step \) cannot be represented in iterator-type in any of the \( i_j + step \) computations for any \( 0 \leq j < N \) in C/C++ or \( 0 < j \leq N \) in Fortran.

\textbf{Restrictions}

Restrictions to iterators are as follows:

- An expression that contains an iterator identifier can only appear in clauses that explicitly allow expressions that contain iterators.
- The iterator-type must not declare a new type.
  \begin{itemize}
  \item The iterator-type must be an integral or pointer type.
  \item The iterator-type must not be \texttt{const} qualified.
  \item The iterator-type must be an integer type.
  \item If the step expression of a range-specification equals zero, the behavior is unspecified.
  \item Each iterator identifier can only be defined once in an iterators-definition.
  \item Iterators cannot appear in the range-specification.
\end{itemize}
2.2 Conditional Compilation

In implementations that support a preprocessor, the _OPENMP macro name is defined to have the decimal value yyyymm where yyyy and mm are the year and month designations of the version of the OpenMP API that the implementation supports.

If a #define or a #undef preprocessing directive in user code defines or undefines the _OPENMP macro name, the behavior is unspecified.

The OpenMP API requires Fortran lines to be compiled conditionally, as described in the following sections.

2.2.1 Fixed Source Form Conditional Compilation Sentinels

The following conditional compilation sentinels are recognized in fixed form source files:

| !$ | *$ | c$ |

To enable conditional compilation, a line with a conditional compilation sentinel must satisfy the following criteria:

- The sentinel must start in column 1 and appear as a single word with no intervening white space;
- After the sentinel is replaced with two spaces, initial lines must have a space or zero in column 6 and only white space and numbers in columns 1 through 5;
- After the sentinel is replaced with two spaces, continuation lines must have a character other than a space or zero in column 6 and only white space in columns 1 through 5.

If these criteria are met, the sentinel is replaced by two spaces. If these criteria are not met, the line is left unchanged.

Note – In the following example, the two forms for specifying conditional compilation in fixed source form are equivalent (the first line represents the position of the first 9 columns):

```
c23456789
!$ 10 iam = omp_get_thread_num() +
!$ & index

#ifdef _OPENMP
10 iam = omp_get_thread_num() +
 & index
#endif
```
### 2.2.2 Free Source Form Conditional Compilation Sentinel

The following conditional compilation sentinel is recognized in free form source files:

```plaintext
!$
```

To enable conditional compilation, a line with a conditional compilation sentinel must satisfy the following criteria:

- The sentinel can appear in any column but must be preceded only by white space;
- The sentinel must appear as a single word with no intervening white space;
- Initial lines must have a space after the sentinel;
- Continued lines must have an ampersand as the last non-blank character on the line, prior to any comment appearing on the conditionally compiled line.

Continuation lines can have an ampersand after the sentinel, with optional white space before and after the ampersand. If these criteria are met, the sentinel is replaced by two spaces. If these criteria are not met, the line is left unchanged.

Note – In the following example, the two forms for specifying conditional compilation in free source form are equivalent (the first line represents the position of the first 9 columns):

```plaintext
c23456789
!$ iam = omp_get_thread_num() + &
!$& index

#define _OPENMP
  iam = omp_get_thread_num() + &
  index
#define
```

### 2.3 Variant Directives

#### 2.3.1 OpenMP Context

At any point in a program, an OpenMP context exists that defines traits that describe the active OpenMP constructs, the execution devices, functionality supported by the implementation and available dynamic values. The traits are grouped into trait sets. The following trait sets exist: `construct`, `device`, `target_device`, `implementation` and `dynamic`. Traits are categorized as name-list
traits, clause-list traits, non-property traits and extension traits. This categorization determines the syntax that is used to match the trait, as defined in Section 2.3.2.

The *construct* set is composed of the directive names, each being a trait, of all enclosing constructs at that point in the program up to a *target* construct. Combined and composite constructs are added to the set as distinct constructs in the same nesting order specified by the original construct. Whether the *dispatch* construct is added to the *construct* set is implementation defined. If it is added, it will only be added for the *target-call* of the associated code. The set is ordered by nesting level in ascending order. Specifically, the ordering of the set of constructs is \(c_1, \ldots, c_N\), where \(c_1\) is the construct at the outermost nesting level and \(c_N\) is the construct at the innermost nesting level. In addition, if the point in the program is not enclosed by a *target* construct, the following rules are applied in order:

1. For procedures with a *declare simd* directive, the *simd* trait is added to the beginning of the set as \(c_1\) for any generated SIMD versions so the total size of the set is increased by 1.

2. For procedures that are determined to be function variants by a *declare variant* directive, the selectors \(c_1, \ldots, c_M\) of the *construct* selector set are added in the same order to the beginning of the set as \(c_1, \ldots, c_M\) so the total size of the set is increased by \(M\).

3. For *device routines*, the *target* trait is added to the beginning of the set as \(c_1\) for any versions of the procedure that are generated for *target* regions so the total size of the set is increased by 1.

The *simd* trait is a clause-list trait that is defined with properties that match the clauses accepted by the *declare simd* directive with the same name and semantics. The *simd* trait defines at least the *simdlen* property and one of the *inbranch* or *notinbranch* properties. Traits in the *construct* set other than *simd* are non-property traits.

The *device* set includes traits that define the characteristics of the device being targeted by the compiler at that point in the program. For each *target device* that the implementation supports, a *target_device* set exists that defines the characteristics of that device. At least the following traits must be defined for the *device* and all *target_device* sets:

- The *kind(kind-name-list)* trait specifies the general kind of the device. The following *kind-name* values are defined:
  - *host*, which specifies that the device is the host device;
  - *nohost*, which specifies that the devices is not the host device; and
  - the values defined in the *OpenMP Additional Definitions* document.
- The *isa(isa-name-list)* trait specifies the Instruction Set Architectures supported by the device. The accepted *isa-name* values are implementation defined.
- The *arch(arch-name-list)* trait specifies the architectures supported by the device. The accepted *arch-name* values are implementation defined.

The *kind*, *isa* and *arch* traits in the *device* and *target_device* sets are name-list traits.
Additionally, the target_device set defines the following trait:

- The device_num trait specifies the device number of the device.

The implementation set includes traits that describe the functionality supported by the OpenMP implementation at that point in the program. At least the following traits can be defined:

- The vendor(vendor-name-list) trait, which specifies the vendor identifiers of the implementation. OpenMP defined values for vendor-name are defined in the OpenMP Additional Definitions document.

- The extension(extension-name-list) trait, which specifies vendor specific extensions to the OpenMP specification. The accepted extension-name values are implementation defined.

- A trait with a name that is identical to the name of any clause that was supplied to the requires directive prior to the program point. Such traits other than the atomic_default_mem_order trait are non-property traits. The presence of these traits has been deprecated.

- A requires(requires-clause-list) trait, which is a clause-list trait for which the properties are the clauses that have been supplied to the requires directive prior to the program point as well as implementation defined implicit requirements.

The vendor and extension traits in the implementation set are name-list traits.

Implementations can define additional traits in the device, target_device and implementation sets; these traits are extension traits.

The dynamic trait set includes traits that define the dynamic properties of a program at a point in its execution. The data state trait in the dynamic trait set refers to the complete data state of the program that may be accessed at runtime.

### 2.3.2 Context Selectors

Context selectors are used to define the properties that can match an OpenMP context. OpenMP defines different sets of selectors, each containing different selectors.

The syntax for a context selector is context-selector-specification as described in the following grammar:

```plaintext
context-selector-specification : 
  trait-set-selector[, trait-set-selector[,...]] 

trait-set-selector : 
  trait-set-selector-name={trait-selector[, trait-selector[,...]]} 

trait-selector : 
  trait-selector-name {trait-score : } trait-property[, trait-property[,...]] 

trait-property : 
```

CHAPTER 2. DIRECTIVES 55
trait-property-name
or
trait-property-clause
or
trait-property-expression
or
trait-property-extension

trait-property-clause :
  clause

trait-property-name :
  identifier
or
  string-literal

trait-property-expression
  scalar-expression (for C/C++)
or
  scalar-logical-expression (for Fortran)
or
  scalar-integer-expression (for Fortran)

trait-score :
  score(score-expression)

trait-property-extension :
  trait-property-name
or
  identifier(trait-property-extension[, trait-property-extension[, ...]])
or
  constant integer expression

For trait selectors that correspond to name-list traits, each \textit{trait-property} should be \textit{trait-property-name} and for any value that is a valid identifier both the identifier and the corresponding string literal (for C/C++) and the corresponding \textit{char-literal-constant} (for Fortran) representation are considered representations of the same value.

For trait selectors that correspond to clause-list traits, each \textit{trait-property} should be \textit{trait-property-clause}. The syntax is the same as for the matching OpenMP clause.

The \texttt{construct} selector set defines the \texttt{construct} traits that should be active in the OpenMP context. The following selectors can be defined in the \texttt{construct} set: \texttt{target}, \texttt{teams}, \texttt{parallel}, \texttt{for} (in C/C++); \texttt{do} (in Fortran); \texttt{simd} and \texttt{dispatch}. Each \textit{trait-property} of the \texttt{simd} selector is a \textit{trait-property-clause}. The syntax is the same as for a valid clause of the
declare simd directive and the restrictions on the clauses from that directive apply. The construct selector is an ordered list \( c_1, \ldots, c_N \).

The device and implementation selector sets define the traits that should be active in the corresponding trait set of the OpenMP context. The target_device selector set defines the traits that should be active in the target_device trait set for the device that the specified device_num selector identifies. The same traits that are defined in the corresponding traits sets can be used as selectors with the same properties. The kind selector of the device and target_device selector sets can also specify the value any, which is as if no kind selector was specified. If a device_num selector does not appear in the target_device selector set then a device_num selector that specifies the value of the default-device-var ICV is implied. For the device_num selector of the target_device selector set, a single trait-property-expression must be specified. For the atomic_default_mem_order selector of the implementation set, a single trait-property must be specified as an identifier equal to one of the valid arguments to the atomic_default_mem_order clause on the requires directive. For the requires selector of the implementation set, each trait-property is a trait-property-clause. The syntax is the same as for a valid clause of the requires directive and the restrictions on the clauses from that directive apply.

The user selector set defines the condition selector that provides additional user-defined conditions.

The condition selector contains a single trait-property-expression that must evaluate to true for the selector to be true.

Any non-constant expression that is evaluated to determine the suitability of a variant is evaluated according to the data state trait in the dynamic trait set of the OpenMP context.

The user selector set is dynamic if the condition selector is present and the expression in the condition selector is not a constant expression; otherwise, it is static.

All parts of a context selector define the static part of the context selector except the following parts, which define the dynamic part of a context selector:

- Its user selector set if it is dynamic; and
- Its target_device selector set.

For the match clause of a declare variant directive, any argument of the base function that is referenced in an expression that appears in the context selector is treated as a reference to the expression that is passed into that argument at the call to the base function. Otherwise, a variable or procedure reference in an expression that appears in a context selector is a reference to the variable or procedure of that name that is visible at the location of the directive on which the selector appears.
Each occurrence of the `this` pointer in an expression in a context selector that appears in the `match` clause of a `declare variant` directive is treated as an expression that is the address of the object on which the associated base function is invoked.

Implementations can allow further selectors to be specified. Each specified `trait-property` for these implementation-defined selectors should be `trait-property-extension`. Implementations can ignore specified selectors that are not those described in this section.

**Restrictions**

Restrictions to context selectors are as follows:

- Each `trait-property` can only be specified once in a `trait-selector` other than the `construct` selector set.
- Each `trait-set-selector-name` can only be specified once.
- Each `trait-selector-name` can only be specified once.
- A `trait-score` cannot be specified in traits from the `construct`, `device` or `target_device` trait-selector-sets.
- A `score-expression` must be a non-negative constant integer expression.
- The expression of a `device_num` trait must evaluate to a non-negative integer value that is less than or equal to the value of `omp_get_num_devices()`.
- A variable or procedure that is referenced in an expression that appears in a context selector must be visible at the location of the directive on which the selector appears unless the directive is a `declare variant` directive and the variable is an argument of the associated base function.
- If `trait-property` any is specified in the `kind` trait-selector of the `device` or `target_device` selector set, no other `trait-property` may be specified in the same selector.
- For a `trait-selector` that corresponds to a name-list trait, at least one `trait-property` must be specified.
- For a `trait-selector` that corresponds to a non-property trait, no `trait-property` may be specified.
- For the `requires` selector of the `implementation` selector set, at least one `trait-property` must be specified.
2.3.3 Matching and Scoring Context Selectors

A given context selector is compatible with a given OpenMP context if the following conditions are satisfied:

- All selectors in the **user** set of the context selector are true;

- All traits and trait properties that are defined by selectors in the **target_device** set of the context selector are active in the **target_device** trait set for the device that is identified by the **device_num** selector;

- All traits and trait properties that are defined by selectors in the **construct**, **device** and **implementation** sets of the context selector are active in the corresponding trait sets of the OpenMP context;

- For each selector in the context selector, its properties are a subset of the properties of the corresponding trait of the OpenMP context;

- Selectors in the **construct** set of the context selector appear in the same relative order as their corresponding traits in the **construct** trait set of the OpenMP context; and

- No specified implementation-defined selector is ignored by the implementation.

Some properties of the **simd** selector have special rules to match the properties of the **simd** trait:

- The **simdlen**(N) property of the selector matches the **simdlen**(M) trait of the OpenMP context if M%N equals zero; and

- The **aligned**(list:N) property of the selector matches the **aligned**(list:M) trait of the OpenMP context if N%M equals zero.

Among compatible context selectors, a score is computed using the following algorithm:

1. Each trait selector for which the corresponding trait appears in the **construct** trait set in the OpenMP context is given the value 2^{p-1} where p is the position of the corresponding trait, c_p, in the context **construct** trait set; if the traits that correspond to the **construct** selector set appear multiple times in the OpenMP context, the highest valued subset of context traits that contains all selectors in the same order are used;

2. The **kind**, **arch**, and **isa** selectors, if specified, are given the values 2^l, 2^{l+1} and 2^{l+2}, respectively, where l is the number of traits in the **construct** set;

3. Trait selectors for which a **trait-score** is specified are given the value specified by the **trait-score** score-expression;

4. The values given to any additional selectors allowed by the implementation are implementation defined;

5. Other selectors are given a value of zero; and

6. A context selector that is a strict subset of another context selector has a score of zero. For other context selectors, the final score is the sum of the values of all specified selectors plus 1.
2.3.4 Metadirectives

Summary
A metadirective is a directive that can specify multiple directive variants of which one may be conditionally selected to replace the metadirective based on the enclosing OpenMP context.

Syntax

C / C++

The syntax of a metadirective is as follows:

```c
#pragma omp metadirective [clause [, clause] ... ] new-line
```

or

```c
#pragma omp begin metadirective [clause [, clause] ... ] new-line stmt(s)
#pragma omp end metadirective
```

where clause is one of the following:

```c
when (context-selector-specification: [directive-variant])
default ([directive-variant])
```

Fortran

In the when clause, context-selector-specification specifies a context selector (see Section 2.3.2). In the when and default clauses, directive-variant has the following form and specifies a directive variant that specifies an OpenMP directive with clauses that apply to it.

```c
directive-name [clause [, clause] ... ]
```

```fortran
!$omp metadirective [clause [, clause] ... ]

or

 !$omp begin metadirective [clause [, clause] ... ] stmt(s)
 !$omp end metadirective
```

where clause is one of the following:

```fortran
when (context-selector-specification: [directive-variant])
default ([directive-variant])
```
Description

A metadirective is replaced by a **nothing** directive or one of the directive variants specified by the **when** clauses or the **default** clause.

If no **default** clause is specified the effect is as if a **default** clause without an associated directive variant was specified.

The **default** clause is treated as a **when** clause with the specified directive variant, if any, and an always compatible static context selector that has a score lower than the scores associated with any other clause.

If a **when** clause does not explicitly specify a directive variant it implicitly specifies a **nothing** directive as the directive variant.

The OpenMP context for a given metadirective is defined according to Section 2.3.1.

The directive variant specified by a **when** clause is a candidate to replace the metadirective if the static part of the corresponding context selector is compatible with the OpenMP context according to the matching rules defined in Section 2.3.3.

Replacement candidates are ordered according to the following rules in decreasing precedence:

- A candidate is before another one if the score associated with the context selector of the corresponding **when** clause is higher.
- A candidate that was explicitly specified is before one that was implicitly specified.
- Candidates are ordered according to the order in which they lexically appear on the metadirective.

The list of dynamic replacement candidates is the prefix of the sorted list of replacement candidates up to and including the first candidate for which the corresponding **when** clause has a static context selector.

The first dynamic replacement candidate for which the corresponding **when** clause has a compatible context selector, according to the matching rules defined in Section 2.3.3, replaces the metadirective.

Expressions that appear in the context selector of a **when** clause are evaluated if no prior dynamic replacement candidate has a compatible context selector, and the number of times each expression is evaluated is implementation defined. All variables referenced by these expressions are considered to be referenced by the metadirective.

A directive variant that is associated with a **when** clause may only affect the program if the directive variant is a dynamic replacement candidate.

The **begin metadirective** directive behaves identically to the **metadirective** directive, except that the directive syntax for the specified directive variants other than the **nothing** directive must accept a paired **end directive**. For any directive variant that is selected to replace the **begin metadirective** directive, the **end metadirective** directive will be implicitly replaced by its paired **end directive** to demarcate the statements that are affected by or are
associated with the directive variant. If the `nothing` directive is selected to replace the `begin metadirective` directive, its paired `end metadirective` directive is ignored.

**Restrictions**

Restrictions to metadirectives are as follows:

- The directive variant appearing in a `when` or `default` clause must not specify a `metadirective`, `begin metadirective`, or `end metadirective` directive.
  
- The directive variant that appears in a `when` or `default` clause must not specify a `begin declare variant` or `end declare variant`.

- The context selector that appears in a `when` clause must not specify any properties for the `simd` selector.

- Replacement of the metadirective with the directive variant associated with any of the dynamic replacement candidates must result in a conforming OpenMP program.

- Insertion of user code at the location of a metadirective must be allowed if the first dynamic replacement candidate does not have a static context selector.

- All items must be executable directives if the first dynamic replacement candidate does not have a static context selector.

- Any directive variant that is specified by a `when` or `default` clause on a `begin metadirective` directive must be an OpenMP directive that has a paired `end directive` or must be the `nothing` directive, and the `begin metadirective` directive must have a paired `end metadirective` directive.

- The `default` clause may appear at most once on a metadirective.
2.3.5 Declare Variant Directive

Summary
The declare variant directive declares a specialized variant of a base function and specifies the context in which that specialized variant is used. The declare variant directive is a declarative directive.

Syntax
The syntax of the declare variant directive is as follows:

```
#pragma omp declare variant(variant-func-id) clause [...] new-line
```

or

```
#pragma omp begin declare variant clause new-line
declaration-definition-seq
#pragma omp end declare variant new-line
```

where clause is one of the following:

```
match(context-selector-specification)
adjust_args(adjust-op: argument-list)
append_args(append-op[[, append-op]...])
```

where adjust-op is one of the following:

```
nothing
need_device_ptr
```

where append-op is one of the following:

```
interop(interop-type[[, interop-type]...])
```

and where variant-func-id is the name of a function variant that is either a base language identifier or, for C++, a template-id.
The syntax of the `declare variant` directive is as follows:

```fortran
!$omp declare variant ([base-proc-name : variant-proc-name] clause [[ , ] clause] ...)
```

where `clause` is one of the following:

- `match (context-selector-specification)`
- `adjust_args (adjust-op : argument-list)`
- `append_args (append-op [ , append-op]...)`

where `adjust-op` is one of the following:

- `nothing`
- `need_device_ptr`

where `append-op` is one of the following:

- `interop (interop-type [ , interop-type]...)`

and where `variant-proc-name` is the name of a function variant that is a base language identifier.

**Description**

The declare variant directive declares a base function to have the specified function variant. The context selector in the `match` clause is associated with the variant.

The `begin declare variant` directive associates the context selector in the `match` clause with each function definition in `declaration-definition-seq`.

For the purpose of call resolution, each function definition that appears between a `begin declare variant` directive and its paired `end declare variant` directive is a function variant for an assumed base function, with the same name and a compatible prototype, that is declared elsewhere without an associated declare variant directive.

If a declare variant directive appears between a `begin declare variant` directive and its paired `end declare variant` directive the effective context selectors of the outer directive are appended to the context selector of the inner directive to form the effective context selector of the inner directive. If a `trait-set-selector` is present on both directives, the `trait-selector` list of the outer directive is appended to the `trait-selector` list of the inner directive after equivalent `trait-selectors` have been removed from the outer list. Restrictions that apply to explicitly specified context selectors also apply to effective context selectors constructed through this process.

The symbol name of a function definition that appears between a `begin declare variant` directive and its paired `end declare variant` directive shall be determined through the base
language rules after the name of the function has been augmented with a string that shall be
determined according to the effective context selector of the **begin declare variant**
directive. The symbol names of two definitions of a function shall be equal if and only if their
effective context selectors are equivalent.

If the context selector of a **begin declare variant** directive contains traits in the `device` or
`implementation` set that are known never to be compatible with an OpenMP context during the
current compilation, the preprocessed code that follows the **begin declare variant** directive
up to the matching **end declare variant** directive shall be elided.

The OpenMP context for a direct call to a given base function is defined according to Section 2.3.1.
If a declare variant directive for the base function is visible at the call site and the static part of the
context selector that is associated with the declared function variant is compatible with the
OpenMP context of the call according to the matching rules defined in Section 2.3.3 then the
variant is a replacement candidate to be called instead of the base function. Replacement
candidates are ordered in decreasing order of the score associated with the context selector. If two
replacement candidates have the same score then their order is implementation defined.

The list of dynamic replacement candidates is the prefix of the sorted list of replacement candidates
up to and including the first candidate for which the corresponding context selector is static.

The first dynamic replacement candidate for which the corresponding context selector is
compatible, according to the matching rules defined in Section 2.3.3, is called instead of the base
function. If no compatible candidate exists then the base function is called.

Expressions that appear in the context selector of a **match** clause are evaluated if no prior dynamic
replacement candidate has a compatible context selector, and the number of times each expression
is evaluated is implementation defined. All variables referenced by these expressions are
considered to be referenced at the call site.

For calls to **constexpr** base functions that are evaluated in constant expressions, whether any
variant replacement occurs is implementation defined.

For indirect function calls that can be determined to call a particular base function, whether any
variant replacement occurs is unspecified.

For each **adjust args** clause that is present on the selected variant the adjustment operation
specified by **adjust-op** will be applied to each of the arguments specified in the clause before being
passed to the selected variant.

If the **adjust-op** modifier is **need_device_ptr**, the arguments are converted to corresponding
device pointers of the default device. If an argument has the **is_device_ptr** property in its
**interoperability requirement set** then the argument will be passed as is. Otherwise, the argument
will be converted in the same manner that a **use_device_ptr** clause on a **target data**
construct converts its pointer list items into device pointers. If the argument cannot be converted into a device pointer then the NULL value will be passed as the argument.

If the adjust-op modifier is nothing, the argument is passed to the selected variant without being modified.

If an append_args clause is present on the matching directive then additional arguments are passed in the call. The arguments are constructed according to any specified append-op modifiers and are passed in the same order in which they are specified in the append_args clause.

The interop operation constructs an argument of type omp_interop_t from the interoperability requirement set of the encountering task.

The interop operation constructs an argument of type omp_interop_kind from the interoperability requirement set of the encountering task.

The argument is constructed as if an interop construct with an init clause of interop-types was specified. If the interoperability requirement set contains one or more properties that could be used as clauses for an interop construct of the interop-type type, the behavior is as if the corresponding clauses would also be part of the aforementioned interop construct and those properties will be removed from the interoperability requirement set.

This argument is destroyed after the call to the selected variant returns, as if an interop construct with a destroy clause was used with the same clauses that were used to initialize the argument.

Any differences that the specific OpenMP context requires in the prototype of the variant from the base function prototype are implementation defined.

For the declare variant directive, any expressions in the match clause are interpreted as if they appeared in the scope of arguments of the base function.

Different declare variant directives may be specified for different declarations of the same base function.

The function variant is determined by base language standard name lookup rules ([basic.lookup]) of variant-func-id using the argument types at the call site after implementation-defined changes have been made according to the OpenMP context.

For the declare variant directive, the variant-func-id and any expressions in the match clause are interpreted as if they appeared at the scope of the trailing return type of the base function.
For the `begin declare variant` directive, any expressions in the `match` clause are interpreted at the location of the directive.

The procedure to which `base-proc-name` refers is resolved at the location of the directive according to the establishment rules for procedure names in the base language.

**Restrictions**

Restrictions to the declare variant directive are as follows:

- Calling functions that a declare variant directive determined to be a function variant directly in an OpenMP context that is different from the one that the `construct` selector set of the context selector specifies is non-conforming.

- If a function is determined to be a function variant through more than one declare variant directive then the `construct` selector set of their context selectors must be the same.

- A function determined to be a function variant may not be specified as a base function in another declare variant directive.

- All variables that are referenced in an expression that appears in the context selector of a `match` clause must be accessible at a call site to the base function according to the base language rules.

- At most one `match` clause can appear on a `declare variant` directive.

- At most one `append_args` clause can appear on a `declare variant` directive.

- Each argument can only appear in a single `adjust_args` clause for each `declare variant` directive.

- An `adjust_args` clause or `append_args` clause can only be specified if the `dispatch` selector of the `construct` selector set appears in the `match` clause.

- The type of the function variant must be compatible with the type of the base function after the implementation-defined transformation for its OpenMP context.

- Only the `match` clause can appear on a `begin declare variant` directive.

- The `match` clause of a `begin declare variant` directive may not contain a `simd` `trait-selector-name`.

- Matching pairs of `begin declare variant` and `end declare variant` directives shall either encompass disjoint source ranges or they shall be perfectly nested.
• The declare variant directive cannot be specified for a virtual, defaulted or deleted function.

• The declare variant directive cannot be specified for a constructor or destructor.

• The declare variant directive cannot be specified for an immediate function.

• The function that a declare variant directive determined to be a function variant may not be an immediate function.

• A match clause that appears on a `begin declare target` directive must not contain a dynamic context selector that references the `this` pointer.

• If an expression in the context selector that appears in a match clause references the `this` pointer, the base function must be a non-static member function.

• `base-proc-name` must not be a generic name, an entry name, the name of a procedure pointer, a dummy procedure or a statement function.

• If `base-proc-name` is omitted then the declare variant directive must appear in an interface block or the specification part of a procedure.

• Any declare variant directive must appear in the specification part of a subroutine subprogram, function subprogram, or interface body to which it applies.

• If the directive is specified for a procedure that is declared via a procedure declaration statement, the `base-proc-name` must be specified.

• The procedure `base-proc-name` must have an accessible explicit interface at the location of the directive.

• Each argument that appears in a `need_device_ptr adjust-op` must be of type `C_PTR` in the dummy argument declaration.

Cross References

• OpenMP Context Specification, see Section 2.3.1.

• Context Selectors, see Section 2.3.2.
2.3.6 dispatch Construct

Summary
The dispatch construct controls whether variant substitution occurs for a given call.

Syntax

The syntax of the dispatch construct is as follows:

```
#pragma omp dispatch [clause [ , ] clause ... ] new-line
expression-stmt
```

where expression-stmt is an expression statement with one of the following forms:

```
ext = target-call ( [expression-list] );
target-call ( [expression-list] );
```

and where clause is one of the following:

```
device(integer-expression)
depend([depend-modifier, ] dependence-type : locator-list)
nowait
novariants(scalar-expression)
nocontext(scalar-expression)
is_device_ptr(list)
```

The syntax of the dispatch construct is as follows:

```
!$omp dispatch [clause [ , ] clause ... ]
stmt
```

where stmt is an expression statement with one of the following forms:

```
ext = target-call ( [arguments] )
CALL target-call [ ( [arguments] ) ]
```
and where *clause* is one of the following:

- **device** *(scalar-integer-expression)*
- **depend** *(depend-modifier, [dependence-type : locator-list]*)
- **nowait**
- **novariants** *(scalar-logical-expression)*
- **nocontext** *(scalar-logical-expression)*
- **is_device_ptr** *(list)*

**Fortran**

### Binding

The binding task set for a **dispatch** region is the generating task. The **dispatch** region binds to the region of the generating task.

### Description

When a **novariants** clause is present on the **dispatch** construct, and the **novariants** clause expression evaluates to *true*, no function variant will be selected for the *target-call* even if one would be selected normally. The use of a variable in a **novariants** clause expression of a **dispatch** construct causes an implicit reference to the variable in all enclosing constructs.

The **novariants** clause expression is evaluated in the enclosing context.

When a **nocontext** clause is present on the **dispatch** construct, and the **nocontext** clause expression evaluates to *true*, the **dispatch** construct is not added to the **construct** set of the OpenMP context. The use of a variable in a **nocontext** clause expression of a **dispatch** construct causes an implicit reference to the variable in all enclosing constructs.

The **nocontext** clause expression is evaluated in the enclosing context.

The **is_device_ptr** clause indicates that its list items are device pointers. For each list item specified in the clause, an **is_device_ptr** property for that list item is added to the **interoperability requirement set**. Support for device pointers created outside of OpenMP, specifically outside of any OpenMP mechanism that returns a device pointer, is implementation defined.

If one or more **depend** clauses are present on the **dispatch** construct, they are added as **depend** properties of the **interoperability requirement set**. If a **nowait** clause is present on the **dispatch** construct the **nowait** property is added to the **interoperability requirement set**.

This construct creates an explicit task, as if the **task** construct was used, that surrounds the associated code. Properties added to the **interoperability requirement set** can be removed by the effect of other directives (see Section 2.15.2) before the task is created. If the **interoperability requirement set** contains one or more **depend** properties, the behavior is as if those properties were applied to the **task** construct as **depend** clauses. If the **interoperability requirement set** does not contain the **nowait** property then the task will also be an included task.
If the `device` clause is present, the value of the `default-device-var` ICV of the generated task is set to the value of the expression in the clause.

**Restrictions**

Restrictions to the `dispatch` construct are as follows:

- At most one `novariants` clause can appear on a `dispatch` directive.
- At most one `nocontext` clause can appear on a `dispatch` directive.
- At most one `nowait` clause can appear on a `dispatch` directive.
- A list item that appears in an `is_device_ptr` clause must be a valid device pointer for the device data environment.

- The `target-call` expression can only be a direct call.

- `target-call` must be a procedure name.
- `target-call` must not be a procedure pointer.
- A list item that appears in an `is_device_ptr` clause must be of type `C_PTR`.

**Cross References**

- `declare variant` directive, see Section 2.3.5.
- Interoperability requirement set, see Section 2.15.2.

### 2.4 Internal Control Variables

An OpenMP implementation must act as if internal control variables (ICVs) control the behavior of an OpenMP program. These ICVs store information such as the number of threads to use for future parallel regions, the schedule to use for worksharing loops and whether nested parallelism is enabled or not. The ICVs are given values at various times (described below) during the execution of the program. They are initialized by the implementation itself and may be given values through OpenMP environment variables and through calls to OpenMP API routines. The program can retrieve the values of these ICVs only through OpenMP API routines.

For purposes of exposition, this document refers to the ICVs by certain names, but an implementation is not required to use these names or to offer any way to access the variables other than through the ways shown in Section 2.4.2.
2.4.1 ICV Descriptions

The following ICVs store values that affect the operation of parallel regions.

- dyn-var - controls whether dynamic adjustment of the number of threads is enabled for encountered parallel regions. One copy of this ICV exists per data environment.
- nthreads-var - controls the number of threads requested for encountered parallel regions. One copy of this ICV exists per data environment.
- thread-limit-var - controls the maximum number of threads that participate in the contention group. One copy of this ICV exists per data environment.
- max-active-levels-var - controls the maximum number of nested active parallel regions when the innermost parallel region is generated by a given task. One copy of this ICV exists per data environment.
- place-partition-var - controls the place partition available to the execution environment for encountered parallel regions. One copy of this ICV exists per implicit task.
- active-levels-var - the number of nested active parallel regions that enclose a given task such that all of the parallel regions are enclosed by the outermost initial task region on the device on which the task executes. One copy of this ICV exists per data environment.
- levels-var - the number of nested parallel regions that enclose a given task such that all of the parallel regions are enclosed by the outermost initial task region on the device on which the task executes. One copy of this ICV exists per data environment.
- bind-var - controls the binding of OpenMP threads to places. When binding is requested, the variable indicates that the execution environment is advised not to move threads between places. The variable can also provide default thread affinity policies. One copy of this ICV exists per data environment.

The following ICVs store values that affect the operation of worksharing-loop regions.

- run-sched-var - controls the schedule that is used for worksharing-loop regions when the runtime schedule kind is specified. One copy of this ICV exists per data environment.
- def-sched-var - controls the implementation defined default scheduling of worksharing-loop regions. One copy of this ICV exists per device.

The following ICVs store values that affect program execution.

- stacksize-var - controls the stack size for threads that the OpenMP implementation creates. One copy of this ICV exists per device.
- wait-policy-var - controls the desired behavior of waiting threads. One copy of this ICV exists per device.
- display-affinity-var - controls whether to display thread affinity. One copy of this ICV exists for the whole program.
• **affinity-format-var** - controls the thread affinity format when displaying thread affinity. One copy of this ICV exists per device.

• **cancel-var** - controls the desired behavior of the `cancel` construct and cancellation points. One copy of this ICV exists for the whole program.

• **default-device-var** - controls the default target device. One copy of this ICV exists per data environment.

• **target-offload-var** - controls the offloading behavior. One copy of this ICV exists for the whole program.

• **max-task-priority-var** - controls the maximum priority value that can be specified in the `priority` clause of the `task` construct. One copy of this ICV exists for the whole program.

The following ICVs store values that affect the operation of the OMPT tool interface.

• **tool-var** - controls whether an OpenMP implementation will try to register a tool. One copy of this ICV exists for the whole program.

• **tool-libraries-var** - specifies a list of absolute paths to tool libraries for OpenMP devices. One copy of this ICV exists for the whole program.

• **tool-verbose-init-var** - controls whether an OpenMP implementation will verbosely log the registration of a tool. One copy of this ICV exists for the whole program.

The following ICVs store values that affect the operation of the OMPD tool interface.

• **debug-var** - controls whether an OpenMP implementation will collect information that an OMPD library can access to satisfy requests from a tool. One copy of this ICV exists for the whole program.

The following ICVs store values that may be queried by interface routines.

• **num-procs-var** - the number of processors that are available to the device. One copy of this ICV exists per device.

• **thread-num-var** - the thread number of an implicit task within its binding team. One copy of this ICV exists per data environment.

• **final-task-var** - whether a given task is a final task. One copy of this ICV exists per data environment.

• **implicit-task-var** - whether a given task is an implicit task. One copy of this ICV exists per data environment.

• **team-size-var** - the size of the current team. One copy of this ICV exists per data environment.

The following ICV stores values that affect default memory allocation.
def-allocator-var – controls the memory allocator to be used by memory allocation routines, directives and clauses when a memory allocator is not specified by the user. One copy of this ICV exists per implicit task.

The following ICVs store values that affect the operation of teams regions.

- nteams-var – controls the number of teams requested for encountered teams regions. One copy of this ICV exists per device.
- teams-thread-limit-var – controls the maximum number of threads that participate in each contention group created by a teams construct. One copy of this ICV exists per device.

### 2.4.2 ICV Initialization

Table 2.1 shows the ICVs, associated environment variables, and initial values.

**TABLE 2.1: ICV Initial Values**

<table>
<thead>
<tr>
<th>ICV</th>
<th>Environment Variable</th>
<th>Initial value</th>
</tr>
</thead>
<tbody>
<tr>
<td>dyn-var</td>
<td>OMP_DYNAMIC</td>
<td>See description below</td>
</tr>
<tr>
<td>nthreads-var</td>
<td>OMP_NUM_THREADS</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>run-sched-var</td>
<td>OMP_SCHEDULE</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>def-sched-var</td>
<td>(none)</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>bind-var</td>
<td>OMP_PROC_BIND</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>stacksize-var</td>
<td>OMP_STACKSIZE</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>wait-policy-var</td>
<td>OMP_WAIT_POLICY</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>thread-limit-var</td>
<td>OMP_THREAD_LIMIT</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>max-active-levels-var</td>
<td>OMP_MAX_ACTIVE_LEVELS, OMP_NESTED, OMP_NUM_THREADS, OMP_PROC_BIND</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>active-levels-var</td>
<td>(none)</td>
<td>zero</td>
</tr>
<tr>
<td>levels-var</td>
<td>(none)</td>
<td>zero</td>
</tr>
<tr>
<td>place-partition-var</td>
<td>OMP_PLACES</td>
<td>Implementation defined</td>
</tr>
</tbody>
</table>

*table continued on next page*
<table>
<thead>
<tr>
<th>ICV</th>
<th>Environment Variable</th>
<th>Initial value</th>
</tr>
</thead>
<tbody>
<tr>
<td>cancel-var</td>
<td>OMP_CANCELLATION</td>
<td>false</td>
</tr>
<tr>
<td>display-affinity-var</td>
<td>OMP_DISPLAY_AFFINITY</td>
<td>false</td>
</tr>
<tr>
<td>affinity-format-var</td>
<td>OMP_AFFINITY_FORMAT</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>default-device-var</td>
<td>OMP_DEFAULT_DEVICE</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>target-offload-var</td>
<td>OMP_TARGET_OFFLOAD</td>
<td>DEFAULT</td>
</tr>
<tr>
<td>max-task-priority-var</td>
<td>OMP_MAX_TASK_PRIORITY</td>
<td>zero</td>
</tr>
<tr>
<td>tool-var</td>
<td>OMP_TOOL</td>
<td>enabled</td>
</tr>
<tr>
<td>tool-libraries-var</td>
<td>OMP_TOOL_LIBRARIES</td>
<td>empty string</td>
</tr>
<tr>
<td>tool-verbose-init-var</td>
<td>OMP_TOOL_VERBOSE_INIT</td>
<td>disabled</td>
</tr>
<tr>
<td>debug-var</td>
<td>OMP_DEBUG</td>
<td>disabled</td>
</tr>
<tr>
<td>num-procs-var</td>
<td>(none)</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>thread-num-var</td>
<td>(none)</td>
<td>zero</td>
</tr>
<tr>
<td>final-task-var</td>
<td>(none)</td>
<td>false</td>
</tr>
<tr>
<td>implicit-task-var</td>
<td>(none)</td>
<td>true</td>
</tr>
<tr>
<td>team-size-var</td>
<td>(none)</td>
<td>one</td>
</tr>
<tr>
<td>def-allocator-var</td>
<td>OMP_ALLOCATOR</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>nteams-var</td>
<td>OMP_NUM_TEAMS</td>
<td>zero</td>
</tr>
<tr>
<td>teams-thread-limit-var</td>
<td>OMP_TEAMS_THREAD_LIMIT</td>
<td>zero</td>
</tr>
</tbody>
</table>

Each ICV that does not have global scope (see Table 2.3) has a set of device-specific environment variables that extend the variables defined in Table 2.1 with the following syntax:

\[<ENVIRONMENT VARIABLE>_DEV[_<device>]\]

where \(<ENVIRONMENT VARIABLE>\) is one of the variables from Table 2.1 and \(<device>\) is the device number as specified in the \texttt{device} clause (see Section 2.14).

**Description**

- Each device has its own ICVs.
- The initial value of \texttt{dyn-var} is implementation defined if the implementation supports dynamic adjustment of the number of threads; otherwise, the initial value is \texttt{false}.
• The value of the *nthreads*-var ICV is a list.

• The value of the *bind*-var ICV is a list.

The host and non-host device ICVs are initialized before any OpenMP API construct or OpenMP API routine executes. After the initial values are assigned, the values of any OpenMP environment variables that were set by the user are read and the associated ICVs are modified accordingly. If no `<device>` number is specified on the device-specific environment variable then the value is applied to all non-host devices.

Cross References

• *OMP_SCHEDULE* environment variable, see Section 6.1.

• *OMP_NUM_THREADS* environment variable, see Section 6.2.

• *OMP_DYNAMIC* environment variable, see Section 6.3.

• *OMP_PROC_BIND* environment variable, see Section 6.4.

• *OMP_PLACES* environment variable, see Section 6.5.

• *OMP_STACKSIZE* environment variable, see Section 6.6.

• *OMP_WAIT_POLICY* environment variable, see Section 6.7.

• *OMP_MAX_ACTIVE_LEVELS* environment variable, see Section 6.8.

• *OMP_NESTED* environment variable, see Section 6.9.

• *OMP_THREAD_LIMIT* environment variable, see Section 6.10.

• *OMP_CANCELLATION* environment variable, see Section 6.11.

• *OMP_DISPLAY AFFINITY* environment variable, see Section 6.13.

• *OMP AFFINITY_FORMAT* environment variable, see Section 6.14.

• *OMP_DEFAULT_DEVICE* environment variable, see Section 6.15.

• *OMP_MAX_TASK_PRIORITY* environment variable, see Section 6.16.

• *OMP_TARGET_OFFLOAD* environment variable, see Section 6.17.

• *OMP_TOOL* environment variable, see Section 6.18.

• *OMP_TOOL_LIBRARIES* environment variable, see Section 6.19.

• *OMP_DEBUG* environment variable, see Section 6.21.

• *OMP_ALLOCATOR* environment variable, see Section 6.22.

• *OMP_NUM_TEAMS* environment variable, see Section 6.23.

• *OMP_TEAMS THREAD LIMIT* environment variable, see Section 6.24.
2.4.3 Modifying and Retrieving ICV Values

Table 2.2 shows the method for modifying and retrieving the values of ICVs through OpenMP API routines. If an ICV is not listed in this table, no OpenMP API routine modifies or retrieves this ICV.

**Table 2.2:** Ways to Modify and to Retrieve ICV Values

<table>
<thead>
<tr>
<th>ICV</th>
<th>Ways to Modify Value</th>
<th>Ways to Retrieve Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>dyn-var</td>
<td>omp_set_dynamic</td>
<td>omp_get_dynamic</td>
</tr>
<tr>
<td>nthreads-var</td>
<td>omp_set_num_threads</td>
<td>omp_get_max_threads</td>
</tr>
<tr>
<td>run-sched-var</td>
<td>omp_set_schedule</td>
<td>omp_get_schedule</td>
</tr>
<tr>
<td>bind-var</td>
<td>(none)</td>
<td>omp_get_proc_bind</td>
</tr>
<tr>
<td>thread-limit-var</td>
<td>target construct, teams construct</td>
<td>omp_get_thread_limit</td>
</tr>
<tr>
<td>max-active-levels-var</td>
<td>omp_set_max_active_levels, omp_set_nested</td>
<td>omp_get_max_active_levels</td>
</tr>
<tr>
<td>active-levels-var</td>
<td>(none)</td>
<td>omp_get_active_level</td>
</tr>
<tr>
<td>levels-var</td>
<td>(none)</td>
<td>omp_get_level</td>
</tr>
<tr>
<td>place-partition-var</td>
<td>(none)</td>
<td>See description below</td>
</tr>
<tr>
<td>cancel-var</td>
<td>(none)</td>
<td>omp_get_cancellation</td>
</tr>
<tr>
<td>affinity-format-var</td>
<td>omp_set_affinity_format</td>
<td>omp_get_affinity_format</td>
</tr>
<tr>
<td>default-device-var</td>
<td>omp_set_default_device</td>
<td>omp_get_default_device</td>
</tr>
<tr>
<td>max-task-priority-var</td>
<td>(none)</td>
<td>omp_get_max_task_priority</td>
</tr>
<tr>
<td>num-procs-var</td>
<td>(none)</td>
<td>omp_get_num_procs</td>
</tr>
<tr>
<td>thread-num-var</td>
<td>(none)</td>
<td>omp_get_thread_num</td>
</tr>
<tr>
<td>final-task-var</td>
<td>(none)</td>
<td>omp_in_final</td>
</tr>
<tr>
<td>team-size-var</td>
<td>(none)</td>
<td>omp_get_num_threads</td>
</tr>
<tr>
<td>def-allocator-var</td>
<td>omp_set_default_allocator</td>
<td>omp_get_default_allocator</td>
</tr>
<tr>
<td>nteams-var</td>
<td>omp_set_num_teams</td>
<td>omp_get_max_teams</td>
</tr>
<tr>
<td>teams-thread-limit-var</td>
<td>omp_set_teams_thread_limit</td>
<td>omp_get_teams_thread_limit</td>
</tr>
</tbody>
</table>

**Description**

- The value of the *nthreads-var* ICV is a list. The runtime call `omp_set_num_threads` sets the value of the first element of this list, and `omp_get_max_threads` retrieves the value of the first element of this list.
• The value of the \textit{bind-var} ICV is a list. The runtime call \texttt{omp\_get\_proc\_bind} retrieves the value of the first element of this list.

• Detailed values in the \textit{place-partition-var} ICV are retrieved using the runtime calls \texttt{omp\_get\_partition\_num\_places}, \texttt{omp\_get\_partition\_place\_nums}, \texttt{omp\_get\_place\_num\_procs}, and \texttt{omp\_get\_place\_proc\_ids}.

\textbf{Cross References}

• \texttt{thread\_limit} clause of the \texttt{teams} construct, see Section 2.7.

• \texttt{thread\_limit} clause of the \texttt{target} construct, see Section 2.14.5.

• \texttt{omp\_set\_num\_threads} routine, see Section 3.2.1.

• \texttt{omp\_get\_num\_threads} routine, see Section 3.2.2.

• \texttt{omp\_get\_max\_threads} routine, see Section 3.2.3.

• \texttt{omp\_get\_thread\_num} routine, see Section 3.2.4.

• \texttt{omp\_set\_dynamic} routine, see Section 3.2.6.

• \texttt{omp\_get\_dynamic} routine, see Section 3.2.7.

• \texttt{omp\_get\_cancellation} routine, see Section 3.2.8.

• \texttt{omp\_set\_nested} routine, see Section 3.2.9.

• \texttt{omp\_get\_schedule} routine, see Section 3.2.11.

• \texttt{omp\_get\_schedule} routine, see Section 3.2.12.

• \texttt{omp\_get\_thread\_limit} routine, see Section 3.2.13.

• \texttt{omp\_get\_supported\_active\_levels}, see Section 3.2.14.

• \texttt{omp\_set\_max\_active\_levels} routine, see Section 3.2.15.

• \texttt{omp\_get\_max\_active\_levels} routine, see Section 3.2.16.

• \texttt{omp\_get\_level} routine, see Section 3.2.17.

• \texttt{omp\_get\_active\_level} routine, see Section 3.2.20.

• \texttt{omp\_get\_proc\_bind} routine, see Section 3.3.1.

• \texttt{omp\_get\_place\_num\_procs} routine, see Section 3.3.3.

• \texttt{omp\_get\_place\_proc\_ids} routine, see Section 3.3.4.

• \texttt{omp\_get\_partition\_num\_places} routine, see Section 3.3.6.

• \texttt{omp\_get\_partition\_place\_nums} routine, see Section 3.3.7.

• \texttt{omp\_set\_affinity\_format} routine, see Section 3.3.8.
• `omp_get_affinity_format` routine, see Section 3.3.9.
• `omp_set_num_teams` routine, see Section 3.4.3.
• `omp_get_max_teams` routine, see Section 3.4.4.
• `omp_set_teams_thread_limit` routine, see Section 3.4.5.
• `omp_get_teams_thread_limit` routine, see Section 3.4.6.
• `omp_get_max_task_priority` routine, see Section 3.5.1.
• `omp_in_final` routine, see Section 3.5.2.
• `omp_get_num_procs` routine, see Section 3.7.1.
• `omp_set_default_device` routine, see Section 3.7.2.
• `omp_get_default_device` routine, see Section 3.7.3.
• `omp_set_default_allocator` routine, see Section 3.13.4.
• `omp_get_default_allocator` routine, see Section 3.13.5.

2.4.4 How ICVs are Scoped

Table 2.3 shows the ICVs and their scope.

**Table 2.3:** Scopes of ICVs

<table>
<thead>
<tr>
<th>ICV</th>
<th>Scope</th>
</tr>
</thead>
<tbody>
<tr>
<td>dyn-var</td>
<td>data environment</td>
</tr>
<tr>
<td>nthreads-var</td>
<td>data environment</td>
</tr>
<tr>
<td>run-sched-var</td>
<td>data environment</td>
</tr>
<tr>
<td>def-sched-var</td>
<td>device</td>
</tr>
<tr>
<td>bind-var</td>
<td>data environment</td>
</tr>
<tr>
<td>stacksize-var</td>
<td>device</td>
</tr>
<tr>
<td>wait-policy-var</td>
<td>device</td>
</tr>
<tr>
<td>thread-limit-var</td>
<td>data environment</td>
</tr>
<tr>
<td>max-active-levels-var</td>
<td>data environment</td>
</tr>
</tbody>
</table>

*table continued on next page*
ICV    | Scope
---|---
active-levels-var  | data environment
levels-var        | data environment
place-partition-var | implicit task
cancel-var         | global
display-affinity-var | global
affinity-format-var | device
default-device-var | data environment
target-offload-var | global
max-task-priority-var | global
tool-var           | global
tool-libraries-var | global
tool-verbose-init-var | global
display-affinity-var | global
debug-var          | global
num-procs-var      | device
thread-num-var      | implicit task
final-task-var      | data environment
implicit-task-var   | data environment
team-size-var       | team
def-allocon-ator-var | implicit task
nteams-var          | device
teams-thread-limit-var | device

**Description**

- One copy of each ICV with device scope exists per device.
- Each data environment has its own copies of ICVs with data environment scope.
- Each implicit task has its own copy of ICVs with implicit task scope.

Calls to OpenMP API routines retrieve or modify data environment scoped ICVs in the data environment of their binding tasks.
2.4.4.1 How the Per-Data Environment ICVs Work

When a `task` construct, a `parallel` construct or a `teams` construct is encountered, each
generated task inherits the values of the data environment scoped ICVs from each generating task’s
ICV values.

When a `parallel` construct is encountered, the value of each ICV with implicit task scope is
inherited from the implicit binding task of the generating task unless otherwise specified.

When a `task` construct is encountered, the generated task inherits the value of `nthreads-var` from
the generating task’s `nthreads-var` value. When a `parallel` construct is encountered, and the
generating task’s `nthreads-var` list contains a single element, the generated implicit tasks inherit
that list as the value of `nthreads-var`. When a `parallel` construct is encountered, and the
generating task’s `nthreads-var` list contains multiple elements, the generated implicit tasks inherit
the value of `nthreads-var` as the list obtained by deletion of the first element from the generating
task’s `nthreads-var` value. The `bind-var` ICV is handled in the same way as the `nthreads-var` ICV.

When a `target task` executes an active `target` region, the generated initial task uses the values of
the data environment scoped ICVs from the device data environment ICV values of the device that
will execute the region.

When a `target task` executes an inactive `target` region, the generated initial task uses the values
of the data environment scoped ICVs from the data environment of the task that encountered the
`target` construct.

If a `target` construct with a `thread_limit` clause is encountered, the `thread-limit-var` ICV
from the data environment of the generated initial task is instead set to an implementation defined
value between one and the value specified in the clause.

If a `target` construct with no `thread_limit` clause is encountered, the `thread-limit-var` ICV
from the data environment of the generated initial task is set to an implementation defined value
that is greater than zero.

If a `teams` construct with a `thread_limit` clause is encountered, the `thread-limit-var` ICV
from the data environment of the initial task for each team is instead set to an implementation defined value
between one and the value specified in the clause.

If a `teams` construct with no `thread_limit` clause is encountered, the `thread-limit-var` ICV
from the data environment of the initial task of each team is set to an implementation defined value
that is greater than zero and does not exceed `teams-thread-limit-var`, if `teams-thread-limit-var` is
greater than zero.

When encountering a worksharing-loop region for which the `runtime` schedule kind is specified,
all implicit task regions that constitute the binding parallel region must have the same value for
`run-sched-var` in their data environments. Otherwise, the behavior is unspecified.
2.4.5 ICV Override Relationships

Table 2.4 shows the override relationships among construct clauses and ICVs. The table only lists ICVs that can be overwritten by a clause.

TABLE 2.4: ICV Override Relationships

<table>
<thead>
<tr>
<th>ICV</th>
<th>construct clause, if used</th>
</tr>
</thead>
<tbody>
<tr>
<td>nthreads-var</td>
<td>num_threads</td>
</tr>
<tr>
<td>run-sched-var</td>
<td>schedule</td>
</tr>
<tr>
<td>def-sched-var</td>
<td>schedule</td>
</tr>
<tr>
<td>bind-var</td>
<td>proc_bind</td>
</tr>
<tr>
<td>def-allocator-var</td>
<td>allocator</td>
</tr>
<tr>
<td>nteams-var</td>
<td>num_teams</td>
</tr>
<tr>
<td>teams-thread-limit-var</td>
<td>thread_limit</td>
</tr>
</tbody>
</table>

Description

- The `num_threads` clause overrides the value of the first element of the `nthreads-var` ICV.
- If a `schedule` clause specifies a modifier then that modifier overrides any modifier that is specified in the `run-sched-var` ICV.
- If `bind-var` is not set to `false` then the `proc_bind` clause overrides the value of the first element of the `bind-var` ICV; otherwise, the `proc_bind` clause has no effect.

Cross References

- `parallel` construct, see Section 2.6.
- `proc_bind` clause, Section 2.6.
- `num_threads` clause, see Section 2.6.1.
- `num_teams` clause, see Section 2.7.
- `thread_limit` clause, see Section 2.7.
- Worksharing-loop construct, see Section 2.11.4.
- `schedule` clause, see Section 2.11.4.1.
2.5 Informational and Utility Directives

This section covers all directives that may function as an informational directive or a utility directive in an OpenMP program.

2.5.1 requires Directive

Summary

The requires directive specifies the features that an implementation must provide in order for the code to compile and to execute correctly. The requires directive is an informational directive.

Syntax

C / C++

The syntax of the requires directive is as follows:

```
#pragma omp requires clause[ [ , ] clause] ... ] new-line
```

where clause is either a clause of the form ext_implementation-defined-requirement for an implementation defined requirement clause or one of the requirement clauses listed below:

```
reverse_offload
unified_address
unified_shared_memory
atomic_default_mem_order(seq_cst | acq_rel | relaxed)
dynamic_allocators
```

Fortran

The syntax of the requires directive is as follows:

```
!$omp requires clause[ [ , ] clause] ...
```

where clause is either a clause of the form ext_implementation-defined-requirement for an implementation defined requirement clause or one of the requirement clauses listed below:

```
reverse_offload
unified_address
unified_shared_memory
atomic_default_mem_order(seq_cst | acq_rel | relaxed)
dynamic_allocators
```
Description

The **requires** directive specifies features that an implementation must support for correct execution. The behavior that a requirement clause specifies may override the normal behavior specified elsewhere in this document. Whether an implementation supports the feature that a given requirement clause specifies is implementation defined.

The **requires** directive specifies requirements for the execution of all code in the current compilation unit.

Note – Use of this directive makes your code less portable. Users should be aware that not all devices or implementations support all requirements.

When the **reverse_offload** clause appears on a **requires** directive, the implementation guarantees that a **target** region, for which the **target** construct specifies a **device** clause in which the **ancestor** modifier appears, can execute on the parent device of an enclosing **target** region.

When the **unified_address** clause appears on a **requires** directive, the implementation guarantees that all devices accessible through OpenMP API routines and directives use a unified address space. In this address space, a pointer will always refer to the same location in memory from all devices accessible through OpenMP. Any OpenMP mechanism that returns a device pointer is guaranteed to return a device address that supports pointer arithmetic, and the **is_device_ptr** clause is not necessary to obtain device addresses from device pointers for use inside **target** regions. Host pointers may be passed as device pointer arguments to device memory routines and device pointers may be passed as host pointer arguments to device memory routines. Non-host devices may still have discrete memories and dereferencing a device pointer on the host device or a host pointer on a non-host device remains unspecified behavior.

When the **unified_address** clause appears on a **requires** directive, the base pointer of a zero-length array section that is mapped on a **target** construct and for which a matching mapped list item cannot be found is not initialized to **NULL** but instead retains its original value.

Memory local to a specific execution context may be exempt from the **unified_address** requirement, following the restrictions of locality to a given execution context, thread or contention group.

The **unified_shared_memory** clause implies the **unified_address** requirement, inheriting all of its behaviors. Additionally, storage locations in memory are accessible to threads on all available devices that are supported by the implementation, except for memory that is local to a specific execution context as defined in the description of **unified_address** above. Every device address that refers to storage allocated through OpenMP device memory routines is a valid host pointer that may be dereferenced.
The `unified_shared_memory` clause makes the `map` clause optional on `target` constructs and the declare target directive optional for variables with static storage duration that are accessed inside functions to which a declare target directive is applied. Scalar variables are still firstprivate by default when referenced inside `target` constructs. Values stored into memory by one device may not be visible to another device until those two devices synchronize with each other or both devices synchronize with the host.

The `atomic_default_mem_order` clause specifies the default memory ordering behavior for `atomic` constructs that must be provided by an implementation. Its parameter appears as a clause on any `atomic` construct that does not already specify a memory order clause.

The `dynamic Allocators` clause removes certain restrictions on the use of memory allocators in `target` regions. It makes the `uses Allocators` clause optional on `target` constructs for the purpose of using allocators in the corresponding `target` regions. It allows calls to the `omp_init_allocator` and `omp_destroy_allocator` API routines in `target` regions. Finally, it allows default allocators to be used by `allocate` directives, `allocate` clauses, and `omp_alloc` API routines in `target` regions.

Implementers are allowed to include additional implementation-defined requirement clauses. All implementation defined requirements should begin with `ext_`. Requirement names that do not start with `ext_` are reserved.

The clauses of a `requires` directive are added to the `requires` trait in the OpenMP context for all program points that follow the directive.

**Restrictions**

The restrictions to the `requires` directive are as follows:

- Each of the clauses can appear at most once on the directive.
- All `atomic_default_mem_order` clauses that appear on a `requires` directive in the same compilation unit must specify the same parameter.
- A `requires` directive with a `unified_address`, `unified_shared_memory`, or `reverse_offload` clause must appear lexically before any device constructs or device routines.
- A `requires` directive may not appear lexically after a context selector in which any clause of the `requires` directive is used.
- A `requires` directive with any of the following clauses must appear in all `compilation units` of a program that contain device constructs or device routines or in none of them:
  - `reverse_offload`
  - `unified_address`
  - `unified_shared_memory`
• The `requires` directive with `atomic_default_mem_order` clause may not appear lexically after any `atomic` construct on which `memory-order-clause` is not specified.

• The `requires` directive may only appear at file scope.

• The `requires` directive may only appear at file or namespace scope.

• The `requires` directive must appear in the specification part of a `program unit`, after any `USE` statement, any `IMPORT` statement, and any `IMPLICIT` statement, unless the directive appears by referencing a module and each clause already appeared with the same parameters in the specification part of the `program unit`.

2.5.2 Assume Directive

Summary
The assume directive provides invariants to the implementation that may be used for optimization purposes. If the invariants do not hold at runtime, the behavior is unspecified. The assume directive is an informational directive.

Syntax
The syntax of the assume directive is as follows:

```
#pragma omp assumes clause[ [ , ] clause] ... ] new-line
```

or

```
#pragma omp begin assumes clause[ [ , ] clause] ... ] new-line
declaration-definition-seq
#pragma omp end assumes new-line
```

or

```
#pragma omp assume clause[ [ , ] clause] ... ] new-line
structured-block
```

where `clause` is either `assumption-clause` or a clause of the form `ext_implementation-defined-assumption` for an implementation-defined assumption clause, and where `assumption-clause` is one of the following:
The syntax of the assume directive is as follows:

```
!$omp assumes clause[[,] clause] ...
```

or

```
loosely-structured-block
!!$omp end assume
```

or

```
strictly-structured-block
[!!$omp end assume]
```

where `clause` is either `assumption-clause` or a clause of the form `ext_implementation-defined-assumption` for an implementation-defined assumption clause, where `assumption-clause` is one of the following:

- `absent(directive-name [[, directive-name]..])`
- `contains(directive-name [[, directive-name]..])`
- `holds(scalar-expression)`
- `no_openmp`
- `no_openmp_routines`
- `no_parallelism`
Description

The assume directive gives the implementation additional information about the expected properties of the program that can optionally be used to optimize the implementation. An implementation may ignore this information without altering the behavior of the program.

The scope of the assumes directive is the code executed and reached from the current compilation unit. The scope of the assume directive is the code executed in the corresponding region or in any region that is nested in the corresponding region.

The scope of the begin assumes directive is the code that is executed and reached from any of the declared functions in declaration-definition-seq.

The absent and contains clauses accept a list of one or more directive names that may match a construct that is encountered within the scope of the directive. An encountered construct matches the directive name if it has the same name as one of the specified directive names or if it is a combined or composite construct for which a constituent construct has the same name as one of the specified directive names.

When the absent clause appears on an assume directive, the application guarantees that no constructs that match a listed directive name are encountered in the scope of the assume directive.

When the contains clause appears on an assume directive, the application provides a hint that constructs that match the listed directive names are likely to be encountered in the scope of the assume directive.

When the holds clause appears on an assume directive, the application guarantees that the listed expression evaluates to true in the scope of the directive. The effect of the clause does not include an evaluation of the expression that is observable.

The no_openmp clause guarantees that no OpenMP related code is executed in the scope of the directive.

The no_openmp_routines clause guarantees that no explicit OpenMP runtime library calls are executed in the scope of the directive.

The no_parallelism clause guarantees that no OpenMP tasks (explicit or implicit) will be generated and that no SIMD constructs will be executed in the scope of the directive.

Implementers are allowed to include additional implementation-defined assumption clauses. All implementation-defined assumptions should begin with ext_. Assumption names that do not start with ext_ are reserved.
Restrictions

The restrictions to the assume directive are as follows:

- Each clause except the absent, contains or holds clause can appear at most once on the directive.
- Each directive-name listed in the clauses can appear at most once on the directive.
- A directive-name that appears in an absent or contains clause may not be a combined or composite directive name.
- A directive-name that appears in an absent or contains clause may not be a directive that is not associated with the execution of user or implementation code, i.e., a nothing directive, a declarative directive, a metadirective, or a loop transformation directive.
- The assumes directive may only appear at file scope.
- The assumes directive may only appear at file or namespace scope.
- The assumes directive may only appear in the specification part of a module or subprogram, after any USE statement, any IMPORT statement, and any IMPLICIT statement.

2.5.3 nothing Directive

Summary

The nothing directive has no effect. It can be used in a metadirective and other contexts to indicate explicitly that the intent is to have no effect on the execution of the OpenMP program. The nothing directive is a utility directive.

Syntax

The syntax of the nothing directive is as follows:

| #pragma omp nothing new-line |
|-------------------------------|---|
| !$omp nothing                 |

The syntax of the nothing directive is as follows:
Description
The `nothing` directive has no effect on the execution of the OpenMP program.

Cross References
- Metadirectives, see Section 2.3.4.

2.5.4 error Directive

Summary
The `error` directive instructs the compiler or runtime to display a message and to perform an error action. The `error` directive is a utility directive unless the `at` clause is specified with the `execution` value, in which case the `error` directive is a stand-alone directive.

Syntax

The syntax of the `error` directive is as follows:

```c
#pragma omp error [clause [, clause] ... ]
```

where `clause` is one of the following:

- `at(compilation | execution)`
- `severity(fatal | warning)`
- `message(msg-string)`

where `msg-string` is a string of const char * type.

The syntax of the `error` directive as follows:

```c
!$omp error [clause [, clause] ... ]
```

where `clause` is one of the following:

- `at(compilation | execution)`
- `severity(fatal | warning)`
- `message(msg-string)`

where `msg-string` is a character string of character(len=*) type
Description
The error directive performs an error action. The error action includes the display of an implementation defined message. The severity clause determines whether the error action includes anything other than the message display.

The at clause determines when the implementation performs the error action. When the at clause specifies compilation, the error action is performed during compilation if the error directive appears in a declarative context or in an executable context that is reachable at runtime. When the at clause specifies compilation and the error directive appears in an executable context that is not reachable at runtime, the error action may or may not be performed. When the at clause specifies execution, the error action is performed during program execution when a thread encounters the directive. If the at clause is not specified then the error directive behaves as if the at clause specifies compilation.

The severity clause determines the action that the implementation performs. When the severity clause specifies warning, the implementation takes no action besides displaying the message. When the severity clause specifies fatal and the at clause specifies compile then the message is displayed and compilation of the current compilation unit is aborted. When the severity clause specifies fatal and the at clause specifies execution then the message is displayed and program execution is aborted. If no severity clause is specified then the error directive behaves as if the severity clause specifies fatal.

If the message clause is specified then msg-string is included in the implementation defined message.

Execution Model Events
The runtime-error event occurs when a thread encounters an error directive for which the at clause specifies execution.

Tool Callbacks
A thread dispatches a registered ompt_callback_error callback for each occurrence of a runtime-error event in the context of the encountering task. This callback has the type signature ompt_callback_error_t.

Restrictions
Restrictions to the error directive are as follows:

- At most one at clause can appear on the directive.
- At most one severity clause can appear on the directive.
- At most one message clause can appear on the directive.

Cross References
- Stand-alone directives, see Section 2.1.3.
- ompt_callback_error_t, see Section 4.5.2.30.
2.6 parallel Construct

Summary
The **parallel** construct creates a team of OpenMP threads that execute the region.

Syntax

C / C++
The syntax of the **parallel** construct is as follows:

```c
#pragma omp parallel [clause[ , ] clause] ... ] new-line
structured-block
```

where **clause** is one of the following:

- **if**(*parallel : scalar-expression*)
- **num_threads**(*integer-expression*)
- **default**(*data-sharing-attribute*)
- **private**(*list*)
- **firstprivate**(*list*)
- **shared**(*list*)
- **copyin**(*list*)
- **reduction**(*reduction-modifier , reduction-identifier : list*)
- **proc_bind**(*affinity-policy*)
- **allocate**(*allocator : list*)

where **affinity-policy** is one of the following:

- **primary**
- **master** [deprecated]
- **close**
- **spread**

Fortran
The syntax of the **parallel** construct is as follows:

```fortran
!$omp parallel [clause[ , ] clause] ... ]
loosely-structured-block
!$omp end parallel
```
or

```fortran
!$omp parallel [clause[ [, ] clause] ... ]
strictly-structured-block
[!$omp end parallel]
```

where `clause` is one of the following:

```fortran
if(/ parallel : ] scalar-logical-expression)
num_threads(scalar-integer-expression)
default(data-sharing-attribute)
private(list)
firstprivate(list)
shared(list)
copyin(list)
reduction([reduction-modifier , ] reduction-identifier : list)
proc_bind(affinity-policy)
allocate([allocator : ] list)
```

where `affinity-policy` is one of the following:

```fortran
primary
master [deprecated]
close
spread
```

### Binding

The binding thread set for a `parallel` region is the encountering thread. The encountering thread becomes the primary thread of the new team.

### Description

When a thread encounters a `parallel` construct, a team of threads is created to execute the `parallel` region (see Section 2.6.1 for more information about how the number of threads in the team is determined, including the evaluation of the `if` and `num_threads` clauses). The thread that encountered the `parallel` construct becomes the primary thread of the new team, with a thread number of zero for the duration of the new `parallel` region. All threads in the new team, including the primary thread, execute the region. Once the team is created, the number of threads in the team remains constant for the duration of that `parallel` region.
The optional \texttt{proc_bind} clause, described in Section 2.6.2, specifies the mapping of OpenMP threads to places within the current place partition, that is, within the places listed in the \texttt{place-partition-var} ICV for the implicit task of the encountering thread.

Within a \texttt{parallel} region, thread numbers uniquely identify each thread. Thread numbers are consecutive whole numbers ranging from zero for the primary thread up to one less than the number of threads in the team. A thread may obtain its own thread number by a call to the \texttt{omp_get_thread_num} library routine.

A set of implicit tasks, equal in number to the number of threads in the team, is generated by the encountering thread. The structured block of the \texttt{parallel} construct determines the code that will be executed in each implicit task. Each task is assigned to a different thread in the team and becomes tied. The task region of the task that the encountering thread is executing is suspended and each thread in the team executes its implicit task. Each thread can execute a path of statements that is different from that of the other threads.

The implementation may cause any thread to suspend execution of its implicit task at a task scheduling point, and to switch to execution of any explicit task generated by any of the threads in the team, before eventually resuming execution of the implicit task (for more details see Section 2.12).

An implicit barrier occurs at the end of a \texttt{parallel} region. After the end of a \texttt{parallel} region, only the primary thread of the team resumes execution of the enclosing task region.

If a thread in a team that is executing a \texttt{parallel} region encounters another \texttt{parallel} directive, it creates a new team, according to the rules in Section 2.6.1, and it becomes the primary thread of that new team.

If execution of a thread terminates while inside a \texttt{parallel} region, execution of all threads in all teams terminates. The order of termination of threads is unspecified. All work done by a team prior to any barrier that the team has passed in the program is guaranteed to be complete. The amount of work done by each thread after the last barrier that it passed and before it terminates is unspecified.

**Execution Model Events**

The \texttt{parallel-begin} event occurs in a thread that encounters a \texttt{parallel} construct before any implicit task is created for the corresponding \texttt{parallel} region.

Upon creation of each implicit task, an \texttt{implicit-task-begin} event occurs in the thread that executes the implicit task after the implicit task is fully initialized but before the thread begins to execute the structured block of the \texttt{parallel} construct.

If the \texttt{parallel} region creates a native thread, a \texttt{native-thread-begin} event occurs as the first event in the context of the new thread prior to the \texttt{implicit-task-begin} event.

Events associated with implicit barriers occur at the end of a \texttt{parallel} region. Section 2.19.3 describes events associated with implicit barriers.

When a thread finishes an implicit task, an \texttt{implicit-task-end} event occurs in the thread after events associated with implicit barrier synchronization in the implicit task.
The parallel-end event occurs in the thread that encounters the parallel construct after the thread executes its implicit-task-end event but before the thread resumes execution of the encountering task.

If a native thread is destroyed at the end of a parallel region, a native-thread-end event occurs in the thread as the last event prior to destruction of the thread.

**Tool Callbacks**

A thread dispatches a registered ompt_callback_parallel_begin callback for each occurrence of a parallel-begin event in that thread. The callback occurs in the task that encounters the parallel construct. This callback has the type signature ompt_callback_parallel_begin_t. In the dispatched callback, (flags & ompt_parallel_team) evaluates to true.

A thread dispatches a registered ompt_callback_implicit_task callback with ompt_scope_begin as its endpoint argument for each occurrence of an implicit-task-begin event in that thread. Similarly, a thread dispatches a registered ompt_callback_implicit_task callback with ompt_scope_end as its endpoint argument for each occurrence of an implicit-task-end event in that thread. The callbacks occur in the context of the implicit task and have type signature ompt_callback_implicit_task_t.

In the dispatched callback, (flags & ompt_task_implicit) evaluates to true.

A thread dispatches a registered ompt_callback_parallel_end callback for each occurrence of a parallel-end event in that thread. The callback occurs in the task that encounters the parallel construct. This callback has the type signature ompt_callback_parallel_end_t.

Restrictions

Restrictions to the parallel construct are as follows:

- A program must not depend on any ordering of the evaluations of the clauses of the parallel directive, or on any side effects of the evaluations of the clauses.
- At most one if clause can appear on the directive.
- At most one proc_bind clause can appear on the directive.
- At most one num_threads clause can appear on the directive. The num_threads expression must evaluate to a positive integer value.
C++

- A `throw` executed inside a `parallel` region must cause execution to resume within the same `parallel` region, and the same thread that threw the exception must catch it.

Cross References
- OpenMP execution model, see Section 1.3.
- `num_threads` clause, see Section 2.6.
- `proc_bind` clause, see Section 2.6.2.
- `allocate` clause, see Section 2.13.4.
- `if` clause, see Section 2.18.
- `default`, `shared`, `private`, `firstprivate`, and `reduction` clauses, see Section 2.21.4.
- `copyin` clause, see Section 2.21.6.
- `omp_get_thread_num` routine, see Section 3.2.4.
- `ompt_scope_begin` and `ompt_scope_end`, see Section 4.4.4.11.
- `ompt_callback_thread_begin_t`, see Section 4.5.2.1.
- `ompt_callback_thread_end_t`, see Section 4.5.2.2.
- `ompt_callback_parallel_begin_t`, see Section 4.5.2.3.
- `ompt_callback_parallel_end_t`, see Section 4.5.2.4.
- `ompt_callback_implicit_task_t`, see Section 4.5.2.11.

2.6.1 Determining the Number of Threads for a `parallel` Region

When execution encounters a `parallel` directive, the value of the `if` clause or `num_threads` clause (if any) on the directive, the current parallel context, and the values of the `nthreads-var`, `dyn-var`, `thread-limit-var`, and `max-active-levels-var` ICVs are used to determine the number of threads to use in the region.

Using a variable in an `if` or `num_threads` clause expression of a `parallel` construct causes an implicit reference to the variable in all enclosing constructs. The `if` clause expression and the `num_threads` clause expression are evaluated in the context outside of the `parallel` construct, and no ordering of those evaluations is specified. In what order or how many times any side effects of the evaluation of the `num_threads` or `if` clause expressions occur is also unspecified.
When a thread encounters a `parallel` construct, the number of threads is determined according to Algorithm 2.1.

**Algorithm 2.1**

```
let ThreadsBusy be the number of OpenMP threads currently executing in this contention group;
if an if clause exists
then let IfClauseValue be the value of the if clause expression;
else let IfClauseValue = true;
if a num_threads clause exists
then let ThreadsRequested be the value of the num_threads clause expression;
else let ThreadsRequested = value of the first element of nthreads-var;
let ThreadsAvailable = (thread-limit-var - ThreadsBusy + 1);
if (IfClauseValue = false)
then number of threads = 1;
else if (active-levels-var ≥ max-active-levels-var)
then number of threads = 1;
else if (dyn-var = true) and (ThreadsRequested ≤ ThreadsAvailable)
then 1 ≤ number of threads ≤ ThreadsRequested;
else if (dyn-var = true) and (ThreadsRequested > ThreadsAvailable)
then 1 ≤ number of threads ≤ ThreadsAvailable;
else if (dyn-var = false) and (ThreadsRequested ≤ ThreadsAvailable)
then number of threads = ThreadsRequested;
else if (dyn-var = false) and (ThreadsRequested > ThreadsAvailable)
then behavior is implementation defined;
```
Note – Since the initial value of the dyn-var ICV is implementation defined, programs that depend on a specific number of threads for correct execution should explicitly disable dynamic adjustment of the number of threads.

Cross References
- nthreads-var, dyn-var, thread-limit-var, and max-active-levels-var ICVs, see Section 2.4.
- parallel construct, see Section 2.6.
- num_threads clause, see Section 2.6.
- if clause, see Section 2.18.

2.6.2 Controlling OpenMP Thread Affinity

When a thread encounters a parallel directive without a proc_bind clause, the bind-var ICV is used to determine the policy for assigning OpenMP threads to places within the current place partition, that is, within the places listed in the place-partition-var ICV for the implicit task of the encountering thread. If the parallel directive has a proc_bind clause then the binding policy specified by the proc_bind clause overrides the policy specified by the first element of the bind-var ICV. Once a thread in the team is assigned to a place, the OpenMP implementation should not move it to another place.

The primary thread affinity policy instructs the execution environment to assign every thread in the team to the same place as the primary thread. The place partition is not changed by this policy, and each implicit task inherits the place-partition-var ICV of the parent implicit task. The master thread-affinity policy, which has been deprecated, has identical semantics to the primary thread affinity policy.

The close thread affinity policy instructs the execution environment to assign the threads in the team to places close to the place of the parent thread. The place partition is not changed by this policy, and each implicit task inherits the place-partition-var ICV of the parent implicit task. If $T$ is the number of threads in the team, and $P$ is the number of places in the parent’s place partition, then the assignment of threads in the team to places is as follows:

- $T \leq P$: The primary thread executes on the place of the parent thread. The thread with the next smallest thread number executes on the next place in the place partition, and so on, with wrap around with respect to the place partition of the primary thread.

- $T > P$: Each place $p$ will contain $S_p$ threads with consecutive thread numbers where $\lfloor T/P \rfloor \leq S_p \leq \lceil T/P \rceil$. The first $S_0$ threads (including the primary thread) are assigned to the place of the parent thread. The next $S_1$ threads are assigned to the next place in the place partition, and so on, with wrap around with respect to the place partition of the primary thread.
When $P$ does not divide $T$ evenly, the exact number of threads in a particular place is implementation defined.

The purpose of the \texttt{spread} thread affinity policy is to create a sparse distribution for a team of $T$ threads among the $P$ places of the parent's place partition. A sparse distribution is achieved by first subdividing the parent partition into $T$ subpartitions if $T \leq P$, or $P$ subpartitions if $T > P$. Then one thread ($T \leq P$) or a set of threads ($T > P$) is assigned to each subpartition. The \texttt{place-partition-var} ICV of each implicit task is set to its subpartition. The subpartitioning is not only a mechanism for achieving a sparse distribution, it also defines a subset of places for a thread to use when creating a nested \texttt{parallel} region. The assignment of threads to places is as follows:

- $T \leq P$: The parent thread's place partition is split into $T$ subpartitions, where each subpartition contains $\lfloor P/T \rfloor$ or $\lceil P/T \rceil$ consecutive places. A single thread is assigned to each subpartition. The primary thread executes on the place of the parent thread and is assigned to the subpartition that includes that place. The thread with the next smallest thread number is assigned to the first place in the next subpartition, and so on, with wrap around with respect to the original place partition of the primary thread.

- $T > P$: The parent thread's place partition is split into $P$ subpartitions, each consisting of a single place. Each subpartition is assigned $S_p$ threads with consecutive thread numbers, where $\lfloor T/P \rfloor \leq S_p \leq \lceil T/P \rceil$. The first $S_0$ threads (including the primary thread) are assigned to the subpartition that contains the place of the parent thread. The next $S_1$ threads are assigned to the next subpartition, and so on, with wrap around with respect to the original place partition of the primary thread. When $P$ does not divide $T$ evenly, the exact number of threads in a particular subpartition is implementation defined.

The determination of whether the affinity request can be fulfilled is implementation defined. If the affinity request cannot be fulfilled, then the affinity of threads in the team is implementation defined.

\begin{flushleft}
\textbf{Note} – Wrap around is needed if the end of a place partition is reached before all thread assignments are done. For example, wrap around may be needed in the case of \texttt{close} and $T \leq P$, if the primary thread is assigned to a place other than the first place in the place partition. In this case, thread 1 is assigned to the place after the place of the primary thread, thread 2 is assigned to the place after that, and so on. The end of the place partition may be reached before all threads are assigned. In this case, assignment of threads is resumed with the first place in the place partition.
\end{flushleft}
2.7 teams Construct

Summary

The `teams` construct creates a league of initial teams and the initial thread in each team executes the region.

Syntax

```C / C++
#pragma omp teams [clause[ , ] clause] ... ]
structured-block
```

where `clause` is one of the following:

- `num_teams([lower-bound: ]upper-bound)`
- `thread_limit(integer-expression)`
- `default(data-sharing-attribute)`
- `private(list)`
- `firstprivate(list)`
- `shared(list)`
- `reduction([default , ] reduction-identifier : list)`
- `allocate([allocator : ] list)`

and where `lower-bound` and `upper-bound` are scalar integer expressions.

```C / C++
!$omp teams [clause[ , ] clause] ... ]
loosely-structured-block
```

or

```C / C++
!$omp teams [clause[ , ] clause] ... ]
strictly-structured-block
```

```C / C++
/$omp end teams/
```
where \textit{clause} is one of the following:

\begin{verbatim}
num_teams([lower-bound:]upper-bound)
thread_limit(scalar-integer-expression)
default(data-sharing-attribute)
private(list)
firstprivate(list)
shared(list)
reduction([default],[reduction-identifier:]list)
allocation([allocator:]list)
\end{verbatim}

and where \textit{lower-bound} and \textit{upper-bound} are scalar integer expressions.

\section*{Binding}

The binding thread set for a \texttt{teams} region is the encountering thread.

\section*{Description}

When a thread encounters a \texttt{teams} construct, a league of teams is created. Each team is an initial team, and the initial thread in each team executes the \texttt{teams} region.

If the \texttt{num_teams} clause is present, \textit{lower-bound} is the specified lower bound and \textit{upper-bound} is the specified upper bound on the number of teams requested. If a lower bound is not specified, the lower bound is set to the specified upper bound. The number of teams created is implementation defined, but it will be greater than or equal to the lower bound and less than or equal to the upper bound.

If the \texttt{num_teams} clause is not specified and the value of the \texttt{nteams-var ICV} is greater than zero, the number of teams created is less or equal to the value of the \texttt{nteams-var ICV}. Otherwise, the number of teams created is implementation defined, but it will be greater than or equal to 1.

A thread may obtain the number of teams created by the construct with a call to the \texttt{omp_get_num_teams} routine.

If a \texttt{thread_limit} clause is not present on the \texttt{teams} construct, but the construct is closely nested inside a \texttt{target} construct on which the \texttt{thread_limit} clause is specified, the behavior is as if that \texttt{thread_limit} clause is also specified for the \texttt{teams} construct.

As described in Section 2.4.4.1, the \texttt{teams} construct limits the number of threads that may participate in a contention group initiated by each team by setting the value of the \texttt{thread-limit-var ICV} for the initial task to an implementation defined value greater than zero. If the \texttt{thread_limit} clause is specified, the number of threads will be less than or equal to the value specified in the clause. Otherwise, if the \texttt{teams-thread-limit-var ICV} is greater than zero, the number of threads will be less than or equal to that value.
On a combined or composite construct that includes `target` and `teams` constructs, the expressions in `num_teams` and `thread_limit` clauses are evaluated on the host device on entry to the `target` construct.

Once the teams are created, the number of initial teams remains constant for the duration of the `teams` region.

Within a `teams` region, initial team numbers uniquely identify each initial team. Initial team numbers are consecutive whole numbers ranging from zero to one less than the number of initial teams. A thread may obtain its own initial team number by a call to the `omp_get_team_num` library routine.

The place list, given by the `place-partition-var` ICV of the encountering thread, is split into subpartitions in an implementation-defined manner, and each team is assigned to a subpartition by setting the `place-partition-var` of its initial thread to the subpartition.

The `teams` construct sets the `default-device-var` ICV for each initial thread to an implementation-defined value.

After the teams have completed execution of the `teams` region, the encountering task resumes execution of the enclosing task region.

**Execution Model Events**

The `teams-begin` event occurs in a thread that encounters a `teams` construct before any initial task is created for the corresponding `teams` region.

Upon creation of each initial task, an `initial-task-begin` event occurs in the thread that executes the initial task after the initial task is fully initialized but before the thread begins to execute the structured block of the `teams` construct.

If the `teams` region creates a native thread, a `native-thread-begin` event occurs as the first event in the context of the new thread prior to the `initial-task-begin` event.

When a thread finishes an initial task, an `initial-task-end` event occurs in the thread.

The `teams-end` event occurs in the thread that encounters the `teams` construct after the thread executes its `initial-task-end` event but before it resumes execution of the encountering task.

If a native thread is destroyed at the end of a `teams` region, a `native-thread-end` event occurs in the thread as the last event prior to destruction of the thread.

**Tool Callbacks**

A thread dispatches a registered `omp_callback_parallel_begin` callback for each occurrence of a `teams-begin` event in that thread. The callback occurs in the task that encounters the `teams` construct. This callback has the type signature `omp_callback_parallel_begin_t`. In the dispatched callback, `(flags & omp_parallel_league)` evaluates to true.
A thread dispatches a registered `ompt_callback_implicit_task` callback with `ompt_scope_begin` as its *endpoint* argument for each occurrence of an *initial-task-begin* in that thread. Similarly, a thread dispatches a registered `ompt_callback_implicit_task` callback with `ompt_scope_end` as its *endpoint* argument for each occurrence of an *initial-task-end* event in that thread. The callbacks occur in the context of the initial task and have type signature `ompt_callback_implicit_task_t`. In the dispatched callback, `(flags & ompt_task_initial)` evaluates to `true`.

A thread dispatches a registered `ompt_callback_parallel_end` callback for each occurrence of a *teams-end* event in that thread. The callback occurs in the task that encounters the *teams* construct. This callback has the type signature `ompt_callback_parallel_end_t`.

A thread dispatches a registered `ompt_callback_thread_begin` callback for the *native-thread-begin* event in that thread. The callback occurs in the context of the thread. The callback has type signature `ompt_callback_thread_begin_t`.

A thread dispatches a registered `ompt_callback_thread_end` callback for the *native-thread-end* event in that thread. The callback occurs in the context of the thread. The callback has type signature `ompt_callback_thread_end_t`.

**Restrictions**

Restrictions to the *teams* construct are as follows:

- A program that branches into or out of a *teams* region is non-conforming.
- A program must not depend on any ordering of the evaluations of the clauses of the *teams* directive, or on any side effects of the evaluation of the clauses.
- At most one *thread_limit* clause can appear on the directive. The *thread_limit* expression must evaluate to a positive integer value.
- At most one *num_teams* clause can appear on the directive. The *lower-bound* and *upper-bound* specified in the *num_teams* clause must evaluate to positive integer values.
- A *teams* region must be strictly nested within the implicit parallel region that surrounds the whole OpenMP program or a *target* region. If a *teams* region is nested inside a *target* region, the corresponding *target* construct must not contain any statements, declarations or directives outside of the corresponding *teams* construct.
- *distribute*, *distribute simd*, distribute parallel worksharing-loop, distribute parallel worksharing-loop SIMD, *parallel* regions, including any *parallel* regions arising from combined constructs, `omp_get_num_teams()` regions, and `omp_get_team_num()` regions are the only OpenMP regions that may be strictly nested inside the *teams* region.
Cross References

- **parallel** construct, see Section 2.6.
- **distribute** construct, see Section 2.11.6.1.
- **distribute simd** construct, see Section 2.11.6.2.
- **allocate** clause, see Section 2.13.4.
- **target** construct, see Section 2.14.5.
- Data-sharing attribute clauses, see Section 2.21.4.
- **omp_get_num_teams** routine, see Section 3.4.1.
- **omp_get_team_num** routine, see Section 3.4.2.
- ompt_callback_thread_begin_t, see Section 4.5.2.1.
- ompt_callback_thread_end_t, see Section 4.5.2.2.
- ompt_callback_parallel_begin_t, see Section 4.5.2.3.
- ompt_callback_parallel_end_t, see Section 4.5.2.4.
- ompt_callback_implicit_task_t, see Section 4.5.2.11.

2.8 masked Construct

Summary

The **masked** construct specifies a structured block that is executed by a subset of the threads of the current team.

Syntax

The syntax of the **masked** construct is as follows:

```
C / C++

#pragma omp masked [ filter(integer-expression) ]\new-line
structured-block
```

```
C / C++

Fortran
```

The syntax of the **masked** construct is as follows:

```
!$omp masked [ filter(scalar-integer-expression) ]
loosely-structured-block
```

```
!$omp end masked
```
or

```fortran
!$omp masked [ filter(scalar-integer-expression) ]
strictly-structured-block
[/!$omp end masked]
```

The **master** construct, which has been deprecated, has the same syntax as the **masked** construct other than the use of **master** as the directive name and that the **filter** clause may not be specified for the **master** construct.

**Binding**

The binding thread set for a **masked** region is the current team. A **masked** region binds to the innermost enclosing parallel region.

**Description**

Only the threads of the team that executes the binding parallel region that the **filter** clause selects participate in the execution of the structured block of a **masked** region. Other threads in the team do not execute the associated structured block. No implied barrier occurs either on entry to, or exit from, the **masked** construct.

If a **filter** clause is present on the construct and the parameter specifies the thread number of the current thread in the current team then the current thread executes the associated structured block. If the **filter** clause is not present, the construct behaves as if the parameter is a constant integer expression that evaluates to zero, so that only the primary thread executes the associated structured block. The use of a variable in a **filter** clause expression causes an implicit reference to the variable in all enclosing constructs. The result of evaluating the parameter of the **filter** clause may vary across threads.

If more than one thread in the team executes the structured block of a **masked** region, the structured block must include any synchronization required to ensure that data races do not occur.

The **master** construct, which has been deprecated, has identical semantics to the **masked** construct with no **filter** clause present.

**Execution Model Events**

The **masked-begin** event occurs in any thread of a team that executes the **masked** region on entry to the region.

The **masked-end** event occurs in any thread of a team that executes the **masked** region on exit from the region.
Tool Callbacks
A thread dispatches a registered `ompt_callback_masked` callback with
`ompt_scope_begin` as its `endpoint` argument for each occurrence of a `masked-begin` event in
that thread. Similarly, a thread dispatches a registered `ompt_callback_masked` callback with
`ompt_scope_end` as its `endpoint` argument for each occurrence of a `masked-end` event in that
thread. These callbacks occur in the context of the task executed by the current thread and have the
type signature `ompt_callback_masked_t`.

Restrictions
Restrictions to the `masked` construct are as follows:

- A throw executed inside a `masked` region must cause execution to resume within the same
  `masked` region, and the same thread that threw the exception must catch it.

Cross References
- `parallel` construct, see Section 2.6.
- `ompt_scope_begin` and `ompt_scope_end`, see Section 4.4.11.
- `ompt_callback_masked_t`, see Section 4.5.2.12.

2.9 scope Construct

Summary
The `scope` construct defines a structured block that is executed by all threads in a team but where
additional OpenMP operations can be specified.

Syntax
The syntax of the scope construct is as follows:

```
#pragma omp scope [clause[ [, ] clause] ... ] new-line
structured-block
```

where `clause` is one of the following:

- `private(list)`
- `reduction([reduction-modifier , ] reduction-identifier : list)`
- `nowait`
The syntax of the `scope` construct is as follows:

```fortran
!$omp scope [clause [ , ] clause] ... ]
loosely-structured-block
|$omp end scope /nowait/
```

or

```fortran
!$omp scope [clause [ , ] clause] ... ]
strictly-structured-block
|$omp end scope /nowait]/
```

where `clause` is one of the following:

```fortran
private(list)
reduction([reduction-modifier , ] reduction-identifier : list)
```

### Binding

The binding thread set for a `scope` region is the current team. A `scope` region binds to the innermost enclosing parallel region. Only the threads of the team that executes the binding parallel region participate in the execution of the structured block and the implied barrier of the `scope` region if the barrier is not eliminated by a `nowait` clause.

### Description

All encountering threads will execute the structured block associated with the `scope` construct. An implicit barrier occurs at the end of a `scope` region if the `nowait` clause is not specified.

### Execution Model Events

The `scope-begin` event occurs after an implicit task encounters a `scope` construct but before the task starts to execute the structured block of the `scope` region.

The `scope-end` event occurs after an implicit task finishes execution of a `scope` region but before it resumes execution of the enclosing region.

### Tool Callbacks

A thread dispatches a registered `ompt_callback_work` callback with `ompt_scope_begin` as its `endpoint` argument and `ompt_work_scope` as its `wstype` argument for each occurrence of a `scope-begin` event in that thread. Similarly, a thread dispatches a registered `ompt_callback_work` callback with `ompt_scope_end` as its `endpoint` argument and `ompt_work_scope` as its `wstype` argument for each occurrence of a `scope-end` event in that thread. The callbacks occur in the context of the implicit task. The callbacks have type signature `ompt_callback_work_t`. 
Restrictions
Restrictions to the `scope` construct are as follows:

- Each `scope` region must be encountered by all threads in a team or by none at all, unless cancellation has been requested for the innermost enclosing parallel region.
- The sequence of worksharing regions, `scope` regions and `barrier` regions encountered must be the same for every thread in a team.
- At most one `nowait` clause can appear on a `scope` construct.

\[ C++ \]

- A throw executed inside a `scope` region must cause execution to resume within the same `scope` region, and the same thread that threw the exception must catch it.

Cross References
- `reduction` clause, see Section 2.21.4.
- `ompt_scope_begin` and `ompt_scope_end`, see Section 4.4.4.11.
- `ompt_work_scope`, see Section 4.4.4.15.
- `ompt_callback_work_t`, see Section 4.5.2.5.

2.10 Worksharing Constructs

A worksharing construct distributes the execution of the corresponding region among the members of the team that encounters it. Threads execute portions of the region in the context of the implicit tasks that each one is executing. If the team consists of only one thread then the worksharing region is not executed in parallel.

A worksharing region has no barrier on entry; however, an implied barrier exists at the end of the worksharing region, unless a `nowait` clause is specified. If a `nowait` clause is present, an implementation may omit the barrier at the end of the worksharing region. In this case, threads that finish early may proceed straight to the instructions that follow the worksharing region without waiting for the other members of the team to finish the worksharing region, and without performing a flush operation.

The OpenMP API defines the worksharing constructs that are described in this section as well as the worksharing-loop construct, which is described in Section 2.11.4.
Restrictions

The following restrictions apply to worksharing constructs:

- Each worksharing region must be encountered by all threads in a team or by none at all, unless
cancellation has been requested for the innermost enclosing parallel region.

- The sequence of worksharing regions, scope regions and barrier regions encountered must
be the same for every thread in a team.

2.10.1 sections Construct

Summary

The sections construct is a non-iterative worksharing construct that contains a set of structured
blocks that are to be distributed among and executed by the threads in a team. Each structured
block is executed once by one of the threads in the team in the context of its implicit task.

Syntax

```c
#pragma omp sections [clause[ , ] clause] ... ] new-line
{
    [#pragma omp section new-line]
    structured-block-sequence
    [#pragma omp section new-line
     structured-block-sequence]
    ...
}
```

where clause is one of the following:

```c
private(list)
firstprivate(list)
lastprivate([lastprivate-modifier:] list)
reduction([reduction-modifier , ] reduction-identifier : list)
allocate([allocator :] list)
nowait
```
The syntax of the `sections` construct is as follows:

```
!$omp sections [clause[ [, ] clause] ... ]
  !$omp section
  structured-block-sequence
  !$omp section
  structured-block-sequence

... 

!$omp end sections [nowait]
```

where `clause` is one of the following:

```
private(list)
firstprivate(list)
lastprivate([ lastprivate-modifier: ] list)
reduction([reduction-modifier , ] reduction-identifier : list)
allocate([allocator : ] list)
```

**Binding**

The binding thread set for a `sections` region is the current team. A `sections` region binds to the innermost enclosing `parallel` region. Only the threads of the team that executes the binding `parallel` region participate in the execution of the structured block sequences and the implied barrier of the `sections` region if the barrier is not eliminated by a `nowait` clause.

**Description**

Each structured block sequence in the `sections` construct is preceded by a `section` directive except possibly the first sequence, for which a preceding `section` directive is optional.

The method of scheduling the structured block sequences among the threads in the team is implementation defined.

An implicit barrier occurs at the end of a `sections` region if the `nowait` clause is not specified.

**Execution Model Events**

The `section-begin` event occurs after an implicit task encounters a `sections` construct but before the task executes any structured block sequences of the `sections` region.

The `sections-end` event occurs after an implicit task finishes execution of a `sections` region but before it resumes execution of the enclosing context.

The `section-begin` event occurs before an implicit task starts to execute a structured block sequence in the `sections` construct for each of those structured block sequences that the task executes.
Tool Callbacks

A thread dispatches a registered `ompt_callback_work` callback with `ompt_scope_begin` as its `endpoint` argument and `ompt_work_sections` as its `wstype` argument for each occurrence of a `section-begin` event in that thread. Similarly, a thread dispatches a registered `ompt_callback_work` callback with `ompt_scope_end` as its `endpoint` argument and `ompt_work_sections` as its `wstype` argument for each occurrence of a `sections-end` event in that thread. The callbacks occur in the context of the implicit task. The callbacks have type signature `ompt_callback_work_t`.

A thread dispatches a registered `ompt_callback_dispatch` callback for each occurrence of a `section-begin` event in that thread. The callback occurs in the context of the implicit task. The callback has type signature `ompt_callback_dispatch_t`.

Restrictions

Restrictions to the `sections` construct are as follows:

- Orphaned `section` directives are prohibited. That is, the `section` directives must appear within the `sections` construct and must not be encountered elsewhere in the `sections` region.

- The code enclosed in a `sections` construct must be a structured block sequence.

- Only a single `nowait` clause can appear on a `sections` directive.

- A throw executed inside a `sections` region must cause execution to resume within the same section of the `sections` region, and the same thread that threw the exception must catch it.

Cross References

- `allocate` clause, see Section 2.13.4.

- `private`, `firstprivate`, `lastprivate`, and `reduction` clauses, see Section 2.21.4.

- `ompt_scope_begin` and `ompt_scope_end`, see Section 4.4.4.11.

- `ompt_work_sections`, see Section 4.4.4.15.

- `ompt_callback_work_t`, see Section 4.5.2.5.

- `ompt_callback_dispatch_t`, see Section 4.5.2.6.
2.10.2 single Construct

Summary
The single construct specifies that the associated structured block is executed by only one of the threads in the team (not necessarily the primary thread), in the context of its implicit task. The other threads in the team, which do not execute the block, wait at an implicit barrier at the end of a single region unless a nowait clause is specified.

Syntax

C / C++

The syntax of the single construct is as follows:

```c
#pragma omp single [clause[ , ] clause] ... ] new-line
  structured-block
```

where clause is one of the following:

- `private(list)`
- `firstprivate(list)`
- `copyprivate(list)`
- `allocate([allocator :] list)`
- `nowait`

Fortran

The syntax of the single construct is as follows:

```fortran
!$omp single [clause[ , ] clause] ... ]
  loosely-structured-block

!$omp end single [end_clause[ , ] end_clause] ... ]
```

or

```fortran
!$omp single [clause[ , ] clause] ... ]
  strictly-structured-block

[!$omp end single [end_clause[ , ] end_clause] ... ]]
```

where clause is one of the following:

- `private(list)`
- `firstprivate(list)`
- `allocate([allocator :] list)`
and \textit{end\_clause} is one of the following:

\begin{itemize}
  \item \texttt{copyprivate(list)}
  \item \texttt{nowait}
\end{itemize}

\section*{Fortran}

\textbf{Binding}

The binding thread set for a \texttt{single} region is the current team. A \texttt{single} region binds to the innermost enclosing \texttt{parallel} region. Only the threads of the team that executes the binding \texttt{parallel} region participate in the execution of the structured block and the implied barrier of the \texttt{single} region if the barrier is not eliminated by a \texttt{nowait} clause.

\section*{Description}

Only one of the encountering threads will execute the structured block associated with the \texttt{single} construct. The method of choosing a thread to execute the structured block each time the team encounters the construct is implementation defined. An implicit barrier occurs at the end of a \texttt{single} region if the \texttt{nowait} clause is not specified.

\section*{Execution Model Events}

The \texttt{single-begin} event occurs after an implicit task encounters a \texttt{single} construct but before the task starts to execute the structured block of the \texttt{single} region.

The \texttt{single-end} event occurs after an implicit task finishes execution of a \texttt{single} region but before it resumes execution of the enclosing region.

\section*{Tool Callbacks}

A thread dispatches a registered \texttt{ompt\_callback\_work} callback with \texttt{ompt\_scope\_begin} as its \textit{endpoint} argument for each occurrence of a \texttt{single-begin} event in that thread. Similarly, a thread dispatches a registered \texttt{ompt\_callback\_work} callback with \texttt{ompt\_scope\_end} as its \textit{endpoint} argument for each occurrence of a \texttt{single-end} event in that thread. For each of these callbacks, the \textit{wstype} argument is \texttt{ompt\_work\_single\_executor} if the thread executes the structured block associated with the \texttt{single} region; otherwise, the \textit{wstype} argument is \texttt{ompt\_work\_single\_other}. The callback has type signature \texttt{ompt\_callback\_work\_t}.

\section*{Restrictions}

Restrictions to the \texttt{single} construct are as follows:

\begin{itemize}
  \item The \texttt{copyprivate} clause must not be used with the \texttt{nowait} clause.
  \item At most one \texttt{nowait} clause can appear on a \texttt{single} construct.
  \item A throw executed inside a \texttt{single} region must cause execution to resume within the same \texttt{single} region, and the same thread that threw the exception must catch it.
\end{itemize}
Cross References

• allocate clause, see Section 2.13.4.
• private and firstprivate clauses, see Section 2.21.4.
• copyprivate clause, see Section 2.21.6.2.
• ompt_scope_begin and ompt_scope_end, see Section 4.4.4.11.
• ompt_work_single_executor and ompt_work_single_other, see Section 4.4.4.15.
• ompt_callback_work_t, Section 4.5.2.5.

2.10.3 workshare Construct

Summary
The workshare construct divides the execution of the enclosed structured block into separate units of work, and causes the threads of the team to share the work such that each unit is executed only once by one thread, in the context of its implicit task.

Syntax
The syntax of the workshare construct is as follows:

```
!$omp workshare
  loosely-structured-block
!$omp end workshare /nowait/
```
or

```
!$omp workshare
  strictly-structured-block
[/(!$omp end workshare /nowait/]
```

Binding
The binding thread set for a workshare region is the current team. A workshare region binds to the innermost enclosing parallel region. Only the threads of the team that executes the binding parallel region participate in the execution of the units of work and the implied barrier of the workshare region if the barrier is not eliminated by a nowait clause.
Description
An implicit barrier occurs at the end of a `workshare` region if a `nowait` clause is not specified.

An implementation of the `workshare` construct must insert any synchronization that is required to maintain standard Fortran semantics. For example, the effects of one statement within the structured block must appear to occur before the execution of succeeding statements, and the evaluation of the right hand side of an assignment must appear to complete prior to the effects of assigning to the left hand side.

The statements in the `workshare` construct are divided into units of work as follows:

- For array expressions within each statement, including transformational array intrinsic functions that compute scalar values from arrays:
  - Evaluation of each element of the array expression, including any references to elemental functions, is a unit of work.
  - Evaluation of transformational array intrinsic functions may be freely subdivided into any number of units of work.
- For an array assignment statement, the assignment of each element is a unit of work.
- For a scalar assignment statement, the assignment operation is a unit of work.
- For a `WHERE` statement or construct, the evaluation of the mask expression and the masked assignments are each a unit of work.
- For a `FORALL` statement or construct, the evaluation of the mask expression, expressions occurring in the specification of the iteration space, and the masked assignments are each a unit of work.
- For an `atomic` construct, the atomic operation on the storage location designated as `x` is a unit of work.
- For a `critical` construct, the construct is a single unit of work.
- For a `parallel` construct, the construct is a unit of work with respect to the `workshare` construct. The statements contained in the `parallel` construct are executed by a new thread team.
- If none of the rules above apply to a portion of a statement in the structured block, then that portion is a unit of work.

The transformational array intrinsic functions are `MATMUL`, `DOT_PRODUCT`, `SUM`, `PRODUCT`, `MAXVAL`, `MINVAL`, `COUNT`, `ANY`, `ALL`, `SPREAD`, `PACK`, `UNPACK`, `RESHAPE`, `TRANSPOSE`, `EOSHIFT`, `CSHIFT`, `MINLOC`, and `MAXLOC`.

It is unspecified how the units of work are assigned to the threads that execute a `workshare` region.
If an array expression in the block references the value, association status, or allocation status of private variables, the value of the expression is undefined, unless the same value would be computed by every thread.

If an array assignment, a scalar assignment, a masked array assignment, or a `FORALL` assignment assigns to a private variable in the block, the result is unspecified.

The `workshare` directive causes the sharing of work to occur only in the `workshare` construct, and not in the remainder of the `workshare` region.

**Execution Model Events**
The `workshare-begin` event occurs after an implicit task encounters a `workshare` construct but before the task starts to execute the structured block of the `workshare` region.

The `workshare-end` event occurs after an implicit task finishes execution of a `workshare` region but before it resumes execution of the enclosing context.

**Tool Callbacks**
A thread dispatches a registered `ompt_callback_work` callback with `ompt_scope_begin` as its endpoint argument and `ompt_work_workshare` as its wstype argument for each occurrence of a `workshare-begin` event in that thread. Similarly, a thread dispatches a registered `ompt_callback_work` callback with `ompt_scope_end` as its endpoint argument and `ompt_work_workshare` as its wstype argument for each occurrence of a `workshare-end` event in that thread. The callbacks occur in the context of the implicit task. The callbacks have type signature `ompt_callback_work_t`.

**Restrictions**
Restrictions to the `workshare` construct are as follows:

- The only OpenMP constructs that may be closely nested inside a `workshare` construct are the `atomic`, `critical`, and `parallel` constructs.
- Base language statements that are encountered inside a `workshare` construct but that are not enclosed within a `parallel` construct that is nested inside the `workshare` construct must consist of only the following:
  - array assignments
  - scalar assignments
  - `FORALL` statements
  - `FORALL` constructs
  - `WHERE` statements
  - `WHERE` constructs
• All array assignments, scalar assignments, and masked array assignments that are encountered inside a `workshare` construct but are not nested inside a `parallel` construct that is nested inside the `workshare` construct must be intrinsic assignments.

• The construct must not contain any user-defined function calls unless the function is `ELEMENTAL` or the function call is contained inside a `parallel` construct that is nested inside the `workshare` construct.

Cross References
• `parallel` construct, see Section 2.6.
• `critical` construct, see Section 2.19.1.
• `atomic` construct, see Section 2.19.7.
• `ompt_scope_begin` and `ompt_scope_end`, see Section 4.4.11.
• `ompt_work_workshare`, see Section 4.4.15.
• `ompt_callback_work_t`, see Section 4.5.2.5.

2.11 Loop-Related Directives

2.11.1 Canonical Loop Nest Form

A loop nest has *canonical loop nest form* if it conforms to `loop-nest` in the following grammar:

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>loop-nest</code></td>
<td>One of the following:</td>
</tr>
<tr>
<td><code>C / C++</code></td>
<td><code>for (init-expr; test-expr; incr-expr)</code></td>
</tr>
<tr>
<td><code>loop-body</code></td>
<td></td>
</tr>
<tr>
<td><code>C / C++</code></td>
<td>or</td>
</tr>
<tr>
<td><code>C++</code></td>
<td><code>for (range-decl: range-expr)</code></td>
</tr>
<tr>
<td><code>loop-body</code></td>
<td></td>
</tr>
</tbody>
</table>

A range-based `for` loop is equivalent to a regular `for` loop using iterators, as defined in the base language. A range-based `for` loop has no iteration variable.
### Fortran

```fortran
DO [ label ] var = lb , ub [ , incr ]
   [intervening-code]
   loop-body
   [intervening-code]
   [ label ] END DO
```

If the loop-nest is a nonblock-do-construct, it is treated as a block-do-construct for each DO construct.

The value of incr is the increment of the loop. If not specified, its value is assumed to be 1.

### Fortran

or

```fortran
loop-transformation-construct
```

or

```fortran
generated-canonical-loop
```

### C / C++

```c
{ [intervening-code]
   loop-body
   [intervening-code]
}
```

### C / C++

```
```
Symbol | Meaning
----- | ------
1 | or

```
Fortran

**BLOCK**

[intervening-code]

loop-body

[intervening-code]

END BLOCK
```

Fortran

or if none of the previous productions match

```
final-loop-body
```

7

```
loop-transformation-construct
```

A loop transformation construct.

9

```
generated-canonical-loop
```

A generated loop from a loop transformation construct that has canonical loop nest form and for which the loop body matches `loop-body`.

10

```
intervening-code
```

A structured block sequence that does not contain OpenMP directives or calls to the OpenMP runtime API in its corresponding region, referred to as intervening code. If intervening code is present, then a loop at the same depth within the loop nest is not a perfectly nested loop.

11

```
C / C++
```

It must not contain iteration statements, `continue` statements or `break` statements that apply to the enclosing loop.

12

```
C / C++
```

Fortran

```
C / C++
```

```
Fortran
```

It must not contain loops, array expressions, `CYCLE` statements or `EXIT` statements.

13

```
final-loop-body
```

A structured block that terminates the scope of loops in the loop nest. If the loop nest is associated with a loop-associated directive, loops in this structured block cannot be associated with that directive.

14

```
C / C++
```

Fortran

```
Fortran
```

```
C / C++
```

```
Fortran
```

```
Fortran
```

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<table>
<thead>
<tr>
<th>Symbol</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>C / C++</strong></td>
<td></td>
</tr>
<tr>
<td><code>init-expr</code></td>
<td>One of the following:</td>
</tr>
<tr>
<td></td>
<td><code>var = lb</code></td>
</tr>
<tr>
<td></td>
<td><code>integer-type var = lb</code></td>
</tr>
<tr>
<td><code>C</code></td>
<td></td>
</tr>
<tr>
<td><code>pointer-type var = lb</code></td>
<td></td>
</tr>
<tr>
<td><code>C</code></td>
<td></td>
</tr>
<tr>
<td><code>C++</code></td>
<td></td>
</tr>
<tr>
<td><code>random-access-iterator-type var = lb</code></td>
<td></td>
</tr>
<tr>
<td><code>C++</code></td>
<td></td>
</tr>
<tr>
<td><code>test-expr</code></td>
<td>One of the following:</td>
</tr>
<tr>
<td></td>
<td><code>var relational-op ub</code></td>
</tr>
<tr>
<td></td>
<td><code>ub relational-op var</code></td>
</tr>
<tr>
<td><code>C / C++</code></td>
<td></td>
</tr>
<tr>
<td><code>relational-op</code></td>
<td>One of the following:</td>
</tr>
<tr>
<td></td>
<td><code>&lt;</code></td>
</tr>
<tr>
<td></td>
<td><code>&lt;=</code></td>
</tr>
<tr>
<td></td>
<td><code>&gt;</code></td>
</tr>
<tr>
<td></td>
<td><code>&gt;=</code></td>
</tr>
<tr>
<td></td>
<td><code>!=</code></td>
</tr>
<tr>
<td><code>C / C++</code></td>
<td></td>
</tr>
<tr>
<td><code>incr-expr</code></td>
<td>One of the following:</td>
</tr>
<tr>
<td></td>
<td><code>++var</code></td>
</tr>
<tr>
<td></td>
<td><code>var++</code></td>
</tr>
<tr>
<td></td>
<td><code>- - var</code></td>
</tr>
<tr>
<td></td>
<td><code>var - -</code></td>
</tr>
<tr>
<td></td>
<td><code>var += incr</code></td>
</tr>
<tr>
<td></td>
<td><code>var -= incr</code></td>
</tr>
<tr>
<td></td>
<td><code>var = var + incr</code></td>
</tr>
<tr>
<td></td>
<td><code>var = incr + var</code></td>
</tr>
<tr>
<td></td>
<td><code>var = var - incr</code></td>
</tr>
<tr>
<td></td>
<td>The value of <code>incr</code>, respectively 1 and -1 for the increment and decrement operators, is the increment of the loop.</td>
</tr>
<tr>
<td><code>C / C++</code></td>
<td></td>
</tr>
</tbody>
</table>
Symbol	Meaning

1

\[ \text{\textit{var}} \]	One of the following:

\[ \begin{align*}
\text{C / C++} \\
A \text{ variable of a signed or unsigned integer type.}
\end{align*} \]

\[ \text{C} \]	A variable of a pointer type.

\[ \text{C++} \]	A variable of a random access iterator type.

\[ \text{Fortran} \]	A variable of integer type.

\textit{var} is the iteration variable of the loop. It must not be modified during the execution of \textit{intervening-code} or \textit{loop-body} in the loop.

\[ \text{\textit{lb}, \textit{ub}} \]	One of the following:

Expressions of a type compatible with the type of \textit{var} that are loop invariant with respect to the outermost loop.

or

One of the following:

\begin{align*}
\text{\textit{var-outer}} \\
\textit{var-outer} + a2 \\
a2 + \textit{var-outer} \\
\textit{var-outer} - a2
\end{align*}

where \textit{var-outer} is of a type compatible with the type of \textit{var}.

or

If \textit{var} is of an integer type, one of the following:

\begin{align*}
a2 - \textit{var-outer} \\
a1 * \textit{var-outer} \\
a1 * \textit{var-outer} + a2 \\
a2 + a1 * \textit{var-outer}
\end{align*}
Symbol Meaning

1. $a_1 \times \text{var-outer} - a_2$
2. $a_2 - a_1 \times \text{var-outer}$
3. $\text{var-outer} \times a_1$
4. $\text{var-outer} \times a_1 + a_2$
5. $a_2 + \text{var-outer} \times a_1$
6. $\text{var-outer} \times a_1 - a_2$
7. $a_2 - \text{var-outer} \times a_1$

where var-outer is of an integer type.

$lb$ and $ub$ are loop bounds. A loop for which $lb$ or $ub$ refers to var-outer is a non-rectangular loop. If $var$ is of an integer type, var-outer must be of an integer type with the same signedness and bit precision as the type of var.

The coefficient in a loop bound is 0 if the bound does not refer to var-outer. If a loop bound matches a form in which $a_1$ appears, the coefficient is $-a_1$ if the product of var-outer and $a_1$ is subtracted from $a_2$, and otherwise the coefficient is $a_1$. For other matched forms where $a_1$ does not appear, the coefficient is $-1$ if var-outer is subtracted from $a_2$, and otherwise the coefficient is $1$.

$a_1, a_2, incr$ Integer expressions that are loop invariant with respect to the outermost loop of the loop nest.

If the loop is associated with a loop-associated directive, the expressions are evaluated before the construct formed from that directive.

var-outer The loop iteration variable of a surrounding loop in the loop nest.

range-decl A declaration of a variable as defined by the base language for range-based for loops.

range-expr An expression that is valid as defined by the base language for range-based for loops. It must be invariant with respect to the outermost loop of the loop nest and the iterator derived from it must be a random access iterator.
A loop transformation construct that appears inside a loop nest is replaced according to its semantics before any loop can be associated with a loop-associated directive that is applied to the loop nest. The depth of the loop nest is determined according to the loops in the loop nest, after any such replacements have taken place. A loop counts towards the depth of the loop nest if it is a base language loop statement or generated loop and it matches `loop-nest` while applying the production rules for canonical loop nest form to the loop nest.

A loop-associated directive controls some number of the outermost loops of an associated loop nest, called the associated loops, in accordance with its specified clauses. The canonical loop nest form allows the iteration count of all associated loops to be computed before executing the outermost loop.

For any associated loop, the iteration count is computed as follows:

- If `var` has a signed integer type and the `var` operand of `test-expr` after usual arithmetic conversions has an unsigned integer type then the loop iteration count is computed from `lb`, `test-expr` and `incr` using an unsigned integer type corresponding to the type of `var`.

- Otherwise, if `var` has an integer type then the loop iteration count is computed from `lb`, `test-expr` and `incr` using the type of `var`.

- If `var` has a pointer type then the loop iteration count is computed from `lb`, `test-expr` and `incr` using the type `ptrdiff_t`.

- If `var` has a random access iterator type then the loop iteration count is computed from `lb`, `test-expr` and `incr` using the type `std::iterator_traits<random-access-iterator-type>::difference_type`.

- For range-based for loops, the loop iteration count is computed from `range-expr` using the type `std::iterator_traits<random-access-iterator-type>::difference_type` where `random-access-iterator-type` is the iterator type derived from `range-expr`.

- The loop iteration count is computed from `lb`, `ub` and `incr` using the type of `var`. 
The behavior is unspecified if any intermediate result required to compute the iteration count cannot be represented in the type determined above.

No synchronization is implied during the evaluation of the \( lb, \) \( ub, \) \( incr \) or \( range-expr \) expressions. Whether, in what order, or how many times any side effects within the \( lb, \) \( ub, \) \( incr, \) or \( range-expr \) expressions occur is unspecified.

The iterations of some number of associated loops can be collapsed into one larger iteration space that is called the logical iteration space. The particular integer type used to compute the iteration count for the collapsed loop is implementation defined.

For directives that result in the execution of a collapsed logical iteration space, the number of times that any intervening code between any two loops of the same logical iteration space will be executed is unspecified but will be the same for all intervening code at the same depth, at least once per iteration of the loop enclosing the intervening code and at most once per logical iteration. If the iteration count of any loop is zero and that loop does not enclose the intervening code, the behavior is unspecified.

**Restrictions**

Restrictions to canonical loop nests are as follows:

For \( C / C++ \):

- If \( test-expr \) is of the form \( var \ relational-op \ b \) and \( relational-op \) is \(<\) or \( <=\) then \( incr-expr \) must cause \( var \) to increase on each iteration of the loop. If \( test-expr \) is of the form \( var \ relational-op \ b \) and \( relational-op \) is \(>\) or \( >=\) then \( incr-expr \) must cause \( var \) to decrease on each iteration of the loop. Increase and decrease are using the order induced by \( relational-op \).

- If \( test-expr \) is of the form \( ub \ relational-op \ var \) and \( relational-op \) is \(<\) or \( <=\) then \( incr-expr \) must cause \( var \) to decrease on each iteration of the loop. If \( test-expr \) is of the form \( ub \ relational-op \ var \) and \( relational-op \) is \(>\) or \( >=\) then \( incr-expr \) must cause \( var \) to increase on each iteration of the loop. Increase and decrease are using the order induced by \( relational-op \).

- If \( relational-op \) is \( !=\) then \( incr-expr \) must cause \( var \) to always increase by 1 or always decrease by 1 and the increment must be a constant expression.

- \( final-loop-body \) must not contain any \textbf{break} statement that would cause the termination of the innermost loop.

For \( Fortran \):

- \( final-loop-body \) must not contain any \textbf{EXIT} statement that would cause the termination of the innermost loop.
• A *loop-nest* must also be a structured block.

• For a non-rectangular loop, if *var-outer* is referenced in *lb* and *ub* then they must both refer to the same iteration variable.

• For a non-rectangular loop, let $a_1lb$ and $a_1ub$ be the respective coefficients in *lb* and *ub*, $incr_{inner}$ the increment of the non-rectangular loop and $incr_{outer}$ the increment of the loop referenced by *var-outer*. $incr_{inner}(a_1ub - a_1lb)$ must be a multiple of $incr_{outer}$.

• The loop iteration variable may not appear in a **threadprivate** directive.

**Cross References**

• Loop transformation constructs, see Section 2.11.9.

• **threadprivate** directive, see Section 2.21.2.

### 2.11.2 Consistent Loop Schedules

For constructs formed from loop-associated directives that have consistent schedules, the implementation will guarantee that memory effects of a logical iteration in the first loop nest happen before the execution of the same logical iteration in the second loop nest.

Two constructs formed from loop-associated directives have consistent schedules if all of the following conditions hold:

• The constructs are formed from directives with the same directive name;

• The regions that correspond to the two constructs have the same binding region;

• The constructs have the same reproducible schedule;

• The associated loop nests have identical logical iteration vector spaces; and

• The associated loop nests are either both rectangular or both non-rectangular.

### 2.11.3 order Clause

**Summary**

The *order* clause specifies an expected order of execution for the iterations of the associated loops of a loop-associated directive.

**Syntax**

The syntax of the *order* clause is as follows:

```
order([ order-modifier : ] concurrent)
```
where order-modifier is one of the following:

- reproducible
- unconstrained

**Description**

The order clause specifies an expected order of execution for the iterations of the associated loops of a loop-associated directive. The specified order must be concurrent.

The order clause is part of the schedule specification for the purpose of determining its consistency with other schedules (see Section 2.11.2).

If the order clause specifies concurrent, the logical iterations of the associated loops may execute in any order, including concurrently.

If order-modifier is not unconstrained, the behavior is as if the reproducible modifier is present.

The specified schedule is reproducible if the reproducible modifier is present.

**Restrictions**

Restrictions to the order clause are as follows:

- The only constructs that may be encountered inside a region that corresponds to a construct with an order clause that specifies concurrent are the loop construct, the parallel construct, the simd construct, and combined constructs for which the first construct is a parallel construct.

- A region that corresponds to a construct with an order clause that specifies concurrent may not contain calls to procedures that contain OpenMP directives.

- A region that corresponds to a construct with an order clause that specifies concurrent may not contain calls to the OpenMP Runtime API.

- If a threadprivate variable is referenced inside a region that corresponds to a construct with an order clause that specifies concurrent, the behavior is unspecified.

- At most one order clause may appear on a construct.

**2.11.4 Worksharing-Loop Construct**

**Summary**

The worksharing-loop construct specifies that the iterations of one or more associated loops will be executed in parallel by threads in the team in the context of their implicit tasks. The iterations are distributed across threads that already exist in the team that is executing the parallel region to which the worksharing-loop region binds.
The syntax of the worksharing-loop construct is as follows:

```c
#pragma omp for [clause[ , clause] ... ] loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` is one of the following:

- `private(list)`
- `firstprivate(list)`
- `lastprivate([[lastprivate-modifier:]list])`
- `linear(list[:linear-step])`
- `reduction([reduction-modifier,]reduction-identifier:list)`
- `schedule([modifier [, modifier]:]kind[, chunk_size])`
- `collapse(n)`
- `ordered/(n)`
- `nowait`
- `allocate([allocator:list])`
- `order([order-modifier:]concurrent)`

The syntax of the worksharing-loop construct is as follows:

```fortran
!$omp do [clause[ , clause] ... ] loop-nest
[!$omp end do [nowait]]
```

where `loop-nest` is a canonical loop nest and `clause` is one of the following:

- `private(list)`
- `firstprivate(list)`
- `lastprivate([[lastprivate-modifier:]list])`
- `linear(list[:linear-step])`
- `reduction([reduction-modifier,]reduction-identifier:list)`
- `schedule([modifier [, modifier]:]kind[, chunk_size])`
- `collapse(n)`
ordered\((n)\)
allocate\([\text{allocator}:\text{list}]\)
order\([\text{order-modifier}:\text{concurrent}]\)

If an \texttt{end do} directive is not specified, an \texttt{end do} directive is assumed at the end of the \texttt{do-loops}.

\textbf{Binding}

The binding thread set for a worksharing-loop region is the current team. A worksharing-loop region binds to the innermost enclosing \texttt{parallel} region. Only the threads of the team executing the binding \texttt{parallel} region participate in the execution of the loop iterations and the implied barrier of the worksharing-loop region when that barrier is not eliminated by a \texttt{nowait} clause.

\textbf{Description}

An implicit barrier occurs at the end of a worksharing-loop region if a \texttt{nowait} clause is not specified.

The \texttt{collapse} and \texttt{ordered} clauses may be used to specify the number of loops from the loop nest that are associated with the worksharing-loop construct. If specified, their parameters must be constant positive integer expressions.

The \texttt{collapse} clause specifies the number of loops that are collapsed into a logical iteration space that is then divided according to the \texttt{schedule} clause. If the \texttt{collapse} clause is omitted, the behavior is as if a \texttt{collapse} clause with a parameter value of one was specified.

If the \texttt{ordered} clause is specified with parameter \(n\) then the \(n\) outer loops from the associated loop nest form a \texttt{doacross loop nest}. The parameter of the \texttt{ordered} clause does not affect how the logical iteration space is divided.

At the beginning of each logical iteration, the loop iteration variable or the variable declared by \texttt{range-decl} of each associated loop has the value that it would have if the set of the associated loops was executed sequentially. The \texttt{schedule} clause specifies how iterations of these associated loops are divided into contiguous non-empty subsets, called chunks, and how these chunks are distributed among threads of the team. Each thread executes its assigned chunks in the context of its implicit task. The iterations of a given chunk are executed in sequential order by the assigned thread. The \texttt{chunk_size} expression is evaluated using the original list items of any variables that are made private in the worksharing-loop construct. Whether, in what order, or how many times, any side effects of the evaluation of this expression occur is unspecified. The use of a variable in a \texttt{schedule} clause expression of a worksharing-loop construct causes an implicit reference to the variable in all enclosing constructs.

See Section 2.11.4.1 for details of how the schedule for a worksharing-loop region is determined.

The schedule \texttt{kind} can be one of those specified in Table 2.5.
The schedule modifier can be one of those specified in Table 2.6. If the static schedule kind is specified or if the ordered clause is specified, and if the nonmonotonic modifier is not specified, the effect is as if the monotonic modifier is specified. Otherwise, unless the monotonic modifier is specified, the effect is as if the nonmonotonic modifier is specified. If a schedule clause specifies a modifier then that modifier overrides any modifier that is specified in the run-sched-var ICV.

If an order clause is present then the semantics are as described in Section 2.11.3.

The schedule is reproducible if one of the following conditions is true:

- The order clause is present and uses the reproducible modifier; or
- The schedule clause is specified with static as the kind parameter and the simd modifier is not present.

Programs can only depend on which thread executes a particular iteration if the schedule is reproducible. Schedule reproducibility is also used for determining its consistency with other schedules (see Section 2.11.2).

**TABLE 2.5: schedule Clause kind Values**

<table>
<thead>
<tr>
<th>kind</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>static</td>
<td>When kind is static, iterations are divided into chunks of size chunk_size, and the chunks are assigned to the threads in the team in a round-robin fashion in the order of the thread number. Each chunk contains chunk_size iterations, except for the chunk that contains the sequentially last iteration, which may have fewer iterations. When no chunk_size is specified, the iteration space is divided into chunks that are approximately equal in size, and at most one chunk is distributed to each thread. The size of the chunks is unspecified in this case.</td>
</tr>
<tr>
<td>dynamic</td>
<td>When kind is dynamic, the iterations are distributed to threads in the team in chunks. Each thread executes a chunk of iterations, then requests another chunk, until no chunks remain to be distributed. Each chunk contains chunk_size iterations, except for the chunk that contains the sequentially last iteration, which may have fewer iterations. When no chunk_size is specified, it defaults to 1.</td>
</tr>
<tr>
<td>guided</td>
<td>When kind is guided, the iterations are assigned to threads in the team in chunks. Each thread executes a chunk of iterations, then requests another chunk, until no chunks remain to be assigned.</td>
</tr>
</tbody>
</table>
For a chunk_size of 1, the size of each chunk is proportional to the number of unassigned iterations divided by the number of threads in the team, decreasing to 1. For a chunk_size with value \( k \) (greater than 1), the size of each chunk is determined in the same way, with the restriction that the chunks do not contain fewer than \( k \) iterations (except for the chunk that contains the sequentially last iteration, which may have fewer than \( k \) iterations).

When no chunk_size is specified, it defaults to 1.

**auto**

When kind is auto, the decision regarding scheduling is delegated to the compiler and/or runtime system. The programmer gives the implementation the freedom to choose any possible mapping of iterations to threads in the team.

**runtime**

When kind is runtime, the decision regarding scheduling is deferred until run time, and the schedule and chunk size are taken from the run-sched-var ICV. If the ICV is set to auto, the schedule is implementation defined.

---

**Note** – For a team of \( p \) threads and a loop of \( n \) iterations, let \( \left\lceil \frac{n}{p} \right\rceil \) be the integer \( q \) that satisfies \( n = p \times q - r \), with \( 0 <= r < p \). One compliant implementation of the static schedule (with no specified chunk_size) would behave as though chunk_size had been specified with value \( q \). Another compliant implementation would assign \( q \) iterations to the first \( p - r \) threads, and \( q - 1 \) iterations to the remaining \( r \) threads. This illustrates why a conforming program must not rely on the details of a particular implementation.

A compliant implementation of the guided schedule with a chunk_size value of \( k \) would assign \( q = \left\lceil \frac{n}{p} \right\rceil \) iterations to the first available thread and set \( n \) to the larger of \( n - q \) and \( p \times k \). It would then repeat this process until \( q \) is greater than or equal to the number of remaining iterations, at which time the remaining iterations form the final chunk. Another compliant implementation could use the same method, except with \( q = \left\lceil \frac{n}{2p} \right\rceil \), and set \( n \) to the larger of \( n - q \) and \( 2 \times p \times k \).
TABLE 2.6: schedule Clause modifier Values

<table>
<thead>
<tr>
<th>Modifier</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>monotonic</td>
<td>When the monotonic modifier is specified then each thread executes the chunks that it is assigned in increasing logical iteration order.</td>
</tr>
<tr>
<td>nonmonotonic</td>
<td>When the nonmonotonic modifier is specified then chunks are assigned to threads in any order and the behavior of an application that depends on any execution order of the chunks is unspecified.</td>
</tr>
<tr>
<td>simd</td>
<td>When the simd modifier is specified and the loop is associated with a SIMD construct, the chunk_size for all chunks except the first and last chunks is ( new_chunk_size = \left\lceil \frac{\text{chunk}_size}{\text{simd}_width} \right\rceil \times \text{simd}_width ) where simd_width is an implementation-defined value. The first chunk will have at least new_chunk_size iterations except if it is also the last chunk. The last chunk may have fewer iterations than new_chunk_size. If the simd modifier is specified and the loop is not associated with a SIMD construct, the modifier is ignored.</td>
</tr>
</tbody>
</table>

Execution Model Events

The ws-loop-begin event occurs after an implicit task encounters a worksharing-loop construct but before the task starts execution of the structured block of the worksharing-loop region.

The ws-loop-end event occurs after a worksharing-loop region finishes execution but before resuming execution of the encountering task.

The ws-loop-iteration-begin event occurs once for each iteration of a worksharing-loop before the iteration is executed by an implicit task.

Tool Callbacks

A thread dispatches a registered ompt_callback_work callback with ompt_scope_begin as its endpoint argument and work_loop as its wstype argument for each occurrence of a ws-loop-begin event in that thread. Similarly, a thread dispatches a registered ompt_callback_work callback with ompt_scope_end as its endpoint argument and work_loop as its wstype argument for each occurrence of a ws-loop-end event in that thread. The callbacks occur in the context of the implicit task. The callbacks have type signature ompt_callback_work_t.

A thread dispatches a registered ompt_callback_dispatch callback for each occurrence of a ws-loop-iteration-begin event in that thread. The callback occurs in the context of the implicit task. The callback has type signature ompt_callback_dispatch_t.
Restrictions
Restrictions to the worksharing-loop construct are as follows:

- If the ordered clause with a parameter is present, all associated loops must be perfectly nested.
- If a reduction clause with the inscan modifier is specified, neither the ordered nor schedule clause may appear on the worksharing-loop directive.
- The values of the loop control expressions of the loops associated with the worksharing-loop construct must be the same for all threads in the team.
- At most one schedule clause can appear on a worksharing-loop directive.
- If the schedule or ordered clause is present then none of the associated loops may be non-rectangular loops.
- The ordered clause must not appear on the worksharing-loop directive if the associated loops include the generated loops of a tile directive.
- At most one collapse clause can appear on a worksharing-loop directive.
- chunk_size must be a loop invariant integer expression with a positive value.
- The value of the chunk_size expression must be the same for all threads in the team.
- The value of the run-sched-var ICV must be the same for all threads in the team.
- When schedule(runtime) or schedule(auto) is specified, chunk_size must not be specified.
- A modifier may not be specified on a linear clause.
- At most one ordered clause can appear on a worksharing-loop directive.
- The ordered clause must be present on the worksharing-loop construct if any ordered region ever binds to a worksharing-loop region arising from the worksharing-loop construct.
- The nonmonotonic modifier cannot be specified if an ordered clause is specified.
- Each schedule clause modifier may be specified at most once on the same schedule clause.
- Either the monotonic modifier or the nonmonotonic modifier can be specified but not both.
- If both the collapse and ordered clause with a parameter are specified, the parameter of the ordered clause must be greater than or equal to the parameter of the collapse clause.
- The values of the parameters specified by the collapse and ordered clauses must not exceed the depth of the associated loop nest.
- A linear clause or an ordered clause with a parameter can be specified on a worksharing-loop directive but not both.
C / C++

1. At most one `nowait` clause can appear on a `for` directive.

C++

2. If an `ordered` clause with a parameter is specified, none of the associated loops may be a range-based `for` loop.

C++

Cross References

- Canonical loop nest form, see Section 2.11.1.
- `order` clause, see Section 2.11.3.
- `tile` construct, see Section 2.11.9.1.
- `ordered` construct, see Section 2.19.9.
- `depend` clause, see Section 2.19.11.
- Data-sharing attribute clauses, see Section 2.21.4.
- `omp_t_scope_begin` and `omp_t_scope_end`, see Section 4.4.4.11.
- `omp_t_work_loop`, see Section 4.4.4.15.
- `omp_t_callback_work_t`, see Section 4.5.2.5.
- `OMP_SCHEDULE` environment variable, see Section 6.1.

2.11.4.1 Determining the Schedule of a Worksharing-Loop

When execution encounters a worksharing-loop directive, the `schedule` clause (if any) on the directive, and the `run-sched-var` and `def-sched-var` ICVs are used to determine how loop iterations are assigned to threads. See Section 2.4 for details of how the values of the ICVs are determined. If the worksharing-loop directive does not have a `schedule` clause then the current value of the `def-sched-var` ICV determines the schedule. If the worksharing-loop directive has a `schedule` clause that specifies the `runtime` schedule kind then the current value of the `run-sched-var` ICV determines the schedule. Otherwise, the value of the `schedule` clause determines the schedule. Figure 2.1 describes how the schedule for a worksharing-loop is determined.

Cross References

- ICVs, see Section 2.4.
2.11.5 SIMD Directives

2.11.5.1 simd Construct

Summary
The simd construct can be applied to a loop to indicate that the loop can be transformed into a SIMD loop (that is, multiple iterations of the loop can be executed concurrently by using SIMD instructions).

Syntax

C / C++

The syntax of the simd construct is as follows:

```c
#pragma omp simd [clause [ , clause] ... ] new-line
```

FIGURE 2.1: Determining the schedule for a Worksharing-Loop
where loop-nest is a canonical loop nest and clause is one of the following:

```c
if([ simd : ] scalar-expression)
safelen(length)
simdlen(length)
linear(list[ : linear-step])
aligned(list[ : alignment])
nontemporal(list)
private(list)
lastprivate([ lastprivate-modifier : ] list)
reduction([ reduction-modifier, ]reduction-identifier : list)
collapse(n)
order([ order-modifier : ]concurrent)
```

C / C++

Fortran

The syntax of the simd construct is as follows:

```c
!$omp simd [clause[ , ] clause] ... ]
loop-nest
[$omp end simd]
```

where loop-nest is a canonical loop nest and clause is one of the following:

```c
if([ simd : ] scalar-logical-expression)
safelen(length)
simdlen(length)
linear(list[ : linear-step])
aligned(list[ : alignment])
nontemporal(list)
private(list)
lastprivate([ lastprivate-modifier : ] list)
reduction([ reduction-modifier, ]reduction-identifier : list)
collapse(n)
order([ order-modifier : ]concurrent)
```

Fortran

If an end simd directive is not specified, an end simd directive is assumed at the end of the do-loops.
**Binding**

A **simd** region binds to the current task region. The binding thread set of the **simd** region is the current team.

**Description**

The **simd** construct enables the execution of multiple iterations of the associated loops concurrently by using SIMD instructions.

The **collapse** clause may be used to specify how many loops are associated with the **simd** construct. The **collapse** clause specifies the number of loops that are collapsed into a logical iteration space that is then executed with SIMD instructions. The parameter of the **collapse** clause must be a constant positive integer expression. If the **collapse** clause is omitted, the behavior is as if a **collapse** clause with a parameter value of one was specified.

At the beginning of each logical iteration, the loop iteration variable or the variable declared by **range-decl** of each associated loop has the value that it would have if the set of the associated loops was executed sequentially. The number of iterations that are executed concurrently at any given time is implementation defined. Each concurrent iteration will be executed by a different SIMD lane. Each set of concurrent iterations is a SIMD chunk. Lexical forward dependences in the iterations of the original loop must be preserved within each SIMD chunk, unless an **order** clause that specifies **concurrent** is present.

The **safelen** clause specifies that no two concurrent iterations within a SIMD chunk can have a distance in the logical iteration space that is greater than or equal to the value given in the clause. The parameter of the **safelen** clause must be a constant positive integer expression. The **simdlen** clause specifies the preferred number of iterations to be executed concurrently, unless an **if** clause is present and evaluates to **false**, in which case the preferred number of iterations to be executed concurrently is one. The parameter of the **simdlen** clause must be a constant positive integer expression.

If an **order** clause is present then the semantics are as described in Section 2.11.3.

```c
C / C++

The **aligned** clause declares that the object to which each list item points is aligned to the number of bytes expressed in the optional parameter of the **aligned** clause.
```

```fortran
Fortran

The optional parameter of the **aligned** clause, **alignment**, must be a constant positive integer expression. If no optional parameter is specified, implementation-defined default alignments for SIMD instructions on the target platforms are assumed.
```

The **nontemporal** clause specifies that accesses to the storage locations to which the list items refer have low temporal locality across the iterations in which those storage locations are accessed.
Restrictions

Restrictions to the simd construct are as follows:

• At most one collapse clause can appear on a simd directive.
• A list-item cannot appear in more than one aligned clause.
• A list-item cannot appear in more than one nontemporal clause.
• At most one safelen clause can appear on a simd directive.
• At most one simdlen clause can appear on a simd directive.
• At most one if clause can appear on a simd directive.
• If both simdlen and safelen clauses are specified, the value of the simdlen parameter must be less than or equal to the value of the safelen parameter.
• A modifier may not be specified on a linear clause.
• The only OpenMP constructs that can be encountered during execution of a simd region are the atomic construct, the loop construct, the simd construct, and the ordered construct with the simd clause.
• If an order clause that specifies concurrent appears on a simd directive, the safelen clause may not also appear.

- The simd region cannot contain calls to the longjmp or setjmp functions.

- The type of list items appearing in the aligned clause must be array or pointer.

- The type of list items appearing in the aligned clause must be array, pointer, reference to array, or reference to pointer.

- No exception can be raised in the simd region.

- The only random access iterator types that are allowed for the associated loops are pointer types.
• If a list item on the `aligned` clause has the `ALLOCATABLE` attribute, the allocation status must be allocated.

• If a list item on the `aligned` clause has the `POINTER` attribute, the association status must be associated.

• If the type of a list item on the `aligned` clause is either `C_PTR` or Cray pointer, the list item must be defined. Cray pointer support has been deprecated.

---

### Cross References

- Canonical loop nest form, see Section 2.11.1.
- `order` clause, see Section 2.11.3.
- `if` clause, see Section 2.18.
- Data-sharing attribute clauses, see Section 2.21.4.

---

#### 2.11.5.2 Worksharing-Loop SIMD Construct

**Summary**

The worksharing-loop SIMD construct specifies that the iterations of one or more associated loops will be distributed across threads that already exist in the team and that the iterations executed by each thread can also be executed concurrently using SIMD instructions. The worksharing-loop SIMD construct is a composite construct.

**Syntax**

```c
#pragma omp for simd [clause [ [, ] clause ] ... ] new-line
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `for` or `simd` directives with identical meanings and restrictions.
The syntax of the worksharing-loop SIMD construct is as follows:

```
!$omp do simd [clause [ , ] clause] ... ]
  loop-nest
[ !$omp end do simd /nowait/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `simd` or `do` directives, with identical meanings and restrictions.

If an `end do simd` directive is not specified, an `end do simd` directive is assumed at the end of the `do-loops`.

### Description
The worksharing-loop SIMD construct will first distribute the logical iterations of the associated loops across the implicit tasks of the parallel region in a manner consistent with any clauses that apply to the worksharing-loop construct. Each resulting chunk of iterations will then be converted to a SIMD loop in a manner consistent with any clauses that apply to the `simd` construct.

### Execution Model Events
This composite construct generates the same events as the worksharing-loop construct.

### Tool Callbacks
This composite construct dispatches the same callbacks as the worksharing-loop construct.

### Restrictions
All restrictions to the worksharing-loop construct and the `simd` construct apply to the worksharing-loop SIMD construct. In addition, the following restrictions apply:

- No `ordered` clause with a parameter can be specified.
- A list item may appear in a `linear` or `firstprivate` clause, but not in both.

### Cross References
- Canonical loop nest form, see Section 2.11.1.
- Worksharing-loop construct, see Section 2.11.4.
- `simd` construct, see Section 2.11.5.1.
- Data-sharing attribute clauses, see Section 2.21.4.
2.11.5.3 declare simd Directive

Summary
The declare simd directive can be applied to a function (C, C++, and Fortran) or a subroutine (Fortran) to enable the creation of one or more versions that can process multiple arguments using SIMD instructions from a single invocation in a SIMD loop. The declare simd directive is a declarative directive. Multiple declare simd directives may be specified for a function (C, C++, and Fortran) or subroutine (Fortran).

Syntax

C / C++

The syntax of the declare simd directive is as follows:

```
#pragma omp declare simd [clause [ , ] clause] ... ] new-line
```

where clause is one of the following:

```
simdlen(length)
linear(linear-list[ : linear-step])
aligned(argument-list[ : alignment])
uniform(argument-list)
inbranch
notinbranch
```

Fortran

The syntax of the declare simd directive is as follows:

```
!$omp declare simd [ (proc-name) ] [clause [ , ] clause] ... ]
```

where clause is one of the following:

```
simdlen(length)
linear(linear-list[ : linear-step])
aligned(argument-list[ : alignment])
uniform(argument-list)
inbranch
notinbranch
```
The use of one or more `declare simd` directives on a function declaration or definition enables the creation of corresponding SIMD versions of the associated function that can be used to process multiple arguments from a single invocation in a SIMD loop concurrently.

The expressions appearing in the clauses of each directive are evaluated in the scope of the arguments of the function declaration or definition.

The use of one or more `declare simd` directives in a subroutine or function enables the creation of corresponding SIMD versions of the subroutine or function that can be used to process multiple arguments from a single invocation in a SIMD loop concurrently.

If a SIMD version is created, the number of concurrent arguments for the function is determined by the `simdlen` clause. If the `simdlen` clause is used, its value corresponds to the number of concurrent arguments of the function. The parameter of the `simdlen` clause must be a constant positive integer expression. Otherwise, the number of concurrent arguments for the function is implementation defined.

The special `this` pointer can be used as if it was one of the arguments to the function in any of the `linear`, `aligned`, or `uniform` clauses.

The `uniform` clause declares one or more arguments to have an invariant value for all concurrent invocations of the function in the execution of a single SIMD loop.

The `aligned` clause declares that the object to which each list item points is aligned to the number of bytes expressed in the optional parameter of the `aligned` clause.

The `aligned` clause declares that the target of each list item is aligned to the number of bytes expressed in the optional parameter of the `aligned` clause.
The optional parameter of the `aligned` clause, `alignment`, must be a constant positive integer expression. If no optional parameter is specified, implementation-defined default alignments for SIMD instructions on the target platforms are assumed.

The `inbranch` clause specifies that the SIMD version of the function will always be called from inside a conditional statement of a SIMD loop. The `notinbranch` clause specifies that the SIMD version of the function will never be called from inside a conditional statement of a SIMD loop. If neither clause is specified, then the SIMD version of the function may or may not be called from inside a conditional statement of a SIMD loop.

**Restrictions**

Restrictions to the `declare simd` directive are as follows:

- Each argument can appear in at most one `uniform` or `linear` clause.
- At most one `simdlen` clause can appear in a `declare simd` directive.
- Either `inbranch` or `notinbranch` may be specified, but not both.
- When a `linear-step` expression is specified in a `linear` clause it must be either a constant integer expression or an integer-typed parameter that is specified in a `uniform` clause on the directive.
- The function or subroutine body must be a structured block.
- The execution of the function or subroutine, when called from a SIMD loop, cannot result in the execution of an OpenMP construct except for an `ordered` construct with the `simd` clause or an `atomic` construct.
- The execution of the function or subroutine cannot have any side effects that would alter its execution for concurrent iterations of a SIMD chunk.
- A program that branches into or out of the function is non-conforming.

```
C / C++
```

- If the function has any declarations, then the `declare simd` directive for any declaration that has one must be equivalent to the one specified for the definition. Otherwise, the result is unspecified.

```
C / C++
```

- The function cannot contain calls to the `longjmp` or `setjmp` functions.

```
C / C++
```

```
C
```

- The type of list items appearing in the `aligned` clause must be array or pointer.
The function cannot contain any calls to `throw`.

The type of list items appearing in the `aligned` clause must be array, pointer, reference to array, or reference to pointer.

- `proc-name` must not be a generic name, procedure pointer, or entry name.
- If `proc-name` is omitted, the `declare simd` directive must appear in the specification part of a subroutine subprogram or a function subprogram for which creation of the SIMD versions is enabled.
- Any `declare simd` directive must appear in the specification part of a subroutine subprogram, function subprogram, or interface body to which it applies.
- If a `declare simd` directive is specified in an interface block for a procedure, it must match a `declare simd` directive in the definition of the procedure.
- If a procedure is declared via a procedure declaration statement, the procedure `proc-name` should appear in the same specification.
- If a `declare simd` directive is specified for a procedure name with explicit interface and a `declare simd` directive is also specified for the definition of the procedure then the two `declare simd` directives must match. Otherwise the result is unspecified.
- Procedure pointers may not be used to access versions created by the `declare simd` directive.
- The type of list items appearing in the `aligned` clause must be `C_PTR` or Cray pointer, or the list item must have the `POINTER` or `ALLOCATABLE` attribute. Cray pointer support has been deprecated.

Cross References
- `linear` clause, see Section 2.21.4.6.
- `reduction` clause, see Section 2.21.5.4.

### 2.11.6 `distribute` Loop Constructs

#### 2.11.6.1 `distribute` Construct

**Summary**

The `distribute` construct specifies that the iterations of one or more loops will be executed by the initial teams in the context of their implicit tasks. The iterations are distributed across the initial threads of all initial teams that execute the `teams` region to which the `distribute` region binds.
The syntax of the `distribute` construct is as follows:

```c
#pragma omp distribute [clause [ , clause] ... ] new-line
```

where `loop-nest` is a canonical loop nest and `clause` is one of the following:

- `private(list)`
- `firstprivate(list)`
- `lastprivate(list)`
- `collapse(n)`
- `dist_schedule(kind[, chunk_size])`
- `allocate([allocator :]list)`
- `order([ order-modifier :]concurrent)`

If an `end distribute` directive is not specified, an `end distribute` directive is assumed at the end of the `do-loops`. 

The syntax of the `distribute` construct is as follows:

```fortran
!$omp distribute [clause [ , clause] ... ]
```

where `loop-nest` is a canonical loop nest and `clause` is one of the following:

- `private(list)`
- `firstprivate(list)`
- `lastprivate(list)`
- `collapse(n)`
- `dist_schedule(kind[, chunk_size])`
- `allocate([allocator :]list)`
- `order([ order-modifier :]concurrent)`

If an `end distribute` directive is not specified, an `end distribute` directive is assumed at the end of the `do-loops`. 

---

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Binding
The binding thread set for a `distribute` region is the set of initial threads executing an enclosing `teams` region. A `distribute` region binds to this `teams` region.

Description
The `distribute` construct is associated with a loop nest consisting of one or more loops that follow the directive.

The `collapse` clause may be used to specify how many loops are associated with the `distribute` construct. The parameter of the `collapse` clause must be a constant positive integer expression. If the `collapse` clause is omitted, the behavior is as if a `collapse` clause with a parameter value of one was specified.

No implicit barrier occurs at the end of a `distribute` region. To avoid data races the original list items that are modified due to `lastprivate` or `linear` clauses should not be accessed between the end of the `distribute` construct and the end of the `teams` region to which the `distribute` binds.

At the beginning of each logical iteration, the loop iteration variable or the variable declared by `range-decl` of each associated loop has the value that it would have if the set of the associated loops was executed sequentially.

If the `dist_schedule` clause is specified, `kind` must be `static`. If specified, iterations are divided into chunks of size `chunk_size`. These chunks are assigned to the initial teams of the league in a round-robin fashion in the order of the initial team number. When `chunk_size` is not specified, the iteration space is divided into chunks that are approximately equal in size, and at most one chunk is distributed to each initial team of the league.

When `dist_schedule` clause is not specified, the schedule is implementation defined.

If an `order` clause is present then the semantics are as described in Section 2.11.3.

The schedule is reproducible if one of the following conditions is true:

- The `order` clause is present and uses the `reproducible` modifier; or
- The `dist_schedule` clause is specified with `static` as the `kind` parameter.

Programs can only depend on which team executes a particular iteration if the schedule is reproducible. Schedule reproducibility is also used for determining its consistency with other schedules (see Section 2.11.2).

Execution Model Events
The `distribute-begin` event occurs after an implicit task encounters a `distribute` construct but before the task starts to execute the structured block of the `distribute` region.

The `distribute-end` event occurs after an implicit task finishes execution of a `distribute` region but before it resumes execution of the enclosing context.
Tool Callbacks
A thread dispatches a registered `ompt_callback_work` callback with `ompt_scope_begin` as its `endpoint` argument and `ompt_work_distribute` as its `wstype` argument for each occurrence of a `distribute-begin` event in that thread. Similarly, a thread dispatches a registered `ompt_callback_work` callback with `ompt_scope_end` as its `endpoint` argument and `ompt_work_distribute` as its `wstype` argument for each occurrence of a `distribute-end` event in that thread. The callbacks occur in the context of the implicit task. The callbacks have type signature `ompt_callback_work_t`.

Restrictions
Restrictions to the `distribute` construct are as follows:

- The `distribute` construct inherits the restrictions of the worksharing-loop construct.
- Each `distribute` region must be encountered by the initial threads of all initial teams in a league or by none at all.
- The sequence of the `distribute` regions encountered must be the same for every initial thread of every initial team in a league.
- The region that corresponds to the `distribute` construct must be strictly nested inside a `teams` region.
- A list item may appear in a `firstprivate` or `lastprivate` clause, but not in both.
- At most one `dist_schedule` clause can appear on the directive.
- If the `dist_schedule` is present then none of the associated loops may be non-rectangular loops.

Cross References
- `teams` construct, see Section 2.7
- Canonical loop nest form, see Section 2.11.1.
- `order` clause, see Section 2.11.3.
- Worksharing-loop construct, see Section 2.11.4.
- `tile` construct, see Section 2.11.9.1.
- `ompt_work_distribute`, see Section 4.4.4.15.
- `ompt_callback_work_t`, see Section 4.5.2.5.
2.11.6.2 distribute simd Construct

Summary
The distribute simd construct specifies a loop that will be distributed across the primary threads of the teams region and executed concurrently using SIMD instructions. The distribute simd construct is a composite construct.

Syntax

C / C++

#pragma omp distribute simd [clause[ [, ] clause] ... ]

loop-nest

The syntax of the distribute simd construct is as follows:

- where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the distribute or simd directives with identical meanings and restrictions.

Fortran

!$omp distribute simd [clause[ [, ] clause] ... ]

loop-nest

|$omp end distribute simd]

The syntax of the distribute simd construct is as follows:

- where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the distribute or simd directives with identical meanings and restrictions.

If an end distribute simd directive is not specified, an end distribute simd directive is assumed at the end of the do-loops.

Description
The distribute simd construct will first distribute the logical iterations of the associated loops across the initial tasks of the teams region in a manner consistent with any clauses that apply to the distribute construct. Each resulting chunk of iterations will then be converted to a SIMD loop in a manner consistent with any clauses that apply to the simd construct.

Execution Model Events
This composite construct generates the same events as the distribute construct.

Tool Callbacks
This composite construct dispatches the same callbacks as the distribute construct.
Restrictions

All restrictions to the `distribute` and `simd` constructs apply to the `distribute simd` construct. In addition, the following restrictions apply:

- A list item may not appear in a `linear` clause unless it is the loop iteration variable of a loop that is associated with the construct.

- The `conditional` modifier may not appear in a `lastprivate` clause.

Cross References

- Canonical loop nest form, see Section 2.11.1.
- `simd` construct, see Section 2.11.5.1.
- `distribute` construct, see Section 2.11.6.1.
- Data-sharing attribute clauses, see Section 2.21.4.

2.11.6.3 Distribute Parallel Worksharing-Loop Construct

Summary

The distribute parallel worksharing-loop construct specifies a loop that can be executed in parallel by multiple threads that are members of multiple teams. The distribute parallel worksharing-loop construct is a composite construct.

Syntax

```
#pragma omp distribute parallel for [clause[ [, ] clause] ... ] new-line
    loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `distribute` or parallel worksharing-loop directives with identical meanings and restrictions.

```
!$omp distribute parallel do [clause[ [, ] clause] ... ]
    loop-nest

/$omp end distribute parallel do/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `distribute` or parallel worksharing-loop directives with identical meanings and restrictions.

If an `end distribute parallel do` directive is not specified, an `end distribute parallel do` directive is assumed at the end of the `do-loops`. 
Description
The distribute parallel worksharing-loop construct will first distribute the logical iterations of the
associated loops across the initial tasks of the \textit{teams} region in a manner consistent with any
clauses that apply to the \texttt{distribute} construct. Each resulting chunk of iterations will then
execute as if part of a parallel worksharing-loop region in a manner consistent with any clauses that
apply to the parallel worksharing-loop construct.

Execution Model Events
This composite construct generates the same events as the \texttt{distribute} and parallel
worksharing-loop constructs.

Tool Callbacks
This composite construct dispatches the same callbacks as the \texttt{distribute} and parallel
worksharing-loop constructs.

Restrictions
All restrictions to the \texttt{distribute} and parallel worksharing-loop constructs apply to the
distribute parallel worksharing-loop construct. In addition, the following restrictions apply:

\begin{itemize}
  \item No \texttt{ordered} clause can be specified.
  \item No \texttt{linear} clause can be specified.
  \item The \texttt{conditional} modifier must not appear in a \texttt{lastprivate} clause.
\end{itemize}

Cross References
\begin{itemize}
  \item Canonical loop nest form, see Section 2.11.1.
  \item \texttt{distribute} construct, see Section 2.11.6.1.
  \item Parallel worksharing-loop construct, see Section 2.16.1.
  \item Data-sharing attribute clauses, see Section 2.21.4.
\end{itemize}

2.11.6.4 Distribute Parallel Worksharing-Loop SIMD Construct

Summary
The distribute parallel worksharing-loop SIMD construct specifies a loop that can be executed
concurrently using SIMD instructions in parallel by multiple threads that are members of multiple
teams. The distribute parallel worksharing-loop SIMD construct is a composite construct.
The syntax of the distribute parallel worksharing-loop SIMD construct is as follows:

```c
#pragma omp distribute parallel for simd \
    [clause [, clause] ... ] new-line
loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `distribute` or parallel worksharing-loop SIMD directives with identical meanings and restrictions.

The syntax of the distribute parallel worksharing-loop SIMD construct is as follows:

```fortran
!$omp distribute parallel do simd [clause [, clause] ... ]
loop-nest

/$omp end distribute parallel do simd/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `distribute` or parallel worksharing-loop SIMD directives with identical meanings and restrictions.

If an `end distribute parallel do simd` directive is not specified, an `end distribute parallel do simd` directive is assumed at the end of the `do-loops`.

Description

The distribute parallel worksharing-loop SIMD construct will first distribute the logical iterations of the associated loops across the initial tasks of the `teams` region in a manner consistent with any clauses that apply to the `distribute` construct. Each resulting chunk of iterations will then execute as if part of a parallel worksharing-loop SIMD region in a manner consistent with any clauses that apply to the parallel worksharing-loop SIMD construct.

Execution Model Events

This composite construct generates the same events as the `distribute` and parallel worksharing-loop SIMD constructs.

Tool Callbacks

This composite construct dispatches the same callbacks as the `distribute` and parallel worksharing-loop SIMD constructs.
Restrictions
All restrictions to the `distribute` and parallel worksharing-loop SIMD constructs apply to the `distribute` parallel worksharing-loop SIMD construct. In addition, the following restrictions apply:

- No `ordered` clause can be specified.
- A list item may not appear in a `linear` clause unless it is the loop iteration variable of a loop that is associated with the construct.
- The `conditional` modifier may not appear in a `lastprivate` clause.
- If any `if` clause on the directive includes a `directive-name-modifier` then all `if` clauses on the directive must include a `directive-name-modifier`.
- At most one `if` clause without a `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `parallel` directive-name-modifier can appear on the directive.
- At most one `if` clause with the `simd` directive-name-modifier can appear on the directive.

Cross References
- Canonical loop nest form, see Section 2.11.1.
- `distribute` construct, see Section 2.11.6.1.
- Parallel worksharing-loop SIMD construct, see Section 2.16.5.
- Data-sharing attribute clauses, see Section 2.21.4.

2.11.7 Loop Construct

Summary
A `loop` construct specifies that the logical iterations of the associated loops may execute concurrently and permits the encountering threads to execute the loop accordingly.

Syntax
The syntax of the `loop` construct is as follows:

```c
#pragma omp loop [clause[ [, clause] ... ] new-line
loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` is one of the following:

```c
bind(binding)
collapse(n)
order([ order-modifier :]concurrent)
private(list)
```
lastprivate(list)
reduction([default, reduction-identifier : list])

where binding is one of the following:

teams
parallel
thread

C / C++

Fortran

The syntax of the loop construct is as follows:

!$omp loop [clause[ [, clause] ... ]]

loop-nest

[!$omp end loop]

where loop-nest is a canonical loop nest and clause is one of the following:

bind(binding)
collapse(n)

order([order-modifier : ]concurrent)

private(list)

lastprivate(list)

reduction([default, reduction-identifier : list])

where binding is one of the following:

teams
parallel
thread

If an end loop directive is not specified, an end loop directive is assumed at the end of the do-loops.
Binding
If the **bind** clause is present on the construct, the binding region is determined by **binding**.
Specifically, if **binding** is **teams** and an innermost enclosing **teams** region exists then the binding region is that **teams** region; if **binding** is **parallel** then the binding region is the innermost enclosing parallel region, which may be an implicit parallel region; and if **binding** is **thread** then the binding region is not defined. If the **bind** clause is not present on the construct and the **loop** construct is closely nested inside a **teams** or **parallel** construct, the binding region is the corresponding **teams** or **parallel** region. If none of those conditions hold, the binding region is not defined.

If the binding region is a **teams** region, then the binding thread set is the set of initial threads that are executing that region. If the binding region is a parallel region, then the binding thread set is the team of threads that are executing that region. If the binding region is not defined, then the binding thread set is the encountering thread.

Description
The **loop** construct is associated with a loop nest that consists of one or more loops that follow the directive. The directive asserts that the iterations may execute in any order, including concurrently.

The **collapse** clause may be used to specify how many loops are associated with the **loop** construct. The parameter of the **collapse** clause must be a constant positive integer expression. If the **collapse** clause is omitted, the behavior is as if a **collapse** clause with a parameter value of one was specified. The **collapse** clause specifies the number of loops that are collapsed into a logical iteration space.

At the beginning of each logical iteration, the loop iteration variable or the variable declared by **range-decl** of each associated loop has the value that it would have if the set of the associated loops was executed sequentially.

Each logical iteration is executed once per instance of the **loop** region that is encountered by the binding thread set.

If an **order** clause is present then the semantics are as described in Section 2.11.3. If the **order** clause is not present, the behavior is as if an **order** clause that specifies **concurrent** appeared on the construct.

The set of threads that may execute the iterations of the **loop** region is the binding thread set. Each iteration is executed by one thread from this set.

If the **loop** region binds to a **teams** region, the threads in the binding thread set may continue execution after the **loop** region without waiting for all logical iterations of the associated loops to complete. The iterations are guaranteed to complete before the end of the **teams** region.

If the **loop** region does not bind to a **teams** region, all logical iterations of the associated loops must complete before the encountering threads continue execution after the **loop** region.

For the purpose of determining its consistency with other schedules (see Section 2.11.2), the schedule is defined by the implicit **order** clause.
The schedule is reproducible if the schedule specified through the implicit `order` clause is reproducible.

**Restrictions**

Restrictions to the `loop` construct are as follows:

- At most one `collapse` clause can appear on a `loop` directive.
- A list item may not appear in a `lastprivate` clause unless it is the loop iteration variable of a loop that is associated with the construct.
- If a `loop` construct is not nested inside another OpenMP construct and it appears in a procedure, the `bind` clause must be present.
- If a `loop` region binds to a `teams` or parallel region, it must be encountered by all threads in the binding thread set or by none of them.
- At most one `bind` clause can appear on a `loop` directive.
- If the `bind` clause is present on the `loop` construct and `binding` is `teams` then the corresponding `loop` region must be strictly nested inside a `teams` region.
- If the `bind` clause, with `teams` specified as `binding`, is present on the `loop` construct and the corresponding `loop` region executes on a non-host device then the behavior of a `reduction` clause that appears on the construct is unspecified if the construct is not nested inside a `teams` construct.
- If the `bind` clause is present on the `loop` construct and `binding` is `parallel` then the behavior is unspecified if the corresponding `loop` region is closely nested inside a `simd` region.

**Cross References**

- The `single` construct, see Section 2.10.2.
- Canonical loop nest form, see Section 2.11.1.
- `order` clause, see Section 2.11.3.
- The Worksharing-Loop construct, see Section 2.11.4.
- SIMD directives, see Section 2.11.5.
- `distribute` construct, see Section 2.11.6.1.

**2.11.8 scan Directive**

**Summary**

The `scan` directive specifies that scan computations update the list items on each iteration of an enclosing loop nest that is associated with a worksharing-loop, worksharing-loop SIMD, or `simd` directive.
Syntax

The syntax of the **scan** directive and the loop body that contains it is as follows:

```c
{  
    structured-block-sequence
    #pragma omp scan clause new-line  
    structured-block-sequence
}
```

where **clause** is one of the following:

- **inclusive(list)**
- **exclusive(list)**

and where the containing loop body belongs to the innermost loop that is associated with the directive of an enclosing **for**, **for simd**, or **simd** construct.

The syntax of the **scan** directive and the loop body that contains it is as follows:

```fortran
structured-block-sequence
$omp scan clause  
structured-block-sequence
```

where **clause** is one of the following:

- **inclusive(list)**
- **exclusive(list)**

and where the containing loop body belongs to the innermost loop that is associated with the directive of an enclosing **do**, **do simd**, or **simd** construct.


**Description**

A scan directive is associated with the same worksharing-loop, worksharing-loop SIMD, or simd directive as the associated loop to which its containing loop body belongs. The directive specifies that a scan computation updates each list item on each logical iteration of the associated loops controlled by its associated directive. The directive specifies that either an inclusive scan computation is to be performed for each list item that appears in an inclusive clause on the directive, or an exclusive scan computation is to be performed for each list item that appears in an exclusive clause on the directive. For each list item for which a scan computation is specified, statements that lexically precede or follow the directive constitute one of two phases for a given logical iteration of the loop — an input phase or a scan phase.

If the list item appears in an inclusive clause, all statements in the structured block sequence that lexically precede the directive constitute the input phase and all statements in the structured block sequence that lexically follow the directive constitute the scan phase. If the list item appears in an exclusive clause, all statements in the structured block sequence that lexically precede the directive constitute the scan phase and all statements in the structured block sequence that lexically follow the directive constitute the input phase. The input phase contains all computations that update the list item in the iteration, and the scan phase ensures that any statement that reads the list item uses the result of the scan computation for that iteration.

The list items that appear in an inclusive or exclusive clause may include array sections.

The result of a scan computation for a given iteration is calculated according to the last generalized prefix sum \( \text{PRESUM}_{\text{last}} \) applied over the sequence of values given by the original value of the list item prior to the loop and all preceding updates to the list item in the logical iteration space of the loop. The operation \( \text{PRESUM}_{\text{last}}( \text{op}, a_1, \ldots, a_N ) \) is defined for a given binary operator \( \text{op} \) and a sequence of \( N \) values \( a_1, \ldots, a_N \) as follows:

- if \( N = 1 \), \( a_1 \)
- if \( N > 1 \), \( \text{op}( \text{PRESUM}_{\text{last}}( \text{op}, a_1, \ldots, a_K ), \text{PRESUM}_{\text{last}}( \text{op}, a_L, \ldots, a_N ) ) \), where \( 1 \leq K + 1 = L \leq N \).

At the beginning of the input phase of each iteration, the list item is initialized with the initializer value of the reduction-identifier specified by the reduction clause on the innermost enclosing construct. The update value of a list item is, for a given iteration, the value of the list item on completion of its input phase.

Let \( \text{orig-val} \) be the value of the original list item on entry to the enclosing worksharing-loop, worksharing-loop SIMD, or simd construct. Let \( \text{combiner} \) be the combiner for the reduction-identifier specified by the reduction clause on the construct. And let \( u_I \) be the update value of a list item for iteration \( I \). For list items that appear in an inclusive clause on the scan directive, at the beginning of the scan phase for iteration \( I \) the list item is assigned the result of the operation \( \text{PRESUM}_{\text{last}}( \text{combiner}, \text{orig-val}, u_0, \ldots, u_I ) \). For list items that appear in an exclusive clause on the scan directive, at the beginning of the scan phase for iteration \( I = 0 \) the list item is assigned the value \( \text{orig-val} \), and at the beginning of the scan phase for iteration \( I > 0 \) the list item is assigned the result of the operation \( \text{PRESUM}_{\text{last}}( \text{combiner}, \text{orig-val}, u_0, \ldots, u_{I-1} ) \).
For list items that appear in an **inclusive** clause, at the end of the enclosing worksharing-loop, worksharing-loop SIMD, or **simd** construct, the original list item is assigned the private copy from the last logical iteration of the loops associated with the enclosing construct. For list items that appear in an **exclusive** clause, let \( L \) be the last logical iteration of the loops associated with the enclosing construct. At the end of the enclosing construct, the original list item is assigned the result of the operation \( \text{PRESUM}_{last}(\text{combiner}, \text{orig-val}, u_0, \ldots, u_L) \).

**Restrictions**

Restrictions to the **scan** directive are as follows:

- Exactly one **scan** directive must be associated with a given worksharing-loop, worksharing-loop SIMD, or **simd** directive on which a **reduction** clause with the **inscan** modifier is present.
- The loops that are associated with the directive to which the **scan** directive is associated must all be perfectly nested.
- A list item that appears in the **inclusive** or **exclusive** clause must appear in a **reduction** clause with the **inscan** modifier on the associated worksharing-loop, worksharing-loop SIMD, or **simd** construct.
- Cross-iteration dependences across different logical iterations must not exist, except for dependences for the list items specified in an **inclusive** or **exclusive** clause.
- Intra-iteration dependences from a statement in the structured block sequence that precede a **scan** directive to a statement in the structured block sequence that follows a **scan** directive must not exist, except for dependences for the list items specified in an **inclusive** or **exclusive** clause.
- The private copy of list items that appear in the **inclusive** or **exclusive** clause may not be modified in the **scan phase**.

**Cross References**

- Worksharing-loop construct, see Section 2.11.4.
- **simd** construct, see Section 2.11.5.1.
- Worksharing-loop SIMD construct, see Section 2.11.5.2.
- **reduction** clause, see Section 2.21.5.4.

**2.11.9 Loop Transformation Constructs**

A loop transformation construct replaces itself, including its associated loop nest, with a structured block that may be another loop nest. If the loop transformation construct is nested inside another loop nest, its replacement becomes part of that loop nest and therefore its generated loops may become associated with another loop-associated directive that forms an enclosing construct. A loop...
transformation construct that is closely nested within another loop transformation construct applies before the enclosing loop transformation construct.

The associated loop nest of a loop transformation construct must have *canonical loop nest form* (see Section 2.11.1). All generated loops have canonical loop nest form, unless otherwise specified. Loop iteration variables of generated loops are always private in the enclosing *teams*, *parallel*, *simd*, or task generating construct.

**Cross References**
- Canonical loop nest form, see Section 2.11.1.

### 2.11.9.1 tile Construct

**Summary**
The *tile* construct tiles one or more loops.

**Syntax**

```c
#pragma omp tile sizes(size-list) new-line loop-nest
```

where `loop-nest` is a canonical loop nest and `size-list` is a list \( s_1, \ldots, s_n \) of positive integer expressions.

```fortran
!$omp tile sizes(size-list) loop-nest
[
  !$omp end tile/
]
```

where `loop-nest` is a canonical loop nest and `size-list` is a list \( s_1, \ldots, s_n \) of positive integer expressions.

If an *end tile* directive is not specified, an *end tile* directive is assumed at the end of the *do-loops*. 

Description

The tile construct controls the outer $n$ loops of the associated loop nest, where $n$ is the number of items in size-list. Let $\ell_1, \ldots, \ell_n$ be the associated loops, from outermost to innermost, which the construct replaces with a loop nest that consists of $2n$ perfectly nested loops. Let $f_1, \ldots, f_n, t_1, \ldots, t_n$ be the generated loops, from outermost to innermost. The loops $f_1, \ldots, f_n$ are the floor loops and the loops $t_1, \ldots, t_n$ are the tile loops. The tile loops do not have canonical loop nest form.

Let $\Omega$ be the logical iteration vector space of the associated loops. For any $(\alpha_1, \ldots, \alpha_n) \in \mathbb{N}^n$, define a tile $T_{\alpha_1, \ldots, \alpha_n}$ as the set of iterations

$$\{(i_1, \ldots, i_n) \in \Omega \mid \forall k \in \{1, \ldots, n\} : s_k \alpha_k \leq i_k < s_k \alpha_k + s_k\}$$

$F = \{T_{\alpha_1, \ldots, \alpha_n} \mid T_{\alpha_1, \ldots, \alpha_n} \neq \emptyset\}$ as the set of tiles with at least one iteration. Tiles that contain $\prod_{k=1}^n s_k$ iterations are complete tiles. Otherwise, they are partial tiles.

The floor loops iterate over all tiles $\{T_{\alpha_1, \ldots, \alpha_n} \in F\}$ in lexicographic order with respect to their indices $(\alpha_1, \ldots, \alpha_n)$ and the tile loops iterate over the iterations in $T_{\alpha_1, \ldots, \alpha_n}$ in the lexicographic order of the corresponding iteration vectors. An implementation may reorder the sequential execution of two iterations if at least one is from a partial tile and if their respective logical iteration vectors in loop-nest do not have a product order relation.

Restrictions

Restrictions to the tile construct are as follows:

- The depth of the associated loop nest must be greater than or equal to $n$.
- All loops that are associated with the construct must be perfectly nested.
- No loop that is associated with the construct may be a non-rectangular loop.

Cross References

- Canonical loop nest form, see Section 2.11.1.
- Worksharing-loop construct, see Section 2.11.4.
- distribute construct, see Section 2.11.6.1.
- taskloop construct, see Section 2.12.2.
2.11.9.2 unroll Construct

Summary
The \texttt{unroll} construct fully or partially unrolls a loop.

Syntax

\begin{verbatim}
C / C++
#pragma omp unroll [clause] new-line loop-nest
\end{verbatim}

The syntax of the \texttt{unroll} construct is as follows:

where \textit{loop-nest} is a canonical loop nest and \textit{clause} is one of the following:

\begin{verbatim}
full
partial[(unroll-factor)]
\end{verbatim}

where \textit{unroll-factor} is a positive integer expression that is a compile-time constant.

\begin{verbatim}
C / C++
#pragma omp unroll [clause]
loop-nest

Fortran
!$omp unroll [clause]
loop-nest

/$omp end unroll/
\end{verbatim}

The syntax of the \texttt{unroll} construct is as follows:

where \textit{loop-nest} is a canonical loop nest and \textit{clause} is one of the following:

\begin{verbatim}
full
partial[(unroll-factor)]
\end{verbatim}

where \textit{unroll-factor} is a positive integer expression that is a compile-time constant.

If an \texttt{end unroll} directive is not specified, an \texttt{end unroll} directive is assumed at the end of the \texttt{do-loop}.

Description
The \texttt{unroll} construct controls the outermost loop of the loop nest.

When the \texttt{full} clause is specified, the associated loop is \textit{fully unrolled} – it is replaced with \textit{n} instances of its loop body, one for each logical iteration of the associated loop and in the order of its logical iterations. The construct is replaced by a structured block that only contains the \textit{n} loop body instances.
When the `partial` clause is specified, the associated loop is first tiled with a tile size of `unroll-factor`. Then, the generated tile loop is fully unrolled. If the `partial` clause is used without an `unroll-factor` argument then the unroll factor is a positive integer that is implementation defined.

When neither the `full` nor the `partial` clauses are specified, if and how the loop is unrolled is implementation defined.

The `unroll` construct results in a generated loop that has canonical loop nest form if and only if the `partial` clause is specified.

**Restrictions**

Restrictions to the `unroll` construct are as follows:

- If the `full` clause is specified, the iteration count of the loop must be a compile-time constant.

**Cross References**

- Canonical loop nest form, see Section 2.11.4.
- `tile` construct, see Section 2.11.9.1.

## 2.12 Tasking Constructs

### 2.12.1 task Construct

**Summary**

The `task` construct defines an explicit task.

**Syntax**

C / C++

The syntax of the `task` construct is as follows:

```
#pragma omp task [clause[ [, ] clause] ... ] new-line
structured-block
```

where `clause` is one of the following:

- `if([ task : ] scalar-expression)`
- `final (scalar-expression)`
- `untied`
- `default (data-sharing-attribute)`
- `mergeable`
- `private (list)`
- `firstprivate (list)`
shared(list)

in_reduction(reduction-identifier : list)

depend([depend-modifier , ] dependence-type : locator-list)

priority(priority-value)

allocate([allocator : ] list)

affinity([aff-modifier : ] locator-list)

detach(event-handle)

where event-handle is a variable of omp_event_handle_t type and aff-modifier is one of the following:

iterator(iterators-definition)

The syntax of the task construct is as follows:

!$omp task [clause[ [, ] clause] ... ]

loosely-structured-block

!$omp end task

or

!$omp task [clause[ [, ] clause] ... ]

strictly-structured-block

/ !$omp end task/

where clause is one of the following:

if( [ task : ] scalar-logical-expression)

final(scalar-logical-expression)

untied

default(data-sharing-attribute)

mergeable

private(list)

firstprivate(list)

shared(list)

in_reduction(reduction-identifier : list)

depend([depend-modifier , ] dependence-type : locator-list)

priority(priority-value)
allocate([allocator : ] list)
affinity([aff-modifier : ] locator-list)
detach(event-handle)

where event-handle is an integer variable of omp_event_handle_kind kind and aff-modifier is one of the following:

iterator(iterators-definition)

Binding
The binding thread set of the task region is the current team. A task region binds to the innermost enclosing parallel region.

Description
The task construct is a task generating construct. When a thread encounters a task construct, an explicit task is generated from the code for the associated structured block. The data environment of the task is created according to the data-sharing attribute clauses on the task construct, per-data environment ICVs, and any defaults that apply. The data environment of the task is destroyed when the execution code of the associated structured block is completed.

The encountering thread may immediately execute the task, or defer its execution. In the latter case, any thread in the team may be assigned the task. Completion of the task can be guaranteed using task synchronization constructs and clauses. If a task construct is encountered during execution of an outer task, the generated task region that corresponds to this construct is not a part of the outer task region unless the generated task is an included task.

If a detach clause is present on a task construct a new allow-completion event is created and connected to the completion of the associated task region. The original event-handle is updated to represent that allow-completion event before the task data environment is created. The event-handle is considered as if it was specified on a firstprivate clause. The use of a variable in a detach clause expression of a task construct causes an implicit reference to the variable in all enclosing constructs.

If no detach clause is present on a task construct the generated task is completed when the execution of its associated structured block is completed. If a detach clause is present on a task construct, the task is completed when the execution of its associated structured block is completed and the allow-completion event is fulfilled.

When an if clause is present on a task construct and the if clause expression evaluates to false, an undeferred task is generated, and the encountering thread must suspend the current task region, for which execution cannot be resumed until execution of the structured block that is associated with the generated task is completed. The use of a variable in an if clause expression of a task construct causes an implicit reference to the variable in all enclosing constructs.
When a **final** clause is present on a **task** construct and the **final** clause expression evaluates to **true**, the generated task is a final task. All **task** constructs that are encountered during execution of a final task generate final and included tasks. The use of a variable in a **final** clause expression of a **task** construct causes an implicit reference to the variable in all enclosing constructs. Encountering a **task** construct with the **detach** clause during the execution of a final task results in unspecified behavior.

The **if** clause expression and the **final** clause expression are evaluated in the context outside of the **task** construct, and no ordering of those evaluations is specified.

A thread that encounters a task scheduling point within the **task** region may temporarily suspend the **task** region. By default, a task is tied and its suspended **task** region can only be resumed by the thread that started its execution. If the **untied** clause is present on a **task** construct, any thread in the team can resume the **task** region after a suspension. The **untied** clause is ignored if the task is a final or an included task.

The **task** construct includes a task scheduling point in the task region of its generating task, immediately following the generation of the explicit task. Each explicit **task** region includes a task scheduling point at the end of its associated structured block.

When the **mergeable** clause is present on a **task** construct, the generated task is a **mergeable task**.

The **priority** clause is a hint for the priority of the generated task. The **priority-value** is a non-negative integer expression that provides a hint for task execution order. Among all tasks ready to be executed, higher priority tasks (those with a higher numerical value in the **priority** clause expression) are recommended to execute before lower priority ones. The default **priority-value** when no **priority** clause is specified is zero (the lowest priority). If a value is specified in the **priority** clause that is higher than the **max-task-priority-var** ICV then the implementation will use the value of that ICV. A program that relies on the task execution order being determined by the **priority-value** may have unspecified behavior.

The **affinity** clause is a hint to indicate data affinity of the generated task. The task is recommended to execute close to the location of the list items. A program that relies on the task execution location being determined by this list may have unspecified behavior.

The list items that appear in the **affinity** clause may reference iterators defined by an **iterators-definition** that appears in the same clause. The list items that appear in the **affinity** clause may include array sections.

```
C / C++

The list items that appear in the **affinity** clause may use shape-operators.
```

If a list item appears in an **affinity** clause then data affinity refers to the original list item.
Note – When storage is shared by an explicit task region, the programmer must ensure, by adding proper synchronization, that the storage does not reach the end of its lifetime before the explicit task region completes its execution.

**Execution Model Events**

The task-create event occurs when a thread encounters a construct that causes a new task to be created. The event occurs after the task is initialized but before it begins execution or is deferred.

**Tool Callbacks**

A thread dispatches a registered ompt_callback_task_create callback for each occurrence of a task-create event in the context of the encountering task. This callback has the type signature ompt_callback_task_create_t and the flags argument indicates the task types shown in Table 2.7.

**TABLE 2.7: ompt_callback_task_create Callback Flags Evaluation**

<table>
<thead>
<tr>
<th>Operation</th>
<th>Evaluates to true</th>
</tr>
</thead>
<tbody>
<tr>
<td>(flags &amp; ompt_task_explicit)</td>
<td>Always in the dispatched callback</td>
</tr>
<tr>
<td>(flags &amp; ompt_task_undeferred)</td>
<td>If the task is an undeferred task</td>
</tr>
<tr>
<td>(flags &amp; ompt_task_final)</td>
<td>If the task is a final task</td>
</tr>
<tr>
<td>(flags &amp; ompt_task_untied)</td>
<td>If the task is an untied task</td>
</tr>
<tr>
<td>(flags &amp; ompt_task_mergeable)</td>
<td>If the task is a mergeable task</td>
</tr>
<tr>
<td>(flags &amp; ompt_task_merged)</td>
<td>If the task is a merged task</td>
</tr>
</tbody>
</table>

**Restrictions**

Restrictions to the task construct are as follows:

- A program must not depend on any ordering of the evaluations of the clauses of the task directive, or on any side effects of the evaluations of the clauses.
- At most one if clause can appear on the directive.
- At most one final clause can appear on the directive.
- At most one priority clause can appear on the directive.
- At most one detach clause can appear on the directive.
- If a detach clause appears on the directive, then a mergeable clause cannot appear on the same directive.
A throw executed inside a `task` region must cause execution to resume within the same `task` region, and the same thread that threw the exception must catch it.

Cross References
- Task scheduling constraints, see Section 2.12.6.
- `allocate` clause, see Section 2.13.4.
- `if` clause, see Section 2.18.
- `depend` clause, see Section 2.19.11.
- Data-sharing attribute clauses, Section 2.21.4.
- `in_reduction` clause, see Section 2.21.5.6.
- `omp_fulfill_event`, see Section 3.11.1.
- `ompt_callback_task_create_t`, see Section 4.5.2.7.

2.12.2 `taskloop` Construct

Summary
The `taskloop` construct specifies that the iterations of one or more associated loops will be executed in parallel using explicit tasks. The iterations are distributed across tasks generated by the construct and scheduled to be executed.

Syntax
```
#pragma omp taskloop [clause[, ] clause] ... ] new-line
loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` is one of the following:
```
if([ taskloop : ] scalar-expression)
shared(list)
private(list)
firstprivate(list)
lastprivate(list)
reduction([default , ] reduction-identifier : list)
in_reduction(reduction-identifier : list)
```
The syntax of the `taskloop` construct is as follows:

```
!$omp taskloop [clause[ [, ] clause] ...]
  loop-nest
/ !$omp end taskloop/
```

where `loop-nest` is a canonical loop nest and `clause` is one of the following:

```
if(/ taskloop : ) scalar-logical-expression)
shared(list)
private(list)
firstprivate(list)
lastprivate(list)
reduction([/default , ]reduction-identifier : list)
in_reduction(reduction-identifier : list)
default(data-sharing-attribute)
grainsize([/strict: ]grain-size)
num_tasks([/strict: ]num-tasks)
collapse(n)
final(scalar-logical-exp)
priority(priority-value)
untied
mergeable
```
allocate([allocator : list])

If an \texttt{end taskloop} directive is not specified, an \texttt{end taskloop} directive is assumed at the end of the do-loops.

\textbf{Binding}

The binding thread set of the \texttt{taskloop} region is the current team. A \texttt{taskloop} region binds to the innermost enclosing \texttt{parallel} region.

\textbf{Description}

The \texttt{taskloop} construct is a task generating construct. When a thread encounters a \texttt{taskloop} construct, the construct partitions the iterations of the associated loops into explicit tasks for parallel execution. The data environment of each generated task is created according to the data-sharing attribute clauses on the \texttt{taskloop} construct, per-data environment ICVs, and any defaults that apply. The order of the creation of the loop tasks is unspecified. Programs that rely on any execution order of the logical iterations are non-conforming.

By default, the \texttt{taskloop} construct executes as if it was enclosed in a \texttt{taskgroup} construct with no statements or directives outside of the \texttt{taskloop} construct. Thus, the \texttt{taskloop} construct creates an implicit \texttt{taskgroup} region. If the \texttt{nogroup} clause is present, no implicit \texttt{taskgroup} region is created.

If a \texttt{reduction} clause is present, the behavior is as if a \texttt{task_reduction} clause with the same reduction operator and list items was applied to the implicit \texttt{taskgroup} construct that encloses the \texttt{taskloop} construct. The \texttt{taskloop} construct executes as if each generated task was defined by a \texttt{task} construct on which an \texttt{in_reduction} clause with the same reduction operator and list items is present. Thus, the generated tasks are participants of the reduction defined by the \texttt{task_reduction} clause that was applied to the implicit \texttt{taskgroup} construct.

If an \texttt{in_reduction} clause is present, the behavior is as if each generated task was defined by a \texttt{task} construct on which an \texttt{in_reduction} clause with the same reduction operator and list items is present. Thus, the generated tasks are participants of a reduction previously defined by a reduction scoping clause.

If a \texttt{grainsize} clause is present, the number of logical iterations assigned to each generated task is greater than or equal to the minimum of the value of the \texttt{grain-size} expression and the number of logical iterations, but less than two times the value of the \texttt{grain-size} expression. If the \texttt{grainsize} clause has the \texttt{strict} modifier, the number of logical iterations assigned to each generated task is equal to the value of the \texttt{grain-size} expression, except for the generated task that contains the sequentially last iteration, which may have fewer iterations. The parameter of the \texttt{grainsize} clause must be a positive integer expression.

If \texttt{num_tasks} is specified, the \texttt{taskloop} construct creates as many tasks as the minimum of the \texttt{num-tasks} expression and the number of logical iterations. Each task must have at least one logical
iteration. The parameter of the `num_tasks` clause must be a positive integer expression. If the `num_tasks` clause has the `strict` modifier for a task loop with \( N \) logical iterations, the logical iterations are partitioned in a balanced manner and each partition is assigned, in order, to a generated task. The partition size is \( \lceil N/\text{num-tasks} \rceil \) until the number of remaining iterations divides the number of remaining tasks evenly, at which point the partition size becomes \( \lfloor N/\text{num-tasks} \rfloor \).

If neither a `grainsize` nor `num_tasks` clause is present, the number of loop tasks generated and the number of logical iterations assigned to these tasks is implementation defined.

The `collapse` clause may be used to specify how many loops are associated with the `taskloop` construct. The parameter of the `collapse` clause must be a constant positive integer expression. If the `collapse` clause is omitted, the behavior is as if a `collapse` clause with a parameter value of one was specified. The `collapse` clause specifies the number of loops that are collapsed into a logical iteration space that is then divided according to the `grainsize` and `num_tasks` clauses.

At the beginning of each logical iteration, the loop iteration variable or the variable declared by `range-decl` of each associated loop has the value that it would have if the set of the associated loops was executed sequentially.

The iteration count for each associated loop is computed before entry to the outermost loop. If execution of any associated loop changes any of the values used to compute any of the iteration counts, then the behavior is unspecified.

When an `if` clause is present and the `if` clause expression evaluates to `false`, undeferred tasks are generated. The use of a variable in an `if` clause expression causes an implicit reference to the variable in all enclosing constructs.

When a `final` clause is present and the `final` clause expression evaluates to `true`, the generated tasks are final tasks. The use of a variable in a `final` clause expression of a `taskloop` construct causes an implicit reference to the variable in all enclosing constructs.

When a `priority` clause is present, the generated tasks use the `priority-value` as if it was specified for each individual task. If the `priority` clause is not specified, tasks generated by the `taskloop` construct have the default task priority (zero).

When the `untied` clause is present, each generated task is an untied task.

When the `mergeable` clause is present, each generated task is a mergeable task.

\[
\begin{align*}
\text{\textbf{C++}}
\end{align*}
\]

For `firstprivate` variables of class type, the number of invocations of copy constructors that perform the initialization is implementation defined.
Note – When storage is shared by a taskloop region, the programmer must ensure, by adding proper synchronization, that the storage does not reach the end of its lifetime before the taskloop region and its descendant tasks complete their execution.

Execution Model Events
The taskloop-begin event occurs after a task encounters a taskloop construct but before any other events that may trigger as a consequence of executing the taskloop region. Specifically, a taskloop-begin event for a taskloop region will precede the taskgroup-begin that occurs unless a nogroup clause is present. Regardless of whether an implicit taskgroup is present, a taskloop-begin will always precede any task-create events for generated tasks.

The taskloop-end event occurs after a taskloop region finishes execution but before resuming execution of the encountering task.

The taskloop-iteration-begin event occurs before an explicit task executes each iteration of a taskloop region.

Tool Callbacks
A thread dispatches a registered ompt_callback_work callback for each occurrence of a taskloop-begin and taskloop-end event in that thread. The callback occurs in the context of the encountering task. The callback has type signature ompt_callback_work_t. The callback receives ompt_scope_begin or ompt_scope_end as its endpoint argument, as appropriate, and ompt_work_taskloop as its wstype argument.

A thread dispatches a registered ompt_callback_dispatch callback for each occurrence of a taskloop-iteration-begin event in that thread. The callback occurs in the context of the encountering task. The callback has type signature ompt_callback_dispatch_t.

Restrictions
Restrictions to the taskloop construct are as follows:

• If a reduction clause is present, the nogroup clause must not be specified.
• The same list item cannot appear in both a reduction and an in_reduction clause.
• At most one grainsize clause can appear on the directive.
• At most one num_tasks clause can appear on the directive.
• Neither the grainsize clause nor the num_tasks clause may appear on the directive if any of the associated loops is a non-rectangular loop.
• The grainsize clause and num_tasks clause are mutually exclusive and may not appear on the same taskloop directive.
• At most one collapse clause can appear on the directive.
• At most one if clause can appear on the directive.
• At most one final clause can appear on the directive.
• At most one priority clause can appear on the directive.

Cross References
• Canonical loop nest form, see Section 2.11.1.
• tile construct, see Section 2.11.9.1.
• task construct, Section 2.12.1.
• if clause, see Section 2.18.
• taskgroup construct, Section 2.19.6.
• Data-sharing attribute clauses, Section 2.21.4.
• ompt_scope_begin and ompt_scope_end, see Section 4.4.4.11.
• ompt_work_taskloop, see Section 4.4.4.15.
• ompt_callback_work_t, see Section 4.5.2.5.
• ompt_callback_dispatch_t, see Section 4.5.2.6.

2.12.3 taskloop simd Construct

Summary
The taskloop simd construct specifies a loop that can be executed concurrently using SIMD instructions and that those iterations will also be executed in parallel using explicit tasks. The taskloop simd construct is a composite construct.

Syntax

\begin{verbatim}
#pragma omp taskloop simd [clause[, clause] ...] new-line
\end{verbatim}

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the taskloop or simd directives with identical meanings and restrictions.
The syntax of the `taskloop simd` construct is as follows:

```
!$omp taskloop simd [clause[, clause] ...]
loop-nest
/$omp end taskloop simd/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `taskloop` or `simd` directives with identical meanings and restrictions.

If an `end taskloop simd` directive is not specified, an `end taskloop simd` directive is assumed at the end of the `do-loops`.

**Binding**

The binding thread set of the `taskloop simd` region is the current team. A `taskloop simd` region binds to the innermost enclosing parallel region.

**Description**

The `taskloop simd` construct first distributes the iterations of the associated loops across tasks in a manner consistent with any clauses that apply to the `taskloop` construct. The resulting tasks are then converted to a SIMD loop in a manner consistent with any clauses that apply to the `simd` construct, except for the `collapse` clause. For the purposes of each task’s conversion to a SIMD loop, the `collapse` clause is ignored and the effect of any `in_reduction` clause is as if a `reduction` clause with the same reduction operator and list items is present on the `simd` construct.

**Execution Model Events**

This composite construct generates the same events as the `taskloop` construct.

**Tool Callbacks**

This composite construct dispatches the same callbacks as the `taskloop` construct.

**Restrictions**

Restrictions to the `taskloop simd` construct are as follows:

- The restrictions for the `taskloop` and `simd` constructs apply.
- The `conditional` modifier may not appear in a `lastprivate` clause.
- If any `if` clause on the directive includes a `directive-name-modifier` then all `if` clauses on the directive must include a `directive-name-modifier`.
- At most one `if` clause without a `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `taskloop` `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `simd` `directive-name-modifier` can appear on the directive.
2.12.4 taskyield Construct

Summary
The taskyield construct specifies that the current task can be suspended in favor of execution of a different task. The taskyield construct is a stand-alone directive.

Syntax
The syntax of the taskyield construct is as follows:

C / C++

#pragma omp taskyield new-line

Fortran

!$omp taskyield

Binding
A taskyield region binds to the current task region. The binding thread set of the taskyield region is the current team.

Description
The taskyield region includes an explicit task scheduling point in the current task region.

Cross References
- Task scheduling, see Section 2.12.6.
2.12.5 Initial Task

Execution Model Events

No events are associated with the implicit parallel region in each initial thread.

The *initial-thread-begin* event occurs in an initial thread after the OpenMP runtime invokes the tool initializer but before the initial thread begins to execute the first OpenMP region in the initial task.

The *initial-task-begin* event occurs after an *initial-thread-begin* event but before the first OpenMP region in the initial task begins to execute.

The *initial-task-end* event occurs before an *initial-thread-end* event but after the last OpenMP region in the initial task finishes execution.

The *initial-thread-end* event occurs as the final event in an initial thread at the end of an initial task immediately prior to invocation of the tool finalizer.

Tool Callbacks

A thread dispatches a registered *ompt_callback_thread_begin* callback for the *initial-thread-begin* event in an initial thread. The callback occurs in the context of the initial thread. The callback has type signature *ompt_callback_thread_begin_t*. The callback receives *ompt_thread_initial* as its *thread_type* argument.

A thread dispatches a registered *ompt_callback_implicit_task* callback with *ompt_scope_begin* as its *endpoint* argument for each occurrence of an *initial-task-begin* event in that thread. Similarly, a thread dispatches a registered *ompt_callback_implicit_task* callback with *ompt_scope_end* as its *endpoint* argument for each occurrence of an *initial-task-end* event in that thread. The callbacks occur in the context of the initial task and have type signature *ompt_callback_implicit_task_t*. In the dispatched callback, *(flag & ompt_task_initial)* always evaluates to *true*.

A thread dispatches a registered *ompt_callback_thread_end* callback for the *initial-thread-end* event in that thread. The callback occurs in the context of the thread. The callback has type signature *ompt_callback_thread_end_t*. The implicit parallel region does not dispatch a *ompt_callback_parallel_end* callback; however, the implicit parallel region can be finalized within this *ompt_callback_thread_end* callback.

Cross References

- *ompt_thread_initial*, see Section 4.4.4.10.
- *ompt_task_initial*, see Section 4.4.4.18.
- *ompt_callback_thread_begin_t*, see Section 4.5.2.1.
- *ompt_callback_thread_end_t*, see Section 4.5.2.2.
- *ompt_callback_parallel_begin_t*, see Section 4.5.2.3.
- *ompt_callback_parallel_end_t*, see Section 4.5.2.4.
- *ompt_callback_implicit_task_t*, see Section 4.5.2.11.
2.12.6 Task Scheduling

Whenever a thread reaches a task scheduling point, the implementation may cause it to perform a
task switch, beginning or resuming execution of a different task bound to the current team. Task
scheduling points are implied at the following locations:

• during the generation of an explicit task;
• the point immediately following the generation of an explicit task;
• after the point of completion of the structured block associated with a task;
• in a taskyield region;
• in a taskwait region;
• at the end of a taskgroup region;
• in an implicit barrier region;
• in an explicit barrier region;
• during the generation of a target region;
• the point immediately following the generation of a target region;
• at the beginning and end of a target data region;
• in a target update region;
• in a target enter data region;
• in a target exit data region;
• in the omp_target_memcpy routine;
• in the omp_target_memcpy_async routine;
• in the omp_target_memcpy_rect routine; and
• in the omp_target_memcpy_rect_async routine.

When a thread encounters a task scheduling point it may do one of the following, subject to the
Task Scheduling Constraints (below):

• begin execution of a tied task bound to the current team;
• resume any suspended task region, bound to the current team, to which it is tied;
• begin execution of an untied task bound to the current team; or
• resume any suspended untied task region bound to the current team.

If more than one of the above choices is available, which one is chosen is unspecified.
Task Scheduling Constraints are as follows:

1. Scheduling of new tied tasks is constrained by the set of task regions that are currently tied to the thread and that are not suspended in a barrier region. If this set is empty, any new tied task may be scheduled. Otherwise, a new tied task may be scheduled only if it is a descendant task of every task in the set.

2. A dependent task shall not start its execution until its task dependences are fulfilled.

3. A task shall not be scheduled while any task with which it is mutually exclusive has been scheduled but has not yet completed.

4. When an explicit task is generated by a construct that contains an if clause for which the expression evaluated to false, and the previous constraints are already met, the task is executed immediately after generation of the task.

A program that relies on any other assumption about task scheduling is non-conforming.

Note – Task scheduling points dynamically divide task regions into parts. Each part is executed uninterrupted from start to end. Different parts of the same task region are executed in the order in which they are encountered. In the absence of task synchronization constructs, the order in which a thread executes parts of different schedulable tasks is unspecified.

A program must behave correctly and consistently with all conceivable scheduling sequences that are compatible with the rules above.

For example, if threadprivate storage is accessed (explicitly in the source code or implicitly in calls to library routines) in one part of a task region, its value cannot be assumed to be preserved into the next part of the same task region if another schedulable task exists that modifies it.

As another example, if a lock acquire and release happen in different parts of a task region, no attempt should be made to acquire the same lock in any part of another task that the executing thread may schedule. Otherwise, a deadlock is possible. A similar situation can occur when a critical region spans multiple parts of a task and another schedulable task contains a critical region with the same name.

The use of threadprivate variables and the use of locks or critical sections in an explicit task with an if clause must take into account that when the if clause evaluates to false, the task is executed immediately, without regard to Task Scheduling Constraint 2.

Execution Model Events

The task-schedule event occurs in a thread when the thread switches tasks at a task scheduling point; no event occurs when switching to or from a merged task.
Tool Callbacks

A thread dispatches a registered `ompt_callback_task_schedule` callback for each occurrence of a `task-schedule` event in the context of the task that begins or resumes. This callback has the type signature `ompt_callback_task_schedule_t`. The argument `prior_task_status` is used to indicate the cause for suspending the prior task. This cause may be the completion of the prior task region, the encountering of a `taskyield` construct, or the encountering of an active cancellation point.

Cross References

- `ompt_callback_task_schedule_t`, see Section 4.5.2.10.

2.13 Memory Management Directives

2.13.1 Memory Spaces

OpenMP memory spaces represent storage resources where variables can be stored and retrieved. Table 2.8 shows the list of predefined memory spaces. The selection of a given memory space expresses an intent to use storage with certain traits for the allocations. The actual storage resources that each memory space represents are implementation defined.

**Table 2.8:** Predefined Memory Spaces

<table>
<thead>
<tr>
<th>Memory space name</th>
<th>Storage selection intent</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>omp_default_mem_space</code></td>
<td>Represents the system default storage</td>
</tr>
<tr>
<td><code>omp_large_cap_mem_space</code></td>
<td>Represents storage with large capacity</td>
</tr>
<tr>
<td><code>omp_const_mem_space</code></td>
<td>Represents storage optimized for variables with constant values</td>
</tr>
<tr>
<td><code>omp_high_bw_mem_space</code></td>
<td>Represents storage with high bandwidth</td>
</tr>
<tr>
<td><code>omp_low_lat_mem_space</code></td>
<td>Represents storage with low latency</td>
</tr>
</tbody>
</table>

Variables allocated in the `omp_const_mem_space` memory space may be initialized through the `firstprivate` clause or with compile time constants for static and constant variables. Implementation-defined mechanisms to provide the constant value of these variables may also be supported.

Restrictions

Restrictions to OpenMP memory spaces are as follows:

- Variables in the `omp_const_mem_space` memory space may not be written.
2.13.2 Memory Allocators

OpenMP memory allocators can be used by a program to make allocation requests. When a memory allocator receives a request to allocate storage of a certain size, an allocation of logically consecutive memory in the resources of its associated memory space of at least the size that was requested will be returned if possible. This allocation will not overlap with any other existing allocation from an OpenMP memory allocator.

The behavior of the allocation process can be affected by the allocator traits that the user specifies. Table 2.9 shows the allowed allocator traits, their possible values and the default value of each trait.

**Table 2.9: Allocator Traits**

<table>
<thead>
<tr>
<th>Allocator trait</th>
<th>Allowed values</th>
<th>Default value</th>
</tr>
</thead>
<tbody>
<tr>
<td>sync_hint</td>
<td>contended, uncontended, serialized, private</td>
<td>contended</td>
</tr>
<tr>
<td>alignment</td>
<td>A positive integer value that is a power of 2</td>
<td>1 byte</td>
</tr>
<tr>
<td>access</td>
<td>all, cgroup, pteam, thread</td>
<td>all</td>
</tr>
<tr>
<td>pool_size</td>
<td>Positive integer value</td>
<td>Implementation defined</td>
</tr>
<tr>
<td>fallback</td>
<td>default_mem_fb, null_fb, abort_fb, allocator_fb</td>
<td>default_mem_fb</td>
</tr>
<tr>
<td>fb_data</td>
<td>an allocator handle</td>
<td>(none)</td>
</tr>
<tr>
<td>pinned</td>
<td>true, false</td>
<td>false</td>
</tr>
<tr>
<td>partition</td>
<td>environment, nearest, blocked, interleaved</td>
<td>environment</td>
</tr>
</tbody>
</table>

The `sync_hint` trait describes the expected manner in which multiple threads may use the allocator. The values and their descriptions are:

- **contended**: high contention is expected on the allocator; that is, many threads are expected to request allocations simultaneously.
- **uncontended**: low contention is expected on the allocator; that is, few threads are expected to request allocations simultaneously.
• **serialized**: only one thread at a time will request allocations with the allocator. Requesting two allocations simultaneously when specifying `serialized` results in unspecified behavior.

• **private**: the same thread will request allocations with the allocator every time. Requesting an allocation from different threads, simultaneously or not, when specifying `private` results in unspecified behavior.

Allocated memory will be byte aligned to at least the value specified for the **alignment** trait of the allocator. Some directives and API routines can specify additional requirements on alignment beyond those described in this section.

Memory allocated by allocators with the **access** trait defined to be **all** must be accessible by all threads in the device where the allocation was requested. Memory allocated by allocators with the **access** trait defined to be **cgroup** will be memory accessible by all threads in the same contention group as the thread that requested the allocation. Attempts to access the memory returned by an allocator with the **access** trait defined to be **cgroup** from a thread that is not part of the same contention group as the thread that allocated the memory result in unspecified behavior.

Memory allocated by allocators with the **access** trait defined to be **pteam** will be memory accessible by all threads that bind to the same **parallel** region of the thread that requested the allocation. Attempts to access the memory returned by an allocator with the **access** trait defined to be **pteam** from a thread that does not bind to the same **parallel** region as the thread that allocated the memory result in unspecified behavior. Memory allocated by allocators with the **access** trait defined to be **thread** will be memory accessible by the thread that requested the allocation. Attempts to access the memory returned by an allocator with the **access** trait defined to be **thread** from a thread other than the one that allocated the memory result in unspecified behavior.

The total amount of storage in bytes that an allocator can use is limited by the **pool_size** trait. For allocators with the **access** trait defined to be **all**, this limit refers to allocations from all threads that access the allocator. For allocators with the **access** trait defined to be **cgroup**, this limit refers to allocations from threads that access the allocator from the same contention group. For allocators with the **access** trait defined to be **pteam**, this limit refers to allocations from threads that access the allocator from the same parallel team. For allocators with the **access** trait defined to be **thread**, this limit refers to allocations from each thread that accesses the allocator. Requests that would result in using more storage than **pool_size** will not be fulfilled by the allocator.

The **fallback** trait specifies how the allocator behaves when it cannot fulfill an allocation request. If the **fallback** trait is set to **null_fb**, the allocator returns the value zero if it fails to allocate the memory. If the **fallback** trait is set to **abort_fb**, program execution will be terminated if the allocation fails. If the **fallback** trait is set to **allocator_fb** then when an allocation fails the request will be delegated to the allocator specified in the **fb_data** trait. If the **fallback** trait is set to **default_mem_fb** then when an allocation fails another allocation will be tried in **omp_default_mem_space**, which assumes all allocator traits to be set to their default values except for **fallback** trait, which will be set to **null_fb**.
Allocators with the \texttt{pinned} trait defined to be \texttt{true} ensure that their allocations remain in the same storage resource at the same location for their entire lifetime.

The \texttt{partition} trait describes the partitioning of allocated memory over the storage resources represented by the memory space associated with the allocator. The partitioning will be done in parts with a minimum size that is implementation defined. The values are:

- \texttt{environment}: the placement of allocated memory is determined by the execution environment;
- \texttt{nearest}: allocated memory is placed in the storage resource that is nearest to the thread that requests the allocation;
- \texttt{blocked}: allocated memory is partitioned into parts of approximately the same size with at most one part per storage resource; and
- \texttt{interleaved}: allocated memory parts are distributed in a round-robin fashion across the storage resources.

Table 2.10 shows the list of predefined memory allocators and their associated memory spaces. The predefined memory allocators have default values for their allocator traits unless otherwise specified.

\begin{table}[h]
\centering
\begin{tabular}{|l|l|l|}
\hline
Allocator name & Associated memory space & Non-default trait values \\
\hline
\texttt{omp\_default\_mem\_alloc} & \texttt{omp\_default\_mem\_space} & \texttt{fallback:}\texttt{null\_fb} \\
\texttt{omp\_large\_cap\_mem\_alloc} & \texttt{omp\_large\_cap\_mem\_space} & (none) \\
\texttt{omp\_const\_mem\_alloc} & \texttt{omp\_const\_mem\_space} & (none) \\
\texttt{omp\_high\_bw\_mem\_alloc} & \texttt{omp\_high\_bw\_mem\_space} & (none) \\
\texttt{omp\_low\_lat\_mem\_alloc} & \texttt{omp\_low\_lat\_mem\_space} & (none) \\
\texttt{omp\_cgroup\_mem\_alloc} & Implementation defined & \texttt{access:}\texttt{cgroup} \\
\texttt{omp\_pteam\_mem\_alloc} & Implementation defined & \texttt{access:}\texttt{pteam} \\
\texttt{omp\_thread\_mem\_alloc} & Implementation defined & \texttt{access:}\texttt{thread} \\
\hline
\end{tabular}
\end{table}

If any operation of the base language causes a reallocation of a variable that is allocated with a memory allocator then that memory allocator will be used to deallocate the current memory and to allocate the new memory. For allocated allocatable components of such variables, the allocator that will be used for the deallocation and allocation is unspecified.
Cross References

• **omp_init_allocator** routine, see Section 3.13.2.

• **omp_destroy_allocator** routine, see Section 3.13.3.

• **omp_set_default_allocator** routine, see Section 3.13.4.

• **omp_get_default_allocator** routine, see Section 3.13.5.

• **OMP_ALLOCATOR** environment variable, see Section 6.22.

2.13.3 allocate Directive

Summary

The *allocate* directive specifies how to allocate the specified variables. The *allocate* directive is a declarative directive if it is not associated with an allocation statement.

Syntax

```
C / C++
#pragma omp allocate(list) [clause [, clause] ... ] new-line
```

where *clause* is one of the following:

```
allocator(alloctor)
align(alignment)
```

where *allocator* is an expression of **omp_allocator_handle_t** type and *alignment* is a constant positive integer expression with a value that is a power of two.

```
C / C++
```

```
Fortran
```

The syntax of the *allocate* directive is as follows:

```
$omp allocate(list) [clause [, clause] ... ]
```

or

```
$omp allocate(list) [clause [, clause] ... ]
$omp allocate(list) [clause [, clause] ... ]
allocate-stmt
```

where *allocate-stmt* is a Fortran **ALLOCATE** statement and *clause* is one of the following:

```
allocator(alloctor)
align(alignment)
```
where \textit{allocator} is an integer expression of \texttt{omp_allocator_handle_kind \textit{kind}} and \textit{alignment} is a constant scalar positive integer expression with a value that is a power of two.

\textbf{Description}

The storage for each list item that appears in the \texttt{allocate} directive is provided by an allocation through a memory allocator. If no \texttt{allocator} clause is specified then the memory allocator specified by the \texttt{def-allocator-var} ICV is used. If the \texttt{allocator} clause is specified, the memory allocator specified in the clause is used. If the \texttt{align} clause is specified then the allocation of each list item is byte aligned to at least the maximum of the alignment required by the base language for the type of that list item, the \texttt{alignment} trait of the allocator and the \texttt{alignment} value of the \texttt{align} clause. If the \texttt{align} clause is not specified then the allocation of each list item is byte aligned to at least the maximum of the alignment required by the base language for the type of that list item and the \texttt{alignment} trait of the allocator.

The scope of this allocation is that of the list item in the base language. At the end of the scope for a given list item the memory allocator used to allocate that list item deallocates the storage.

\textbf{Restrictions}

Restrictions to the \texttt{allocate} directive are as follows:

- At most one \texttt{allocator} clause may appear on the directive.
- At most one \texttt{align} clause may appear on the directive.
- A variable that is part of another variable (as an array or structure element) cannot appear in a declarative \texttt{allocate} directive.
- A declarative \texttt{allocate} directive must appear in the same scope as the declarations of each of its list items and must follow all such declarations.
- A declared variable may appear as a list item in at most one declarative \texttt{allocate} directive in a given compilation unit.
- At most one \texttt{allocator} clause can appear on the \texttt{allocate} directive.
• allocate directives that appear in a target region must specify an allocator clause unless a requires directive with the dynamic allocators clause is present in the same compilation unit.

- If a list item has static storage duration, the allocator clause must be specified and the allocator expression in the clause must be a constant expression that evaluates to one of the predefined memory allocator values.

- A variable that is declared in a namespace or global scope may only appear as a list item in an allocate directive if an allocate directive that lists the variable follows a declaration that defines the variable and if all allocate directives that list the variable specify the same allocator.

- After a list item has been allocated, the scope that contains the allocate directive must not end abnormally, such as through a call to the longjmp function.

- After a list item has been allocated, the scope that contains the allocate directive must not end abnormally, such as through a call to the longjmp function, other than through C++ exceptions.

- A variable that has a reference type may not appear as a list item in an allocate directive.

- An allocate directive that is associated with an allocate-stmt and specifies a list must be preceded by an executable statement or OpenMP construct.

- A list item that is specified in a declarative allocate directive must not have the ALLOCATABLE or POINTER attribute.

- If a list item is specified in an allocate directive that is associated with an allocate-stmt, it must appear as one of the data objects in the allocation list of that statement.

- A list item may not be specified more than once in an allocate directive that is associated with an allocate-stmt.

- If multiple directives are associated with an allocate-stmt then at most one directive may specify no list items.

- If a list item has the SAVE attribute, either explicitly or implicitly, or is a common block name then the allocator clause must be specified and only predefined memory allocator parameters can be used in the clause.

- A variable that is part of a common block may not be specified as a list item in an allocate directive, except implicitly via the named common block.
• A named common block may appear as a list item in at most one allocate directive in a given compilation unit.

• If a named common block appears as a list item in an allocate directive, it must appear as a list item in an allocate directive that specifies the same allocator in every compilation unit in which the common block is used.

• An associate name may not appear as a list item in an allocate directive.

• A type parameter inquiry cannot appear in an allocate directive.

Cross References
- def-allocator-var ICV, see Section 2.4.1.
- Memory allocators, see Section 2.13.2.
- omp_allocator_handle_t and omp_allocator_handle_kind, see Section 3.13.1.

2.13.4 allocate Clause

Summary
The allocate clause specifies the memory allocator to be used to obtain storage for private variables of a directive.

Syntax
The syntax of the allocate clause is one of the following:

allocate([allocator:] list)
allocate(allocate-modifier [, allocate-modifier]: list)

where allocate-modifier is one of the following:

allocator(allocator)
align(alignment)

where alignment is a constant positive integer expression with a value that is a power of two; and

where allocator is an expression of the omp_allocator_handle_t type.

where allocator is an integer expression of the omp_allocator_handle_kind kind.
Description
The storage for new list items that arise from list items that appear in the directive is provided through a memory allocator. If an allocator is specified in the clause, that allocator is used for allocations. If no allocator is specified in the clause and the directive is not a target directive then the memory allocator that is specified by the def-allocator-var ICV is used for the list items that are specified in the allocate clause. If no allocator is specified in the clause and the directive is a target directive the behavior is unspecified. If the align allocate-modifier is specified then the allocation of each list item is byte aligned to at least the maximum of the alignment required by the base language for the type of that list item, the alignment trait of the allocator and the alignment value of the align allocate-modifier. If the align allocate-modifier is not specified then the allocation of each list item is byte aligned to at least the maximum of the alignment required by the base language for the type of that list item and the alignment trait of the allocator.

For allocations that arise from this clause the null_fb value of the fallback allocator trait behaves as if the abort_fb had been specified.

Restrictions
Restrictions to the allocate clause are as follows:

- At most one allocator allocate-modifier may be specified on the clause.
- At most one align allocate-modifier may be specified on the clause.
- For any list item that is specified in the allocate clause on a directive, a data-sharing attribute clause that may create a private copy of that list item must be specified on the same directive.
- For task, taskloop or target directives, allocation requests to memory allocators with the trait access set to thread result in unspecified behavior.
- allocate clauses that appear on a target construct or on constructs in a target region must specify an allocator expression unless a requires directive with the dynamic_allocators clause is present in the same compilation unit.

Cross References
- def-allocator-var ICV, see Section 2.4.1.
- Memory allocators, see Section 2.13.2.
- List Item Privatization, see Section 2.21.3.
- omp_allocator_handle_t and omp_allocator_handle_kind, see Section 3.13.1.
2.14 Device Directives

2.14.1 Device Initialization

Execution Model Events

The device-initialize event occurs in a thread that encounters the first target, target data, or target enter data construct or a device memory routine that is associated with a particular target device after the thread initiates initialization of OpenMP on the device and the device’s OpenMP initialization, which may include device-side tool initialization, completes.

The device-load event for a code block for a target device occurs in some thread before any thread executes code from that code block on that target device.

The device-unload event for a target device occurs in some thread whenever a code block is unloaded from the device.

The device-finalize event for a target device that has been initialized occurs in some thread before an OpenMP implementation shuts down.

Tool Callbacks

A thread dispatches a registered ompt_callback_device_initialize callback for each occurrence of a device-initialize event in that thread. This callback has type signature ompt_callback_device_initialize_t.

A thread dispatches a registered ompt_callback_device_load callback for each occurrence of a device-load event in that thread. This callback has type signature ompt_callback_device_load_t.

A thread dispatches a registered ompt_callback_device_unload callback for each occurrence of a device-unload event in that thread. This callback has type signature ompt_callback_device_unload_t.

A thread dispatches a registered ompt_callback_device_finalize callback for each occurrence of a device-finalize event in that thread. This callback has type signature ompt_callback_device_finalize_t.

Restrictions

Restrictions to OpenMP device initialization are as follows:

- No thread may offload execution of an OpenMP construct to a device until a dispatched ompt_callback_device_initialize callback completes.
- No thread may offload execution of an OpenMP construct to a device after a dispatched ompt_callback_device_finalize callback occurs.
Cross References

- ompt_callback_device_initialize_t, see Section 4.5.2.19.
- ompt_callback_device_finalize_t, see Section 4.5.2.20.
- ompt_callback_device_load_t, see Section 4.5.2.21.
- ompt_callback_device_unload_t, see Section 4.5.2.22.

2.14.2 target data Construct

Summary

The target data construct maps variables to a device data environment for the extent of the region.

Syntax

C / C++

The syntax of the target data construct is as follows:

```
#pragma omp target data clause[ [ , ] clause]... ] new-line
  structured-block
```

where clause is one of the following:

```
if([ target data : ] scalar-expression)
device(integer-expression)
map([ [map-type-modifier[, ] [map-type-modifier[, ] ...]] map-type : ] locator-list)
use_device_ptr(list)
use_device_addr(list)
```

Fortran

The syntax of the target data construct is as follows:

```
!$omp target data clause[ [ , ] clause]... ]
  loosely-structured-block
!$omp end target data
```

or

```
!$omp target data clause[ [ , ] clause]... ]
  strictly-structured-block
!$omp end target data/
```
where clause is one of the following:

```
if(/ target data :/ scalar-logical-expression)
device(scalar-integer-expression)
map(/[[map-type-modifier[, ] [map-type-modifier[, ] ...]] map-type : ] locator-list)
use_device_ptr(list)
use_device_addr(list)
```

### Fortran

**Binding**

The binding task set for a target data region is the generating task. The target data region binds to the region of the generating task.

**Description**

When a target data construct is encountered, the encountering task executes the region. If no device clause is present, the behavior is as if the device clause appeared with an expression equal to the value of the default-device-var ICV. When an if clause is present and the if clause expression evaluates to false, the target device is the host. Variables are mapped for the extent of the region, according to any data-mapping attribute clauses, from the data environment of the encountering task to the device data environment.

If a list item that appears in a use_device_addr clause has corresponding storage in the device data environment, references to the list item in the associated structured block are converted into references to the corresponding list item. If the list item is not a mapped list item, it is assumed to be accessible on the target device. Inside the structured block, the list item has a device address and its storage may not be accessible from the host device. The list items that appear in a use_device_addr clause may include array sections.

**C / C++**

If a list item in a use_device_addr clause is an array section that has a base pointer, the effect of the clause is to convert the base pointer to a pointer that is local to the structured block and that contains the device address. This conversion may be elided if the list item was not already mapped.

If a list item that appears in a use_device_ptr clause is a pointer to an object that is mapped to the device data environment, references to the list item in the associated structured block are converted into references to a device pointer that is local to the structured block and that refers to the device address of the corresponding object. If the list item does not point to a mapped object, it must contain a valid device address for the target device, and the list item references are instead converted to references to a local device pointer that refers to this device address.
If a list item that appears in a use_device_ptr clause is of type C_PTR and points to a data entity that is mapped to the device data environment, references to the list item in the associated structured block are converted into references to a device pointer that is local to the structured block and that refers to the device address of the corresponding entity. If a list item of type C_PTR does not point to a mapped object, it must contain a valid device address for the target device, and the list item references are instead converted to references to a local device pointer that refers to this device address.

If a list item in a use_device_ptr clause is not of type C_PTR, the behavior is as if the list item appeared in a use_device_addr clause. Support for such list items in a use_device_ptr clause is deprecated.

If one or more map clauses are present, the list item conversions that are performed for any use_device_ptr or use_device_addr clause occur after all variables are mapped on entry to the region according to those map clauses.

**Execution Model Events**
The events associated with entering a target data region are the same events as associated with a target enter data construct, as described in Section 2.14.3.

The events associated with exiting a target data region are the same events as associated with a target exit data construct, as described in Section 2.14.4.

**Tool Callbacks**
The tool callbacks dispatched when entering a target data region are the same as the tool callbacks dispatched when encountering a target enter data construct, as described in Section 2.14.3.

The tool callbacks dispatched when exiting a target data region are the same as the tool callbacks dispatched when encountering a target exit data construct, as described in Section 2.14.4.

**Restrictions**
Restrictions to the target data construct are as follows:

- A program must not depend on any ordering of the evaluations of the clauses of the target data directive, except as explicitly stated for map clauses and for map clauses relative to use_device_ptr and use_device_addr clauses, or on any side effects of the evaluations of the clauses.

- At most one device clause can appear on the directive. The device clause expression must evaluate to a non-negative integer value that is less than or equal to the value of omp_get_num_devices().
• At most one `if` clause can appear on the directive.

• A `map-type` in a `map` clause must be `to`, `from`, `tofrom` or `alloc`.

• At least one `map`, `use_device_addr` or `use_device_ptr` clause must appear on the directive.

• A list item may not be specified more than once in `use_device_ptr` clauses that appear on the directive.

• A list item may not be specified more than once in `use_device_addr` clauses that appear on the directive.

• A list item may not be specified in both a `use_device_addr` clause and a `use_device_ptr` clause on the directive.

• A list item in a `use_device_addr` clause must have a corresponding list item in the device data environment or be accessible on the target device.

• A list item that appears in a `use_device_ptr` or `use_device_addr` clause must not be a structure element.

• A list item in a `use_device_ptr` clause must be a pointer for which the value is the address of an object that has corresponding storage in the device data environment or is accessible on the target device.

• If a list item in a `use_device_addr` clause is an array section, the base expression must be a base language identifier.

• The value of a list item in a `use_device_ptr` clause that is of type `C_PTR` must be the address of a data entity that has corresponding storage in the device data environment or is accessible on the target device.

• If a list item in a `use_device_addr` clause is an array section, the designator of the base expression must be a name without any selectors.

**Cross References**

• `default-device-var`, see Section 2.4.

• `if` clause, see Section 2.18.

• `map` clause, see Section 2.21.7.1.

• `omp_get_num_devices` routine, see Section 3.7.4.
2.14.3 target enter data Construct

Summary
The target enter data directive specifies that variables are mapped to a device data environment. The target enter data directive is a stand-alone directive.

Syntax

C / C++

The syntax of the target enter data construct is as follows:

```
#pragma omp target enter data [clause [ , ] clause] ...
```

where clause is one of the following:

```
if([ target enter data : ] scalar-expression)
device(integer-expression)
map([map-type-modifier [ , ] [map-type-modifier [ , ] ...]] map-type : locator-list)
depend([depend-modifier [ , ] dependence-type : locator-list])
nowait
```

Fortran

The syntax of the target enter data is as follows:

```
!$omp target enter data [clause [ , ] clause] ...
```

where clause is one of the following:

```
if([ target enter data : ] scalar-logical-expression)
device(scalar-integer-expression)
map([map-type-modifier [ , ] [map-type-modifier [ , ] ...]] map-type : locator-list)
depend([depend-modifier [ , ] dependence-type : locator-list])
nowait
```

Binding
The binding task set for a target enter data region is the generating task, which is the target task generated by the target enter data construct. The target enter data region binds to the corresponding target task region.
Description

When a \texttt{target enter data} construct is encountered, the list items are mapped to the device data environment according to the \texttt{map} clause semantics.

The \texttt{target enter data} construct is a task generating construct. The generated task is a \texttt{target task}. The generated task region encloses the \texttt{target enter data} region.

All clauses are evaluated when the \texttt{target enter data} construct is encountered. The data environment of the \texttt{target task} is created according to the data-mapping attribute clauses on the \texttt{target enter data} construct, per-data environment ICVs, and any default data-sharing attribute rules that apply to the \texttt{target enter data} construct. If a variable or part of a variable is mapped by the \texttt{target enter data} construct, the variable has a default data-sharing attribute of shared in the data environment of the \texttt{target task}.

Assignment operations associated with mapping a variable (see Section 2.21.7.1) occur when the \texttt{target task} executes.

If the \texttt{nowait} clause is present, execution of the \texttt{target task} may be deferred. If the \texttt{nowait} clause is not present, the \texttt{target task} is an included task.

If a \texttt{depend} clause is present, it is associated with the \texttt{target task}.

If no \texttt{device} clause is present, the behavior is as if the \texttt{device} clause appears with an expression equal to the value of the \texttt{default-device-var} ICV.

When an \texttt{if} clause is present and the \texttt{if} clause expression evaluates to \texttt{false}, the target device is the host.

Execution Model Events

Events associated with a \texttt{target task} are the same as for the \texttt{task} construct defined in Section 2.12.1.

The \texttt{target-enter-data-begin} event occurs when a thread enters a \texttt{target enter data} region.

The \texttt{target-enter-data-end} event occurs when a thread exits a \texttt{target enter data} region.

Tool Callbacks

Callbacks associated with events for \texttt{target tasks} are the same as for the \texttt{task} construct defined in Section 2.12.1; \texttt{(flags & ompt_task_target)} always evaluates to \texttt{true} in the dispatched callback.

A thread dispatches a registered \texttt{ompt_callback_target} or \texttt{ompt_callback_target_emis} callback with \texttt{ompt_scope_begin} as its \texttt{endpoint} argument and \texttt{ompt_target_enter_data} or \texttt{ompt_target_enter_data_nowait} if the \texttt{nowait} clause is present as its \texttt{kind} argument for each occurrence of a \texttt{target-enter-data-begin} event in that thread in the context of the target task on the host. Similarly, a thread dispatches a registered \texttt{ompt_callback_target} or \texttt{ompt_callback_target_emis} callback with \texttt{ompt_scope_end} as its \texttt{endpoint} argument and \texttt{ompt_target_enter_data} or
ompt_target_enter_data_nowait if the nowait clause is present as its kind argument for each occurrence of a target-enter-data-end event in that thread in the context of the target task on the host. These callbacks have type signature ompt_callback_target_t or ompt_callback_target_emi_t, respectively.

Restrictions
Restrictions to the target enter data construct are as follows:

- A program must not depend on any ordering of the evaluations of the clauses of the target enter data directive, or on any side effects of the evaluations of the clauses.
- At least one map clause must appear on the directive.
- At most one device clause can appear on the directive. The device clause expression must evaluate to a non-negative integer value that is less than or equal to the value of omp_get_num_devices().
- At most one if clause can appear on the directive.
- A map-type must be specified in all map clauses and must be either to or alloc.
- At most one nowait clause can appear on the directive.

Cross References
- default-device-var, see Section 2.4.1.
- task, see Section 2.12.1.
- task scheduling constraints, see Section 2.12.6.
- target data, see Section 2.14.2.
- target exit data, see Section 2.14.4.
- if clause, see Section 2.18.
- map clause, see Section 2.21.7.1.
- omp_get_num_devices routine, see Section 3.7.4.
- ompt_callback_target_t and ompt_callback_target_emi_t callback type, see Section 4.5.2.26.

2.14.4 target exit data Construct

Summary
The target exit data directive specifies that list items are unmapped from a device data environment. The target exit data directive is a stand-alone directive.
Syntax

The syntax of the **target exit data** construct is as follows:
```
#pragma omp target exit data [clause [, clause] ... ] new-line
```
where `clause` is one of the following:
```
if([ target exit data : ] scalar-expression)
device(integer-expression)
map([map-type-modifier, ] [map-type-modifier[ , ] ...] map-type : locator-list)
depend([depend-modifier, ] dependence-type : locator-list)
nowait
```

C / C++

The syntax of the **target exit data** is as follows:
```
!$omp target exit data [clause [, clause] ... ]
```
where clause is one of the following:
```
if([ target exit data : ] scalar-logical-expression)
device(scalar-integer-expression)
map([map-type-modifier[ ] [map-type-modifier[ , ] ...] map-type : locator-list)
depend([depend-modifier, ] dependence-type : locator-list)
nowait
```

Fortran

Binding

The binding task set for a **target exit data** region is the generating task, which is the **target task** generated by the **target exit data** construct. The **target exit data** region binds to the corresponding **target task** region.

Description

When a **target exit data** construct is encountered, the list items in the **map** clauses are unmapped from the device data environment according to the **map** clause semantics.

The **target exit data** construct is a task generating construct. The generated task is a **target task**. The generated task region encloses the **target exit data** region.
All clauses are evaluated when the **target exit data** construct is encountered. The data environment of the **target task** is created according to the data-mapping attribute clauses on the **target exit data** construct, per-data environment ICVs, and any default data-sharing attribute rules that apply to the **target exit data** construct. If a variable or part of a variable is mapped by the **target exit data** construct, the variable has a default data-sharing attribute of shared in the data environment of the **target task**.

Assignment operations associated with mapping a variable (see Section 2.21.7.1) occur when the **target task** executes.

If the **nowait** clause is present, execution of the **target task** may be deferred. If the **nowait** clause is not present, the **target task** is an included task.

If a **depend** clause is present, it is associated with the **target task**.

If no **device** clause is present, the behavior is as if the **device** clause appears with an expression equal to the value of the **default-device-var** ICV.

When an **if** clause is present and the **if** clause expression evaluates to **false**, the target device is the host.

**Execution Model Events**

Events associated with a **target task** are the same as for the **task** construct defined in Section 2.12.1.

The **target-exit-data-begin** event occurs when a thread enters a **target exit data** region.

The **target-exit-data-end** event occurs when a thread exits a **target exit data** region.

**Tool Callbacks**

Callbacks associated with events for **target tasks** are the same as for the **task** construct defined in Section 2.12.1; (**flags & ompt_task_target**) always evaluates to **true** in the dispatched callback.

A thread dispatches a registered **ompt_callback_target** or **ompt_callback_target_emii** callback with **ompt_scope_begin** as its **endpoint** argument and **ompt_target_exit_data** or **ompt_target_exit_data_nowait** if the **nowait** clause is present as its **kind** argument for each occurrence of a **target-exit-data-begin** event in that thread in the context of the target task on the host. Similarly, a thread dispatches a registered **ompt_callback_target** or **ompt_callback_target_emii** callback with **ompt_scope_end** as its **endpoint** argument and **ompt_target_exit_data** or **ompt_target_exit_data_nowait** if the **nowait** clause is present as its **kind** argument for each occurrence of a **target-exit-data-end** event in that thread in the context of the target task on the host. These callbacks have type signature **ompt_callback_target_t** or **ompt_callback_target_emii_t**, respectively.
**Restrictions**

Restrictions to the `target exit data` construct are as follows:

- A program must not depend on any ordering of the evaluations of the clauses of the `target exit data` directive, or on any side effects of the evaluations of the clauses.
- At least one `map` clause must appear on the directive.
- At most one `device` clause can appear on the directive. The `device` clause expression must evaluate to a non-negative integer value that is less than or equal to the value of `omp_get_num_devices()`.
- At most one `if` clause can appear on the directive.
- A `map-type` must be specified in all `map` clauses and must be either `from`, `release`, or `delete`.
- At most one `nowait` clause can appear on the directive.

**Cross References**

- `default-device-var`, see Section 2.4.1.
- `task`, see Section 2.12.1.
- `task scheduling constraints`, see Section 2.12.6.
- `target data`, see Section 2.14.2.
- `target enter data`, see Section 2.14.3.
- `if` clause, see Section 2.18.
- `map` clause, see Section 2.21.7.1.
- `omp_get_num_devices` routine, see Section 3.7.4.
- `ompt_callback_target_t` and `ompt_callback_target_emi_t` callback type, see Section 4.5.2.26.
2.14.5 target Construct

Summary

The target construct maps variables to a device data environment and executes the construct on that device.

Syntax

The syntax of the target construct is as follows:

```c
#pragma omp target [clause[ [, ] clause] ...] new-line
structured-block
```

where clause is one of the following:

- `if([ target : ] scalar-expression)
- `device([ device-modifier : ] integer-expression)
- `thread_limit(integer-expression)
- `private(list)
- `firstprivate(list)
- `in_reduction(reduction-identifier : list)
- `map([map-type-modifier[, ] [map-type-modifier[, ] ...]] map-type : ] locator-list)
- `is_device_ptr(list)
- `has_device_addr(list)
- `defaultmap(in_reduction-implicit-behavior[:variable-category])
- `nowait
- `depend([depend-modifier, ] dependence-type : locator-list)
- `allocate([allocator : ] list)
- `uses_allocators(allocator[ (allocator-traits-array) ]
  [, allocator[ (allocator-traits-array) ] ...])
```

and where device-modifier is one of the following:

- `ancestor
- `device_num

and where allocator is an identifier of omp_allocator_handle_t type and allocator-traits-array is an identifier of const omp_alloctrait_t type.
The syntax of the `target` construct is as follows:

```fortran
!$omp target [clause [ , clause] ... ]
loosely-structured-block
!

or

!$omp target [clause [ , clause] ... ]
strictly-structured-block
!

where `clause` is one of the following:

```fortran
if([ target : ] scalar-logical-expression)
device([ device-modifier : ] scalar-integer-expression)
thread_limit(scalar-integer-expression)
private(list)
firstprivate(list)
in_reduction(reduction-identifier : list)
map([ [map-type-modifier[ , ] [map-type-modifier[ , ] ...]] map-type : ] locator-list)
is_device_ptr(list)
has_device_addr(list)
defaultmap(implicit-behavior[:variable-category])
nowait
depend([[depend-modifier, ] dependence-type : locator-list])
allocate([[[allocator : ] list])
uses Allocators(allocator[ (allocator-traits-array) ]

[,] allocator[ (allocator-traits-array) ] ...])
```

and where `device-modifier` is one of the following:

```fortran
ancestor
device_num
```

and where `allocator` is an integer expression of `omp_allocator_handle_kind kind` and `allocator-traits-array` is an array of `type(omp_alloctrait)` type.
Binding

The binding task set for a target region is the generating task, which is the target task generated by the target construct. The target region binds to the corresponding target task region.

Description

The target construct provides a superset of the functionality provided by the target data directive, except for the use_device_ptr and use_device_addr clauses.

The functionality added to the target directive is the inclusion of an executable region to be executed on a device. That is, the target directive is an executable directive.

The target construct is a task generating construct. The generated task is a target task. The generated task region encloses the target region. The device clause determines the device on which the target region executes.

All clauses are evaluated when the target construct is encountered. The data environment of the target task is created according to the data-sharing and data-mapping attribute clauses on the target construct, per-data environment ICVs, and any default data-sharing attribute rules that apply to the target construct. If a variable or part of a variable is mapped by the target construct and does not appear as a list item in an in_reduction clause on the construct, the variable has a default data-sharing attribute of shared in the data environment of the target task.

Assignment operations associated with mapping a variable (see Section 2.21.7.1) occur when the target task executes.

As described in Section 2.4.4.1, the target construct limits the number of threads that may participate in a contention group initiated by the initial thread by setting the value of the thread-limit-var ICV for the initial task to an implementation defined value greater than zero. If the thread_limit clause is specified, the number of threads will be less than or equal to the value specified in the clause.

If a device clause in which the device_num device-modifier appears is present on the construct, the device clause expression specifies the device number of the target device. If device-modifier does not appear in the clause, the behavior of the clause is as if device-modifier is device_num.

If a device clause in which the ancestor device-modifier appears is present on the target construct and the device clause expression evaluates to 1, execution of the target region occurs on the parent device of the enclosing target region. If the target construct is not encountered in a target region, the current device is treated as the parent device. The encountering thread waits for completion of the target region on the parent device before resuming. For any list item that appears in a map clause on the same construct, if the corresponding list item exists in the device data environment of the parent device, it is treated as if it has a reference count of positive infinity.

If no device clause is present, the behavior is as if the device clause appears without a device-modifier and with an expression equal to the value of the default-device-var ICV.
If the `nowait` clause is present, execution of the target task may be deferred. If the `nowait` clause is not present, the target task is an included task.

If a `depend` clause is present, it is associated with the target task.

When an `if` clause is present and the `if` clause expression evaluates to `false`, the target region is executed by the host device.

The `has_device_addr` clause indicates that its list items already have device addresses and therefore they may be directly accessed from a target device. If the device address of a list item is not for the device on which the target region executes, accessing the list item inside the region results in unspecified behavior. The list items may include array sections.

---

### C / C++

The `is_device_ptr` clause indicates that its list items are device pointers. Inside the target construct, each list item is privatized and the new list item is initialized to the device address to which the original list item refers. Support for device pointers created outside of OpenMP, specifically outside of any OpenMP mechanism that returns a device pointer, is implementation defined.

---

### Fortran

The `is_device_ptr` clause indicates that its list items are device pointers if they are of type `C_PTR`. Inside the target construct, each list item of type `C_PTR` is privatized and the new list item is initialized to the device address to which the original list item refers. Support for device pointers created outside of OpenMP, specifically outside of any OpenMP mechanism that returns a device pointer, is implementation defined. For any list item in an `is_device_ptr` clause that is not of type `C_PTR`, the behavior is as if the list item appeared in a `has_device_addr` clause. Support for such list items in an `is_device_ptr` clause is deprecated.

---

If a function (C, C++, Fortran) or subroutine (Fortran) is referenced in a target construct that does not specify a `device` clause in which the ancestor device-modifier appears then that function or subroutine is treated as if its name had appeared in a `to` clause on a declare target directive.

If a variable with static storage duration is declared in a target construct that does not specify a `device` clause in which the ancestor device-modifier appears then the named variable is treated as if it had appeared in a `to` clause on a declare target directive.

Each memory allocator specified in the `uses_allocators` clause will be made available in the target region. For each non-predefined allocator that is specified, a new allocator handle will be associated with an allocator that is created with the specified traits as if by a call to `omp_init_allocator` at the beginning of the target region. Each non-predefined allocator will be destroyed as if by a call to `omp_destroy_allocator` at the end of the target region.
If a list item in a `map` clause has a base pointer and it is a scalar variable with a predetermined data-sharing attribute of firstprivate (see Section 2.21.1.1), then on entry to the `target` region:

- If the list item is not a zero-length array section, the corresponding private variable is initialized such that the corresponding list item in the device data environment can be accessed through the pointer in the `target` region.
- If the list item is a zero-length array section, the corresponding private variable is initialized according to Section 2.21.7.2.

When an internal procedure is called in a `target` region, any references to variables that are host associated in the procedure have unspecified behavior.

**Execution Model Events**

Events associated with a `target` task are the same as for the `task` construct defined in Section 2.12.1.

Events associated with the initial task that executes the `target` region are defined in Section 2.12.5.

The `target-begin` event occurs when a thread enters a `target` region.

The `target-end` event occurs when a thread exits a `target` region.

The `target-submit-begin` event occurs prior to initiating creation of an initial task on a target device for a `target` region.

The `target-submit-end` event occurs after initiating creation of an initial task on a target device for a `target` region.

**Tool Callbacks**

Callbacks associated with events for `target` tasks are the same as for the `task` construct defined in Section 2.12.1; `(flags & ompt_task_target)` always evaluates to `true` in the dispatched callback.

A thread dispatches a registered `ompt_callback_target` or `ompt_callback_target_emi` callback with `ompt_scope_begin` as its `endpoint` argument and `ompt_target` or `ompt_target_nowait` if the `nowait` clause is present as its `kind` argument for each occurrence of a `target-begin` event in that thread in the context of the target task on the host. Similarly, a thread dispatches a registered `ompt_callback_target` or `ompt_callback_target_emi` callback with `ompt_scope_end` as its `endpoint` argument and `ompt_target` or `ompt_target_nowait` if the `nowait` clause is present as its `kind` argument.
argument for each occurrence of a target-end event in that thread in the context of the target task on
the host. These callbacks have type signature ompt_callback_target_t or
ompt_callback_target_emi_t, respectively.

A thread dispatches a registered ompt_callback_target_submit_emi callback with
ompt_scope_begin as its endpoint argument for each occurrence of a target-submit-begin
event in that thread. Similarly, a thread dispatches a registered
ompt_callback_target_submit_emi callback with ompt_scope_end as its endpoint
argument for each occurrence of a target-submit-end event in that thread. These callbacks have type
signature ompt_callback_target_submit_emi_t.

A thread dispatches a registered ompt_callback_target_submit callback for each
occurrence of a target-submit-begin event in that thread. The callback occurs in the context of the
target task and has type signature ompt_callback_target_submit_t.

Restrictions
Restrictions to the target construct are as follows:

- If a target update, target data, target enter data, or target exit data
  construct is encountered during execution of a target region, the behavior is unspecified.

- The result of an omp_set_default_device, omp_get_default_device, or
  omp_get_num_devices routine called within a target region is unspecified.

- The effect of an access to a threadprivate variable in a target region is unspecified.

- If a list item in a map clause is a structure element, any other element of that structure that is
  referenced in the target construct must also appear as a list item in a map clause.

- A list item cannot appear in both a map clause and a data-sharing attribute clause on the same
  target construct.

- A variable referenced in a target region but not the target construct that is not declared in
  the target region must appear in a declare target directive.

- At most one defaultmap clause for each category can appear on the directive.

- At most one nowait clause can appear on the directive.

- At most one if clause can appear on the directive.

- A map-type in a map clause must be to, from, tofrom or alloc.

- A list item that appears in an is_device_ptr clause must be a valid device pointer for the
device data environment.

- A list item that appears in a has_device_addr clause must have a valid device address for
  the device data environment.

- A list item may not be specified in both an is_device_ptr clause and a
  has_device_addr clause on the directive.
• A list item that appears in an `is_device_ptr` or a `has_device_addr` clause must not be specified in any data-sharing attribute clause on the same `target` construct.

• At most one `device` clause can appear on the directive. The `device` clause expression must evaluate to a non-negative integer value that is less than or equal to the value of `omp_get_num_devices()`.

• If a `device` clause in which the `ancestor` device-modifier appears is present on the construct, then the following restrictions apply:
  – A `requires` directive with the `reverse_offload` clause must be specified;
  – The `device` clause expression must evaluate to 1;
  – Only the `device, firstprivate, private, defaultmap, and map` clauses may appear on the construct;
  – No OpenMP constructs or calls to OpenMP API runtime routines are allowed inside the corresponding `target` region.

• Memory allocators that do not appear in a `uses Allocators` clause cannot appear as an allocator in an `allocate` clause or be used in the `target` region unless a `requires` directive with the `dynamic Allocators` clause is present in the same compilation unit.

• Memory allocators that appear in a `uses Allocators` clause cannot appear in other data-sharing attribute clauses or data-mapping attribute clauses in the same construct.

• Predefined allocators appearing in a `uses Allocators` clause cannot have `traits` specified.

• Non-predefined allocators appearing in a `uses Allocators` clause must have `traits` specified.

• Arrays that contain allocator traits that appear in a `uses Allocators` clause must be constant arrays, have constant values and be defined in the same scope as the construct in which the clause appears.

• Any IEEE floating-point exception status flag, halting mode, or rounding mode set prior to a `target` region is unspecified in the region.

• Any IEEE floating-point exception status flag, halting mode, or rounding mode set in a `target` region is unspecified upon exiting the region.

• An attached pointer must not be modified in a `target` region.

• A list item that appears in an `is_device_ptr` clause must have a type of pointer or array.
• A list item that appears in an `is_device_ptr` clause must have a type of pointer, array, reference to pointer or reference to array.

• A `throw` executed inside a `target` region must cause execution to resume within the same `target` region, and the same thread that threw the exception must catch it.

• The run-time type information (RTTI) of an object can only be accessed from the device on which it was constructed.

• An attached pointer that is associated with a given pointer target must not become associated with a different pointer target in a `target` region.

• If a list item in a `map` clause is an array section, and the array section is derived from a variable with a `POINTER` or `ALLOCATABLE` attribute then the behavior is unspecified if the corresponding list item’s variable is modified in the region.

• A reference to a coarray that is encountered on a non-host device must not be coindexed or appear as an actual argument to a procedure where the corresponding dummy argument is a coarray.

• If the allocation status of a mapped variable that has the `ALLOCATABLE` attribute is unallocated on entry to a `target` region, the allocation status of the corresponding variable in the device data environment must be unallocated upon exiting the region.

• If the allocation status of a mapped variable that has the `ALLOCATABLE` attribute is allocated on entry to a `target` region, the allocation status and shape of the corresponding variable in the device data environment may not be changed, either explicitly or implicitly, in the region after entry to it.

• If the association status of a list item with the `POINTER` attribute that appears in a `map` clause on the construct is associated upon entry to the `target` region, the list item must be associated with the same pointer target upon exit from the region.

• If the association status of a list item with the `POINTER` attribute that appears in a `map` clause on the construct is disassociated upon entry to the `target` region, the list item must be disassociated upon exit from the region.

• If the association status of a list item with the `POINTER` attribute that appears in a `map` clause on the construct is disassociated or undefined on entry to the `target` region and if the list item is associated with a pointer target inside the `target` region, the pointer association status must become disassociated before the end the region.
Cross References
- default-device-var, see Section 2.4.
- task construct, see Section 2.12.1.
- task scheduling constraints, see Section 2.12.6
- Memory allocators, see Section 2.13.2.
- target data construct, see Section 2.14.2.
- if clause, see Section 2.18.
- private and firstprivate clauses, see Section 2.21.4.
- Data-Mapping Attribute Rules and Clauses, see Section 2.21.7.
- omp_get_num_devices routine, see Section 3.7.4.
- omp_alloctrait_t and omp_alloctrait types, see Section 3.13.1.
- omp_set_default_allocator routine, see Section 3.13.4.
- omp_get_default_allocator routine, see Section 3.13.5.
- ompt_callback_target_t or ompt_callback_target_emi_t callback type, see Section 4.5.2.26.
- ompt_callback_target_submit_t or ompt_callback_target_submit_emi_t callback type, Section 4.5.2.28.

2.14.6 target update Construct

Summary
The target update directive makes the corresponding list items in the device data environment consistent with their original list items, according to the specified motion clauses. The target update construct is a stand-alone directive.

Syntax

```
C / C++

#pragma omp target update clause[ [ , ] clause ] ... ] new-line
```

where clause is either motion-clause or one of the following:

```
if(/ target update :/) scalar-expression
device(integer-expression)
nowait
depend([depend-modifier, ] dependence-type : locator-list)
```
and motion-clause is one of the following:

\[
\begin{align*}
to & (\{[\text{motion-modifier}[, ] [\text{motion-modifier}[, ] \ldots]: ] \text{ locator-list} \}) \\
from & (\{[\text{motion-modifier}[, ] [\text{motion-modifier}[, ] \ldots]: ] \text{ locator-list} \})
\end{align*}
\]

where motion-modifier is one of the following:

- present
- mapper \((\text{mapper-identifier})\)
- iterator \((\text{iterators-definition})\)

The syntax of the `target update` construct is as follows:

\[
!\$omp \text{ target update} \quad \text{clause[ [ [ ] clause] \ldots ]}
\]

where clause is either motion-clause or one of the following:

- `if` \((/ \text{ target update } : ] \text{ scalar-logical-expression})\)
- `device` \((\text{scalar-integer-expression})\)
- `nowait`
- `depend` \((/\text{depend-modifier}, ] \text{ dependence-type} : ] \text{ locator-list})\)

and motion-clause is one of the following:

\[
\begin{align*}
to & (\{[\text{motion-modifier}[, ] [\text{motion-modifier}[, ] \ldots]: ] \text{ locator-list} \}) \\
from & (\{[\text{motion-modifier}[, ] [\text{motion-modifier}[, ] \ldots]: ] \text{ locator-list} \})
\end{align*}
\]

where motion-modifier is one of the following:

- present
- mapper \((\text{mapper-identifier})\)
- iterator \((\text{iterators-definition})\)

### Binding

The binding task set for a `target update` region is the generating task, which is the `target task` generated by the `target update` construct. The `target update` region binds to the corresponding `target task` region.
Description

A corresponding list item and an original list item exist for each list item in a to or from clause. If the corresponding list item is not present in the device data environment and the present modifier is not specified in the clause then no assignment occurs to or from the original list item. Otherwise, each corresponding list item in the device data environment has an original list item in the current task’s data environment.

If the mapper motion-modifier is not present, the behavior is as if the mapper (default) modifier was specified. The effect of a motion clause on a list item is modified by a visible user-defined mapper (see Section 2.21.7.4) if the mapper has the same mapper-identifier as the mapper motion-modifier and is specified for a type that matches the type of the list item. For a to clause, each list item is replaced with the list items that the given mapper specifies are to be mapped with a to or tofrom map type. For a from clause, each list item is replaced with the list items that the given mapper specifies are to be mapped with a from or tofrom map type.

If a list item is an array or array section of a type for which a user-defined mapper exists, each array element is updated according to the mapper as if it appeared as a list item in the clause.

If a present modifier appears in the clause and the corresponding list item is not present in the device data environment then an error occurs and the program terminates. For a list item that is replaced with a set of list items as a result of a user-defined mapper, the present modifier only applies to those mapper list items that share storage with the original list item.

The list items that appear in a to or from may reference iterators defined by an iterators-definition appearing on an iterator modifier.

Fortran

If a list item or a subobject of a list item in a motion clause has the ALLOCATABLE attribute, its update is performed only if its allocation status is allocated and only with respect to the allocated storage.

If a list item in a motion clause has the POINTER attribute and its association status is associated, the effect is as if the update is performed with respect to the pointer target.

Fortran

For each list item in a from or a to clause:

- For each part of the list item that is an attached pointer:

C / C++

- On exit from the region that part of the original list item will have the value it had on entry to the region;

C / C++

- On exit from the region that part of the corresponding list item will have the value it had on entry to the region;
Fortran

- On exit from the region that part of the original list item, if associated, will be associated with the same pointer target with which it was associated on entry to the region;
- On exit from the region that part of the corresponding list item, if associated, will be associated with the same pointer target with which it was associated on entry to the region.


- For each part of the list item that is not an attached pointer:
  - If the clause is from, the value of that part of the corresponding list item is assigned to that part of the original list item;
  - If the clause is to, the value of that part of the original list item is assigned to that part of the corresponding list item.

- To avoid data races:
  - Concurrent reads or updates of any part of the original list item must be synchronized with the update of the original list item that occurs as a result of the from clause;
  - Concurrent reads or updates of any part of the corresponding list item must be synchronized with the update of the corresponding list item that occurs as a result of the to clause.

C / C++

The list items that appear in the to or from clauses may use shape-operators.

The list items that appear in the to or from clauses may include array sections with stride expressions.

The target update construct is a task generating construct. The generated task is a target task. The generated task region encloses the target update region.

All clauses are evaluated when the target update construct is encountered. The data environment of the target task is created according to the motion clauses on the target update construct, per-data environment ICVs, and any default data-sharing attribute rules that apply to the target update construct. If a variable or part of a variable is a list item in a motion clause on the target update construct, the variable has a default data-sharing attribute of shared in the data environment of the target task.

Assignment operations associated with mapping a variable (see Section 2.21.7.1) occur when the target task executes.

If the nowait clause is present, execution of the target task may be deferred. If the nowait clause is not present, the target task is an included task.

If a depend clause is present, it is associated with the target task.
The device is specified in the `device` clause. If no `device` clause is present, the device is determined by the `default-device-var` ICV. When an `if` clause is present and the `if` clause expression evaluates to `false`, no assignments occur.

**Execution Model Events**

Events associated with a `target task` are the same as for the `task` construct defined in Section 2.12.1.

The `target-update-begin` event occurs when a thread enters a `target update` region.

The `target-update-end` event occurs when a thread exits a `target update` region.

**Tool Callbacks**

Callbacks associated with events for `target tasks` are the same as for the `task` construct defined in Section 2.12.1; `(flags & ompt_task_target)` always evaluates to `true` in the dispatched callback.

A thread dispatches a registered `ompt_callback_target` or `ompt_callback_target_emi` callback with `ompt_scope_begin` as its `endpoint` argument and `ompt_target_update` or `ompt_target_update_nowait` if the `nowait` clause is present as its `kind` argument for each occurrence of a `target-update-begin` event in that thread in the context of the target task on the host. Similarly, a thread dispatches a registered `ompt_callback_target` or `ompt_callback_target_emi` callback with `ompt_scope_end` as its `endpoint` argument and `ompt_target_update` or `ompt_target_update_nowait` if the `nowait` clause is present as its `kind` argument for each occurrence of a `target-update-end` event in that thread in the context of the target task on the host. These callbacks have type signature `ompt_callback_target_t` or `ompt_callback_target_emi_t`, respectively.

**Restrictions**

Restrictions to the `target update` construct are as follows:

- A program must not depend on any ordering of the evaluations of the clauses of the `target update` directive, or on any side effects of the evaluations of the clauses.
- Each of the `motion-modifier` modifiers can appear at most once on a motion clause.
- At least one `motion-clause` must be specified.
- A list item can only appear in a `to` or `from` clause, but not in both.
- A list item in a `to` or `from` clause must have a mappable type.
- At most one `device` clause can appear on the directive. The `device` clause expression must evaluate to a non-negative integer value that is less than or equal to the value of `omp_get_num_devices()`.
- At most one `if` clause can appear on the directive.
- At most one `nowait` clause can appear on the directive.
2.14.7 Declare Target Directive

Summary
The declare target directive specifies that variables, functions (C, C++ and Fortran), and subroutines (Fortran) are mapped to a device. The declare target directive is a declarative directive.

Syntax

```
#pragma omp declare target new-line
declaration-definition-seq
#pragma omp end declare target new-line
```

or

```
#pragma omp declare target (extended-list) new-line
```

or

```
#pragma omp declare target clause[ [ , ] clause] ... ] new-line
```
or

```
#pragma omp begin declare target [clause [ , ] clause] ... ] new-line
declaration-definition-seq
#pragma omp end declare target new-line
```

where `clause` is one of the following:

```
to(extended-list)
link(list)
device_type(host | nohost | any)
indirect[(invoked-by-fptr)]
```

where `invoked-by-fptr` is a constant boolean expression that evaluates to true or false at compile time.

The syntax of the `declare target` directive is as follows:

```
!$omp declare target (extended-list)
```

or

```
!$omp declare target [clause [ , ] clause] ...
```

where `clause` is one of the following:

```
to(extended-list)
link(list)
device_type(host | nohost | any)
indirect[(invoked-by-fptr)]
```

where `invoked-by-fptr` is a constant boolean expression that evaluates to true or false at compile time.
**Description**

The declare target directive ensures that procedures and global variables can be executed or accessed on a device. Variables are mapped for all device executions, or for specific device executions through a `link` clause.

If an `extended-list` is present with no clause then the `to` clause is assumed.

The `device_type` clause specifies if a version of the procedure or variable should be made available on the host, device or both. If `host` is specified only a host version of the procedure or variable is made available. If `any` is specified then both device and host versions of the procedure or variable are made available. If `nohost` is specified for a procedure then only a device version of the procedure is made available. If `nohost` is specified for a variable then that variable is not available on the host. If no `device_type` clause is present then the behavior is as if the `device_type` clause appears with `any` specified.

If a variable with static storage duration is declared in a device routine then the named variable is treated as if it had appeared in a `to` clause on a declare target directive.

```c

// C / C++
```

If a function appears in a `to` clause in the same compilation unit in which the definition of the function occurs then a device-specific version of the function is created.

If a variable appears in a `to` clause in the same compilation unit in which the definition of the variable occurs then the original list item is allocated a corresponding list item in the device data environment of all devices.

```c

// C / C++
```

```fortran

// Fortran
```

If a procedure appears in a `to` clause in the same compilation unit in which the definition of the procedure occurs then a device-specific version of the procedure is created.

If a variable that is host associated appears in a `to` clause then the original list item is allocated a corresponding list item in the device data environment of all devices.

```fortran

// Fortran
```

If a variable appears in a `to` clause then the corresponding list item in the device data environment of each device is initialized once, in the manner specified by the program, but at an unspecified point in the program prior to the first reference to that list item. The list item is never removed from those device data environments as if its reference count was initialized to positive infinity.

Including list items in a `link` clause supports compilation of functions called in a `target` region that refer to the list items. The list items are not mapped by the declare target directive. Instead, they are mapped according to the data mapping rules described in Section 2.21.7.
If a function is referenced in a function that appears as a list item in a `to` clause on a declare target directive that does not specify a `device_type` clause with `host` and the function reference is not enclosed in a `target` construct that specifies a `device` clause in which the `ancestor` `device-modifier` appears then the name of the referenced function is treated as if it had appeared in a `to` clause on a declare target directive.

If a variable with static storage duration or a function (except `lambda` for C++) is referenced in the initializer expression list of a variable with static storage duration that appears as a list item in a `to` clause on a declare target directive then the name of the referenced variable or function is treated as if it had appeared in a `to` clause on a declare target directive.

The form, preceded by either the `declare target` directive that has no clauses and no `extended-list` or the `begin declare target` directive and followed by a matching `end declare target` directive, defines an implicit `extended-list`. The implicit `extended-list` consists of the variable names of any variable declarations at file or namespace scope that appear between the two directives and of the function names of any function declarations at file, namespace or class scope that appear between the two directives. The implicit `extended-list` is converted to an implicit `to` clause.

The `declaration-definition-seq` preceded by either `begin declare target` directive or a `declare target` directive without any clauses or an `extended-list` and followed by an `end declare target` directive may contain declare target directives. If a `device_type` clause is present on the contained declare target directive, then its argument determines which versions are made available. If a list item appears both in an implicit and explicit list, the explicit list determines which versions are made available.

If a procedure is referenced in a procedure that appears as a list item in a `to` clause on a `declare target` directive that does not specify a `device_type` clause with `host` and the procedure reference is not enclosed in a `target` construct that specifies a `device` clause in which the `ancestor` `device-modifier` appears then the name of the referenced procedure is treated as if it had appeared in a `to` clause on a `declare target` directive.

If a `declare target` does not have any clauses and does not have an `extended-list` then an implicit `to` clause with one item is formed from the name of the enclosing subroutine subprogram, function subprogram or interface body to which it applies.

If a `declare target` directive has a `device_type` clause then any enclosed internal procedures cannot contain any `declare target` directives. The enclosing `device_type` clause implicitly applies to internal procedures.
If the `indirect` clause is present and `invoked-by-fptr` is not specified, the effect of the clause is as if `invoked-by-fptr` evaluates to true.

If the `indirect` clause is present and `invoked-by-fptr` evaluates to true, any procedures that appear in a `to` clause on the directive may be called with an indirect device invocation. If the `indirect` clause is present and `invoked-by-fptr` does not evaluate to true, any procedures that appear in a `to` clause on the directive may not be called with an indirect device invocation. Unless otherwise specified by an `indirect` clause, procedures may not be called with an indirect device invocation.

If a function appears in the `to` clause of a `begin declare target` directive and in the `to` clause of a declare target directive that is contained in the `declaration-definition-seq` associated with the `begin declare target` directive, and if an `indirect` clause appears on both directives, then the `indirect` clause on the `begin declare target` directive has no effect for that function.

**Execution Model Events**
The `target-global-data-op` event occurs when an original variable is associated with a corresponding variable on a device as a result of a declare target directive; the event occurs before the first access to the corresponding variable.

**Tool Callbacks**
A thread dispatches a registered `ompt_callback_target_data_op` callback, or a registered `ompt_callback_target_data_op_emi` callback with `ompt_scope_beginend` as its endpoint argument for each occurrence of a `target-global-data-op` event in that thread. These callbacks have type signature `ompt_callback_target_data_op_t` or `ompt_callback_target_data_op_emi_t`, respectively.

**Restrictions**
Restrictions to the declare target directive are as follows:

- A thread private variable cannot appear in the directive.
- A variable declared in the directive must have a mappable type.
- The same list item must not appear multiple times in clauses on the same directive.
- The same list item must not explicitly appear in both a `to` clause on one declare target directive and a `link` clause on another declare target directive.
- If the directive has a clause, it must contain at least one `to` clause or at least one `link` clause.
- A variable for which `nohost` is specified may not appear in a `link` clause.
- At most one `indirect` clause can be specified on the directive.
- At most one `device_type` clause can be specified on the directive.
• If an indirect clause is present and invoked-by-fptr evaluates to true then the only permitted
  device_type clause is device_type(any).

• A to clause or a link clause cannot appear in a begin declare target directive.

• The function names of overloaded functions or template functions may only be specified within
  an implicit extended-list.

• If a lambda declaration and definition appears between a declare target directive and the paired
  end declare target directive, all variables that are captured by the lambda expression must
  also appear in a to clause.

• A module export or import statement cannot appear between a declare target directive and the
  paired end declare target directive.

• If a list item is a procedure name, it must not be a generic name, procedure pointer, entry name,
  or statement function name.

• If the directive does not have any clause or has a device_type clause it must appear in a
  specification part of a subroutine subprogram, function subprogram or interface body.

• If a list item is a procedure name, the directive must be in the specification part of that subroutine
  or function subprogram or in the specification part of that subroutine or function in an interface
  body.

• If the directive has a variable name in extended-list, it must appear in the specification part of a
  subroutine subprogram, function subprogram, program or module.

• If the directive is specified in an interface block for a procedure, it must match a
  declare target directive in the definition of the procedure, including the device_type
  clause if present.

• If an external procedure is a type-bound procedure of a derived type and the directive is specified
  in the definition of the external procedure, it must appear in the interface block that is accessible
  to the derived-type definition.

• If any procedure is declared via a procedure declaration statement that is not in the type-bound
  procedure part of a derived-type definition, any declare target with the procedure name
  must appear in the same specification part.

• A variable that is part of another variable (as an array, structure element or type parameter
  inquiry) cannot appear in the directive.

• The directive must appear in the declaration section of a scoping unit in which the common block
  or variable is declared.
• If a `declare target` directive that specifies a common block name appears in one program unit, then such a directive must also appear in every other program unit that contains a `COMMON` statement that specifies the same name, after the last such `COMMON` statement in the program unit.

• If a list item is declared with the `BIND` attribute, the corresponding C entities must also be specified in a `declare target` directive in the C program.

• A variable can only appear in a `declare target` directive in the scope in which it is declared. It must not be an element of a common block or appear in an `EQUIVALENCE` statement.

• A variable that appears in a `declare target` directive must be declared in the Fortran scope of a module or have the `SAVE` attribute, either explicitly or implicitly.

Cross References

• `target data` construct, see Section 2.14.2.

• `target` construct, see Section 2.14.5.

• `ompt_callback_target_data_op_t` or `ompt_callback_target_data_op_emi_t` callback type, see Section 4.5.2.25.

2.15 Interoperability

An OpenMP implementation may interoperate with one or more foreign runtime environments through the use of the `interop` construct that is described in this section, the `interop` operation for a declared variant function and the interoperability routines that are available through the OpenMP Runtime API.

The implementation must provide `foreign-runtime-id` values that are enumerators of type `omp_interop_fr_t` and that correspond to the supported foreign runtime environments.

The implementation must provide `foreign-runtime-id` values that are named integer constants with kind `omp_interop_fr_kind` and that correspond to the supported foreign runtime environments.

Each `foreign-runtime-id` value provided by an implementation will be available as `omp_ifr_name`, where `name` is the name of the foreign runtime environment. Available names include those that are listed in the `OpenMP Additional Definitions` document; implementation-defined names may also be supported. The value of `omp_ifr_last` is defined as one greater than the value of the highest supported `foreign-runtime-id` value that is listed in the aforementioned document.
Cross References

• `declare variant` directive, see Section 2.3.5.
• Interoperability routines, see Section 3.12.

2.15.1 interop Construct

Summary
The `interop` construct retrieves interoperability properties from the OpenMP implementation to enable interoperability with foreign execution contexts. The `interop` construct is a stand-alone directive.

Syntax
In the following syntax, `interop-type` is the type of interoperability information being requested or used by the `interop` construct, and `action-clause` is a clause that indicates the action to take with respect to that interop object.

```
C / C++
```

The syntax of the `interop` construct is as follows:

```
#pragma omp interop clause[ ... ] new-line
```

where `clause` is `action-clause` or one of the following:

```
device(integer-expression)
depend([depend-modifier,] dependence-type : locator-list)
```

where `action-clause` is one of the following:

```
init([interop-modifier,] interop-type[ , interop-type] ... : interop-var)
destroy(interop-var)
use(interop-var)
ownait
```

where `interop-var` is a variable of type `omp_interop_t`, and `interop-type` is one of the following:

```
target
targetsync
```

and `interop-modifier` is one of the following:

```
prefer_type(preference-list)
```
where preference-list is a comma separated list for which each item is a foreign-runtime-id, which is a base language string literal or a compile-time constant integral expression. Allowed values for foreign-runtime-id include the names (as string literals) and integer values specified in the OpenMP Additional Definitions document and the corresponding omp_ifr_name constants of omp_interop_fr_t type; implementation-defined values may also be supported.

The syntax of the interop construct is as follows:

```language-plaintext
!$omp interop clause[ [ [, ] clause] ... ]
```

where clause is action-clause or one of the following:

- `device(integer-expression)`
- `depend([depend-modifier, ] dependence-type : locator-list)`

where action-clause is one of the following:

- `init([interop-modifier, ]interop-type[, interop-type] ... ):interop-var)`
- `destroy(interop-var)`
- `use(interop-var)`
- `nowait`

where interop-var is a scalar integer variable of kind omp_interop_kind, and interop-type is one of the following:

- `target`
- `targetsync`

and interop-modifier is one of the following:

- `prefer_type(preference-list)`

where preference-list is a comma separated list for which each item is a foreign-runtime-id, which is a base language string literal or a compile-time constant integral expression. Allowed values for foreign-runtime-id include the names (as string literals) and integer values specified in the OpenMP Additional Definitions document and the corresponding omp_ifr_name integer constants of kind omp_interop_fr_kind; implementation-defined values may also be supported.

**Binding**

The binding task set for an interop region is the generating task. The interop region binds to the region of the generating task.
When an `interop` construct is encountered, the encountering task executes the region. If no `device` clause is present, the behavior is as if the `device` clause appears with an expression equal to the value of the `default-device-var` ICV.

If the `init` action-clause is specified, the `interop-var` is initialized to refer to the list of properties associated with the given `interop-type`. For any `interop-type`, the properties `type`, `type_name`, `vendor`, `vendor_name` and `device_num` will be available. If the implementation is unable to initialize the `interop-var`, it will be initialized to the value of `omp_interop_none`, which is defined to be zero.

The `targetsync` `interop-type` will additionally provide the `targetsync` property, which is the handle to a foreign synchronization object for enabling synchronization between OpenMP tasks and foreign tasks that execute in the foreign execution context.

The `target` `interop-type` will additionally provide the following properties:

- `device`, which will be a foreign device handle;
- `device_context`, which will be a foreign device context handle; and
- `platform`, which will be a handle to a foreign platform of the device.

If the `destroy` action-clause is specified, the `interop-var` is set to the value of `omp_interop_none` after resources associated with `interop-var` are released. The object referred to by the `interop-var` will be unusable after destruction and the effect of using values associated with it is unspecified until `interop-var` is initialized again by another `interop` construct.

If the `use` action-clause is specified, the `interop-var` is used for other effects of this directive but is not initialized, destroyed or otherwise modified.

For the `destroy` or `use` action-clause, the `interop-type` is inferred based on the `interop-type` used to initialize the `interop-var`.

If the `interop-type` specified is `targetsync`, or the `interop-var` was initialized with `targetsync`, an empty `mergeable task` is generated. If the `nowait` clause is not present on the construct then the task is also an `included task`. Any `depend` clauses that are present on the construct apply to the generated task. The `interop` construct ensures an ordered execution of the generated task relative to foreign tasks executed in the foreign execution context through the foreign synchronization object accessible through the `targetsync` property of `interop-var`. When the creation of the foreign task precedes the encountering of an `interop` construct in happens before order (see Section 1.4.5), the foreign task must complete execution before the generated task begins execution. Similarly, when the creation of a foreign task follows the encountering of an `interop` construct in happens before order, the foreign task must not begin execution until the generated task completes execution. No ordering is imposed between the encountering thread and either foreign tasks or OpenMP tasks by the `interop` construct.
If the `prefer_type` `interop-modifier` clause is specified, the first supported `foreign-runtime-id` in `preference-list` in left-to-right order is used. The `foreign-runtime-id` that is used if the implementation does not support any of the items in `preference-list` is implementation defined.

**Restrictions**

Restrictions to the `interop` construct are as follows:

- At least one `action-clause` must appear on the directive.
- Each `interop-type` may be specified on an `action-clause` at most once.
- The `interop-var` passed to `init` or `destroy` must be non-const.
- A `depend` clause can only appear on the directive if a `targetsync` `interop-type` is present or the `interop-var` was initialized with the `targetsync` `interop-type`.
- Each `interop-var` may be specified for at most one `action-clause` of each `interop` construct.
- At most one `device` clause can appear on the directive. The `device` clause expression must evaluate to a non-negative integer value less than or equal to the value returned by `omp_get_num_devices()`.
- At most one `nowait` clause can appear on the directive.

**Cross References**

- `depend` clause, see Section 2.19.11.
- Interoperability routines, see Section 3.12.

**2.15.2 Interoperability Requirement Set**

The `interoperability requirement set` of each task is a logical set of properties that can be added or removed by different directives. These properties can be queried by other constructs that have interoperability semantics.

A construct can add the following properties to the set:

- `depend`, which specifies that the construct requires enforcement of the synchronization relationship expressed by the `depend` clause;
- `nowait`, which specifies that the construct is asynchronous; and
- `is_device_ptr(list-item)`, which specifies that the `list-item` is a device pointer in the construct.

The following directives may add properties to the set:

- `dispatch`.

The following directives may remove properties from the set:

- `declare variant`.
Cross References

- `declare variant` directive, see Section 2.3.5.
- `dispatch` construct, see Section 2.3.6.

2.16 Combined Constructs

Combined constructs are shortcuts for specifying one construct immediately nested inside another construct. The semantics of the combined constructs are identical to that of explicitly specifying the first construct containing one instance of the second construct and no other statements.

For combined constructs, tool callbacks are invoked as if the constructs were explicitly nested.

2.16.1 Parallel Worksharing-Loop Construct

Summary

The parallel worksharing-loop construct is a shortcut for specifying a `parallel` construct containing a worksharing-loop construct with a canonical loop nest and no other statements.

Syntax

### C / C++

The syntax of the parallel worksharing-loop construct is as follows:

```c
#pragma omp parallel for [clause [, clause] ... ] new-line loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `parallel` or `for` directives, except the `nowait` clause, with identical meanings and restrictions.

### Fortran

The syntax of the parallel worksharing-loop construct is as follows:

```fortran
!$omp parallel do [clause [, clause] ... ]
loop-nest

[$!$omp end parallel do/]
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `parallel` or `do` directives, with identical meanings and restrictions.

If an `end parallel do` directive is not specified, an `end parallel do` directive is assumed at the end of the `loop-nest`. 
Description
The semantics are identical to explicitly specifying a `parallel` directive immediately followed by a worksharing-loop directive.

Restrictions
The restrictions for the `parallel` construct and the worksharing-loop construct apply.

Cross References
- `parallel` construct, see Section 2.6.
- Canonical loop nest form, see Section 2.11.1.
- Worksharing-loop construct, see Section 2.11.4.
- Data attribute clauses, see Section 2.21.4.

2.16.2 parallel loop Construct

Summary
The `parallel loop` construct is a shortcut for specifying a `parallel` construct containing a `loop` construct with a canonical loop nest and no other statements.

Syntax
C / C++

```c
#pragma omp parallel loop [clause[ [, ] clause] ... ] new-line
loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `parallel` or `loop` directives, with identical meanings and restrictions.

Fortran

```fortran
!$omp parallel loop [clause[ [, ] clause] ... ]
loop-nest
/$omp end parallel loop/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `parallel` or `loop` directives, with identical meanings and restrictions.

If an `end parallel loop` directive is not specified, an `end parallel loop` directive is assumed at the end of the `loop-nest`. 
The semantics are identical to explicitly specifying a parallel directive immediately followed by a loop directive.

The restrictions for the parallel construct and the loop construct apply.

The restrictions for the parallel construct and the loop construct apply.

The parallel sections construct is a shortcut for specifying a parallel construct containing a sections construct and no other statements.

The syntax of the parallel sections construct is as follows:

```c
#pragma omp parallel sections [clause[ [, ] clause] ... ] new-line
{
    [#pragma omp section new-line]
    structured-block-sequence
    [#pragma omp section new-line
    structured-block-sequence]
    ...
}
```

where clause can be any of the clauses accepted by the parallel or sections directives, except the nowait clause, with identical meanings and restrictions.
The syntax of the **parallel sections** construct is as follows:

```fortran
!$omp parallel sections [clause [, clause] ... ]
   [ !$omp section/
     structured-block-sequence
   ]
   [$omp section]
   structured-block-sequence
   [ ... ]
!$omp end parallel sections
```

where `clause` can be any of the clauses accepted by the parallel or sections directives, with identical meanings and restrictions.

### Description

The semantics are identical to explicitly specifying a parallel directive immediately followed by a sections directive.

The semantics are identical to explicitly specifying a parallel directive immediately followed by a sections directive, and an end sections directive immediately followed by an end parallel directive.

### Restrictions

The restrictions for the parallel construct and the sections construct apply.

### Cross References

- parallel construct, see Section 2.6.
- sections construct, see Section 2.10.1.
- Data attribute clauses, see Section 2.21.4.

### 2.16.4 parallel workshare Construct

#### Summary

The parallel workshare construct is a shortcut for specifying a parallel construct containing a workshare construct and no other statements.
Syntax
The syntax of the parallel workshare construct is as follows:

```c
!$omp parallel workshare [clause [ , clause] ... ]
   loosely-structured-block
!$omp end parallel workshare
```

or

```c
!$omp parallel workshare [clause [ , clause] ... ]
   strictly-structured-block
!$omp end parallel workshare
```

where clause can be any of the clauses accepted by the parallel directive, with identical meanings and restrictions.

Description
The semantics are identical to explicitly specifying a parallel directive immediately followed by a workshare directive, and an end workshare directive immediately followed by an end parallel directive.

Restrictions
The restrictions for the parallel construct and the workshare construct apply.

Cross References
• parallel construct, see Section 2.6.
• workshare construct, see Section 2.10.3.
• Data attribute clauses, see Section 2.21.4.

2.16.5 Parallel Worksharing-Loop SIMD Construct

Summary
The parallel worksharing-loop SIMD construct is a shortcut for specifying a parallel construct containing a worksharing-loop SIMD construct and no other statements.

Syntax
The syntax of the parallel worksharing-loop SIMD construct is as follows:

```c
#pragma omp parallel for simd [clause [ , clause] ... ] new-line
   loop-nest
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the parallel or for simd directives, except the nowait clause, with identical meanings and restrictions.
The syntax of the parallel worksharing-loop SIMD construct is as follows:

```fortran
!$omp parallel do simd [clause[ [, ] clause] ... ]
loop-nest
/$omp end parallel do simd/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `parallel` or `do simd` directives, with identical meanings and restrictions.

If an `end parallel do simd` directive is not specified, an `end parallel do simd` directive is assumed at the end of the `loop-nest`.

**Description**

The semantics of the parallel worksharing-loop SIMD construct are identical to explicitly specifying a `parallel` directive immediately followed by a worksharing-loop SIMD directive.

**Restrictions**

The restrictions for the `parallel` construct and the worksharing-loop SIMD construct apply except for the following explicit modifications:

- If any `if` clause on the directive includes a `directive-name-modifier` then all `if` clauses on the directive must include a `directive-name-modifier`.
- At most one `if` clause without a `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `parallel directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `simd directive-name-modifier` can appear on the directive.

**Cross References**

- `parallel` construct, see Section 2.6.
- Canonical loop nest form, see Section 2.11.1.
- Worksharing-loop SIMD construct, see Section 2.11.5.2.
- `if` clause, see Section 2.18.
- Data attribute clauses, see Section 2.21.4.

**2.16.6 parallel masked Construct**

**Summary**

The `parallel masked` construct is a shortcut for specifying a `parallel` construct containing a `masked` construct, and no other statements.
Syntax

The syntax of the parallel masked construct is as follows:

```
#pragma omp parallel masked [clause [ , ] clause] ... ] new-line
```

structured-block

where clause can be any of the clauses accepted by the parallel or masked directives, with identical meanings and restrictions.

```
!$omp parallel masked [clause [ , ] clause] ...
```

loosely-structured-block

```
!$omp end parallel masked
```

or

```
!$omp parallel masked [clause [ , ] clause] ... ]
```

strictly-structured-block

```
/$omp end parallel masked/
```

where clause can be any of the clauses accepted by the parallel or masked directives, with identical meanings and restrictions.

The parallel master construct, which has been deprecated, has identical syntax to the parallel masked construct other than the use of parallel master as the directive name.

Description

The semantics are identical to explicitly specifying a parallel directive immediately followed by a masked directive.

Restrictions

The restrictions for the parallel construct and the masked construct apply.

Cross References

- parallel construct, see Section 2.6.
- masked construct, see Section 2.8.
- Data attribute clauses, see Section 2.21.4.
### 2.16.7 masked taskloop Construct

#### Summary
The masked taskloop construct is a shortcut for specifying a masked construct containing a taskloop construct and no other statements.

#### Syntax

**C / C++**

The syntax of the masked taskloop construct is as follows:

```
#pragma omp masked taskloop [clause[ , clause] ... ] new-line
   loop-nest
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the masked or taskloop directives with identical meanings and restrictions.

**Fortran**

The syntax of the masked taskloop construct is as follows:

```
!$omp masked taskloop [clause[ , clause] ... ]
   loop-nest
!$omp end masked taskloop/
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the masked or taskloop directives with identical meanings and restrictions.

If an end masked taskloop directive is not specified, an end masked taskloop directive is assumed at the end of the loop-nest.

The master taskloop construct, which has been deprecated, has identical syntax to the masked taskloop construct other than the use of master taskloop as the directive name.

#### Description
The semantics are identical to explicitly specifying a masked directive immediately followed by a taskloop directive.

#### Restrictions
The restrictions for the masked and taskloop constructs apply.

#### Cross References
- masked construct, see Section 2.8.
- Canonical loop nest form, see Section 2.11.1.
- taskloop construct, see Section 2.12.2.
- Data attribute clauses, see Section 2.21.4.
2.16.8 masked taskloop simd Construct

Summary
The masked taskloop simd construct is a shortcut for specifying a masked construct containing a taskloop simd construct and no other statements.

Syntax

```
  C / C++
  #pragma omp masked taskloop simd [clause[ , ] clause] ... ] new-line
    loop-nest

  Fortran
  !$omp masked taskloop simd [clause[ , ] clause] ... ]
    loop-nest

  !$omp end masked taskloop simd/
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the masked or taskloop simd directives with identical meanings and restrictions.

Description
The semantics are identical to explicitly specifying a masked directive immediately followed by a taskloop simd directive.

Restrictions
The restrictions for the masked and taskloop simd constructs apply.

Cross References
* masked construct, see Section 2.8.
* Canonical loop nest form, see Section 2.11.1.
* taskloop simd construct, see Section 2.12.3.
* Data attribute clauses, see Section 2.21.4.
2.16.9 parallel masked taskloop Construct

Summary
The parallel masked taskloop construct is a shortcut for specifying a parallel construct containing a masked taskloop construct and no other statements.

The syntax of the parallel masked taskloop construct is as follows:

```
#pragma omp parallel masked taskloop [clause[ [, ] clause] ... ] new-line
  loop-nest
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the parallel or masked taskloop directives, except the in_reduction clause, with identical meanings and restrictions.

The syntax of the parallel masked taskloop construct is as follows:

```
!$omp parallel masked taskloop [clause[ [, ] clause] ... ]
  loop-nest
$/omp end parallel masked taskloop/
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the parallel or masked taskloop directives, except the in_reduction clause, with identical meanings and restrictions.

If an end parallel masked taskloop directive is not specified, an end parallel masked taskloop directive is assumed at the end of the loop-nest.

The parallel master taskloop construct, which has been deprecated, has identical syntax to the parallel masked taskloop construct other than the use of parallel master taskloop as the directive name.

Description
The semantics are identical to explicitly specifying a parallel directive immediately followed by a masked taskloop directive.
Restrictions
The restrictions for the parallel construct and the masked taskloop construct apply except for the following explicit modifications:

- If any if clause on the directive includes a directive-name-modifier then all if clauses on the directive must include a directive-name-modifier.
- At most one if clause without a directive-name-modifier can appear on the directive.
- At most one if clause with the parallel directive-name-modifier can appear on the directive.
- At most one if clause with the taskloop directive-name-modifier can appear on the directive.

Cross References
- parallel construct, see Section 2.6.
- Canonical loop nest form, see Section 2.11.1.
- masked taskloop construct, see Section 2.16.7.
- if clause, see Section 2.18.
- Data attribute clauses, see Section 2.21.4.
- in_reduction clause, see Section 2.21.5.6.

2.16.10 parallel masked taskloop simd Construct
Summary
The parallel masked taskloop simd construct is a shortcut for specifying a parallel construct containing a masked taskloop simd construct and no other statements.

Syntax
The syntax of the parallel masked taskloop simd construct is as follows:

```
#pragma omp parallel masked taskloop simd [clause[, clause]... ] new-line
loop-nest
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the parallel or masked taskloop simd directives, except the in_reduction clause, with identical meanings and restrictions.
The syntax of the `parallel masked taskloop simd` construct is as follows:

```
!$omp parallel masked taskloop simd [clause[ [, clause] ... ]] loop-nest

/$omp end parallel masked taskloop simd/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `parallel` or `masked taskloop simd` directives, except the `in_reduction` clause, with identical meanings and restrictions.

If an `end parallel masked taskloop simd` directive is not specified, an `end parallel masked taskloop simd` directive is assumed at the end of the `loop-nest`.

The `parallel master taskloop simd` construct, which has been deprecated, has identical syntax to the `parallel masked taskloop simd` construct other than the use of `parallel master taskloop simd` as the directive name.

**Description**

The semantics are identical to explicitly specifying a `parallel` directive immediately followed by a `masked taskloop simd` directive.

**Restrictions**

The restrictions for the `parallel` construct and the `masked taskloop simd` construct apply except for the following explicit modifications:

- If any `if` clause on the directive includes a `directive-name-modifier` then all `if` clauses on the directive must include a `directive-name-modifier`.
- At most one `if` clause without a `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `parallel` `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `taskloop` `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `simd` `directive-name-modifier` can appear on the directive.

**Cross References**

- `parallel` construct, see Section 2.6.
- Canonical loop nest form, see Section 2.11.1.
- `masked taskloop simd` construct, see Section 2.16.8.
- `if` clause, see Section 2.18.
- Data attribute clauses, see Section 2.21.4.
- `in_reduction` clause, see Section 2.21.5.6.
2.16.11 teams distribute Construct

Summary
The teams distribute construct is a shortcut for specifying a teams construct containing a distribute construct and no other statements.

Syntax

C / C++

The syntax of the teams distribute construct is as follows:

```
#pragma omp teams distribute [clause [, clause] ... ] new-line
loop-nest
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the teams or distribute directives with identical meanings and restrictions.

Fortran

The syntax of the teams distribute construct is as follows:

```
!$omp teams distribute [clause [, clause] ... ]
loop-nest

!$omp end teams distribute/
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the teams or distribute directives with identical meanings and restrictions.

If an end teams distribute directive is not specified, an end teams distribute directive is assumed at the end of the loop-nest.

Description
The semantics are identical to explicitly specifying a teams directive immediately followed by a distribute directive.

Restrictions
The restrictions for the teams and distribute constructs apply.

Cross References
- teams construct, see Section 2.7.
- Canonical loop nest form, see Section 2.11.1.
- distribute construct, see Section 2.11.6.1.
- Data attribute clauses, see Section 2.21.4.
2.16.12 teams distribute simd Construct

Summary
The **teams distribute simd** construct is a shortcut for specifying a **teams** construct containing a **distribute simd** construct and no other statements.

Syntax

```c
#pragma omp teams distribute simd [clause[ [, ] clause] ... ] new-line

loop-nest
```

The syntax of the **teams distribute simd** construct is as follows:

```
!$omp teams distribute simd [clause[ [, ] clause] ... ]

loop-nest

/$omp end teams distribute simd/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the **teams** or **distribute simd** directives with identical meanings and restrictions.

Description
The semantics are identical to explicitly specifying a **teams** directive immediately followed by a **distribute simd** directive.

Restrictions
The restrictions for the **teams** and **distribute simd** constructs apply.

Cross References
- **teams** construct, see Section 2.7.
- Canonical loop nest form, see Section 2.11.1.
- **distribute simd** construct, see Section 2.11.6.2.
- Data attribute clauses, see Section 2.21.4.
2.16.13 Teams Distribute Parallel Worksharing-Loop Construct

Summary
The teams distribute parallel worksharing-loop construct is a shortcut for specifying a `teams` construct containing a distribute parallel worksharing-loop construct and no other statements.

Syntax

The syntax of the teams distribute parallel worksharing-loop construct is as follows:

```c
#pragma omp teams distribute parallel for \\
[clause [, clause] ... ] new-line
loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `teams` or `distribute parallel for` directives with identical meanings and restrictions.

The syntax of the teams distribute parallel worksharing-loop construct is as follows:

```fortran
!$omp teams distribute parallel do [clause [, clause] ... ]
loop-nest

[!$omp end teams distribute parallel do/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `teams` or `distribute parallel do` directives with identical meanings and restrictions.

If an `end teams distribute parallel do` directive is not specified, an `end teams distribute parallel do` directive is assumed at the end of the `loop-nest`.

Description
The semantics are identical to explicitly specifying a `teams` directive immediately followed by a distribute parallel worksharing-loop directive.

Restrictions
The restrictions for the `teams` and distribute parallel worksharing-loop constructs apply.

Cross References
• `teams` construct, see Section 2.7.
• Canonical loop nest form, see Section 2.11.1.
• Distribute parallel worksharing-loop construct, see Section 2.11.6.3.
• Data attribute clauses, see Section 2.21.4.
2.16.14 Teams Distribute Parallel Worksharing-Loop SIMD Construct

Summary
The teams distribute parallel worksharing-loop SIMD construct is a shortcut for specifying a teams construct containing a distribute parallel worksharing-loop SIMD construct and no other statements.

Syntax

The syntax of the teams distribute parallel worksharing-loop SIMD construct is as follows:

C / C++

```c
#pragma omp teams distribute parallel for simd \\
[clause[ , ] clause] ... ] new-line
loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the teams or distribute parallel for simd directives with identical meanings and restrictions.

Fortran

```fortran
!$omp teams distribute parallel do simd [clause[ , ] clause] ... ]
loop-nest

|$omp end teams distribute parallel do simd/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the teams or distribute parallel do simd directives with identical meanings and restrictions.

If an end teams distribute parallel do simd directive is not specified, an end teams distribute parallel do simd directive is assumed at the end of the loop-nest.

Description
The semantics are identical to explicitly specifying a teams directive immediately followed by a distribute parallel worksharing-loop SIMD directive.

Restrictions
The restrictions for the teams and distribute parallel worksharing-loop SIMD constructs apply.
Cross References

- **teams** construct, see Section 2.7.
- Canonical loop nest form, see Section 2.11.1.
- Distribute parallel worksharing-loop SIMD construct, see Section 2.11.6.4.
- Data attribute clauses, see Section 2.21.4.

### 2.16.15 teams loop Construct

**Summary**

The **teams loop** construct is a shortcut for specifying a **teams** construct containing a **loop** construct and no other statements.

**Syntax**

```c
#pragma omp teams loop [clause[ [, ] clause] ... ] new-line
```

where **loop-nest** is a canonical loop nest and **clause** can be any of the clauses accepted by the **teams** or **loop** directives with identical meanings and restrictions.

```fortran
!$omp teams loop [clause[ [, ] clause] ... ]
```

where **loop-nest** is a canonical loop nest and **clause** can be any of the clauses accepted by the **teams** or **loop** directives with identical meanings and restrictions.

If an **end teams loop** directive is not specified, an **end teams loop** directive is assumed at the end of the **loop-nest**.

**Description**

The semantics are identical to explicitly specifying a **teams** directive immediately followed by a **loop** directive.

**Restrictions**

The restrictions for the **teams** and **loop** constructs apply.
Cross References

- teams construct, see Section 2.7.
- Canonical loop nest form, see Section 2.11.1.
- loop construct, see Section 2.11.7.
- Data attribute clauses, see Section 2.21.4.

2.16.16 target parallel Construct

Summary

The target parallel construct is a shortcut for specifying a target construct containing a parallel construct and no other statements.

Syntax

The syntax of the target parallel construct is as follows:

```c
#pragma omp target parallel [clause [ , ] clause] ... ] new-line
structured-block
```

where clause can be any of the clauses accepted by the target or parallel directives, except for copyin, with identical meanings and restrictions.

The syntax of the target parallel construct is as follows:

```fortran
!$omp target parallel [clause [ , ] clause] ... ]
loosely-structured-block
```

or

```fortran
!$omp target parallel [clause [ , ] clause] ... ]
strictly-structured-block
```

where clause can be any of the clauses accepted by the target or parallel directives, except for copyin, with identical meanings and restrictions.

Description

The semantics are identical to explicitly specifying a target directive immediately followed by a parallel directive.
Restrictions
The restrictions for the target and parallel constructs apply except for the following explicit modifications:

- If any if clause on the directive includes a directive-name-modifier then all if clauses on the directive must include a directive-name-modifier.
- At most one if clause without a directive-name-modifier can appear on the directive.
- At most one if clause with the parallel directive-name-modifier can appear on the directive.
- At most one if clause with the target directive-name-modifier can appear on the directive.

Cross References
- parallel construct, see Section 2.6.
- target construct, see Section 2.14.5.
- if clause, see Section 2.18.
- Data attribute clauses, see Section 2.21.4.
- copyin clause, see Section 2.21.6.1.

2.16.17 Target Parallel Worksharing-Loop Construct

Summary
The target parallel worksharing-loop construct is a shortcut for specifying a target construct containing a parallel worksharing-loop construct and no other statements.

Syntax
```
#pragma omp target parallel for [clause[,clause]... ] new-line
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the target or parallel for directives, except for copyin, with identical meanings and restrictions.
The syntax of the target parallel worksharing-loop construct is as follows:

```
!$omp target parallel do [clause[[,] clause] ... ]
loop-nest
/$omp end target parallel do/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `target` or `parallel do` directives, except for `copyin`, with identical meanings and restrictions.

If an `end target parallel do` directive is not specified, an `end target parallel do` directive is assumed at the end of the `loop-nest`.

---

**Description**

The semantics are identical to explicitly specifying a `target` directive immediately followed by a parallel worksharing-loop directive.

**Restrictions**

The restrictions for the `target` and parallel worksharing-loop constructs apply except for the following explicit modifications:

- If any `if` clause on the directive includes a `directive-name-modifier` then all `if` clauses on the directive must include a `directive-name-modifier`.
- At most one `if` clause without a `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `parallel` `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `target` `directive-name-modifier` can appear on the directive.

**Cross References**

- Canonical loop nest form, see Section 2.11.1.
- `target` construct, see Section 2.14.5.
- Parallel Worksharing-Loop construct, see Section 2.16.1.
- `if` clause, see Section 2.18.
- Data attribute clauses, see Section 2.21.4.
- `copyin` clause, see Section 2.21.6.1.
2.16.18 Target Parallel Worksharing-Loop SIMD Construct

Summary
The target parallel worksharing-loop SIMD construct is a shortcut for specifying a `target` construct containing a parallel worksharing-loop SIMD construct and no other statements.

Syntax

C / C++

The syntax of the target parallel worksharing-loop SIMD construct is as follows:

```c
#pragma omp target parallel for simd \\
[clause[, ]clause] ... ] new-line
loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `target` or `parallel for simd` directives, except for `copyin`, with identical meanings and restrictions.

Fortran

The syntax of the target parallel worksharing-loop SIMD construct is as follows:

```fortran
!$omp target parallel do simd [clause[, ]clause] ... ]
loop-nest
/$!omp end target parallel do simd/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `target` or `parallel do simd` directives, except for `copyin`, with identical meanings and restrictions.

If an `end target parallel do simd` directive is not specified, an `end target parallel do simd` directive is assumed at the end of the `loop-nest`.

Description
The semantics are identical to explicitly specifying a `target` directive immediately followed by a parallel worksharing-loop SIMD directive.
Restrictions
The restrictions for the `target` and parallel worksharing-loop SIMD constructs apply except for the following explicit modifications:

- If any `if` clause on the directive includes a `directive-name-modifier` then all `if` clauses on the directive must include a `directive-name-modifier`.
- At most one `if` clause without a `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `parallel` `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `target` `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `simd` `directive-name-modifier` can appear on the directive.

Cross References
- Canonical loop nest form, see Section 2.11.1.
- `target` construct, see Section 2.14.5.
- Parallel worksharing-loop SIMD construct, see Section 2.16.5.
- `if` clause, see Section 2.18.
- Data attribute clauses, see Section 2.21.4.
- `copyin` clause, see Section 2.21.6.1.

2.16.19 target parallel loop Construct

Summary
The `target parallel loop` construct is a shortcut for specifying a `target` construct containing a `parallel loop` construct and no other statements.

Syntax
```
#pragma omp target parallel loop [clause[ [, ] clause] ... ] new-line
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `target` or `parallel loop` directives, except for `copyin`, with identical meanings and restrictions.
The syntax of the `target parallel loop` construct is as follows:

```fortran
!$omp target parallel loop [clause[ , ] clause] ... ]
  loop-nest
}$/omp end target parallel loop/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `target` or `parallel loop` directives, except for `copyin`, with identical meanings and restrictions.

If an `end target parallel loop` directive is not specified, an `end target parallel loop` directive is assumed at the end of the `loop-nest`. `nowait` may not be specified on an `end target parallel loop` directive.

**Description**

The semantics are identical to explicitly specifying a `target` directive immediately followed by a `parallel loop` directive.

**Restrictions**

The restrictions for the `target` and `parallel loop` constructs apply except for the following explicit modifications:

- If any `if` clause on the directive includes a `directive-name-modifier` then all `if` clauses on the directive must include a `directive-name-modifier`.
- At most one `if` clause without a `directive-name-modifier` can appear on the directive.
- At most one `if` clause with the `parallel` directive-name-modifier can appear on the directive.
- At most one `if` clause with the `target` directive-name-modifier can appear on the directive.

**Cross References**

- Canonical loop nest form, see Section 2.11.1.
- `target` construct, see Section 2.14.5.
- `parallel loop` construct, see Section 2.16.2.
- `if` clause, see Section 2.18.
- Data attribute clauses, see Section 2.21.4.
- `copyin` clause, see Section 2.21.6.1.
2.16.20 target simd Construct

Summary
The target simd construct is a shortcut for specifying a target construct containing a simd construct and no other statements.

Syntax

The syntax of the target simd construct is as follows:

```c
#pragma omp target simd [clause[ [, ] clause] ... ] new-line
loop-nest
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the target or simd directives with identical meanings and restrictions.

```fortran
!$omp target simd [clause[ [, ] clause] ... ]
loop-nest
$omp end target simd/
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the target or simd directives with identical meanings and restrictions.

If an end target simd directive is not specified, an end target simd directive is assumed at the end of the loop-nest.

Description
The semantics are identical to explicitly specifying a target directive immediately followed by a simd directive.

Restrictions
The restrictions for the target and simd constructs apply except for the following explicit modifications:

- If any if clause on the directive includes a directive-name-modifier then all if clauses on the directive must include a directive-name-modifier.
- At most one if clause without a directive-name-modifier can appear on the directive.
- At most one if clause with the target directive-name-modifier can appear on the directive.
- At most one if clause with the simd directive-name-modifier can appear on the directive.
Cross References

• Canonical loop nest form, see Section 2.11.1.
• simd construct, see Section 2.11.5.1.
• target construct, see Section 2.14.5.
• if clause, see Section 2.18.
• Data attribute clauses, see Section 2.21.4.

2.16.21 target teams Construct

Summary

The target teams construct is a shortcut for specifying a target construct containing a teams construct and no other statements.

Syntax

The syntax of the target teams construct is as follows:

```
#pragma omp target teams [clause [ , ] clause] ... } new-line
structured-block
```

where clause can be any of the clauses accepted by the target or teams directives with identical meanings and restrictions.

```
!$omp target teams [clause [ , ] clause] ... }
loosely-structured-block

|$omp end target teams
```

or

```
!$omp target teams [clause [ , ] clause] ... }
strictly-structured-block

|$omp end target teams)
```

where clause can be any of the clauses accepted by the target or teams directives with identical meanings and restrictions.

Description

The semantics are identical to explicitly specifying a target directive immediately followed by a teams directive.
Restrictions
The restrictions for the target and teams constructs apply.

Cross References
- teams construct, see Section 2.7.
- target construct, see Section 2.14.5.
- Data attribute clauses, see Section 2.21.4.

2.16.22 target teams distribute Construct

Summary
The target teams distribute construct is a shortcut for specifying a target construct containing a teams distribute construct and no other statements.

Syntax
The syntax of the target teams distribute construct is as follows:
```
#pragma omp target teams distribute [clause[ , ] clause] ... ] new-line
loop-nest
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the target or teams distribute directives with identical meanings and restrictions.

```
!$omp target teams distribute [clause[ , ] clause] ...
loop-nest
[!$omp end target teams distribute/]
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the target or teams distribute directives with identical meanings and restrictions.

If an end target teams distribute directive is not specified, an end target teams distribute directive is assumed at the end of the loop-nest.

Description
The semantics are identical to explicitly specifying a target directive immediately followed by a teams distribute directive.

Restrictions
The restrictions for the target and teams distribute constructs apply.
Cross References

- Canonical loop nest form, see Section 2.11.1.
- `target` construct, see Section 2.14.2.
- `teams distribute` construct, see Section 2.16.11.
- Data attribute clauses, see Section 2.21.4.

2.16.23 `target teams distribute simd` Construct

Summary

The `target teams distribute simd` construct is a shortcut for specifying a `target` construct containing a `teams distribute simd` construct and no other statements.

Syntax

C / C++

```c
#pragma omp target teams distribute simd \\
    [clause[ [,] clause] ... ] new-line
```

Fortran

```fortran
!$omp target teams distribute simd [clause[ [,] clause] ... ]
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `target` or `teams distribute simd` directives with identical meanings and restrictions.

Description

The semantics are identical to explicitly specifying a `target` directive immediately followed by a `teams distribute simd` directive.
Restrictions
The restrictions for the target and teams distribute simd constructs apply except for the following explicit modifications:

- If any if clause on the directive includes a directive-name-modifier then all if clauses on the directive must include a directive-name-modifier.
- At most one if clause without a directive-name-modifier can appear on the directive.
- At most one if clause with the target directive-name-modifier can appear on the directive.
- At most one if clause with the simd directive-name-modifier can appear on the directive.

Cross References
- Canonical loop nest form, see Section 2.11.1.
- target construct, see Section 2.14.2.
- teams distribute simd construct, see Section 2.16.12.
- if clause, see Section 2.18.
- Data attribute clauses, see Section 2.21.4.

2.16.24 target teams loop Construct

Summary
The target teams loop construct is a shortcut for specifying a target construct containing a teams loop construct and no other statements.

Syntax

```
#pragma omp target teams loop [clause[ [, ] clause] ... ] new-line
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the target or teams loop directives with identical meanings and restrictions.
The syntax of the `target teams loop` construct is as follows:

```fortran
!$omp target teams loop [clause[ [, ] clause] ... ]
loop-nest
!$omp end target teams loop/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `target` or `teams loop` directives with identical meanings and restrictions.

If an `end target teams loop` directive is not specified, an `end target teams loop` directive is assumed at the end of the `loop-nest`.

**Description**

The semantics are identical to explicitly specifying a `target` directive immediately followed by a `teams loop` directive.

**Restrictions**

The restrictions for the `target` and `teams loop` constructs apply.

**Cross References**

- Canonical loop nest form, see Section 2.11.1.
- `target` construct, see Section 2.14.5.
- Teams loop construct, see Section 2.16.15.
- Data attribute clauses, see Section 2.21.4.

**Summary**

The target teams distribute parallel worksharing-loop construct is a shortcut for specifying a `target` construct containing a teams distribute parallel worksharing-loop construct and no other statements.
Syntax

The syntax of the target teams distribute parallel worksharing-loop construct is as follows:

```
#pragma omp target teams distribute parallel for \\
[clause[ , clause] ... ] new-line

loop-nest
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `target` or `teams distribute parallel for` directives with identical meanings and restrictions.

```
C / C++

Fortran
```

The syntax of the target teams distribute parallel worksharing-loop construct is as follows:

```
!$omp target teams distribute parallel do [clause[ , clause] ... ]

loop-nest

/$omp end target teams distribute parallel do/
```

where `loop-nest` is a canonical loop nest and `clause` can be any of the clauses accepted by the `target` or `teams distribute parallel do` directives with identical meanings and restrictions.

If an `end target teams distribute parallel do` directive is not specified, an `end target teams distribute parallel do` directive is assumed at the end of the `loop-nest`.

```
Fortran
```

Description

The semantics are identical to explicitly specifying a `target` directive immediately followed by a `teams distribute parallel worksharing-loop` directive.

Restrictions

The restrictions for the `target` and `teams distribute parallel worksharing-loop` constructs apply except for the following explicit modifications:

- If any `if` clause on the directive includes a `directive-name-modifier` then all `if` clauses on the directive must include a `directive-name-modifier`.

- At most one `if` clause without a `directive-name-modifier` can appear on the directive.

- At most one `if` clause with the `parallel` `directive-name-modifier` can appear on the directive.

- At most one `if` clause with the `target` `directive-name-modifier` can appear on the directive.
Cross References

- Canonical loop nest form, see Section 2.11.1.
- target construct, see Section 2.14.5.
- Teams distribute parallel worksharing-loop construct, see Section 2.16.13.
- if clause, see Section 2.18.
- Data attribute clauses, see Section 2.21.4.

2.16.26 Target Teams Distribute Parallel Worksharing-Loop SIMD Construct

Summary

The target teams distribute parallel worksharing-loop SIMD construct is a shortcut for specifying a target construct containing a teams distribute parallel worksharing-loop SIMD construct and no other statements.

Syntax

C / C++

The syntax of the target teams distribute parallel worksharing-loop SIMD construct is as follows:

```
#pragma omp target teams distribute parallel for simd \\
[clause[,clause]... ] new-line
loop-nest
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the target or teams distribute parallel for simd directives with identical meanings and restrictions.

C / C++

Fortran

The syntax of the target teams distribute parallel worksharing-loop SIMD construct is as follows:

```
!$omp target teams distribute parallel do simd [clause[,clause]... ]
loop-nest
```

```
/!$omp end target teams distribute parallel do simd/
```

where loop-nest is a canonical loop nest and clause can be any of the clauses accepted by the target or teams distribute parallel do simd directives with identical meanings and restrictions.

If an end target teams distribute parallel do simd directive is not specified, an end target teams distribute parallel do simd directive is assumed at the end of the loop-nest.
Description
The semantics are identical to explicitly specifying a target directive immediately followed by a teams distribute parallel worksharing-loop SIMD directive.

Restrictions
The restrictions for the target and teams distribute parallel worksharing-loop SIMD constructs apply except for the following explicit modifications:

• If any if clause on the directive includes a directive-name-modifier then all if clauses on the directive must include a directive-name-modifier.

• At most one if clause without a directive-name-modifier can appear on the directive.

• At most one if clause with the parallel directive-name-modifier can appear on the directive.

• At most one if clause with the target directive-name-modifier can appear on the directive.

• At most one if clause with the simd directive-name-modifier can appear on the directive.

Cross References
• Canonical loop nest form, see Section 2.11.1.

• target construct, see Section 2.14.5.

• Teams distribute parallel worksharing-loop SIMD construct, see Section 2.16.14.

• if clause, see Section 2.18.

• Data attribute clauses, see Section 2.21.4.

2.17 Clauses on Combined and Composite Constructs

This section specifies the handling of clauses on combined or composite constructs and the handling of implicit clauses from variables with predetermined data sharing if they are not predetermined only on a particular construct. Some clauses are permitted only on a single leaf construct of the combined or composite construct, in which case the effect is as if the clause is applied to that specific construct. As detailed in this section, other clauses have the effect as if they are applied to one or more leaf constructs.

The collapse clause is applied once to the combined or composite construct.

The effect of the private clause is as if it is applied only to the innermost leaf construct that permits it.
The effect of the `firstprivate` clause is as if it is applied to one or more leaf constructs as follows:

- To the `distribute` construct if it is among the constituent constructs;
- To the `teams` construct if it is among the constituent constructs and the `distribute` construct is not;
- To the worksharing-loop construct if it is among the constituent constructs;
- To the `taskloop` construct if it is among the constituent constructs;
- To the `parallel` construct if it is among the constituent constructs and the worksharing-loop construct or the `taskloop` construct is not;
- To the `target` construct if it is among the constituent constructs and the same list item does not appear in a `lastprivate` or `map` clause.

If the `parallel` construct is among the constituent constructs and the effect is not as if the `firstprivate` clause is applied to it by the above rules, then the effect is as if the `shared` clause with the same list item is applied to the `parallel` construct. If the `teams` construct is among the constituent constructs and the effect is not as if the `firstprivate` clause is applied to it by the above rules, then the effect is as if the `shared` clause with the same list item is applied to the `teams` construct.

The effect of the `lastprivate` clause is as if it is applied to all leaf constructs that permit the clause. If the `parallel` construct is among the constituent constructs and the list item is not also specified in the `firstprivate` clause, then the effect of the `lastprivate` clause is as if the `shared` clause with the same list item is applied to the `parallel` construct. If the `teams` construct is among the constituent constructs and the list item is not also specified in the `firstprivate` clause, then the effect of the `lastprivate` clause is as if the `shared` clause with the same list item is applied to the `teams` construct. If the `target` construct is among the constituent constructs and the list item is not specified in a `map` clause, the effect of the `lastprivate` clause is as if the same list item appears in a `map` clause with a `map-type` of `tofrom`.

The effect of the `shared`, `default`, `order`, or `allocate` clause is as if it is applied to all leaf constructs that permit the clause.

The effect of the `reduction` clause is as if it is applied to all leaf constructs that permit the clause, except for the following constructs:

- The `parallel` construct, when combined with the `sections`, worksharing-loop, `loop`, or `taskloop` construct; and
- The `teams` construct, when combined with the `loop` construct.

For the `parallel` and `teams` constructs above, the effect of the `reduction` clause instead is as if each list item or, for any list item that is an array item, its corresponding base array or base pointer appears in a `shared` clause for the construct. If the `task reduction-modifier` is specified,
the effect is as if it only modifies the behavior of the reduction clause on the innermost leaf construct that accepts the modifier (see Section 2.21.5.4). If the inscan reduction-modifier is specified, the effect is as if it modifies the behavior of the reduction clause on all constructs of the combined construct to which the clause is applied and that accept the modifier. If the target construct is among the constituent constructs and the list item is not specified in a map clause, the effect of the reduction clause is as if the same list item appears in a map clause with a map-type of tofrom.

The in_reduction clause applies to the single leaf construct on which it is permitted. If that construct is a target construct, the effect is as if the same list item also appears in a map clause with a map-type of tofrom and a map-type-modifier of always.

The effect of the if clause is described in Section 2.18.

The effect of the linear clause is as if it is applied to the innermost leaf construct. Additionally, if the list item is not the iteration variable of a simd or worksharing-loop SIMD construct, the effect on the outer leaf constructs is as if the list item was specified in firstprivate and lastprivate clauses on the combined or composite construct, with the rules specified above applied. If a list item of the linear clause is the iteration variable of a simd or worksharing-loop SIMD construct and it is not declared in the construct, the effect on the outer leaf constructs is as if the list item was specified in a lastprivate clause on the combined or composite construct with the rules specified above applied.

The effect of the nowait clause is as if it is applied to the outermost leaf construct that permits it.

If the clauses have expressions on them, such as for various clauses where the argument of the clause is an expression, or lower-bound, length, or stride expressions inside array sections (or subscript and stride expressions in subscript-triplet for Fortran), or linear-step or alignment expressions, the expressions are evaluated immediately before the construct to which the clause has been split or duplicated per the above rules (therefore inside of the outer leaf constructs). However, the expressions inside the num_teams and thread_limit clauses are always evaluated before the outermost leaf construct.

The restriction that a list item may not appear in more than one data sharing clause with the exception of specifying a variable in both firstprivate and lastprivate clauses applies after the clauses are split or duplicated per the above rules.

### 2.18 if Clause

**Summary**
The semantics of an if clause are described in the section on the construct to which it applies. The if clause directive-name-modifier names the associated construct to which an expression applies, and is particularly useful for composite and combined constructs.
The syntax of the if clause is as follows:

| if([ directive-name-modifier : ] scalar-expression) |

The syntax of the if clause is as follows:

| if([ directive-name-modifier : ] scalar-logical-expression) |

Description

The effect of the if clause depends on the construct to which it is applied. For combined or composite constructs, the if clause only applies to the semantics of the construct named in the directive-name-modifier if one is specified. If no directive-name-modifier is specified for a combined or composite construct then the if clause applies to all constructs to which an if clause can apply.

2.19 Synchronization Constructs and Clauses

A synchronization construct orders the completion of code executed by different threads. This ordering is imposed by synchronizing flush operations that are executed as part of the region that corresponds to the construct.

Synchronization through the use of synchronizing flush operations and atomic operations is described in Section 1.4.4 and Section 1.4.6. Section 2.19.8.1 defines the behavior of synchronizing flush operations that are implied at various other locations in an OpenMP program.

2.19.1 critical Construct

Summary

The critical construct restricts execution of the associated structured block to a single thread at a time.

Syntax

```
#pragma omp critical [ (name) ] [ , ] hint(hint-expression) ] ] new-line
structured-block
```

where hint-expression is an integer constant expression that evaluates to a valid synchronization hint (as described in Section 2.19.12).
The syntax of the `critical` construct is as follows:

```fortran
!$omp critical [ (name) [, ] hint (hint-expression) ]  
loosely-structured-block

|$omp end critical [ (name) ]
```

or

```fortran
!$omp critical [ (name) [, ] hint (hint-expression) ]  
strictly-structured-block

|$omp end critical [ (name) ]
```

where `hint-expression` is a constant expression that evaluates to a scalar value with kind `omp_sync_hint_kind` and a value that is a valid synchronization hint (as described in Section 2.19.12).

**Binding**

The binding thread set for a `critical` region is all threads in the contention group.

**Description**

The region that corresponds to a `critical` construct is executed as if only a single thread at a time among all threads in the contention group enters the region for execution, without regard to the teams to which the threads belong. An optional `name` may be used to identify the `critical` construct. All `critical` constructs without a name are considered to have the same unspecified name.

Identifiers used to identify a `critical` construct have external linkage and are in a name space that is separate from the name spaces used by labels, tags, members, and ordinary identifiers.

The names of `critical` constructs are global entities of the program. If a name conflicts with any other entity, the behavior of the program is unspecified.

The threads of a contention group execute the `critical` region as if only one thread of the contention group executes the `critical` region at a time. The `critical` construct enforces these execution semantics with respect to all `critical` constructs with the same name in all threads in the contention group.

If present, the `hint` clause gives the implementation additional information about the expected runtime properties of the `critical` region that can optionally be used to optimize the implementation. The presence of a `hint` clause does not affect the isolation guarantees provided by the `critical` construct. If no `hint` clause is specified, the effect is as if `hint (omp_sync_hint_none)` had been specified.
Execution Model Events
The critical-acquiring event occurs in a thread that encounters the critical construct on entry to the critical region before initiating synchronization for the region.

The critical-acquired event occurs in a thread that encounters the critical construct after it enters the region, but before it executes the structured block of the critical region.

The critical-released event occurs in a thread that encounters the critical construct after it completes any synchronization on exit from the critical region.

Tool Callbacks
A thread dispatches a registered ompt_callback_mutex_acquire callback for each occurrence of a critical-acquiring event in that thread. This callback has the type signature ompt_callback_mutex_acquire_t.

A thread dispatches a registered ompt_callback_mutex_acquired callback for each occurrence of a critical-acquired event in that thread. This callback has the type signature ompt_callback_mutex_t.

A thread dispatches a registered ompt_callback_mutex_released callback for each occurrence of a critical-released event in that thread. This callback has the type signature ompt_callback_mutex_t.

The callbacks occur in the task that encounters the critical construct. The callbacks should receive ompt_mutex_critical as their kind argument if practical, but a less specific kind is acceptable.

Restrictions
Restrictions to the critical construct are as follows:

- Unless the effect is as if hint(omp_sync_hint_none) was specified, the critical construct must specify a name.
- If the hint clause is specified, each of the critical constructs with the same name must have a hint clause for which the hint-expression evaluates to the same value.
- A throw executed inside a critical region must cause execution to resume within the same critical region, and the same thread that threw the exception must catch it.
- If a name is specified on a critical directive, the same name must also be specified on the end critical directive.
- If no name appears on the critical directive, no name can appear on the end critical directive.
Cross References

• Synchronization Hints, see Section 2.19.12.
• ompt_mutex_critical, see Section 4.4.4.16.
• ompt_callback_mutex_acquire_t, see Section 4.5.2.14.
• ompt_callback_mutex_t, see Section 4.5.2.15.

2.19.2 barrier Construct

Summary

The barrier construct specifies an explicit barrier at the point at which the construct appears. The barrier construct is a stand-alone directive.

Syntax

C / C++

The syntax of the barrier construct is as follows:

```
#pragma omp barrier new-line
```

Fortran

The syntax of the barrier construct is as follows:

```
!$omp barrier
```

Binding

The binding thread set for a barrier region is the current team. A barrier region binds to the innermost enclosing parallel region.

Description

All threads of the team that is executing the binding parallel region must execute the barrier region and complete execution of all explicit tasks bound to this parallel region before any are allowed to continue execution beyond the barrier.

The barrier region includes an implicit task scheduling point in the current task region.

Execution Model Events

The explicit-barrier-begin event occurs in each thread that encounters the barrier construct on entry to the barrier region.

The explicit-barrier-wait-begin event occurs when a task begins an interval of active or passive waiting in a barrier region.
The explicit-barrier-wait-end event occurs when a task ends an interval of active or passive waiting and resumes execution in a barrier region.

The explicit-barrier-end event occurs in each thread that encounters the barrier construct after the barrier synchronization on exit from the barrier region.

A cancellation event occurs if cancellation is activated at an implicit cancellation point in a barrier region.

Tool Callbacks
A thread dispatches a registered ompt_callback_sync_region callback with ompt_sync_region_barrier_explicit as its kind argument and ompt_scope_begin as its endpoint argument for each occurrence of an explicit-barrier-begin event. Similarly, a thread dispatches a registered ompt_callback_sync_region callback with ompt_sync_region_barrier_explicit as its kind argument and ompt_scope_end as its endpoint argument for each occurrence of an explicit-barrier-end event. These callbacks occur in the context of the task that encountered the barrier construct and have type signature ompt_callback_sync_region_t.

A thread dispatches a registered ompt_callback_sync_region_wait callback with ompt_sync_region_barrier_explicit as its kind argument and ompt_scope_begin as its endpoint argument for each occurrence of an explicit-barrier-wait-begin event. Similarly, a thread dispatches a registered ompt_callback_sync_region_wait callback with ompt_sync_region_barrier_explicit as its kind argument and ompt_scope_end as its endpoint argument for each occurrence of an explicit-barrier-wait-end event. These callbacks occur in the context of the task that encountered the barrier construct and have type signature ompt_callback_sync_region_t.

A thread dispatches a registered ompt_callback_cancel callback with ompt_cancel_detected as its flags argument for each occurrence of a cancellation event in that thread. The callback occurs in the context of the encountering task. The callback has type signature ompt_callback_cancel_t.

Restrictions
Restrictions to the barrier construct are as follows:

- Each barrier region must be encountered by all threads in a team or by none at all, unless cancellation has been requested for the innermost enclosing parallel region.

- The sequence of worksharing regions and barrier regions encountered must be the same for every thread in a team.
Cross References

- `ompt_scope_begin` and `ompt_scope_end`, see Section 4.4.4.11.
- `ompt_sync_region_barrier`, see Section 4.4.4.13.
- `ompt_callback_sync_region_t`, see Section 4.5.2.13.
- `ompt_callback_cancel_t`, see Section 4.5.2.18.

2.19.3 Implicit Barriers

This section describes the OMPT events and tool callbacks associated with implicit barriers, which occur at the end of various regions as defined in the description of the constructs to which they correspond. Implicit barriers are task scheduling points. For a description of task scheduling points, associated events, and tool callbacks, see Section 2.12.6.

Execution Model Events

The `implicit-barrier-begin` event occurs in each implicit task at the beginning of an implicit barrier region.

The `implicit-barrier-wait-begin` event occurs when a task begins an interval of active or passive waiting in an implicit barrier region.

The `implicit-barrier-wait-end` event occurs when a task ends an interval of active or waiting and resumes execution of an implicit barrier region.

The `implicit-barrier-end` event occurs in each implicit task after the barrier synchronization on exit from an implicit barrier region.

A `cancellation` event occurs if cancellation is activated at an implicit cancellation point in an implicit barrier region.

Tool Callbacks

A thread dispatches a registered `ompt_callback_sync_region` callback for each implicit barrier `begin` and `end` event. Similarly, a thread dispatches a registered `ompt_callback_sync_region_wait` callback for each implicit barrier `wait-begin` and `wait-end` event. All callbacks for implicit barrier events execute in the context of the encountering task and have type signature `ompt_callback_sync_region_t`.

For the implicit barrier at the end of a worksharing construct, the `kind` argument is `ompt_sync_region_barrier_implicit_workshare`. For the implicit barrier at the end of a `parallel` region, the `kind` argument is `ompt_sync_region_barrier_implicit_parallel`. For an extra barrier added by an OpenMP implementation, the `kind` argument is `ompt_sync_region_barrier_implementation`. For a barrier at the end of a `teams` region, the `kind` argument is `ompt_sync_region_barrier_teams`. 

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A thread dispatches a registered `ompt_callback_cancel` callback with `ompt_cancel_detected` as its `flags` argument for each occurrence of a `cancellation` event in that thread. The callback occurs in the context of the encountering task. The callback has type signature `ompt_callback_cancel_t`.

**Restrictions**
Restrictions to implicit barriers are as follows:

- If a thread is in the state `ompt_state_wait_barrier_implicit_parallel`, a call to `omp_get_parallel_info` may return a pointer to a copy of the data object associated with the parallel region rather than a pointer to the associated data object itself. Writing to the data object returned by `omp_get_parallel_info` when a thread is in the `ompt_state_wait_barrier_implicit_parallel` results in unspecified behavior.

**Cross References**
- `ompt_scope_begin` and `ompt_scope_end`, see Section 4.4.4.11.
- `ompt_sync_region_barrier_implementation`, `ompt_sync_region_barrier_implicit_parallel`, `ompt_sync_region_barrier_teams`, and `ompt_sync_region_barrier_implicit_workshare`, see Section 4.4.4.13.
- `ompt_cancel_detected`, see Section 4.4.4.25.
- `ompt_callback_sync_region_t`, see Section 4.5.2.13.
- `ompt_callback_cancel_t`, see Section 4.5.2.18.

### 2.19.4 Implementation-Specific Barriers

An OpenMP implementation can execute implementation-specific barriers that are not implied by the OpenMP specification; therefore, no execution model events are bound to these barriers. The implementation can handle these barriers like implicit barriers and dispatch all events as for implicit barriers. These callbacks are dispatched with `ompt_sync_region_barrier_implementation` — or `ompt_sync_region_barrier`, if the implementation cannot make a distinction — as the `kind` argument.

### 2.19.5 taskwait Construct

**Summary**
The `taskwait` construct specifies a wait on the completion of child tasks of the current task. The `taskwait` construct is a stand-alone directive.
Syntax

The syntax of the `taskwait` construct is as follows:

```c
#pragma omp taskwait [clause[ , clause] ... ]
```

where `clause` is one of the following:

- `depend([depend-modifier, ]dependence-type : locator-list)`
- `nowait`

---

The syntax of the `taskwait` construct is as follows:

```fortran
!$omp taskwait [clause[ , clause] ... ]
```

where `clause` is one of the following:

- `depend([depend-modifier, ]dependence-type : locator-list)`
- `nowait`

---

Binding

The `taskwait` region binds to the current task region. The binding thread set of the `taskwait` region is the current team.

Description

If no `depend` clause is present on the `taskwait` construct, the current task region is suspended at an implicit task scheduling point associated with the construct. The current task region remains suspended until all child tasks that it generated before the `taskwait` region complete execution.

If one or more `depend` clauses are present on the `taskwait` construct and the `nowait` clause is not also present, the behavior is as if these clauses were applied to a `task` construct with an empty associated structured block that generates a `mergeable` and `included task`. Thus, the current task region is suspended until the `predecessor tasks` of this task complete execution.

If one or more `depend` clauses are present on the `taskwait` construct and the `nowait` clause is also present, the behavior is as if these clauses were applied to a `task` construct with an empty associated structured block that generates a task for which execution may be deferred. Thus, all `predecessor tasks` of this task must complete execution before any subsequently generated task that depends on this task starts its execution.
Execution Model Events

The `taskwait-begin` event occurs in a thread when it encounters a `taskwait` construct with no `depend` clause on entry to the `taskwait` region.

The `taskwait-wait-begin` event occurs when a task begins an interval of active or passive waiting in a region corresponding to a `taskwait` construct with no `depend` clause.

The `taskwait-wait-end` event occurs when a task ends an interval of active or passive waiting and resumes execution from a region corresponding to a `taskwait` construct with no `depend` clause.

The `taskwait-end` event occurs in a thread when it encounters a `taskwait` construct with no `depend` clause after the taskwait synchronization on exit from the `taskwait` region.

The `taskwait-init` event occurs in a thread when it encounters a `taskwait` construct with one or more `depend` clauses on entry to the `taskwait` region.

The `taskwait-complete` event occurs on completion of the dependent task that results from a `taskwait` construct with one or more `depend` clauses, in the context of the thread that executes the dependent task and before any subsequently generated task that depends on the dependent task starts its execution.

Tool Callbacks

A thread dispatches a registered `ompt_callback_sync_region` callback with `ompt_sync_region_taskwait` as its `kind` argument and `ompt_scope_begin` as its `endpoint` argument for each occurrence of a `taskwait-begin` event in the task that encounters the `taskwait` construct. Similarly, a thread dispatches a registered `ompt_callback_sync_region` callback with `ompt_sync_region_taskwait` as its `kind` argument and `ompt_scope_end` as its `endpoint` argument for each occurrence of a `taskwait-end` event in the task that encounters the `taskwait` construct. These callbacks occur in the task that encounters the `taskwait` construct and have the type signature `ompt_callback_sync_region_t`.

A thread dispatches a registered `ompt_callback_sync_region_wait` callback with `ompt_sync_region_taskwait` as its `kind` argument and `ompt_scope_begin` as its `endpoint` argument for each occurrence of a `taskwait-wait-begin` event. Similarly, a thread dispatches a registered `ompt_callback_sync_region_wait` callback with `ompt_sync_region_taskwait` as its `kind` argument and `ompt_scope_end` as its `endpoint` argument for each occurrence of a `taskwait-wait-end` event. These callbacks occur in the context of the task that encounters the `taskwait` construct and have type signature `ompt_callback_sync_region_t`.

A thread dispatches a registered `ompt_callback_task_create` callback for each occurrence of a `taskwait-init` event in the context of the encountering task. This callback has the type signature `ompt_callback_task_create_t`. In the dispatched callback, `(flags & ompt_task_taskwait)` always evaluates to `true`. If the `nowait` clause is not present, `(flags & ompt_task_undeferred)` also evaluates to `true`.
A thread dispatches a registered `ompt_callback_task_schedule` callback for each occurrence of a `taskwait-complete` event. This callback has the type signature `ompt_callback_task_schedule_t` with `ompt_taskwait_complete` as its `prior_task_status` argument.

**Restrictions**

Restrictions to the `taskwait` construct are as follows:

- The `mutexinoutset` dependence-type may not appear in a `depend` clause on a `taskwait` construct.
- If the `dependence-type` of a `depend` clause is `depobj` then the dependence objects cannot represent dependences of the `mutexinoutset` dependence type.
- The `nowait` clause may only appear on a `taskwait` directive if the `depend` clause is present.
- At most one `nowait` clause can appear on a `taskwait` directive.

**Cross References**

- `task` construct, see Section 2.12.1.
- Task scheduling, see Section 2.12.6.
- `depend` clause, see Section 2.19.11.
- `ompt_scope_begin` and `ompt_scope_end`, see Section 4.4.4.11.
- `ompt_sync_region_taskwait`, see Section 4.4.4.13.
- `ompt_callback_sync_region_t`, see Section 4.5.2.13.

### 2.19.6 taskgroup Construct

**Summary**

The `taskgroup` construct specifies a wait on completion of child tasks of the current task and their descendnet tasks.

**Syntax**

```
#pragma omp taskgroup [clause[[, ] clause]...] new-line
structured-block
```

where `clause` is one of the following:

- `task_reduction(reduction-identifier : list)`
- `allocate([allocator: ]list)`
The syntax of the `taskgroup` construct is as follows:

```
!$omp taskgroup [clause [ , clause] ...]
  loosely-structured-block
!$omp end taskgroup
```

or

```
!$omp taskgroup [clause [ , clause] ...]
  strictly-structured-block
[!$omp end taskgroup]
```

where `clause` is one of the following:

```
task_reduction(reduction-identifier : list)
allocate([allocator : ]list)
```

**Binding**

The binding task set of a `taskgroup` region is all tasks of the current team that are generated in the region. A `taskgroup` region binds to the innermost enclosing `parallel` region.

**Description**

When a thread encounters a `taskgroup` construct, it starts executing the region. All child tasks generated in the `taskgroup` region and all of their descendants that bind to the same `parallel` region as the `taskgroup` region are part of the `taskgroup set` associated with the `taskgroup` region.

An implicit task scheduling point occurs at the end of the `taskgroup` region. The current task is suspended at the task scheduling point until all tasks in the `taskgroup set` complete execution.

**Execution Model Events**

The `taskgroup-begin` event occurs in each thread that encounters the `taskgroup` construct on entry to the `taskgroup` region.

The `taskgroup-wait-begin` event occurs when a task begins an interval of active or passive waiting in a `taskgroup` region.

The `taskgroup-wait-end` event occurs when a task ends an interval of active or passive waiting and resumes execution in a `taskgroup` region.

The `taskgroup-end` event occurs in each thread that encounters the `taskgroup` construct after the taskgroup synchronization on exit from the `taskgroup` region.
Tool Callbacks
A thread dispatches a registered `ompt_callback_sync_region` callback with
`ompt_sync_region_taskgroup` as its `kind` argument and `ompt_scope_begin` as its
`endpoint` argument for each occurrence of a `taskgroup-begin` event in the task that encounters the
taskgroup construct. Similarly, a thread dispatches a registered
`ompt_callback_sync_region` callback with `ompt_sync_region_taskgroup` as its
`kind` argument and `ompt_scope_end` as its `endpoint` argument for each occurrence of a
`taskgroup-end` event in the task that encounters the taskgroup construct. These callbacks occur
in the task that encounters the taskgroup construct and have the type signature
`ompt_callback_sync_region_t`.
A thread dispatches a registered `ompt_callback_sync_region_wait` callback with
`ompt_sync_region_taskgroup` as its `kind` argument and `ompt_scope_begin` as its
`endpoint` argument for each occurrence of a `taskgroup-wait-begin` event. Similarly, a thread
dispatches a registered `ompt_callback_sync_region_wait` callback with
`ompt_sync_region_taskgroup` as its `kind` argument and `ompt_scope_end` as its
`endpoint` argument for each occurrence of a `taskgroup-wait-end` event. These callbacks occur in the
context of the task that encounters the taskgroup construct and have type signature
`ompt_callback_sync_region_t`.

Cross References
• Task scheduling, see Section 2.12.6.
• `task_reduction` clause, see Section 2.21.5.5.
• `ompt_scope_begin` and `ompt_scope_end`, see Section 4.4.4.11.
• `ompt_sync_region_taskgroup`, see Section 4.4.4.13.
• `ompt_callback_sync_region_t`, see Section 4.5.2.13.

2.19.7 atomic Construct
Summary
The atomic construct ensures that a specific storage location is accessed atomically, rather than
exposing it to the possibility of multiple, simultaneous reading and writing threads that may result
in indeterminate values.
Syntax
In the following syntax, `atomic-clause` is a clause that indicates the semantics for which atomicity is
enforced, and `memory-order-clause` is a clause that indicates the memory ordering behavior of the
construct. Specifically, `atomic-clause` is one of the following:

| read |
| write |
| update |
memory-order-clause is one of the following:

- seq_cst
- acq_rel
- release
- acquire
- relaxed

and clause is either atomic-clause, memory-order-clause or one of the following:

- capture
- compare
- hint (hint-expression)
- fail(seq_cst | acquire | relaxed)
- weak

The syntax of the atomic construct is:

```c
#pragma omp atomic [clause[ , clause] ... ] new-line statement
```

where statement is a statement with one of the following forms:

- If atomic-clause is read then statement is read-expr-stmt, a read expression statement that has the following form:

  ```
  v = x;
  ```

- If atomic-clause is write then statement is write-expr-stmt, a write expression statement that has the following form:

  ```
  x = expr;
  ```

- If atomic-clause is update then statement can be update-expr-stmt, an update expression statement that has one of the following forms:

  ```
  x++; 
  x--; 
  ++x; 
  --x; 
  x binop expr; 
  x = x binop expr; 
  x = expr binop x;
  ```
• If the **compare** clause is present then either **statement** is:

  - **cond-expr-stmt**, a conditional expression statement that has one of the following forms:

    ```
    x = expr ordop x ? expr : x;
    x = x ordop expr ? expr : x;
    x = x == e ? d : x;
    ```

  - or **cond-update-stmt**, a conditional update statement that has one of the following forms:

    ```
    if(expr ordop x) { x = expr; }
    if(x ordop expr) { x = expr; }
    if(x == e) { x = d; }
    ```

• If the **capture** clause is present, **statement** can have one of the following forms:

  ```
  v = expr-stmt
  { v = x; expr-stmt }
  { expr-stmt v = x; }
  ```

  where **expr-stmt** is either **write-expr-stmt**, **update-expr-stmt** or **cond-expr-stmt**.

• If both the **compare** and **capture** clauses are present then the following forms are also valid:

  ```
  { v = x; cond-update-stmt }
  { cond-update-stmt v = x; }
  if(x == e) { x = d; } else { v = x; }
  { r = x == e; if(r) { x = d; } }
  { r = x == e; if(r) { x = d; } else { v = x; } }
  ```

In the preceding expressions:

• **x**, **r** (result), and **v** (as applicable) are **lvalue** expressions with scalar type.

• **e** (expected) is an expression with scalar type, in forms where **e** is assigned it must be an lvalue.

• **d** (desired) is an expression with scalar type.

• **r** must be of integral type.

• During the execution of an **atomic** region, multiple syntactic occurrences of **x** must designate the same storage location.

• None of **v**, **x**, **r**, **d** and **expr** (as applicable) may access the storage location designated by any other in the list.

• **e** and **v** may refer to, or access, the same storage location.

• **expr** is an expression with scalar type.

• The order operation, **ordop**, is one of `<`, `>`, `<=`, `>=`, `==`, `!=`, `&&`, `||`, `&`, `|`, `<<`, or `>>`.
• binop, binop=, ordop, ==, ++, and -- are not overloaded operators.

• The expression $x \text{ binop } expr$ must be numerically equivalent to $x \text{ binop } (expr)$. This requirement is satisfied if the operators in $expr$ have precedence greater than $binop$, or by using parentheses around $expr$ or subexpressions of $expr$.

• The expression $expr \text{ binop } x$ must be numerically equivalent to $(expr) \text{ binop } x$. This requirement is satisfied if the operators in $expr$ have precedence equal to or greater than $binop$, or by using parentheses around $expr$ or subexpressions of $expr$.

• == comparisons are performed by comparing the bits that comprise each object as with memcmp.

• For forms that allow multiple occurrences of $x$, the number of times that $x$ is evaluated is unspecified.

• hint-expression is a constant integer expression that evaluates to a valid synchronization hint.

The syntax of the atomic construct takes any of the following forms:

```
!$omp atomic [clause[...]]
statement
[!$omp end atomic]
```

or

```
!$omp atomic [clause[...]] capture[,...]
capture-statement
[!$omp end atomic]
```

or

```
!$omp atomic [clause[...]] capture[,...]
capture-statement
statement
[!$omp end atomic]
```

where capture-statement has the following form:

```
v = x
```

and statement is a statement with one of the following forms:

• If atomic-clause is read then statement is:

```
v = x;
```

• If atomic-clause is write then statement is:

```
x = expr
```
• If `atomic-clause` is `update` then:
  - `statement` can have one of the following forms:
    ```
    x = x operator expr
    x = expr operator x
    x = intrinsic-procedure-name (x, expr-list)
    x = intrinsic-procedure-name (expr-list, x)
    ```
  - or, if the `capture` clause is present and `statement` is preceded or followed by `capture-statement`, `statement` can also have this form:
    ```
    x = expr
    ```

• If the `compare` clause is present then:
  - `statement` has one of these forms:
    ```
    if (x == e) then
      x = d
    end if
    ```
  - or, if the `capture` clause is also present and `statement` is not preceded or followed by `capture-statement`, `statement` has this form:
    ```
    if (x == e) then
      x = d
    else
      v = x
    end if
    ```

In the preceding statements:
• `x, v, d` and `e` (as applicable) are scalar variables of intrinsic type.
• `x` must not have the `ALLOCATABLE` attribute.
• During the execution of an atomic region, multiple syntactic occurrences of `x` must designate the same storage location.
• None of `v, expr, and expr-list` (as applicable) may access the same storage location as `x`.
• None of `x, expr, and expr-list` (as applicable) may access the same storage location as `v`.
• `expr` is a scalar expression.
• `expr-list` is a comma-separated, non-empty list of scalar expressions. If `intrinsic-procedure-name` refers to `IAND, IOR, or IEO`, exactly one expression must appear in `expr-list`.
• `intrinsic-procedure-name` is one of `MAX, MIN, IAND, IOR, or IEO`.
• `operator` is one of `+, *, -, /, .AND., .OR., .EQV., or .NEQV.`
• The expression \( x \ operator \ expr \) must be numerically equivalent to \( x \ operator (expr) \). This requirement is satisfied if the operators in \( expr \) have precedence greater than \( operator \), or by using parentheses around \( expr \) or subexpressions of \( expr \).

• The expression \( expr \ operator \ x \) must be numerically equivalent to \( (expr) \ operator \ x \). This requirement is satisfied if the operators in \( expr \) have precedence equal to or greater than \( operator \), or by using parentheses around \( expr \) or subexpressions of \( expr \).

• \textit{intrinsic-procedure-name} must refer to the intrinsic procedure name and not to other program entities.

• \textit{operator} must refer to the intrinsic operator and not to a user-defined operator.

• All assignments must be intrinsic assignments.

• For forms that allow multiple occurrences of \( x \), the number of times that \( x \) is evaluated is unspecified.

• \textit{hint-expression} is a constant expression that evaluates to a scalar value with kind \texttt{omp_sync_hint_kind} and a value that is a valid synchronization hint.

\begin{fortran}
Binding
\end{fortran}

If the size of \( x \) is 8, 16, 32, or 64 bits and \( x \) is aligned to a multiple of its size, the binding thread set for the \texttt{atomic} region is all threads on the device. Otherwise, the binding thread set for the \texttt{atomic} region is all threads in the contention group. \texttt{atomic} regions enforce exclusive access with respect to other \texttt{atomic} regions that access the same storage location \( x \) among all threads in the binding thread set without regard to the teams to which the threads belong.

\begin{fortran}
Description
\end{fortran}

If \texttt{atomic-clause} is not present on the construct, the behavior is as if the \texttt{update} clause is specified.

The \texttt{atomic} construct with the \texttt{read} clause results in an atomic read of the location designated by \( x \).

The \texttt{atomic} construct with the \texttt{write} clause results in an atomic write of the location designated by \( x \).

The \texttt{atomic} construct with the \texttt{update} clause results in an atomic update of the location designated by \( x \) using the designated operator or intrinsic. Only the read and write of the location designated by \( x \) are performed mutually atomically. The evaluation of \( expr \) or \( expr-list \) need not be atomic with respect to the read or write of the location designated by \( x \). No task scheduling points are allowed between the read and the write of the location designated by \( x \).

If the \texttt{capture} clause is present, the atomic update is an atomic captured update — an atomic update to the location designated by \( x \) using the designated operator or intrinsic while also capturing the original or final value of the location designated by \( x \) with respect to the atomic update. The original or final value of the location designated by \( x \) is written in the location.
designated by $v$ based on the base language semantics of structured block or statements of the
\texttt{atomic} construct. Only the read and write of the location designated by $x$ are performed mutually
atomically. Neither the evaluation of $\texttt{expr}$ or $\texttt{expr-list}$, nor the write to the location designated by $v$,
need be atomic with respect to the read or write of the location designated by $x$.

If the \texttt{compare} clause is present, the atomic update is an atomic conditional update. For forms
that use an equality comparison, the operation is an atomic compare-and-swap — it atomically
compares the value of $x$ to $e$ and if they are equal writes the value of $d$ into the location designated
by $x$. Based on the base language semantics of the associated structured block of the \texttt{atomic}
construct, the original or final value of the location designated by $x$ is written to the location
designated by $v$, which is allowed to be the same location as designated by $e$, or the result of the
comparison is written to the location designated by $r$. Only the read and write of the location
designated by $x$ are performed mutually atomically. Neither the evaluation of either $e$ or $d$ nor
writes to the locations designated by $v$ and $r$ need be atomic with respect to the read or write of the
location designated by $x$.

\begin{quote}
\textbf{C / C++}
\end{quote}

If the \texttt{compare} clause is present, forms that use \texttt{ordop} are logically an atomic maximum or
minimum, but they may be implemented with a compare-and-swap loop with short-circuiting. For
forms where \texttt{statement} is \texttt{cond-expr-stmt}, if the result of the condition implies that the value of $x$
does not change then the update may not occur.

\begin{quote}
\textbf{C / C++}
\end{quote}

If the \texttt{weak} clause is present, the comparison performed by an atomic compare-and-swap operation
may spuriously fail, evaluating to not equal even when the values are equal.

\begin{quote}
\textbf{Note} – Allowing for spurious failure by specifying a \texttt{weak} clause can result in performance gains
on some systems when using compare-and-swap in a loop. For cases where a single
compare-and-swap would otherwise be sufficient, using a loop over a \texttt{weak} compare-and-swap is
unlikely to improve performance.
\end{quote}

\begin{quote}
If \texttt{memory-order-clause} is present, or implicitly provided by a \texttt{requires} directive, it specifies the
effective memory ordering and otherwise the effective memory ordering is \texttt{relaxed}. If the \texttt{fail}
clause is present, its parameter overrides the effective memory ordering used if the comparison for
an atomic conditional update fails.

The \texttt{atomic} construct may be used to enforce memory consistency between threads, based on the
guarantees provided by Section 1.4.6. A strong flush on the location designated by $x$ is performed
on entry to and exit from the atomic operation, ensuring that the set of all atomic operations applied
to the same location in a race-free program has a total completion order. If the \texttt{write} or \texttt{update}
clause is specified, the atomic operation is not an atomic conditional update for which the
comparison fails, and the effective memory ordering is \texttt{release}, \texttt{acq_rel}, or \texttt{seq_cst}, the
strong flush on entry to the atomic operation is also a release flush. If the \texttt{read} or \texttt{update} clause
is specified and the effective memory ordering is \texttt{acquire}, \texttt{acq_rel}, or \texttt{seq_cst} then the
strong flush on exit from the atomic operation is also an acquire flush. Therefore, if the effective memory ordering is not relaxed, release and/or acquire flush operations are implied and permit synchronization between the threads without the use of explicit flush directives.

For all forms of the atomic construct, any combination of two or more of these atomic constructs enforces mutually exclusive access to the locations designated by $x$ among threads in the binding thread set. To avoid data races, all accesses of the locations designated by $x$ that could potentially occur in parallel must be protected with an atomic construct.

Atomic regions do not guarantee exclusive access with respect to any accesses outside of atomic regions to the same storage location $x$ even if those accesses occur during a critical or ordered region, while an OpenMP lock is owned by the executing task, or during the execution of a reduction clause.

However, other OpenMP synchronization can ensure the desired exclusive access. For example, a barrier that follows a series of atomic updates to $x$ guarantees that subsequent accesses do not form a race with the atomic accesses.

A compliant implementation may enforce exclusive access between atomic regions that update different storage locations. The circumstances under which this occurs are implementation defined.

If the storage location designated by $x$ is not size-aligned (that is, if the byte alignment of $x$ is not a multiple of the size of $x$), then the behavior of the atomic region is implementation defined.

If present, the hint clause gives the implementation additional information about the expected properties of the atomic operation that can optionally be used to optimize the implementation. The presence of a hint clause does not affect the semantics of the atomic construct, and all hints may be ignored. If no hint clause is specified, the effect is as if hint(omp_sync_hint_none) had been specified.

### Execution Model Events

The atomic-acquiring event occurs in the thread that encounters the atomic construct on entry to the atomic region before initiating synchronization for the region.

The atomic-acquired event occurs in the thread that encounters the atomic construct after it enters the region, but before it executes the structured block of the atomic region.

The atomic-released event occurs in the thread that encounters the atomic construct after it completes any synchronization on exit from the atomic region.

### Tool Callbacks

A thread dispatches a registered ompt_callback_mutex_acquire callback for each occurrence of an atomic-acquiring event in that thread. This callback has the type signature ompt_callback_mutex_acquire_t.

A thread dispatches a registered ompt_callback_mutex_acquired callback for each occurrence of an atomic-acquired event in that thread. This callback has the type signature ompt_callback_mutex_t.
A thread dispatches a registered `omp_callback_mutex_released` callback with `omp_mutex_atomic` as the `kind` argument if practical, although a less specific `kind` may be used, for each occurrence of an `atomic-released` event in that thread. This callback has the type signature `omp_callback_mutex_t` and occurs in the task that encounters the atomic construct.

**Restrictions**

Restrictions to the `atomic` construct are as follows:

- OpenMP constructs may not be encountered during execution of an `atomic` region.
- At most one `atomic-clause` may appear on the construct.
- At most one `memory-order-clause` may appear on the construct.
- At most one `hint` clause may appear on the construct.
- At most one `capture` clause may appear on the construct.
- At most one `compare` clause may appear on the construct.
- If a `capture` or `compare` clause appears on the construct then `atomic-clause` must be `update`.
- At most one `fail` clause may appear on the construct.
- At most one `weak` clause may appear on the construct.
- If `atomic-clause` is `read` then `memory-order-clause` must not be `release`.
- If `atomic-clause` is `write` then `memory-order-clause` must not be `acquire`.
- The `weak` clause may only appear if the resulting atomic operation is an atomic conditional update for which the comparison tests for equality.

```
-- C / C++ --
• All atomic accesses to the storage locations designated by `x` throughout the program are required to have a compatible type.
• The `fail` clause may only appear if the resulting atomic operation is an atomic conditional update.

-- C / C++ --
```

```
-- Fortran --
• All atomic accesses to the storage locations designated by `x` throughout the program are required to have the same type and type parameters.
• The `fail` clause may only appear if the resulting atomic operation is an atomic conditional update or an atomic update where `intrinsic-procedure-name` is either `MAX` or `MIN`.

```
Cross References

- `requires` directive, see Section 2.5.1.
- `critical` construct, see Section 2.19.1.
- `barrier` construct, see Section 2.19.2.
- `flush` construct, see Section 2.19.8.
- `ordered` construct, see Section 2.19.9.
- Synchronization hints, see Section 2.19.12.
- `reduction` clause, see Section 2.21.5.4.
- lock routines, see Section 3.9.
- `ompt_mutex_atomic`, see Section 4.4.4.16.
- `ompt_callback_mutex_acquire_t`, see Section 4.5.2.14.
- `ompt_callback_mutex_t`, see Section 4.5.2.15.

2.19.8 flush Construct

Summary

The `flush` construct executes the OpenMP flush operation. This operation makes a thread’s temporary view of memory consistent with memory and enforces an order on the memory operations of the variables explicitly specified or implied. See the memory model description in Section 1.4 for more details. The `flush` construct is a stand-alone directive.

Syntax

```
C / C++

#pragma omp flush [memory-order-clause] [ (list) ] new-line
```

where `memory-order-clause` is one of the following:

```
seq_cst
acq_rel
release
acquire

C / C++
```
The syntax of the `flush` construct is as follows:

```fortran
$omp flush [memory-order-clause] [(list)]
```

where `memory-order-clause` is one of the following:

- `seq_cst`
- `acq_rel`
- `release`
- `acquire`

### Binding

The binding thread set for a `flush` region is all threads in the `device-set` of its flush operation. Execution of a `flush` region affects the memory and it affects the temporary view of memory of the encountering thread. It does not affect the temporary view of other threads. Other threads on devices in the `device-set` must themselves execute a flush operation in order to be guaranteed to observe the effects of the flush operation of the encountering thread.

### Description

If neither `memory-order-clause` nor a list appears on the `flush` construct then the behavior is as if `memory-order-clause` is `seq_cst`.

A `flush` construct with the `seq_cst` clause, executed on a given thread, operates as if all data storage blocks that are accessible to the thread are flushed by a strong flush operation. A `flush` construct with a list applies a strong flush operation to the items in the list, and the flush operation does not complete until the operation is complete for all specified list items. An implementation may implement a `flush` construct with a list by ignoring the list and treating it the same as a `flush` construct with the `seq_cst` clause.

If no list items are specified, the flush operation has the release and/or acquire flush properties:

- If `memory-order-clause` is `seq_cst` or `acq_rel`, the flush operation is both a release flush and an acquire flush.
- If `memory-order-clause` is `release`, the flush operation is a release flush.
- If `memory-order-clause` is `acquire`, the flush operation is an acquire flush.
If a pointer is present in the list, the pointer itself is flushed, not the memory block to which the pointer refers.

A flush construct without a list corresponds to a call to atomic_thread_fence, where the argument is given by the identifier that results from prefixing memory_order_ to memory-order-clause.

For a flush construct without a list, the generated flush region implicitly performs the corresponding call to atomic_thread_fence. The behavior of an explicit call to atomic_thread_fence that occurs in the program and does not have the argument memory_order_consume is as if the call is replaced by its corresponding flush construct.

If the list item or a subobject of the list item has the POINTER attribute, the allocation or association status of the POINTER item is flushed, but the pointer target is not. If the list item is a Cray pointer, the pointer is flushed, but the object to which it points is not. Cray pointer support has been deprecated. If the list item is of type C_PTR, the variable is flushed, but the storage that corresponds to that address is not flushed. If the list item or the subobject of the list item has the ALLOCATABLE attribute and has an allocation status of allocated, the allocated variable is flushed; otherwise the allocation status is flushed.

Note – Use of a flush construct with a list is extremely error prone and users are strongly discouraged from attempting it. The following examples illustrate the ordering properties of the flush operation. In the following incorrect pseudocode example, the programmer intends to prevent simultaneous execution of the protected section by the two threads, but the program does not work properly because it does not enforce the proper ordering of the operations on variables a and b. Any shared data accessed in the protected section is not guaranteed to be current or consistent during or after the protected section. The atomic notation in the pseudocode in the following two examples indicates that the accesses to a and b are atomic write and atomic read operations. Otherwise both examples would contain data races and automatically result in unspecified behavior. The flush operations are strong flushes that are applied to the specified flush lists.
Incorrect example:

\[ a = b = 0 \]

\[ \begin{array}{ll}
\text{thread 1} & \text{thread 2} \\
\text{atomic}(b = 1) & \text{atomic}(a = 1) \\
\text{flush}(b) & \text{flush}(a) \\
\text{flush}(a) & \text{flush}(b) \\
\text{atomic}(\text{tmp} = a) & \text{atomic}(\text{tmp} = b) \\
\text{if (tmp == 0)} & \text{if (tmp == 0)} \\
\text{protected section} & \text{protected section} \\
\text{end if} & \text{end if}
\end{array} \]

The problem with this example is that operations on variables \( a \) and \( b \) are not ordered with respect to each other. For instance, nothing prevents the compiler from moving the flush of \( b \) on thread 1 or the flush of \( a \) on thread 2 to a position completely after the protected section (assuming that the protected section on thread 1 does not reference \( b \) and the protected section on thread 2 does not reference \( a \)). If either re-ordering happens, both threads can simultaneously execute the protected section.

The following pseudocode example correctly ensures that the protected section is executed by only one thread at a time. Execution of the protected section by neither thread is considered correct in this example. This occurs if both flushes complete prior to either thread executing its \( \textbf{if} \) statement.

Correct example:

\[ a = b = 0 \]

\[ \begin{array}{ll}
\text{thread 1} & \text{thread 2} \\
\text{atomic}(b = 1) & \text{atomic}(a = 1) \\
\text{flush}(a,b) & \text{flush}(a,b) \\
\text{atomic}(\text{tmp} = a) & \text{atomic}(\text{tmp} = b) \\
\text{if (tmp == 0)} & \text{if (tmp == 0)} \\
\text{protected section} & \text{protected section} \\
\text{end if} & \text{end if}
\end{array} \]

The compiler is prohibited from moving the flush at all for either thread, ensuring that the respective assignment is complete and the data is flushed before the \( \textbf{if} \) statement is executed.
Execution Model Events

The flush event occurs in a thread that encounters the flush construct.

Tool Callbacks

A thread dispatches a registered ompt_callback_flush callback for each occurrence of a flush event in that thread. This callback has the type signature ompt_callback_flush_t.

Restrictions

Restrictions to the flush construct are as follows:

• If a memory-order-clause is specified, list items must not be specified on the flush directive.

Cross References

• ompt_callback_flush_t, see Section 4.5.2.17.

2.19.8.1 Implicit Flushes

Flush operations implied when executing an atomic region are described in Section 2.19.7.

A flush region that corresponds to a flush directive with the release clause present is implied at the following locations:

• During a barrier region;
• At entry to a parallel region;
• At entry to a teams region;
• At exit from a critical region;
• During an omp_unset_lock region;
• During an omp_unset_nest_lock region;
• Immediately before every task scheduling point;
• At exit from the task region of each implicit task;
• At exit from an ordered region, if a threads clause or a depend clause with a source dependence type is present, or if no clauses are present; and
• During a cancel region, if the cancel-var ICV is true.

For a target construct, the device-set of an implicit release flush that is performed in a target task during the generation of the target region and that is performed on exit from the initial task region that implicitly encloses the target region consists of the devices that execute the target task and the target region.

A flush region that corresponds to a flush directive with the acquire clause present is implied at the following locations:
During a barrier region;

- At exit from a teams region;

- At entry to a critical region;

- If the region causes the lock to be set, during:
  - an omp_set_lock region;
  - an omp_test_lock region;
  - an omp_set_nest_lock region; and
  - an omp_test_nest_lock region;

- Immediately after every task scheduling point;

- At entry to the task region of each implicit task;

- At entry to an ordered region, if a threads clause or a depend clause with a sink dependence type is present, or if no clauses are present; and

- Immediately before a cancellation point, if the cancel-var ICV is true and cancellation has been activated.

For a target construct, the device-set of an implicit acquire flush that is performed in a target task following the generation of the target region or that is performed on entry to the initial task region that implicitly encloses the target region consists of the devices that execute the target task and the target region.

Note – A flush region is not implied at the following locations:

- At entry to worksharing regions; and

- At entry to or exit from masked regions.

The synchronization behavior of implicit flushes is as follows:

- When a thread executes an atomic region for which the corresponding construct has the release, acq_rel, or seq_cst clause and specifies an atomic operation that starts a given release sequence, the release flush that is performed on entry to the atomic operation synchronizes with an acquire flush that is performed by a different thread and has an associated atomic operation that reads a value written by a modification in the release sequence.

- When a thread executes an atomic region for which the corresponding construct has the acquire, acq_rel, or seq_cst clause and specifies an atomic operation that reads a value written by a given modification, a release flush that is performed by a different thread and has an
associated release sequence that contains that modification synchronizes with the acquire flush that is performed on exit from the atomic operation.

- When a thread executes a `critical` region that has a given name, the behavior is as if the release flush performed on exit from the region synchronizes with the acquire flush performed on entry to the next `critical` region with the same name that is performed by a different thread, if it exists.

- When a thread team executes a `barrier` region, the behavior is as if the release flush performed by each thread within the region synchronizes with the acquire flush performed by all other threads within the region.

- When a thread executes a `taskwait` region that does not result in the creation of a dependent task and the task that encounters the corresponding `taskwait` construct has at least one child task, the behavior is as if each thread that executes a child task that is generated before the `taskwait` region performs a release flush upon completion of the child task that synchronizes with an acquire flush performed in the `taskwait` region.

- When a thread executes a `taskgroup` region, the behavior is as if each thread that executes a remaining descendant task performs a release flush upon completion of the descendant task that synchronizes with an acquire flush performed on exit from the `taskgroup` region.

- When a thread executes an `ordered` region that does not arise from a stand-alone `ordered` directive, the behavior is as if the release flush performed on exit from the region synchronizes with the acquire flush performed on entry to an `ordered` region encountered in the next logical iteration to be executed by a different thread, if it exists.

- When a thread executes an `ordered` region that arises from a stand-alone `ordered` directive, the behavior is as if the release flush performed in the `ordered` region from a given source iteration synchronizes with the acquire flush performed in all `ordered` regions executed by a different thread that are waiting for dependences on that iteration to be satisfied.

- When a thread team begins execution of a `parallel` region, the behavior is as if the release flush performed by the primary thread on entry to the `parallel` region synchronizes with the acquire flush performed on entry to each implicit task that is assigned to a different thread.

- When an initial thread begins execution of a `target` region that is generated by a different thread from a target task, the behavior is as if the release flush performed by the generating thread in the target task synchronizes with the acquire flush performed by the initial thread on entry to its initial task region.

- When an initial thread completes execution of a `target` region that is generated by a different thread from a target task, the behavior is as if the release flush performed by the initial thread on exit from its initial task region synchronizes with the acquire flush performed by the generating thread in the target task.

- When a thread encounters a `teams` construct, the behavior is as if the release flush performed by the thread on entry to the `teams` region synchronizes with the acquire flush performed on entry
to each initial task that is executed by a different initial thread that participates in the execution of
the **teams** region.

- When a thread that encounters a **teams** construct reaches the end of the **teams** region, the
  behavior is as if the release flush performed by each different participating initial thread at exit
  from its initial task synchronizes with the acquire flush performed by the thread at exit from the
  **teams** region.

- When a task generates an explicit task that begins execution on a different thread, the behavior is
  as if the thread that is executing the generating task performs a release flush that synchronizes
  with the acquire flush performed by the thread that begins to execute the explicit task.

- When an undeferred task completes execution on a given thread that is different from the thread
  on which its generating task is suspended, the behavior is as if a release flush performed by the
  thread that completes execution of the undeferred task synchronizes with an acquire flush
  performed by the thread that resumes execution of the generating task.

- When a dependent task with one or more predecessor tasks begins execution on a given thread,
  the behavior is as if each release flush performed by a different thread on completion of a
  predecessor task synchronizes with the acquire flush performed by the thread that begins to
  execute the dependent task.

- When a task begins execution on a given thread and it is mutually exclusive with respect to
  another sibling task that is executed by a different thread, the behavior is as if each release flush
  performed on completion of the sibling task synchronizes with the acquire flush performed by
  the thread that begins to execute the task.

- When a thread executes a **cancel** region, the **cancel-var** ICV is **true**, and cancellation is not
  already activated for the specified region, the behavior is as if the release flush performed during
  the **cancel** region synchronizes with the acquire flush performed by a different thread
  immediately before a cancellation point in which that thread observes cancellation was activated
  for the region.

- When a thread executes an **omp_unset_lock** region that causes the specified lock to be unset,
  the behavior is as if a release flush is performed during the **omp_unset_lock** region that
  synchronizes with an acquire flush that is performed during the next **omp_set_lock** or
  **omp_test_lock** region to be executed by a different thread that causes the specified lock to be
  set.

- When a thread executes an **omp_unset_nest_lock** region that causes the specified nested
  lock to be unset, the behavior is as if a release flush is performed during the
  **omp_unset_nest_lock** region that synchronizes with an acquire flush that is performed
  during the next **omp_set_nest_lock** or **omp_test_nest_lock** region to be executed by
  a different thread that causes the specified nested lock to be set.
2.19.9 ordered Construct

Summary

The ordered construct either specifies a structured block in a worksharing-loop, simd, or worksharing-loop SIMD region that will be executed in the order of the loop iterations, or it is a stand-alone directive that specifies cross-iteration dependences in a doacross loop nest. The ordered construct sequentializes and orders the execution of ordered regions while allowing code outside the region to run in parallel.

Syntax

The syntax of the ordered construct is as follows:

**C / C++**

```c
#pragma omp ordered [clause [,] clause] new-line
structured-block
```

where clause is one of the following:

- `threads`
- `simd`

or

```c
#pragma omp ordered clause [([,] clause] ...) new-line
```

where clause is one of the following:

- `depend(source)`
- `depend(sink : vec)`

**C / C++**

**Fortran**

The syntax of the ordered construct is as follows:

```fortran
!$omp ordered clause [,] clause] 
loosely-structured-block
!$omp end ordered
```

or

```fortran
!$omp ordered clause [,] clause] 
strictly-structured-block 
!$omp end ordered/
```

where clause is one of the following:

- `threads`
- `simd`
or

```fortran
!$omp ordered clause \([\ldots]\) ...
```

where \textit{clause} is one of the following:

1. `depend(source)`
2. `depend(sink : vec)`

If the \textit{depend} clause is specified, the \textbf{ordered} construct is a stand-alone directive.

\section*{Binding}

The binding thread set for an \textbf{ordered} region is the current team. An \textbf{ordered} region binds to the innermost enclosing \textbf{simd} or worksharing-loop SIMD region if the \textbf{simd} clause is present, and otherwise it binds to the innermost enclosing worksharing-loop region. \textbf{ordered} regions that bind to different regions execute independently of each other.

\section*{Description}

If no clause is specified, the \textbf{ordered} construct behaves as if the \textbf{threads} clause had been specified. If the \textbf{threads} clause is specified, the threads in the team that is executing the worksharing-loop region execute ordered regions sequentially in the order of the loop iterations. If any \textit{depend} clauses are specified then those clauses specify the order in which the threads in the team execute \textbf{ordered} regions. If the \textbf{simd} clause is specified, the \textbf{ordered} regions encountered by any thread will execute one at a time in the order of the loop iterations.

When the thread that is executing the first iteration of the loop encounters an \textbf{ordered} construct, it can enter the \textbf{ordered} region without waiting. When a thread that is executing any subsequent iteration encounters an \textbf{ordered} construct without a \textit{depend} clause, it waits at the beginning of the \textbf{ordered} region until execution of all \textbf{ordered} regions that belong to all previous iterations has completed. When a thread that is executing any subsequent iteration encounters an \textbf{ordered} construct with one or more `depend(sink: vec)` clauses, it waits until its dependences on all valid iterations specified by the \textit{depend} clauses are satisfied before it completes execution of the \textbf{ordered} region. A specific dependence is satisfied when a thread that is executing the corresponding iteration encounters an \textbf{ordered} construct with a `depend(source)` clause.

\section*{Execution Model Events}

The \textbf{ordered-acquiring} event occurs in the task that encounters the \textbf{ordered} construct on entry to the ordered region before it initiates synchronization for the region.

The \textbf{ordered-acquired} event occurs in the task that encounters the \textbf{ordered} construct after it enters the region, but before it executes the structured block of the \textbf{ordered} region.

The \textbf{ordered-released} event occurs in the task that encounters the \textbf{ordered} construct after it completes any synchronization on exit from the \textbf{ordered} region.
The doacross-sink event occurs in the task that encounters an ordered construct for each depend(sink:vec) clause after the dependence is fulfilled.

The doacross-source event occurs in the task that encounters an ordered construct with a depend(source:vec) clause before signaling the dependence to be fulfilled.

Tool Callbacks
A thread dispatches a registered ompt_callback_mutex_acquire callback for each occurrence of an ordered-acquiring event in that thread. This callback has the type signature ompt_callback_mutex_acquire_t.

A thread dispatches a registered ompt_callback_mutex_acquired callback for each occurrence of an ordered-acquired event in that thread. This callback has the type signature ompt_callback_mutex_t.

A thread dispatches a registered ompt_callback_mutex_released callback with ompt_mutex_ordered as the kind argument if practical, although a less specific kind may be used, for each occurrence of an ordered-released event in that thread. This callback has the type signature ompt_callback_mutex_t and occurs in the task that encounters the atomic construct.

A thread dispatches a registered ompt_callback_dependences callback with all vector entries listed as ompt_dependence_type_sink in the deps argument for each occurrence of a doacross-sink event in that thread. A thread dispatches a registered ompt_callback_dependences callback with all vector entries listed as ompt_dependence_type_source in the deps argument for each occurrence of a doacross-source event in that thread. These callbacks have the type signature ompt_callback_dependences_t.

Restrictions
Restrictions to the ordered construct are as follows:

- At most one threads clause can appear on an ordered construct.
- At most one simd clause can appear on an ordered construct.
- At most one depend(source) clause can appear on an ordered construct.
- The construct that corresponds to the binding region of an ordered region must not specify a reduction clause with the inscan modifier.
- Either depend(sink:vec) clauses or depend(source) clauses may appear on an ordered construct, but not both.
- The worksharing-loop or worksharing-loop SIMD region to which an ordered region corresponding to an ordered construct without a depend clause binds must have an ordered clause without the parameter specified on the corresponding worksharing-loop or worksharing-loop SIMD directive.
• The worksharing-loop region to which an ordered region that corresponds to an ordered construct with any depend clauses binds must have an ordered clause with the parameter specified on the corresponding worksharing-loop directive.

• An ordered construct with the depend clause specified must be closely nested inside a worksharing-loop (or parallel worksharing-loop) construct.

• An ordered region that corresponds to an ordered construct without the simd clause specified must be closely nested inside a worksharing-loop region.

• An ordered region that corresponds to an ordered construct with the simd clause specified must be closely nested inside a simd or worksharing-loop SIMD region.

• An ordered region that corresponds to an ordered construct with both the simd and threads clauses must be closely nested inside a worksharing-loop SIMD region or must be closely nested inside a worksharing-loop and simd region.

• During execution of an iteration of a worksharing-loop or a loop nest within a worksharing-loop, simd, or worksharing-loop SIMD region, a thread must not execute more than one ordered region that corresponds to an ordered construct without a depend clause.

• A throw executed inside a ordered region must cause execution to resume within the same ordered region, and the same thread that threw the exception must catch it.

Cross References
• worksharing-loop construct, see Section 2.11.4.
• simd construct, see Section 2.11.5.1.
• parallel Worksharing-loop construct, see Section 2.16.1.
• depend Clause, see Section 2.19.11
• ompt_mutex_ordered, see Section 4.4.4.16.
• ompt_callback_mutex_acquire_t, see Section 4.5.2.14.
• ompt_callback_mutex_t, see Section 4.5.2.15.

2.19.10 Depend Objects

This section describes constructs that support OpenMP depend objects that can be used to supply user-computed dependences to depend clauses. OpenMP depend objects must be accessed only through the depobj construct or through the depend clause; programs that otherwise access OpenMP depend objects are non-conforming.

An OpenMP depend object can be in one of the following states: uninitialized or initialized.

Initially OpenMP depend objects are in the uninitialized state.
2.19.10.1 depobj Construct

Summary
The depobj construct initializes, updates or destroys an OpenMP depend object. The depobj construct is a stand-alone directive.

Syntax

C / C++

The syntax of the depobj construct is as follows:

```c
#pragma omp depobj(depobj) clause new-line
```

where depobj is an lvalue expression of type omp_depend_t.

where clause is one of the following:

```c
depend(dependence-type : locator)
destroy
update(dependence-type)
```

Fortran

The syntax of the depobj construct is as follows:

```fortran
!$omp depobj(depobj) clause
```

where depobj is a scalar integer variable of the omp_depend_kind kind.

where clause is one of the following:

```fortran
depend(dependence-type : locator)
destroy
update(dependence-type)
```

Binding
The binding thread set for a depobj region is the encountering thread.

Description
A depobj construct with a depend clause present sets the state of depobj to initialized. The depobj is initialized to represent the dependence that the depend clause specifies.

A depobj construct with a destroy clause present changes the state of the depobj to uninitialized.

A depobj construct with an update clause present changes the dependence type of the dependence represented by depobj to the one specified by the update clause.
Restrictions
Restrictions to the depobj construct are as follows:

- A depend clause on a depobj construct must not have source or sink as dependence-type.
- An update clause on a depobj construct must not have source, sink or depobj as dependence-type.
- A depend clause on a depobj construct can only specify one locator.
- The depobj of a depobj construct with the depend clause present must be in the uninitialized state.
- The depobj of a depobj construct with the destroy clause present must be in the initialized state.
- The depobj of a depobj construct with the update clause present must be in the initialized state.

Cross References
- depend clause, see Section 2.19.11.

2.19.11 depend Clause

Summary
The depend clause enforces additional constraints on the scheduling of tasks or loop iterations. These constraints establish dependences only between sibling tasks or between loop iterations.

Syntax
The syntax of the depend clause is as follows:
```plaintext
depend ([depend-modifier, ]dependence-type : locator-list)
```

where dependence-type is one of the following:
```plaintext
in
out
inout
mutexinoutset
inoutset
depobj
```

where depend-modifier is one of the following:
```plaintext
iterator (iterators-definition)
```
or
\[
\text{depend}(\text{dependence-type})
\]

where \(\text{dependence-type}\) is:
\[
\text{source}
\]

or
\[
\text{depend}(\text{dependence-type} : \text{vec})
\]

where \(\text{dependence-type}\) is:
\[
\text{sink}
\]

and where \(\text{vec}\) is the iteration vector, which has the form:

\[
x_1 [\pm d_1], x_2 [\pm d_2], \ldots, x_n [\pm d_n]
\]

where \(n\) is the value specified by the \texttt{ordered} clause in the worksharing-loop directive, \(x_i\) denotes the loop iteration variable of the \(i\)-th nested loop associated with the worksharing-loop directive, and \(d_i\) is a constant non-negative integer.

\section*{Description}

Task dependences are derived from the \texttt{dependence-type} of a \texttt{depend} clause and its list items when \texttt{dependence-type} is \texttt{in}, \texttt{out}, \texttt{inout}, \texttt{mutexinoutset} or \texttt{inoutset}. When the \texttt{dependence-type} is \texttt{depobj}, the task dependences are derived from the dependences represented by the depend objects specified in the \texttt{depend} clause as if the \texttt{depend} clauses of the \texttt{depobj} constructs were specified in the current construct.

The storage location of a list item matches the storage location of another list item if they have the same storage location, or if any of the list items is \texttt{omp_all_memory}.

For the \texttt{in} dependence-type, if the storage location of at least one of the list items matches the storage location of a list item appearing in a \texttt{depend} clause with an \texttt{out}, \texttt{inout}, \texttt{mutexinoutset}, or \texttt{inoutset} \texttt{dependence-type} on a construct from which a sibling task was previously generated, then the generated task will be a dependent task of that sibling task.

For the \texttt{out} and \texttt{inout} dependence-types, if the storage location of at least one of the list items matches the storage location of a list item appearing in a \texttt{depend} clause with an \texttt{in}, \texttt{out}, \texttt{inout}, \texttt{mutexinoutset}, or \texttt{inoutset} \texttt{dependence-type} on a construct from which a sibling task was previously generated, then the generated task will be a dependent task of that sibling task.

For the \texttt{mutexinoutset} dependence-type, if the storage location of at least one of the list items matches the storage location of a list item appearing in a \texttt{depend} clause with an \texttt{in}, \texttt{out}, \texttt{inout}, or \texttt{inoutset} \texttt{dependence-type} on a construct from which a sibling task was previously generated, then the generated task will be a dependent task of that sibling task.

If a list item appearing in a \texttt{depend} clause with a \texttt{mutexinoutset} \texttt{dependence-type} on a task generating construct matches a list item appearing in a \texttt{depend} clause with a \texttt{mutexinoutset}
dependence-type on a different task generating construct, and both constructs generate sibling tasks, the sibling tasks will be mutually exclusive tasks.

For the inoutset dependence-type, if the storage location of at least one of the list items matches the storage location of a list item appearing in a depend clause with an in, out, inout, or mutexinoutset dependence-type on a construct from which a sibling task was previously generated, then the generated task will be a dependent task of that sibling task.

The list items that appear in the depend clause may reference iterators defined by an iterators-definition appearing on an iterator modifier.

The list items that appear in the depend clause may include array sections or the omp_all_memory reserved locator.

Fortran

If a list item has the ALLOCATABLE attribute and its allocation status is unallocated, the behavior is unspecified. If a list item has the POINTER attribute and its association status is disassociated or undefined, the behavior is unspecified.

C / C++

The list items that appear in a depend clause may use shape-operators.

C / C++

Note – The enforced task dependence establishes a synchronization of memory accesses performed by a dependent task with respect to accesses performed by the predecessor tasks. However, it is the responsibility of the programmer to synchronize properly with respect to other concurrent accesses that occur outside of those tasks.

The source dependence-type specifies the satisfaction of cross-iteration dependences that arise from the current iteration.

The sink dependence-type specifies a cross-iteration dependence, where the iteration vector vec indicates the iteration that satisfies the dependence.

If the iteration vector vec does not occur in the iteration space, the depend clause is ignored. If all depend clauses on an ordered construct are ignored then the construct is ignored.

Note – An iteration vector vec that does not indicate a lexicographically earlier iteration may cause a deadlock.
Execution Model Events

The *task-dependences* event occurs in a thread that encounters a task generating construct or a *taskwait* construct with a *depend* clause immediately after the *task-create* event for the new task or the *taskwait-init* event.

The *task-dependence* event indicates an unfulfilled dependence for the generated task. This event occurs in a thread that observes the unfulfilled dependence before it is satisfied.

Tool Callbacks

A thread dispatches the `ompt_callback_dependences` callback for each occurrence of the *task-dependences* event to announce its dependences with respect to the list items in the *depend* clause. This callback has type signature `ompt_callback_dependences_t`.

A thread dispatches the `ompt_callback_task_dependence` callback for a *task-dependence* event to report a dependence between a predecessor task (*src_task_data*) and a dependent task (*sink_task_data*). This callback has type signature `ompt_callback_task_dependence_t`.

Restrictions

Restrictions to the *depend* clause are as follows:

- List items, other than reserved locators, used in *depend* clauses of the same task or sibling tasks must indicate identical storage locations or disjoint storage locations.
- List items used in *depend* clauses cannot be zero-length array sections.
- The `omp_all_memory` reserved locator can only be used in a *depend* clause with an *out* or *inout* dependence-type.
- Array sections cannot be specified in *depend* clauses with the *depobj* dependence type.
- List items used in *depend* clauses with the *depobj* dependence type must be depend objects in the initialized state.

\[
\begin{align*}
\text{C / C++} & \quad \text{Fortran} \\
\end{align*}
\]

- List items used in *depend* clauses with the *depobj* dependence type must be expressions of the `omp_depend_t` type.
- List items that are expressions of the `omp_depend_t` type can only be used in *depend* clauses with the *depobj* dependence type.

\[
\begin{align*}
\text{C / C++} & \quad \text{Fortran} \\
\end{align*}
\]

- A common block name cannot appear in a *depend* clause.
- List items used in *depend* clauses with the *depobj* dependence type must be integer expressions of the `omp_depend_kind` kind.
• For a vec element of sink dependence-type of the form $x_i + d_i$ or $x_i - d_i$ if the loop iteration variable $x_i$ has an integral or pointer type, the expression $x_i + d_i$ or $x_i - d_i$ for any value of the loop iteration variable $x_i$ that can encounter the ordered construct must be computable without overflow in the type of the loop iteration variable.

\begin{Verbatim}[commandchars=\[\]]
C++
\end{Verbatim}

• For a vec element of sink dependence-type of the form $x_i + d_i$ or $x_i - d_i$ if the loop iteration variable $x_i$ is of a random access iterator type other than pointer type, the expression $(x_i - lb_i) + d_i$ or $(x_i - lb_i) - d_i$ for any value of the loop iteration variable $x_i$ that can encounter the ordered construct must be computable without overflow in the type that would be used by std: : distance applied to variables of the type of $x_i$.

\begin{Verbatim}[commandchars=\[\]]
C++
\end{Verbatim}

\begin{Verbatim}[commandchars=\[\]]
C / C++
\end{Verbatim}

• A bit-field cannot appear in a depend clause.

\begin{Verbatim}[commandchars=\[\]]
C / C++
\end{Verbatim}

Cross References

• Array shaping, see Section 2.1.4.

• Array sections, see Section 2.1.5.

• Iterators, see Section 2.1.6.

• task construct, see Section 2.12.1.

• Task scheduling constraints, see Section 2.12.6.

• target enter data construct, see Section 2.14.3.

• target exit data construct, see Section 2.14.4.

• target construct, see Section 2.14.5.

• target update construct, see Section 2.14.6.

• ordered construct, see Section 2.19.9.

• depobj construct, see Section 2.19.10.1.

• ompt_callback_dependencies_t, see Section 4.5.2.8.

• ompt_callback_task_dependence_t, see Section 4.5.2.9.
2.19.12 Synchronization Hints

Hints about the expected dynamic behavior or suggested implementation can be provided by the programmer to locks (by using the `omp_init_lock_with_hint` or `omp_init_nest_lock_with_hint` functions to initialize the lock), and to `atomic` and `critical` directives by using the `hint` clause. The effect of a hint does not change the semantics of the associated construct; if ignoring the hint changes the program semantics, the result is unspecified.

The C/C++ header file (`omp.h`) and the Fortran include file (`omp_lib.h`) and/or Fortran module file (`omp_lib`) define the valid hint constants. The valid constants must include the following, which can be extended with implementation-defined values:

```c
typedef enum omp_sync_hint_t {
  omp_sync_hint_none = 0x0,
  omp_lock_hint_none = omp_sync_hint_none,
  omp_sync_hint_uncontended = 0x1,
  omp_lock_hint_uncontended = omp_sync_hint_uncontended,
  omp_sync_hint_contended = 0x2,
  omp_lock_hint_contended = omp_sync_hint_contended,
  omp_sync_hint_nonspeculative = 0x4,
  omp_lock_hint_nonspeculative = omp_sync_hint_nonspeculative,
  omp_sync_hint_speculative = 0x8,
  omp_lock_hint_speculative = omp_sync_hint_speculative
} omp_sync_hint_t;
```

```c
typedef omp_sync_hint_t omp_lock_hint_t;
```

```fortran
integer, parameter :: omp_lock_hint_kind = omp_sync_hint_kind

integer (kind=omp_sync_hint_kind), &
  parameter :: omp_sync_hint_none = &
  int(Z’0’, kind=omp_sync_hint_kind)
integer (kind=omp_lock_hint_kind), &
  parameter :: omp_lock_hint_none = omp_sync_hint_none
integer (kind=omp_sync_hint_kind), &
  parameter :: omp_sync_hint_uncontended = &
  int(Z’1’, kind=omp_sync_hint_kind)
integer (kind=omp_lock_hint_kind), &
  parameter :: omp_lock_hint_uncontended = &
  omp_sync_hint_uncontended
integer (kind=omp_sync_hint_kind), &
  parameter :: omp_sync_hint_contended
integer (kind=omp_lock_hint_kind), &
  parameter :: omp_lock_hint_contended
integer (kind=omp_sync_hint_kind), &
  parameter :: omp_sync_hint_nonspeculative
integer (kind=omp_lock_hint_kind), &
  parameter :: omp_lock_hint_nonspeculative
integer (kind=omp_sync_hint_kind), &
  parameter :: omp_sync_hint_speculative
integer (kind=omp_lock_hint_kind), &
  parameter :: omp_lock_hint_speculative
```

CHAPTER 2. DIRECTIVES
The hints can be combined by using the + or | operators in C/C++ or the + operator in Fortran. Combining \texttt{omp\_sync\_hint\_none} with any other hint is equivalent to specifying the other hint.

The intended meaning of each hint is:

- \texttt{omp\_sync\_hint\_uncontended}: low contention is expected in this operation, that is, few threads are expected to perform the operation simultaneously in a manner that requires synchronization;

- \texttt{omp\_sync\_hint\_contended}: high contention is expected in this operation, that is, many threads are expected to perform the operation simultaneously in a manner that requires synchronization;

- \texttt{omp\_sync\_hint\_speculative}: the programmer suggests that the operation should be implemented using speculative techniques such as transactional memory; and

- \texttt{omp\_sync\_hint\_nonspeculative}: the programmer suggests that the operation should not be implemented using speculative techniques such as transactional memory.

Note – Future OpenMP specifications may add additional hints to the \texttt{omp\_sync\_hint\_t} type and the \texttt{omp\_sync\_hint\_kind} kind. Implementers are advised to add implementation-defined hints starting from the most significant bit of the \texttt{omp\_sync\_hint\_t} type and \texttt{omp\_sync\_hint\_kind} kind and to include the name of the implementation in the name of the added hint to avoid name conflicts with other OpenMP implementations.
The `omp_sync_hint_t` and `omp_lock_hint_t` enumeration types and the equivalent types in Fortran are synonyms for each other. The type `omp_lock_hint_t` has been deprecated.

**Restrictions**

Restrictions to the synchronization hints are as follows:

- The hints `omp_sync_hint_uncontended` and `omp_sync_hint_contended` cannot be combined.
- The hints `omp_sync_hint_nonspeculative` and `omp_sync_hint_speculative` cannot be combined.

The restrictions for combining multiple values of `omp_sync_hint` apply equally to the corresponding values of `omp_lock_hint`, and expressions that mix the two types.

**Cross References**

- `critical` construct, see Section 2.19.1.
- `atomic` construct, see Section 2.19.7
- `omp_init_lock_with_hint` and `omp_init_nest_lock_with_hint`, see Section 3.9.2.

## 2.20 Cancellation Constructs

### 2.20.1 cancel Construct

**Summary**

The `cancel` construct activates cancellation of the innermost enclosing region of the type specified. The `cancel` construct is a stand-alone directive.

**Syntax**

```c++
#pragma omp cancel construct-type-clause [ [, ] if-clause ] new-line
```

The syntax of the `cancel` construct is as follows:

- `parallel`
- `sections`
- `for`
- `taskgroup`
and if-clause is

```c
if([ cancel : ] scalar-expression)
```

The syntax of the `cancel` construct is as follows:

```fortran
!$omp cancel construct-type-clause [, if-clause]
```

where `construct-type-clause` is one of the following:

- `parallel`
- `sections`
- `do`
- `taskgroup`

and if-clause is

```fortran
if([ cancel : ] scalar-logical-expression)
```

**Binding**

The binding thread set of the `cancel` region is the current team. The binding region of the `cancel` region is the innermost enclosing region of the type corresponding to the `construct-type-clause` specified in the directive (that is, the innermost `parallel`, `sections`, worksharing-loop, or `taskgroup` region).

**Description**

The `cancel` construct activates cancellation of the binding region only if the `cancel-var` ICV is `true`, in which case the `cancel` construct causes the encountering task to continue execution at the end of the binding region if `construct-type-clause` is `parallel`, `for`, `do`, or `sections`. If the `cancel-var` ICV is `true` and `construct-type-clause` is `taskgroup`, the encountering task continues execution at the end of the current task region. If the `cancel-var` ICV is `false`, the `cancel` construct is ignored.

Threads check for active cancellation only at cancellation points that are implied at the following locations:

- `cancel` regions;
- `cancellation point` regions;
- `barrier` regions;
- at the end of a worksharing-loop construct with a `nowait` clause and for which the same list item appears in both `firstprivate` and `lastprivate` clauses; and
- implicit barrier regions.
When a thread reaches one of the above cancellation points and if the cancel-var ICV is true, then:

- If the thread is at a **cancel** or **cancellation point** region and **construct-type-clause** is **parallel**, **for**, **do**, or **sections**, the thread continues execution at the end of the canceled region if cancellation has been activated for the innermost enclosing region of the type specified.

- If the thread is at a **cancel** or **cancellation point** region and **construct-type-clause** is **taskgroup**, the encountering task checks for active cancellation of all of the **taskgroup sets** to which the encountering task belongs, and continues execution at the end of the current task region if cancellation has been activated for any of the **taskgroup sets**.

- If the encountering task is at a barrier region or at the end of a worksharing-loop construct with a **nowait** clause and for which the same list item appears in both **firstprivate** and **lastprivate** clauses, the encountering task checks for active cancellation of the innermost enclosing **parallel** region. If cancellation has been activated, then the encountering task continues execution at the end of the canceled region.

Note – If one thread activates cancellation and another thread encounters a cancellation point, the order of execution between the two threads is non-deterministic. Whether the thread that encounters a cancellation point detects the activated cancellation depends on the underlying hardware and operating system.

When cancellation of tasks is activated through a **cancel** construct with the **taskgroup** **construct-type-clause**, the tasks that belong to the **taskgroup set** of the innermost enclosing **taskgroup** region will be canceled. The task that encountered that construct continues execution at the end of its task region, which implies completion of that task. Any task that belongs to the innermost enclosing **taskgroup** and has already begun execution must run to completion or until a cancellation point is reached. Upon reaching a cancellation point and if cancellation is active, the task continues execution at the end of its task region, which implies the task’s completion. Any task that belongs to the innermost enclosing **taskgroup** and that has not begun execution may be discarded, which implies its completion.

When cancellation is active for a **parallel**, **sections**, or worksharing-loop region, each thread of the binding thread set resumes execution at the end of the canceled region if a cancellation point is encountered. If the canceled region is a **parallel** region, any tasks that have been created by a **task** or a **taskloop** construct and their descendant tasks are canceled according to the above **taskgroup** cancellation semantics. If the canceled region is a **sections**, or worksharing-loop region, no task cancellation occurs.

\[
\text{\textbackslash C++\textbackslash}}
\]

The usual C++ rules for object destruction are followed when cancellation is performed.
All private objects or subobjects with `ALLOCATABLE` attribute that are allocated inside the canceled construct are deallocated.

If the canceled construct contains a `reduction`, `task_reduction` or `lastprivate` clause, the final values of the list items that appeared in those clauses are undefined.

When an `if` clause is present on a `cancel` construct and the `if` expression evaluates to `false`, the `cancel` construct does not activate cancellation. The cancellation point associated with the `cancel` construct is always encountered regardless of the value of the `if` expression.

Note – The programmer is responsible for releasing locks and other synchronization data structures that might cause a deadlock when a `cancel` construct is encountered and blocked threads cannot be canceled. The programmer is also responsible for ensuring proper synchronizations to avoid deadlocks that might arise from cancellation of OpenMP regions that contain OpenMP synchronization constructs.

**Execution Model Events**

If a task encounters a `cancel` construct that will activate cancellation then a `cancel` event occurs. A `discarded-task` event occurs for any discarded tasks.

**Tool Callbacks**

A thread dispatches a registered `ompt_callback_cancel` callback for each occurrence of a `cancel` event in the context of the encountering task. This callback has type signature `ompt_callback_cancel_t; (flags & ompt_cancel_activated) always evaluates to true in the dispatched callback; (flags & ompt_cancel_parallel) evaluates to true in the dispatched callback if `construct-type-clause` is `parallel`; (flags & ompt_cancel_sections) evaluates to true in the dispatched callback if `construct-type-clause` is `sections`; (flags & ompt_cancel_loop) evaluates to true in the dispatched callback if `construct-type-clause` is `for` or `do`; and (flags & ompt_cancel_taskgroup) evaluates to true in the dispatched callback if `construct-type-clause` is `taskgroup`.

A thread dispatches a registered `ompt_callback_cancel` callback with the `ompt_data_t` associated with the discarded task as its `task_data` argument and `ompt_cancel_discarded_task` as its `flags` argument for each occurrence of a `discarded-task` event. The callback occurs in the context of the task that discards the task and has type signature `ompt_callback_cancel_t`. 
Restrictions
Restrictions to the `cancel` construct are as follows:

- The behavior for concurrent cancellation of a region and a region nested within it is unspecified.
- If `construct-type-clause` is `taskgroup`, the `cancel` construct must be closely nested inside a `task` or a `taskloop` construct and the `cancel` region must be closely nested inside a `taskgroup` region.
- If `construct-type-clause` is `sections`, the `cancel` construct must be closely nested inside a `sections` or `section` construct.
- If `construct-type-clause` is neither `sections` nor `taskgroup`, the `cancel` construct must be closely nested inside an OpenMP construct that matches the type specified in `construct-type-clause` of the `cancel` construct.

- A worksharing construct that is canceled must not have a `nowait` clause or a `reduction` clause with a user-defined reduction that uses `omp_orig` in the `initializer-expr` of the corresponding `declare reduction` directive.
- A worksharing-loop construct that is canceled must not have an `ordered` clause or a `reduction` clause with the `inscan` modifier.
- When cancellation is active for a `parallel` region, a thread in the team that binds to that region may not be executing or encounter a worksharing construct with an `ordered` clause, a `reduction` clause with the `inscan` modifier or a `reduction` clause with a user-defined reduction that uses `omp_orig` in the `initializer-expr` of the corresponding `declare reduction` directive.
- When cancellation is active for a `parallel` region, a thread in the team that binds to that region may not be executing or encounter a `scope` construct with a `reduction` clause with a user-defined reduction that uses `omp_orig` in the `initializer-expr` of the corresponding `declare reduction` directive.
- During execution of a construct that may be subject to cancellation, a thread must not encounter an orphaned cancellation point. That is, a cancellation point must only be encountered within that construct and must not be encountered elsewhere in its region.

Cross References
- `cancel-var` ICV, see Section 2.4.1.
- `if` clause, see Section 2.18.
- `cancellation point` construct, see Section 2.20.2.
- `omp_get_cancellation` routine, see Section 3.2.8.
- `omp_cancel_flag_t` enumeration type, see Section 4.4.4.25.
- `ompt_callback_cancel_t`, see Section 4.5.2.18.
2.20.2 cancellation point Construct

Summary
The cancellation point construct introduces a user-defined cancellation point at which implicit or explicit tasks check if cancellation of the innermost enclosing region of the type specified has been activated. The cancellation point construct is a stand-alone directive.

Syntax
The syntax of the cancellation point construct is as follows:

```
#pragma omp cancellation point construct-type-clause new-line
```

where construct-type-clause is one of the following:

- parallel
- sections
- for
- taskgroup

The syntax of the cancellation point construct is as follows:

```
!$omp cancellation point construct-type-clause
```

where construct-type-clause is one of the following:

- parallel
- sections
- do
- taskgroup

Binding
The binding thread set of the cancellation point construct is the current team. The binding region of the cancellation point region is the innermost enclosing region of the type corresponding to the construct-type-clause specified in the directive (that is, the innermost parallel, sections, worksharing-loop, or taskgroup region).
Description

This directive introduces a user-defined cancellation point at which an implicit or explicit task must check if cancellation of the innermost enclosing region of the type specified in the clause has been requested. This construct does not implement any synchronization between threads or tasks.

When an implicit or explicit task reaches a user-defined cancellation point and if the cancel-var ICV is true, then:

- If the construct-type-clause of the encountered cancellation point construct is parallel, for, do, or sections, the thread continues execution at the end of the canceled region if cancellation has been activated for the innermost enclosing region of the type specified.

- If the construct-type-clause of the encountered cancellation point construct is taskgroup, the encountering task checks for active cancellation of all taskgroup sets to which the encountering task belongs and continues execution at the end of the current task region if cancellation has been activated for any of them.

Execution Model Events

The cancellation event occurs if a task encounters a cancellation point and detected the activation of cancellation.

Tool Callbacks

A thread dispatches a registered ompt_callback_cancel callback for each occurrence of a cancel event in the context of the encountering task. This callback has type signature ompt_callback_cancel_t; (flags & ompt_cancel_detected) always evaluates to true in the dispatched callback; (flags & ompt_cancel_parallel) evaluates to true in the dispatched callback if construct-type-clause of the encountered cancellation point construct is parallel; (flags & ompt_cancel_sections) evaluates to true in the dispatched callback if construct-type-clause of the encountered cancellation point construct is sections; (flags & ompt_cancel_loop) evaluates to true in the dispatched callback if construct-type-clause of the encountered cancellation point construct is for or do; and (flags & ompt_cancel_taskgroup) evaluates to true in the dispatched callback if construct-type-clause of the encountered cancellation point construct is taskgroup.

Restrictions

Restrictions to the cancellation point construct are as follows:

- A cancellation point construct for which construct-type-clause is taskgroup must be closely nested inside a task or taskloop construct, and the cancellation point region must be closely nested inside a taskgroup region.

- A cancellation point construct for which construct-type-clause is sections must be closely nested inside a sections or section construct.

- A cancellation point construct for which construct-type-clause is neither sections nor taskgroup must be closely nested inside an OpenMP construct that matches the type specified in construct-type-clause.
Cross References

- `cancel-var` ICV, see Section 2.4.1.
- `cancel` construct, see Section 2.20.1.
- `omp_get_cancellation` routine, see Section 3.2.8.
- `ompt_callback_cancel_t`, see Section 4.5.2.18.

2.21 Data Environment

This section presents directives and clauses for controlling data environments.

2.21.1 Data-Sharing Attribute Rules

This section describes how the data-sharing attributes of variables referenced in data environments are determined. The following two cases are described separately:

- Section 2.21.1.1 describes the data-sharing attribute rules for variables referenced in a construct.
- Section 2.21.1.2 describes the data-sharing attribute rules for variables referenced in a region, but outside any construct.

2.21.1.1 Variables Referenced in a Construct

The data-sharing attributes of variables that are referenced in a construct can be *predetermined*, *explicitly determined*, or *implicitly determined*, according to the rules outlined in this section.

Specifying a variable in a data-sharing attribute clause, except for the `private` clause, or `copyprivate` clause of an enclosed construct, causes an implicit reference to the variable in the enclosing construct. Specifying a variable in a `map` clause of an enclosed construct may cause an implicit reference to the variable in the enclosing construct. Such implicit references are also subject to the data-sharing attribute rules outlined in this section.

Certain variables and objects have *predetermined* data-sharing attributes with respect to the construct in which they are referenced. The first matching rule from the following list of predetermined data-sharing attribute rules applies for variables and objects that are referenced in a construct.

```
C / C++
```

- Variables that appear in `threadprivate` directives or variables with the `_Thread_local` (in C) or `thread_local` (in C++) storage-class specifier are threadprivate.
• Variables with automatic storage duration that are declared in a scope inside the construct are private.

• Variables of non-reference type with automatic storage duration that are declared in a scope inside the construct are private.

• Objects with dynamic storage duration are shared.

• The loop iteration variable in any associated loop of a `for`, `parallel for`, `taskloop`, or `distribute` construct is private.

• The loop iteration variable in the associated loop of a `simd` construct with just one associated loop is linear with a `linear-step` that is the increment of the associated loop.

• The loop iteration variables in the associated loops of a `simd` construct with multiple associated loops are lastprivate.

• The loop iteration variable in any associated loop of a `loop` construct is lastprivate.

• The implicitly declared variables of a range-based `for` loop are private.

• Variables with static storage duration that are declared in a scope inside the construct are shared.

• If a list item in a `map` clause on the `target` construct has a base pointer, and the base pointer is a scalar variable that does not appear in a `map` clause on the construct, the base pointer is firstprivate.

• If a list item in a `reduction` or `in_reduction` clause on a construct has a base pointer then the base pointer is private.

• Static data members are shared.

• The `__func__` variable and similar function-local predefined variables are shared.

• Variables declared within a `BLOCK` construct inside a construct that do not have the `SAVE` attribute are private.

• Variables and common blocks that appear in `threadprivate` directives are threadprivate.

• The loop iteration variable in any associated `do-loop` of a `do`, `parallel do`, `taskloop`, or `distribute` construct is private.

• The loop iteration variable in the associated `do-loop` of a `simd` construct with just one associated `do-loop` is linear with a `linear-step` that is the increment of the associated `do-loop`. 
• The loop iteration variables in the associated *do-loops* of a *simd* construct with multiple
  associated *do-loops* are lastprivate.

• The loop iteration variable in any associated *do-loop* of a *loop* construct is lastprivate.

• Loop iteration variables inside *parallel* or task generating constructs are private in the
  innermost such construct that encloses the loop.

• Implied-do, **FORALL** and **DO CONCURRENT** indices are private.

• Cray pointees have the same data-sharing attribute as the storage with which their Cray pointers
  are associated. Cray pointer support has been deprecated.

• Assumed-size arrays are shared.

• *Named constants* are shared.

• An associate name that may appear in a variable definition context is shared if its association
  occurs outside of the construct and otherwise it has the same data-sharing attribute as the
  selector with which it is associated.

---

**Fortran**

Variables with predetermined data-sharing attributes may not be listed in data-sharing attribute
clauses, except for the cases listed below. For these exceptions only, listing a predetermined
variable in a data-sharing attribute clause is allowed and overrides the variable’s predetermined
data-sharing attributes.

---

**C / C++**

• The loop iteration variable in any associated loop of a *for*, *taskloop*, *distribute*, or
  *loop* construct may be listed in a *private* or *lastprivate* clause.

• If a *simd* construct has just one associated loop then its loop iteration variable may be listed in a
  *private*, *lastprivate*, or *linear* clause with a *linear-step* that is the increment of the
  associated loop.

• If a *simd* construct has more than one associated loop then their loop iteration variables may be
  listed in a *private* or *lastprivate* clause.

• Variables with *const*-qualified type with no mutable members may be listed in a
  *firstprivate* clause, even if they are static data members.

• The **__func__** variable and similar function-local predefined variables may be listed in a
  *shared* or *firstprivate* clause.
• The loop iteration variable in any associated do-loop of a do, taskloop, distribute, or loop construct may be listed in a private or lastprivate clause.

• The loop iteration variable in the associated do-loop of a simd construct with just one associated do-loop may be listed in a private, lastprivate, or linear clause with a linear-step that is the increment of the associated loop.

• The loop iteration variables in the associated do-loops of a simd construct with multiple associated do-loops may be listed in a private or lastprivate clause.

• Loop iteration variables of loops that are not associated with any OpenMP directive may be listed in data-sharing attribute clauses on the surrounding teams, parallel, or task generating construct, and on enclosed constructs, subject to other restrictions.

• Assumed-size arrays may be listed in a shared clause.

• Named constants may be listed in a firstprivate clause.

Additional restrictions on the variables that may appear in individual clauses are described with each clause in Section 2.21.4.

Variables with explicitly determined data-sharing attributes are those that are referenced in a given construct and are listed in a data-sharing attribute clause on the construct.

Variables with implicitly determined data-sharing attributes are those that are referenced in a given construct, do not have predetermined data-sharing attributes, and are not listed in a data-sharing attribute clause on the construct.

Rules for variables with implicitly determined data-sharing attributes are as follows:

• In a parallel, teams, or task generating construct, the data-sharing attributes of these variables are determined by the default clause, if present (see Section 2.21.4.1).

• In a parallel construct, if no default clause is present, these variables are shared.

• For constructs other than task generating constructs, if no default clause is present, these variables reference the variables with the same names that exist in the enclosing context.

• In a target construct, variables that are not mapped after applying data-mapping attribute rules (see Section 2.21.7) are firstprivate.

• In an orphaned task generating construct, if no default clause is present, formal arguments passed by reference are firstprivate.
• In an orphaned task generating construct, if no `default` clause is present, dummy arguments are firstprivate.

• In a task generating construct, if no `default` clause is present, a variable for which the data-sharing attribute is not determined by the rules above and that in the enclosing context is determined to be shared by all implicit tasks bound to the current team is shared.

• In a task generating construct, if no `default` clause is present, a variable for which the data-sharing attribute is not determined by the rules above is firstprivate.

Additional restrictions on the variables for which data-sharing attributes cannot be implicitly determined in a task generating construct are described in Section 2.21.4.4.

### 2.21.1.2 Variables Referenced in a Region but not in a Construct

The data-sharing attributes of variables that are referenced in a region, but not in a construct, are determined as follows:

#### C / C++

• Variables with static storage duration that are declared in called routines in the region are shared.

• File-scope or namespace-scope variables referenced in called routines in the region are shared unless they appear in a `threadprivate` directive.

• Objects with dynamic storage duration are shared.

• Static data members are shared unless they appear in a `threadprivate` directive.

• In C++, formal arguments of called routines in the region that are passed by reference have the same data-sharing attributes as the associated actual arguments.

• Other variables declared in called routines in the region are private.

#### C / C++

• Local variables declared in called routines in the region and that have the `save` attribute, or that are data initialized, are shared unless they appear in a `threadprivate` directive.

• Variables belonging to common blocks, or accessed by host or use association, and referenced in called routines in the region are shared unless they appear in a `threadprivate` directive.

• Dummy arguments of called routines in the region that have the `VALUE` attribute are private.

• Dummy arguments of called routines in the region that do not have the `VALUE` attribute are private if the associated actual argument is not shared.
• Dummy arguments of called routines in the region that do not have the \texttt{VALUE} attribute are shared if the actual argument is shared and it is a scalar variable, structure, an array that is not a pointer or assumed-shape array, or a simply contiguous array section. Otherwise, the data-sharing attribute of the dummy argument is implementation-defined if the associated actual argument is shared.

• Cray pointees have the same data-sharing attribute as the storage with which their Cray pointers are associated. Cray pointer support has been deprecated.

• Implied-do indices, \texttt{DO CONCURRENT} indices, \texttt{FORALL} indices, and other local variables declared in called routines in the region are private.

\section*{2.21.2 \texttt{threadprivate} Directive}

\textbf{Summary}

The \texttt{threadprivate} directive specifies that variables are replicated, with each thread having its own copy. The \texttt{threadprivate} directive is a declarative directive.

\textbf{Syntax}

\begin{verbatim}
C / C++
#pragma omp threadprivate(list)
\end{verbatim}

\begin{verbatim}
C / C++
#pragma omp threadprivate(list) new-line
\end{verbatim}

\begin{verbatim}
Fortran
!$omp threadprivate(list)
\end{verbatim}

where \texttt{list} is a comma-separated list of file-scope, namespace-scope, or static block-scope variables that do not have incomplete types.

\begin{verbatim}
C / C++
#pragma omp threadprivate
\end{verbatim}

\begin{verbatim}
Fortran
#$omp threadprivate(list)
\end{verbatim}

where \texttt{list} is a comma-separated list of named variables and named common blocks. Common block names must appear between slashes.

\textbf{Description}

Unless otherwise specified, each copy of a \texttt{threadprivate} variable is initialized once, in the manner specified by the program, but at an unspecified point in the program prior to the first reference to that copy. The storage of all copies of a \texttt{threadprivate} variable is freed according to how static variables are handled in the base language, but at an unspecified point in the program.
Each copy of a block-scope threadprivate variable that has a dynamic initializer is initialized the first time its thread encounters its definition; if its thread does not encounter its definition, its initialization is unspecified.

The content of a threadprivate variable can change across a task scheduling point if the executing thread switches to another task that modifies the variable. For more details on task scheduling, see Section 1.3 and Section 2.12.

In parallel regions, references by the primary thread will be to the copy of the variable in the thread that encountered the parallel region.

During a sequential part references will be to the initial thread’s copy of the variable. The values of data in the initial thread’s copy of a threadprivate variable are guaranteed to persist between any two consecutive references to the variable in the program provided that no teams construct that is not nested inside of a target construct is encountered between the references and that the initial thread is not nested inside of a teams region. For initial threads nested inside of a teams region, the values of data in the copies of a threadprivate variable of those initial threads are guaranteed to persist between any two consecutive references to the variable inside of that teams region.

The values of data in the threadprivate variables of threads that are not initial threads are guaranteed to persist between two consecutive active parallel regions only if all of the following conditions hold:

• Neither parallel region is nested inside another explicit parallel region;
• The number of threads used to execute both parallel regions is the same;
• The thread affinity policies used to execute both parallel regions are the same;
• The value of the dyn-var internal control variable in the enclosing task region is false at entry to both parallel regions;
• No teams construct that is not nested inside of a target construct is encountered between the parallel regions;
• No construct with an order clause that specifies concurrent is encountered between the parallel regions; and
• Neither the omp_pause_resource nor omp_pause_resource_all routine is called.

If these conditions all hold, and if a threadprivate variable is referenced in both regions, then threads with the same thread number in their respective regions will reference the same copy of that variable.
If the above conditions hold, the storage duration, lifetime, and value of a thread’s copy of a threadprivate variable that does not appear in any `copyin` clause on the second region will span the two consecutive active `parallel` regions. Otherwise, the storage duration, lifetime, and value of a thread’s copy of the variable in the second region is unspecified.

If the above conditions hold, the definition, association, or allocation status of a thread’s copy of a threadprivate variable or a variable in a threadprivate common block that is not affected by any `copyin` clause that appears on the second region (a variable is affected by a `copyin` clause if the variable appears in the `copyin` clause or it is in a common block that appears in the `copyin` clause) will span the two consecutive active `parallel` regions. Otherwise, the definition and association status of a thread’s copy of the variable in the second region are undefined, and the allocation status of an allocatable variable will be implementation defined.

If a threadprivate variable or a variable in a threadprivate common block is not affected by any `copyin` clause that appears on the first `parallel` region in which it is referenced, the thread’s copy of the variable inherits the declared type parameter and the default parameter values from the original variable. The variable or any subobject of the variable is initially defined or undefined according to the following rules:

- If it has the `ALLOCATABLE` attribute, each copy created will have an initial allocation status of unallocated;
- If it has the `POINTER` attribute, each copy will have the same association status as the initial association status.
- If it does not have either the `POINTER` or the `ALLOCATABLE` attribute:
  - If it is initially defined, either through explicit initialization or default initialization, each copy created is so defined;
  - Otherwise, each copy created is undefined.

The address of a threadprivate variable may not be an address constant.

The order in which any constructors for different threadprivate variables of class type are called is unspecified. The order in which any destructors for different threadprivate variables of class type are called is unspecified.
Restrictions

Restrictions to the threadprivate directive are as follows:

• A thread must not reference another thread’s copy of a threadprivate variable.

• A threadprivate variable must not appear in any clause except the copyin, copyprivate, schedule, num_threads, thread_limit, and if clauses.

• A program in which an untied task accesses threadprivate storage is non-conforming.

C / C++

• If the value of a variable referenced in an explicit initializer of a threadprivate variable is modified prior to the first reference to any instance of the threadprivate variable, then the behavior is unspecified.

• A variable that is part of another variable (as an array or structure element) cannot appear in a threadprivate directive unless it is a static data member of a C++ class.

• A threadprivate directive for file-scope variables must appear outside any definition or declaration, and must lexically precede all references to any of the variables in its list.

• A threadprivate directive for namespace-scope variables must appear outside any definition or declaration other than the namespace definition itself, and must lexically precede all references to any of the variables in its list.

• Each variable in the list of a threadprivate directive at file, namespace, or class scope must refer to a variable declaration at file, namespace, or class scope that lexically precedes the directive.

• A threadprivate directive for static block-scope variables must appear in the scope of the variable and not in a nested scope. The directive must lexically precede all references to any of the variables in its list.

• Each variable in the list of a threadprivate directive in block scope must refer to a variable declaration in the same scope that lexically precedes the directive. The variable must have static storage duration.

• If a variable is specified in a threadprivate directive in one translation unit, it must be specified in a threadprivate directive in every translation unit in which it is declared.

C / C++  

C++

• A threadprivate directive for static class member variables must appear in the class definition, in the same scope in which the member variables are declared, and must lexically precede all references to any of the variables in its list.

• A threadprivate variable must not have an incomplete type or a reference type.

• A threadprivate variable with class type must have:
– An accessible, unambiguous default constructor in the case of default initialization without a
  given initializer;

– An accessible, unambiguous constructor that accepts the given argument in the case of direct
  initialization; and

– An accessible, unambiguous copy constructor in the case of copy initialization with an explicit
  initializer.

Fortran

• A variable that is part of another variable (as an array, structure element or type parameter
  inquiry) cannot appear in a threadprivate directive.

• A coarray cannot appear in a threadprivate directive.

• An associate name cannot appear in a threadprivate directive.

• The threadprivate directive must appear in the declaration section of a scoping unit in
  which the common block or variable is declared.

• If a threadprivate directive that specifies a common block name appears in one program
  unit, then such a directive must also appear in every other program unit that contains a COMMON
  statement that specifies the same name. It must appear after the last such COMMON statement in
  the program unit.

• If a threadprivate variable or a threadprivate common block is declared with the BIND
  attribute, the corresponding C entities must also be specified in a threadprivate directive in the C
  program.

• A variable can only appear in a threadprivate directive in the scope in which it is declared.
  It must not be an element of a common block or appear in an EQUIVALENCE statement.

• A variable that appears in a threadprivate directive must be declared in the scope of a
  module or have the SAVE attribute, either explicitly or implicitly.

• The effect of an access to a threadprivate variable in a DO CONCURRENT construct is unspecified.

Cross References

• dyn-var ICV, see Section 2.4.

• Number of threads used to execute a parallel region, see Section 2.6.1.

• order clause, see Section 2.11.3.

• copyin clause, see Section 2.21.6.1.
2.21.3 List Item Privatization

For any construct, a list item that appears in a data-sharing attribute clause, including a reduction clause, may be privatized. Each task that references a privatized list item in any statement in the construct receives at least one new list item if the construct has one or more associated loops, and otherwise each such task receives one new list item. Each SIMD lane used in a simd construct that references a privatized list item in any statement in the construct receives at least one new list item. Language-specific attributes for new list items are derived from the corresponding original list item. Inside the construct, all references to the original list item are replaced by references to a new list item received by the task or SIMD lane.

If the construct has one or more associated loops, within the same logical iteration of the loops, then the same new list item replaces all references to the original list item. For any two logical iterations, if the references to the original list item are replaced by the same list item then the logical iterations must execute in some sequential order.

In the rest of the region, whether references are to a new list item or the original list item is unspecified. Therefore, if an attempt is made to reference the original item, its value after the region is also unspecified. If a task or a SIMD lane does not reference a privatized list item, whether the task or SIMD lane receives a new list item is unspecified.

The value and/or allocation status of the original list item will change only:

- If accessed and modified via pointer;
- If possibly accessed in the region but outside of the construct;
- As a side effect of directives or clauses; or
- If accessed and modified via construct association.

If the construct is contained in a member function, whether accesses anywhere in the region through the implicit this pointer refer to the new list item or the original list item is unspecified.

A new list item of the same type, with automatic storage duration, is allocated for the construct. The storage and thus lifetime of these list items last until the block in which they are created exits. The size and alignment of the new list item are determined by the type of the variable. This allocation occurs once for each task generated by the construct and once for each SIMD lane used by the construct.

The new list item is initialized, or has an undefined initial value, as if it had been locally declared without an initializer.
C++

If the type of a list item is a reference to a type $T$ then the type will be considered to be $T$ for all purposes of this clause.

The order in which any default constructors for different private variables of class type are called is unspecified. The order in which any destructors for different private variables of class type are called is unspecified.

C++

Fortran

If any statement of the construct references a list item, a new list item of the same type and type parameters is allocated. This allocation occurs once for each task generated by the construct and once for each SIMD lane used by the construct. If the type of the list item has default initialization, the new list item has default initialization. Otherwise, the initial value of the new list item is undefined. The initial status of a private pointer is undefined.

For a list item or the subobject of a list item with the **ALLOCATABLE** attribute:

- If the allocation status is unallocated, the new list item or the subobject of the new list item will have an initial allocation status of unallocated;

- If the allocation status is allocated, the new list item or the subobject of the new list item will have an initial allocation status of allocated; and

- If the new list item or the subobject of the new list item is an array, its bounds will be the same as those of the original list item or the subobject of the original list item.

A privatized list item may be storage-associated with other variables when the data-sharing attribute clause is encountered. Storage association may exist because of constructs such as **EQUIVALENCE** or **COMMON**. If $A$ is a variable that is privatized by a construct and $B$ is a variable that is storage-associated with $A$, then:

- The contents, allocation, and association status of $B$ are undefined on entry to the region;

- Any definition of $A$, or of its allocation or association status, causes the contents, allocation, and association status of $B$ to become undefined; and

- Any definition of $B$, or of its allocation or association status, causes the contents, allocation, and association status of $A$ to become undefined.

A privatized list item may be a selector of an **ASSOCIATE** or **SELECT TYPE** construct. If the construct association is established prior to a **parallel** region, the association between the associate name and the original list item will be retained in the region.

Finalization of a list item of a finalizable type or subobjects of a list item of a finalizable type occurs at the end of the region. The order in which any final subroutines for different variables of a finalizable type are called is unspecified.
If a list item appears in both `firstprivate` and `lastprivate` clauses, the update required for the `lastprivate` clause occurs after all initializations for the `firstprivate` clause.

### Restrictions
The following restrictions apply to any list item that is privatized unless otherwise stated for a given data-sharing attribute clause:

- A variable that is part of another variable (as an array or structure element) cannot be privatized.
- A variable that is part of another variable (as an array or structure element) cannot be privatized except if the data-sharing attribute clause is associated with a construct within a class non-static member function and the variable is an accessible data member of the object for which the non-static member function is invoked.
- A variable of class type (or array thereof) that is privatized requires an accessible, unambiguous default constructor for the class type.
- A variable that is privatized must not have a `const`-qualified type unless it is of class type with a `mutable` member. This restriction does not apply to the `firstprivate` clause.
- A variable that is privatized must not have an incomplete type or be a reference to an incomplete type.
- A variable that is part of another variable (as an array or structure element) cannot be privatized.
- Variables that appear in namelist statements, in variable format expressions, and in expressions for statement function definitions, may not be privatized.
- Pointers with the `INTENT(IN)` attribute may not be privatized. This restriction does not apply to the `firstprivate` clause.
- A private variable must not be coindexed or appear as an actual argument to a procedure where the corresponding dummy argument is a coarray.
- Assumed-size arrays may not be privatized in a `target`, `teams`, or `distribute` construct.
2.21.4 Data-Sharing Attribute Clauses

Several constructs accept clauses that allow a user to control the data-sharing attributes of variables referenced in the construct. Not all of the clauses listed in this section are valid on all directives. The set of clauses that is valid on a particular directive is described with the directive.

Most of the clauses accept a comma-separated list of list items (see Section 2.1). All list items that appear in a clause must be visible, according to the scoping rules of the base language. With the exception of the default clause, clauses may be repeated as needed. A list item may not appear in more than one clause on the same directive, except that it may be specified in both firstprivate and lastprivate clauses.

The reduction data-sharing attribute clauses are explained in Section 2.21.5.

C++

If a variable referenced in a data-sharing attribute clause has a type derived from a template, and the program does not otherwise reference that variable then any behavior related to that variable is unspecified.

Fortran

If individual members of a common block appear in a data-sharing attribute clause other than the shared clause, the variables no longer have a Fortran storage association with the common block.

2.21.4.1 default Clause

Summary

The default clause explicitly determines the data-sharing attributes of variables that are referenced in a parallel, teams, or task generating construct and would otherwise be implicitly determined (see Section 2.21.1.1).

Syntax

The syntax of the default clause is as follows:

```
default (data-sharing-attribute)
```

where data-sharing-attribute is one of the following:

```
shared
firstprivate
private
none
```
Description
If data-sharing-attribute is shared or, for Fortran, firstprivate or private, the
data-sharing attribute of all variables referenced in the construct that have implicitly determined
data-sharing attributes will be data-sharing-attribute.

If data-sharing-attribute is firstprivate or private, each variable with static storage
duration that is declared in a namespace or global scope and referenced in the construct, and that
does not have a predetermined data-sharing attribute, must have its data-sharing attribute explicitly
determined by being listed in a data-sharing attribute clause. The data-sharing attribute of all other
variables that are referenced in the construct and that have implicitly determined data-sharing
attributes will be data-sharing-attribute.

The default(none) clause requires that each variable that is referenced in the construct, and
that does not have a predetermined data-sharing attribute, must have its data-sharing attribute
explicitly determined by being listed in a data-sharing attribute clause.

Restrictions
Restrictions to the default clause are as follows:

- Only a single default clause may be specified on a parallel, task, taskloop or
teams directive.

2.21.4.2 shared Clause

Summary
The shared clause declares one or more list items to be shared by tasks generated by a
parallel, teams, or task generating construct.

Syntax
The syntax of the shared clause is as follows:

\[
\text{shared(list)}
\]

Description
All references to a list item within a task refer to the storage area of the original variable at the point
the directive was encountered.

The programmer must ensure, by adding proper synchronization, that storage shared by an explicit
task region does not reach the end of its lifetime before the explicit task region completes its
execution.
The association status of a shared pointer becomes undefined upon entry to and exit from a parallel, teams, or task generating construct if it is associated with a target or a subobject of a target that appears as a privatized list item in a data-sharing attribute clause on the construct.

Note – Passing a shared variable to a procedure may result in the use of temporary storage in place of the actual argument when the corresponding dummy argument does not have the VALUE or CONTIGUOUS attribute and its data-sharing attribute is implementation-defined as per the rules in Section 2.21.1.2. These conditions effectively result in references to, and definitions of, the temporary storage during the procedure reference. Furthermore, the value of the shared variable is copied into the intervening temporary storage before the procedure reference when the dummy argument does not have the INTENT (OUT) attribute, and is copied out of the temporary storage into the shared variable when the dummy argument does not have the INTENT (IN) attribute. Any references to (or definitions of) the shared storage that is associated with the dummy argument by any other task must be synchronized with the procedure reference to avoid possible data races.

Restrictions
Restrictions to the shared clause are as follows:

- A variable that is part of another variable (as an array or structure element) cannot appear in a shared clause.

- A variable that is part of another variable (as an array or structure element) cannot appear in a shared clause except if the shared clause is associated with a construct within a class non-static member function and the variable is an accessible data member of the object for which the non-static member function is invoked.

- A variable that is part of another variable (as an array, structure element or type parameter inquiry) cannot appear in a shared clause.
### 2.21.4.3 private Clause

#### Summary
The `private` clause declares one or more list items to be private to a task or to a SIMD lane.

#### Syntax
The syntax of the `private` clause is as follows:

```
private(list)
```

#### Description
The `private` clause specifies that its list items are to be privatized according to Section 2.21.3. Each task or SIMD lane that references a list item in the construct receives only one new list item, unless the construct has one or more associated loops and an `order` clause that specifies `concurrent` is also present.

#### Restrictions
Restrictions to the `private` clause are as specified in Section 2.21.3.

#### Cross References
- List Item Privatization, see Section 2.21.3.

### 2.21.4.4 firstprivate Clause

#### Summary
The `firstprivate` clause declares one or more list items to be private to a task, and initializes each of them with the value that the corresponding original item has when the construct is encountered.

#### Syntax
The syntax of the `firstprivate` clause is as follows:

```
firstprivate(list)
```

#### Description
The `firstprivate` clause provides a superset of the functionality provided by the `private` clause.

The list items that appear in a `firstprivate` clause may include named constants.
A list item that appears in a `firstprivate` clause is subject to the `private` clause semantics described in Section 2.21.4.3, except as noted. In addition, the new list item is initialized from the original list item that exists before the construct. The initialization of the new list item is done once for each task that references the list item in any statement in the construct. The initialization is done prior to the execution of the construct.

For a `firstprivate` clause on a `parallel`, `task`, `taskloop`, `target`, or `teams` construct, the initial value of the new list item is the value of the original list item that exists immediately prior to the construct in the task region where the construct is encountered unless otherwise specified. For a `firstprivate` clause on a worksharing construct, the initial value of the new list item for each implicit task of the threads that execute the worksharing construct is the value of the original list item that exists in the implicit task immediately prior to the point in time that the worksharing construct is encountered unless otherwise specified.

To avoid data races, concurrent updates of the original list item must be synchronized with the read of the original list item that occurs as a result of the `firstprivate` clause.

```c++
C / C++
```

For variables of non-array type, the initialization occurs by copy assignment. For an array of elements of non-array type, each element is initialized as if by assignment from an element of the original array to the corresponding element of the new array.

```c++
C / C++
```

For each variable of class type:

- If the `firstprivate` clause is not on a `target` construct then a copy constructor is invoked to perform the initialization; and

- If the `firstprivate` clause is on a `target` construct then how many copy constructors, if any, are invoked is unspecified.

If copy constructors are called, the order in which copy constructors for different variables of class type are called is unspecified.

```c++
```

If the original list item does not have the `POINTER` attribute, initialization of the new list items occurs as if by intrinsic assignment unless the original list item has a compatible type-bound defined assignment, in which case initialization of the new list items occurs as if by the defined assignment. If the original list item that does not have the `POINTER` attribute has the allocation status of unallocated, the new list items will have the same status.

If the original list item has the `POINTER` attribute, the new list items receive the same association status of the original list item as if by pointer assignment.

```fortran
```

```fortran
```
Restrictions

Restrictions to the `firstprivate` clause are as follows:

- A list item that is private within a `parallel` region must not appear in a `firstprivate` clause on a worksharing construct if any of the worksharing regions that arise from the worksharing construct ever bind to any of the `parallel` regions that arise from the `parallel` construct.

- A list item that is private within a `teams` region must not appear in a `firstprivate` clause on a `distribute` construct if any of the `distribute` regions that arise from the `distribute` construct ever bind to any of the `teams` regions that arise from the `teams` construct.

- A list item that appears in a `reduction` clause of a `parallel` construct must not appear in a `firstprivate` clause on a worksharing, `task`, or `taskloop` construct if any of the worksharing or `task` regions that arise from the worksharing, `task`, or `taskloop` construct ever bind to any of the `parallel` regions that arise from the `parallel` construct.

- A list item that appears in a `reduction` clause of a `teams` construct must not appear in a `firstprivate` clause on a `distribute` construct if any of the `distribute` regions that arise from the `distribute` construct ever bind to any of the `teams` regions that arise from the `teams` construct.

- A list item that appears in a `reduction` clause of a worksharing construct must not appear in a `firstprivate` clause in a `task` construct encountered during execution of any of the worksharing regions that arise from the worksharing construct.

C++

- A variable of class type (or array thereof) that appears in a `firstprivate` clause requires an accessible, unambiguous copy constructor for the class type.

C / C++

- If a list item in a `firstprivate` clause on a worksharing construct has a reference type then it must bind to the same object for all threads of the team.

Fortran

- If the list item is a polymorphic variable with the `ALLOCATABLE` attribute, the behavior is unspecified.
2.21.4.5 lastprivate Clause

Summary
The lastprivate clause declares one or more list items to be private to an implicit task or to a
SIMD lane, and causes the corresponding original list item to be updated after the end of the region.

Syntax
The syntax of the lastprivate clause is as follows:

```
lastprivate([ lastprivate-modifier:] list)
```

where lastprivate-modifier is:

```
conditional
```

Description
The lastprivate clause provides a superset of the functionality provided by the private
clause.

A list item that appears in a lastprivate clause is subject to the private clause semantics
described in Section 2.21.4.3. In addition, when a lastprivate clause without the
conditional modifier appears on a directive and the list item is not an iteration variable of one
of the associated loops, the value of each new list item from the sequentially last iteration of the
associated loops, or the lexically last section construct, is assigned to the original list item.
When the conditional modifier appears on the clause or the list item is an iteration variable of
one of the associated loops, if sequential execution of the loop nest would assign a value to the list
item then the original list item is assigned the value that the list item would have after sequential
execution of the loop nest.

```
\[ C / C++ \]
```

For an array of elements of non-array type, each element is assigned to the corresponding element
of the original array.

```
\[ C / C++ \] \[ Fortran \]
```

If the original list item does not have the POINTER attribute, its update occurs as if by intrinsic
assignment unless it has a type bound procedure as a defined assignment.

If the original list item has the POINTER attribute, its update occurs as if by pointer assignment.

```
\[ Fortran \]
```

When the conditional modifier does not appear on the lastprivate clause, any list item
that is not an iteration variable of the associated loops and that is not assigned a value by the
sequentially last iteration of the loops, or by the lexically last section construct, has an
unspecified value after the construct. When the conditional modifier does not appear on the
lastprivate clause, a list item that is the iteration variable of an associated loop and that would
not be assigned a value during sequential execution of the loop nest has an unspecified value after
the construct. Unassigned subcomponents also have unspecified values after the construct.
If the `lastprivate` clause is used on a construct to which neither the `nowait` nor the `nogroup` clauses are applied, the original list item becomes defined at the end of the construct. To avoid data races, concurrent reads or updates of the original list item must be synchronized with the update of the original list item that occurs as a result of the `lastprivate` clause.

Otherwise, If the `lastprivate` clause is used on a construct to which the `nowait` or the `nogroup` clauses are applied, accesses to the original list item may create a data race. To avoid this data race, if an assignment to the original list item occurs then synchronization must be inserted to ensure that the assignment completes and the original list item is flushed to memory.

If a list item that appears in a `lastprivate` clause with the `conditional` modifier is modified in the region by an assignment outside the construct or not to the list item then the value assigned to the original list item is unspecified.

**Restrictions**

Restrictions to the `lastprivate` clause are as follows:

- A list item that is private within a `parallel` region, or that appears in the `reduction` clause of a `parallel` construct, must not appear in a `lastprivate` clause on a worksharing construct if any of the corresponding worksharing regions ever binds to any of the corresponding `parallel` regions.

- A list item that appears in a `lastprivate` clause with the `conditional` modifier must be a scalar variable.

- A variable of class type (or array thereof) that appears in a `lastprivate` clause requires an accessible, unambiguous default constructor for the class type, unless the list item is also specified in a `firstprivate` clause.

- A variable of class type (or array thereof) that appears in a `lastprivate` clause requires an accessible, unambiguous copy assignment operator for the class type. The order in which copy assignment operators for different variables of class type are called is unspecified.

- If a list item in a `lastprivate` clause on a worksharing construct has a reference type then it must bind to the same object for all threads of the team.

- A variable that appears in a `lastprivate` clause must be definable.

- If the original list item has the `ALLOCATABLE` attribute, the corresponding list item of which the value is assigned to the original item must have an allocation status of allocated upon exit from the sequentially last iteration or lexically last `section` construct.

- If the list item is a polymorphic variable with the `ALLOCATABLE` attribute, the behavior is unspecified.
2.21.4.6 linear Clause

Summary
The linear clause declares one or more list items to be private and to have a linear relationship with respect to the iteration space of a loop associated with the construct on which the clause appears.

Syntax

The syntax of the linear clause is as follows:

```
linear(linear-list[ : linear-step])
```

where linear-list is one of the following:

```
list
modifier (list)
```

where modifier is one of the following:

```
val
```

The syntax of the linear clause is as follows:

```
linear(linear-list[ : linear-step])
```

where linear-list is one of the following:

```
list
modifier (list)
```

where modifier is one of the following:

```
ref
val
uval
```
The syntax of the `linear` clause is as follows:

```
linear(linear-list [ : linear-step])
```

where `linear-list` is one of the following:

```
list
  modifier (list)
```

where `modifier` is one of the following:

```
ref
val
uval
```

### Description

The `linear` clause provides a superset of the functionality provided by the `private` clause. A list item that appears in a `linear` clause is subject to the `private` clause semantics described in Section 2.21.4.3 except as noted. If `linear-step` is not specified, it is assumed to be 1.

When a `linear` clause is specified on a construct, the value of the new list item on each logical iteration of the associated loops corresponds to the value of the original list item before entering the construct plus the logical number of the iteration times `linear-step`. The value corresponding to the sequentially last logical iteration of the associated loops is assigned to the original list item.

When a `linear` clause is specified on a declarative directive, all list items must be formal parameters (or, in Fortran, dummy arguments) of a function that will be invoked concurrently on each SIMD lane. If no `modifier` is specified or the `val` or `uval` modifier is specified, the value of each list item on each lane corresponds to the value of the list item upon entry to the function plus the logical number of the lane times `linear-step`. If the `uval` modifier is specified, each invocation uses the same storage location for each SIMD lane; this storage location is updated with the final value of the logically last lane. If the `ref` modifier is specified, the storage location of each list item on each lane corresponds to an array at the storage location upon entry to the function indexed by the logical number of the lane times `linear-step`.

### Restrictions

Restrictions to the `linear` clause are as follows:

- The `linear-step` expression must be invariant during the execution of the region that corresponds to the construct.
- Only a loop iteration variable of a loop that is associated with the construct may appear as a `list-item` in a `linear` clause if a `reduction` clause with the `inscan` modifier also appears on the construct.
A list-item that appears in a **linear** clause must be of integral or pointer type.

A list-item that appears in a **linear** clause without the **ref** modifier must be of integral or pointer type, or must be a reference to an integral or pointer type.

The **ref** or **uval** modifier can only be used if the list-item is of a reference type.

If a list item in a **linear** clause on a worksharing construct has a reference type then it must bind to the same object for all threads of the team.

If the list item is of a reference type and the **ref** modifier is not specified and if any write to the list item occurs before any read of the list item then the result is unspecified.

A list-item that appears in a **linear** clause without the **ref** modifier must be of type **integer**.

The **ref** or **uval** modifier can only be used if the list-item is a dummy argument without the **VALUE** attribute.

Variables that have the **POINTER** attribute and Cray pointers may not appear in a **linear** clause. Cray pointer support has been deprecated.

If the list item has the **ALLOCATABLE** attribute and the **ref** modifier is not specified, the allocation status of the list item in the sequentially last iteration must be allocated upon exit from that iteration.

If the **ref** modifier is specified, variables with the **ALLOCATABLE** attribute, assumed-shape arrays and polymorphic variables may not appear in the **linear** clause.

If the list item is a dummy argument without the **VALUE** attribute and the **ref** modifier is not specified, any read of the list item must occur before any write to the list item.

A common block name cannot appear in a **linear** clause.

### 2.21.5 Reduction Clauses and Directives

The reduction clauses are data-sharing attribute clauses that can be used to perform some forms of recurrence calculations in parallel. Reduction clauses include reduction scoping clauses and reduction participating clauses. Reduction scoping clauses define the region in which a reduction is computed. Reduction participating clauses define the participants in the reduction.
2.21.5.1 Properties Common to All Reduction Clauses

Syntax

The syntax of a reduction-identifier is defined as follows:

\[\text{C} \]

A reduction-identifier is either an identifier or one of the following operators: +, −, *, &, |, ^, && and ||.

\[\text{C++} \]

A reduction-identifier is either an id-expression or one of the following operators: +, −, *, &, |, ^, && and ||.

\[\text{C++} \]

A reduction-identifier is either a base language identifier, or a user-defined operator, or one of the following operators: +, −, *, .and., .or., .eqv., .neqv., or one of the following intrinsic procedure names: max, min, iand, ior, ieor.

\[\text{Fortran} \]

A reduction-identifier is either a base language identifier, or a user-defined operator, or one of the following operators: +, −, *, .and., .or., .eqv., .neqv., or one of the following intrinsic procedure names: max, min, iand, ior, ieor.

Table 2.11 lists each reduction-identifier that is implicitly declared at every scope for arithmetic types and its semantic initializer value. The actual initializer value is that value as expressed in the data type of the reduction list item.

**Table 2.11: Implicitly Declared C/C++ reduction-identifiers**

<table>
<thead>
<tr>
<th>Identifier</th>
<th>Initializer</th>
<th>Combiner</th>
</tr>
</thead>
<tbody>
<tr>
<td>+</td>
<td>omp_priv = 0</td>
<td>omp_out += omp_in</td>
</tr>
<tr>
<td>−</td>
<td>omp_priv = 0</td>
<td>omp_out += omp_in</td>
</tr>
<tr>
<td>*</td>
<td>omp_priv = 1</td>
<td>omp_out *= omp_in</td>
</tr>
<tr>
<td>&amp;</td>
<td>omp_priv = ~ 0</td>
<td>omp_out &amp;= omp_in</td>
</tr>
<tr>
<td></td>
<td></td>
<td>omp_priv = 0</td>
</tr>
<tr>
<td>^</td>
<td>omp_priv = 0</td>
<td>omp_out ^= omp_in</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>omp_priv = 1</td>
<td>omp_out = omp_in &amp;&amp; omp_out</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

*table continued on next page*
Table 2.12 lists each reduction-identifier that is implicitly declared for numeric and logical types and its semantic initializer value. The actual initializer value is that value as expressed in the data type of the reduction list item.

**Table 2.12: Implicitly Declared Fortran reduction-identifiers**

<table>
<thead>
<tr>
<th>Identifier</th>
<th>Initializer</th>
<th>Combiner</th>
</tr>
</thead>
<tbody>
<tr>
<td>+</td>
<td>omp_priv = 0</td>
<td>omp_out = omp_in + omp_out</td>
</tr>
<tr>
<td>-</td>
<td>omp_priv = 0</td>
<td>omp_out = omp_in + omp_out</td>
</tr>
<tr>
<td>*</td>
<td>omp_priv = 1</td>
<td>omp_out = omp_in * omp_out</td>
</tr>
<tr>
<td>.and.</td>
<td>omp_priv = .true.</td>
<td>omp_out = omp_in .and. omp_out</td>
</tr>
<tr>
<td>.or.</td>
<td>omp_priv = .false.</td>
<td>omp_out = omp_in .or. omp_out</td>
</tr>
<tr>
<td>.eqv.</td>
<td>omp_priv = .true.</td>
<td>omp_out = omp_in .eqv. omp_out</td>
</tr>
<tr>
<td>.neqv.</td>
<td>omp_priv = .false.</td>
<td>omp_out = omp_in .neqv. omp_out</td>
</tr>
<tr>
<td>max</td>
<td>omp_priv = Minimal representable number in the reduction list item type</td>
<td>omp_out = max(omp_in, omp_out)</td>
</tr>
<tr>
<td>min</td>
<td>omp_priv = Maximal representable number in the reduction list item type</td>
<td>omp_out = min(omp_in, omp_out)</td>
</tr>
</tbody>
</table>
In the above tables, `omp_in` and `omp_out` correspond to two identifiers that refer to storage of the type of the list item. If the list item is an array or array section, the identifiers to which `omp_in` and `omp_out` correspond each refer to an array element. `omp_out` holds the final value of the combiner operation.

Any `reduction-identifier` that is defined with the `declare reduction` directive is also valid. In that case, the initializer and combiner of the `reduction-identifier` are specified by the `initializer-clause` and the `combiner` in the `declare reduction` directive.

**Description**

A reduction clause specifies a `reduction-identifier` and one or more list items.

The `reduction-identifier` specified in a reduction clause must match a previously declared `reduction-identifier` of the same name and type for each of the list items. This match is done by means of a name lookup in the base language.

The list items that appear in a reduction clause may include array sections.

If the type is a derived class, then any `reduction-identifier` that matches its base classes is also a match, if no specific match for the type has been specified.

If the `reduction-identifier` is not an `id-expression`, then it is implicitly converted to one by prepending the keyword `operator` (for example, `+` becomes `operator+`).

If the `reduction-identifier` is qualified then a qualified name lookup is used to find the declaration.

If the `reduction-identifier` is unqualified then an `argument-dependent name lookup` must be performed using the type of each list item.

If a list item is an array or array section, it will be treated as if a reduction clause would be applied to each separate element of the array section.

If a list item is an array section, the elements of any copy of the array section will be stored contiguously.
If the original list item has the **POINTER** attribute, any copies of the list item are associated with private targets.

Any copies of a list item associated with the reduction are initialized with the initializer value of the reduction-identifier.

Any copies are combined using the combiner associated with the reduction-identifier.

### Execution Model Events

The *reduction-begin* event occurs before a task begins to perform loads and stores that belong to the implementation of a reduction and the *reduction-end* event occurs after the task has completed loads and stores associated with the reduction. If a task participates in multiple reductions, each reduction may be bracketed by its own pair of *reduction-begin/reduction-end* events or multiple reductions may be bracketed by a single pair of events. The interval defined by a pair of *reduction-begin/reduction-end* events may not contain a task scheduling point.

### Tool Callbacks

A thread dispatches a registered `ompt_callback_reduction` with `ompt_sync_region_reduction` in its `kind` argument and `ompt_scope_begin` as its `endpoint` argument for each occurrence of a *reduction-begin* event in that thread. Similarly, a thread dispatches a registered `ompt_callback_reduction` with `ompt_sync_region_reduction` in its `kind` argument and `ompt_scope_end` as its `endpoint` argument for each occurrence of a *reduction-end* event in that thread. These callbacks occur in the context of the task that performs the reduction and has the type signature `ompt_callback_sync_region_t`.

### Restrictions

Restrictions common to reduction clauses are as follows:

- Any number of reduction clauses can be specified on the directive, but a list item (or any array element in an array section) can appear only once in reduction clauses for that directive.
- For a *reduction-identifier* declared in a `declare reduction` directive, the directive must appear before its use in a reduction clause.
- If a list item is an array section or an array element, its base expression must be a base language identifier.
- If a list item is an array section, it must specify contiguous storage and it cannot be a zero-length array section.
- If a list item is an array section or an array element, accesses to the elements of the array outside the specified array section or array element result in unspecified behavior.
• A variable that is part of another variable, with the exception of array elements, cannot appear in a reduction clause.

• A variable that is part of another variable, with the exception of array elements, cannot appear in a reduction clause except if the reduction clause is associated with a construct within a class non-static member function and the variable is an accessible data member of the object for which the non-static member function is invoked.

• The type of a list item that appears in a reduction clause must be valid for the reduction-identifier. For a max or min reduction in C, the type of the list item must be an allowed arithmetic data type: char, int, float, double, or _Bool, possibly modified with long, short, signed, or unsigned. For a max or min reduction in C++, the type of the list item must be an allowed arithmetic data type: char, wchar_t, int, float, double, or bool, possibly modified with long, short, signed, or unsigned.

• A list item that appears in a reduction clause must not be const-qualified.

• The reduction-identifier for any list item must be unambiguous and accessible.

• A variable that is part of another variable, with the exception of array elements, cannot appear in a reduction clause.

• A type parameter inquiry cannot appear in a reduction clause.

• The type, type parameters and rank of a list item that appears in a reduction clause must be valid for the combiner and initializer.

• A list item that appears in a reduction clause must be definable.

• A procedure pointer may not appear in a reduction clause.

• A pointer with the INTENT(IN) attribute may not appear in the reduction clause.

• An original list item with the POINTER attribute or any pointer component of an original list item that is referenced in the combiner must be associated at entry to the construct that contains the reduction clause. Additionally, the list item or the pointer component of the list item must not be deallocated, allocated, or pointer assigned within the region.
• An original list item with the ALLOCATABLE attribute or any allocatable component of an
original list item that corresponds to a special variable identifier in the combiner or the initializer
must be in the allocated state at entry to the construct that contains the reduction clause.
Additionally, the list item or the allocatable component of the list item must be neither
deallocated nor allocated, explicitly or implicitly, within the region.

• If the reduction-identifier is defined in a declare reduction directive, the
declare reduction directive must be in the same subprogram, or accessible by host or use
association.

• If the reduction-identifier is a user-defined operator, the same explicit interface for that operator
must be accessible at the location of the declare reduction directive that defines the
reduction-identifier.

• If the reduction-identifier is defined in a declare reduction directive, any procedure
referred in the initializer clause or combiner expression must be an intrinsic function,
or must have an explicit interface where the same explicit interface is accessible as at the
declare reduction directive.

Cross References
• ompt_scope_begin and ompt_scope_end, see Section 4.4.4.11.
• ompt_sync_region_reduction, see Section 4.4.4.13.
• ompt_callback_sync_region_t, see Section 4.5.2.13.

2.21.5.2 Reduction Scoping Clauses

Reduction scoping clauses define the region in which a reduction is computed by tasks or SIMD
lanes. All properties common to all reduction clauses, which are defined in Section 2.21.5.1, apply
to reduction scoping clauses.

The number of copies created for each list item and the time at which those copies are initialized
are determined by the particular reduction scoping clause that appears on the construct.

The time at which the original list item contains the result of the reduction is determined by the
particular reduction scoping clause.

The location in the OpenMP program at which values are combined and the order in which values
are combined are unspecified. Therefore, when comparing sequential and parallel executions, or
when comparing one parallel execution to another (even if the number of threads used is the same),
bitewise-identical results are not guaranteed to be obtained. Similarly, side effects (such as
floating-point exceptions) may not be identical and may not take place at the same location in the
OpenMP program.

To avoid data races, concurrent reads or updates of the original list item must be synchronized with
the update of the original list item that occurs as a result of the reduction computation.
2.21.5.3 Reduction Participating Clauses

A reduction participating clause specifies a task or a SIMD lane as a participant in a reduction defined by a reduction scoping clause. All properties common to all reduction clauses, which are defined in Section 2.21.5.1, apply to reduction participating clauses.

Accesses to the original list item may be replaced by accesses to copies of the original list item created by a region that corresponds to a construct with a reduction scoping clause.

In any case, the final value of the reduction must be determined as if all tasks or SIMD lanes that participate in the reduction are executed sequentially in some arbitrary order.

2.21.5.4 reduction Clause

Summary
The reduction clause specifies a reduction-identifier and one or more list items. For each list item, a private copy is created for each implicit task or SIMD lane and is initialized with the initializer value of the reduction-identifier. After the end of the region, the original list item is updated with the values of the private copies using the combiner associated with the reduction-identifier.

Syntax
The syntax of the reduction clause is as follows:

```
reduction([ reduction-modifier, ] reduction-identifier : list)
```

Where reduction-identifier is defined in Section 2.21.5.1, and reduction-modifier is one of the following:

- inscan
- task
- default

Description
The reduction clause is a reduction scoping clause and a reduction participating clause, as described in Section 2.21.5.2 and Section 2.21.5.3.

If reduction-modifier is not present or the default reduction-modifier is present, the behavior is as follows. For parallel, scope and worksharing constructs, one or more private copies of each list item are created for each implicit task, as if the private clause had been used. For the simd construct, one or more private copies of each list item are created for each SIMD lane, as if the private clause had been used. For the taskloop construct, private copies are created according to the rules of the reduction scoping clauses. For the teams construct, one or more private copies of each list item are created for the initial task of each team in the league, as if the private clause had been used. For the loop construct, private copies are created and used in the
construct according to the description and restrictions in Section 2.21.3. At the end of a region that

corresponds to a construct for which the reduction clause was specified, the original list item is

updated by combining its original value with the final value of each of the private copies, using the

combiner of the specified reduction-identifier.

If the inscan reduction-modifier is present, a scan computation is performed over updates to the
list item performed in each logical iteration of the loop associated with the worksharing-loop,
worksharing-loop SIMD, or simd construct (see Section 2.11.8). The list items are privatized in
the construct according to the description and restrictions in Section 2.21.3. At the end of the
region, each original list item is assigned the value described in Section 2.11.8.

If the task reduction-modifier is present for a parallel, scope, or worksharing construct,
then each list item is privatized according to the description and restrictions in Section 2.21.3, and
an unspecified number of additional private copies may be created to support task reductions. Any
copies associated with the reduction are initialized before they are accessed by the tasks that
participate in the reduction, which include all implicit tasks in the corresponding region and all
participating explicit tasks that specify an in_reduction clause (see Section 2.21.5.6). After
the end of the region, the original list item contains the result of the reduction.

If nowait is not specified for the construct, the reduction computation will be complete at the end
of the region that corresponds to the construct; however, if the reduction clause is used on a
construct to which nowait is also applied, accesses to the original list item will create a race and,
thus, have unspecified effect unless synchronization ensures that they occur after all threads have
executed all of their iterations or section constructs, and the reduction computation has
completed and stored the computed value of that list item. This can be ensured simply through a

barrier synchronization in most cases.

Restrictions

Restrictions to the reduction clause are as follows:

• All restrictions common to all reduction clauses, which are listed in Section 2.21.5.1, apply to
  this clause.

• A list item that appears in a reduction clause of a worksharing construct must be shared in
  the parallel region to which a corresponding worksharing region binds.

• A list item that appears in a reduction clause of a scope construct must be shared in the
  parallel region to which a corresponding scope region binds.

• If an array section or an array element appears as a list item in a reduction clause of a
  worksharing construct, scope construct or loop construct for which the corresponding region
  binds to a parallel region, all threads that participate in the reduction must specify the same
  storage location.

• A list item that appears in a reduction clause with the inscan reduction-modifier must
  appear as a list item in an inclusive or exclusive clause on a scan directive enclosed by
  the construct.
• A reduction clause without the inscan reduction-modifier may not appear on a construct on which a reduction clause with the inscan reduction-modifier appears.

• A reduction clause with the task reduction-modifier may only appear on a parallel construct, a scope construct, a worksharing construct or a combined or composite construct for which any of the aforementioned constructs is a constituent construct and simd or loop are not constituent constructs.

• A reduction clause with the inscan reduction-modifier may only appear on a worksharing-loop construct, a simd construct or a combined or composite construct for which any of the aforementioned constructs is a constituent construct and distribute is not a constituent construct.

• A list item that appears in a reduction clause of the innermost enclosing worksharing, parallel or scope construct may not be accessed in an explicit task generated by a construct for which an in_reduction clause over the same list item does not appear.

• The task reduction-modifier may not appear in a reduction clause if the nowait clause is specified on the same construct.

C / C++

• If a list item in a reduction clause on a worksharing construct, scope construct or loop construct for which the corresponding region binds to a parallel region has a reference type then it must bind to the same object for all threads of the team.

• If a list item in a reduction clause on a worksharing construct, scope or loop construct for which the corresponding region binds to a parallel region is an array section or an array element then the base pointer must point to the same variable for all threads of the team.

• A variable of class type (or array thereof) that appears in a reduction clause with the inscan reduction-modifier requires an accessible, unambiguous default constructor for the class type. The number of calls to the default constructor while performing the scan computation is unspecified.

• A variable of class type (or array thereof) that appears in a reduction clause with the inscan reduction-modifier requires an accessible, unambiguous copy assignment operator for the class type. The number of calls to the copy assignment operator while performing the scan computation is unspecified.

Cross References

• scan directive, see Section 2.11.8.

• List Item Privatization, see Section 2.21.3.

• private clause, see Section 2.21.4.3.
2.21.5.5 *task_reduction* Clause

**Summary**
The *task_reduction* clause specifies a reduction among tasks.

**Syntax**
The syntax of the *task_reduction* clause is as follows:

```plaintext
task_reduction(reduction-identifier : list)
```

where *reduction-identifier* is defined in Section 2.21.5.1.

**Description**
The *task_reduction* clause is a reduction scoping clause, as described in 2.21.5.2.

For each list item, the number of copies is unspecified. Any copies associated with the reduction are initialized before they are accessed by the tasks that participate in the reduction. After the end of the region, the original list item contains the result of the reduction.

**Restrictions**
Restrictions to the *task_reduction* clause are as follows:

- All restrictions common to all reduction clauses, which are listed in Section 2.21.5.1, apply to this clause.

2.21.5.6 *in_reduction* Clause

**Summary**
The *in_reduction* clause specifies that a task participates in a reduction.

**Syntax**
The syntax of the *in_reduction* clause is as follows:

```plaintext
in_reduction(reduction-identifier : list)
```

where *reduction-identifier* is defined in Section 2.21.5.1.

**Description**
The *in_reduction* clause is a reduction participating clause, as described in Section 2.21.5.3.

For a given list item, the *in_reduction* clause defines a task to be a participant in a task reduction that is defined by an enclosing region for a matching list item that appears in a

*task_reduction* clause or a *reduction* clause with *task* as the *reduction-modifier*, where either:

1. The matching list item has the same storage location as the list item in the *in_reduction*
   clause; or

2. A private copy, derived from the matching list item, that is used to perform the task reduction
   has the same storage location as the list item in the *in_reduction* clause.
For the task construct, the generated task becomes the participating task. For each list item, a private copy may be created as if the private clause had been used.

For the target construct, the target task becomes the participating task. For each list item, a private copy may be created in the data environment of the target task as if the private clause had been used. This private copy will be implicitly mapped into the device data environment of the target device, if the target device is not the parent device.

At the end of the task region, if a private copy was created its value is combined with a copy created by a reduction scoping clause or with the original list item.

**Restrictions**

Restrictions to the in_reduction clause are as follows:

- All restrictions common to all reduction clauses, which are listed in Section 2.21.5.1, apply to this clause.

- A list item that appears in a task_reduction clause or a reduction clause with task as the reduction-modifier that is specified on a construct that corresponds to a region in which the region of the participating task is closely nested must match each list item. The construct that corresponds to the innermost enclosing region that meets this condition must specify the same reduction-identifier for the matching list item as the in_reduction clause.

### 2.21.5.7 declare reduction Directive

**Summary**

The following section describes the directive for declaring user-defined reductions. The declare reduction directive declares a reduction-identifier that can be used in a reduction clause. The declare reduction directive is a declarative directive.

**Syntax**

```c
#pragma omp declare reduction(reduction-identifier : typename-list : combiner ) [initializer-clause] new-line
```

where:

- reduction-identifier is either a base language identifier or one of the following operators: +, -, *, &, |, ^, && or ||

- typename-list is a list of type names

- combiner is an expression

- initializer-clause is initializer(initializer-expr) where initializer-expr is omp_priv = initializer or function-name (argument-list)
The syntax of the `declare reduction` directive is as follows:

```c++
#pragma omp declare reduction(reduction-identifier : typename-list :
combiner) [initializer-clause] new-line
```

where:

- `reduction-identifier` is either an `id-expression` or one of the following operators: `+`, `-`, `*`, `&`, `|`, `^`, `&&` or `||`
- `typename-list` is a list of type names
- `combiner` is an expression
- `initializer-clause` is `initializer(initializer-expr)` where `initializer-expr` is `omp_priv initializer` or `function-name(argument-list)`

The syntax of the `declare reduction` directive is as follows:

```fortran
 !$omp declare reduction(reduction-identifier : type-list : combiner)
 [initializer-clause]
```

where:

- `reduction-identifier` is either a base language identifier, or a user-defined operator, or one of the following operators: `+`, `-`, `*`, `.and.`, `.or.`, `.eqv.`, `.neqv.`, or one of the following intrinsic procedure names: `max`, `min`, `iand`, `i-or`, `ieor`.
- `type-list` is a list of type specifiers that must not be `CLASS(*)` or an abstract type
- `combiner` is either an assignment statement or a subroutine name followed by an argument list
- `initializer-clause` is `initializer(initializer-expr)`, where `initializer-expr` is `omp_priv = expression` or `subroutine-name(argument-list)`

**Description**

User-defined reductions can be defined using the `declare reduction` directive. The `reduction-identifier` and the type identify the `declare reduction` directive. The `reduction-identifier` can later be used in a `reduction` clause that uses variables of the type or types specified in the `declare reduction` directive. If the directive specifies several types then the behavior is as if a `declare reduction` directive was specified for each type.
If a type with deferred or assumed length type parameter is specified in a `declare reduction` directive, the `reduction-identifier` of that directive can be used in a reduction clause with any variable of the same type and the same kind parameter, regardless of the length type Fortran parameters with which the variable is declared.

The visibility and accessibility of a user-defined reduction are the same as those of a variable declared at the same location in the program. The enclosing context of the `combiner` and of the `initializer-expr` is that of the `declare reduction` directive. The `combiner` and the `initializer-expr` must be correct in the base language as if they were the body of a function defined at the same location in the program.

If the `reduction-identifier` is the same as the name of a user-defined operator or an extended operator, or the same as a generic name that is one of the allowed intrinsic procedures, and if the operator or procedure name appears in an accessibility statement in the same module, the accessibility of the corresponding `declare reduction` directive is determined by the accessibility attribute of the statement.

If the `reduction-identifier` is the same as a generic name that is one of the allowed intrinsic procedures and is accessible, and if it has the same name as a derived type in the same module, the accessibility of the corresponding `declare reduction` directive is determined by the accessibility of the generic name according to the base language.

The `declare reduction` directive can also appear at the locations in a program where a static data member could be declared. In this case, the visibility and accessibility of the declaration are the same as those of a static data member declared at the same location in the program.

The `combiner` specifies how partial results are combined into a single value. The `combiner` can use the special variable identifiers `omp_in` and `omp_out` that are of the type of the variables that this `reduction-identifier` reduces. Each of the two special variable identifiers denotes one of the values to be combined before executing the `combiner`. The special `omp_out` identifier refers to the storage that holds the resulting combined value after executing the `combiner`.

The number of times that the `combiner` is executed and the order of these executions for any reduction clause are unspecified.

If the `combiner` is a subroutine name with an argument list, the `combiner` is evaluated by calling the subroutine with the specified argument list.

If the `combiner` is an assignment statement, the `combiner` is evaluated by executing the assignment statement.
If a generic name is used in the *combiner* expression and the list item in the corresponding *reduction* clause is an array or array section, it is resolved to the specific procedure that is elemental or only has scalar dummy arguments.

If the *initializer-expr* value of a user-defined reduction is not known *a priori*, the *initializer-clause* can be used to specify one. The content of the *initializer-clause* will be used as the initializer for the private copies of reduction list items where the *omp_priv* identifier will refer to the storage to be initialized. The special identifier *omp_orig* can also appear in the *initializer-clause* and it will refer to the storage of the original variable to be reduced.

The number of times that the *initializer-expr* is evaluated and the order of these evaluations are unspecified.

If the *initializer-expr* is a function name with an argument list, the *initializer-expr* is evaluated by calling the function with the specified argument list. Otherwise, the *initializer-expr* specifies how *omp_priv* is declared and initialized.

If no *initializer-clause* is specified, the private variables will be initialized following the rules for initialization of objects with static storage duration.

If no *initializer-expr* is specified, the private variables will be initialized following the rules for *default-initialization*.

If the *initializer-expr* is a subroutine name with an argument list, the *initializer-expr* is evaluated by calling the subroutine with the specified argument list.

If the *initializer-expr* is an assignment statement, the *initializer-expr* is evaluated by executing the assignment statement.

If no *initializer-clause* is specified, the private variables will be initialized as follows:

- For *complex*, *real*, or *integer* types, the value 0 will be used.
- For *logical* types, the value *.false.* will be used.
- For derived types for which default initialization is specified, default initialization will be used.
- Otherwise, not specifying an *initializer-clause* results in unspecified behavior.
If `reduction-identifier` is used in a **target** region then a declare target directive must be specified for any function that can be accessed through the `combiner` and `initializer-expr`.

If `reduction-identifier` is used in a **target** region then a **declare target** directive must be specified for any function or subroutine that can be accessed through the `combiner` and `initializer-expr`.

### Restrictions

Restrictions to the **declare reduction** directive are as follows:

- The only variables allowed in the `combiner` are `omp_in` and `omp_out`.
- The only variables allowed in the `initializer-clause` are `omp_priv` and `omp_orig`.
- If the variable `omp_orig` is modified in the `initializer-clause`, the behavior is unspecified.
- If execution of the `combiner` or the `initializer-expr` results in the execution of an OpenMP construct or an OpenMP API call, then the behavior is unspecified.
- A `reduction-identifier` may not be re-declared in the current scope for the same type or for a type that is compatible according to the base language rules.
- At most one `initializer-clause` can be specified.
- The `typename-list` must not declare new types.
- A type name in a **declare reduction** directive cannot be a function type, an array type, a reference type, or a type qualified with `const`, `volatile` or `restrict`.
- If the `initializer-expr` is a function name with an argument list, then one of the arguments must be the address of `omp_priv`.
- If the `initializer-expr` is a function name with an argument list, then one of the arguments must be `omp_priv` or the address of `omp_priv`. 
Any selectors in the designator of `omp_in` and `omp_out` can only be `component selectors`

- If the `initializer-expr` is a subroutine name with an argument list, then one of the arguments must be `omp_priv`.

- Any subroutine or function used in the `initializer` clause or `combiner` expression must be an intrinsic function, or must have an accessible interface.

- Any user-defined operator, defined assignment or extended operator used in the `initializer` clause or `combiner` expression must have an accessible interface.

- If any subroutine, function, user-defined operator, defined assignment or extended operator is used in the `initializer` clause or `combiner` expression, it must be accessible to the subprogram in which the corresponding `reduction` clause is specified.

- If the length type parameter is specified for a type, it must be a constant, a colon or an `*`.

- If a type with deferred or assumed length parameter is specified in a `declare reduction` directive, no other `declare reduction` directive with the same type, the same kind parameters and the same `reduction-identifier` is allowed in the same scope.

- Any subroutine used in the `initializer` clause or `combiner` expression must not have any alternate returns appear in the argument list.

- If the list item in the corresponding `reduction` clause is an array or array section, then any procedure used in the `initializer` clause or `combiner` expression must either be elemental or have dummy arguments that are scalar.

- Any procedure called in the region of `initializer-expr` or `combiner` must be pure and may not reference any host-associated variables.

Cross References

- Properties Common to All Reduction Clauses, see Section 2.21.5.1.

2.21.6 Data Copying Clauses

This section describes the `copyin` clause (allowed on the `parallel` construct and combined parallel worksharing constructs) and the `copyprivate` clause (allowed on the `single` construct).

These two clauses support copying data values from private or threadprivate variables of an implicit task or thread to the corresponding variables of other implicit tasks or threads in the team.

The clauses accept a comma-separated list of list items (see Section 2.1). All list items appearing in a clause must be visible, according to the scoping rules of the base language. Clauses may be repeated as needed, but a list item that specifies a given variable may not appear in more than one clause on the same directive.
### 2.21.6.1 copyin Clause

#### Summary

The **copyin** clause provides a mechanism to copy the value of a threadprivate variable of the primary thread to the threadprivate variable of each other member of the team that is executing the **parallel** region.

#### Syntax

The syntax of the **copyin** clause is as follows:

```
copyin(list)
```

#### Description

**C / C++**

The copy is performed after the team is formed and prior to the execution of the associated structured block. For variables of non-array type, the copy is by copy assignment. For an array of elements of non-array type, each element is copied as if by assignment from an element of the array of the primary thread to the corresponding element of the array of all other threads.

**C++**

For class types, the copy assignment operator is invoked. The order in which copy assignment operators for different variables of the same class type are invoked is unspecified.

**Fortran**

The copy is performed, as if by assignment, after the team is formed and prior to the execution of the associated structured block.

On entry to any **parallel** region, each thread’s copy of a variable that is affected by a **copyin** clause for the **parallel** region will acquire the type parameters, allocation, association, and definition status of the copy of the primary thread, according to the following rules:

- If the original list item has the **POINTER** attribute, each copy receives the same association status as that of the copy of the primary thread as if by pointer assignment.

- If the original list item does not have the **POINTER** attribute, each copy becomes defined with the value of the copy of the primary thread as if by intrinsic assignment unless the list item has a type bound procedure as a defined assignment. If the original list item that does not have the **POINTER** attribute has the allocation status of unallocated, each copy will have the same status.

- If the original list item is unallocated or unassociated, each copy inherits the declared type parameters and the default type parameter values from the original list item.
Restrictions
Restrictions to the \texttt{copyin} clause are as follows:

\begin{itemize}
  \item A list item that appears in a \texttt{copyin} clause must be threadprivate.
  \item A variable of class type (or array thereof) that appears in a \texttt{copyin} clause requires an accessible, unambiguous copy assignment operator for the class type.
\end{itemize}

Cross References
\begin{itemize}
  \item parallel construct, see Section 2.6.
  \item threadprivate directive, see Section 2.21.2.
\end{itemize}

2.21.6.2 copyprivate Clause
Summary
The \texttt{copyprivate} clause provides a mechanism to use a private variable to broadcast a value from the data environment of one implicit task to the data environments of the other implicit tasks that belong to the \texttt{parallel} region.

To avoid data races, concurrent reads or updates of the list item must be synchronized with the update of the list item that occurs as a result of the \texttt{copyprivate} clause.

Syntax
The syntax of the \texttt{copyprivate} clause is as follows:

\begin{verbatim}
copyprivate(list)
\end{verbatim}
Description
The effect of the `copyprivate` clause on the specified list items occurs after the execution of the
structured block associated with the `single` construct (see Section 2.10.2), and before any of the
threads in the team have left the barrier at the end of the construct.

\[
\text{C / C++}
\]
In all other implicit tasks that belong to the `parallel` region, each specified list item becomes
defined with the value of the corresponding list item in the implicit task associated with the thread
that executed the structured block. For variables of non-array type, the definition occurs by copy
assignment. For an array of elements of non-array type, each element is copied by copy assignment
from an element of the array in the data environment of the implicit task that is associated with the
thread that executed the structured block to the corresponding element of the array in the data
environment of the other implicit tasks.

\[
\text{C / C++}
\]
\[
\text{C++}
\]
For class types, a copy assignment operator is invoked. The order in which copy assignment
operators for different variables of class type are called is unspecified.

\[
\text{C++}
\]
\[
\text{Fortran}
\]
If a list item does not have the `POINTER` attribute, then in all other implicit tasks that belong to the
`parallel` region, the list item becomes defined as if by intrinsic assignment with the value of the
corresponding list item in the implicit task that is associated with the thread that executed the
structured block. If the list item has a type bound procedure as a defined assignment, the
assignment is performed by the defined assignment.

\[
\text{Fortran}
\]
If the list item has the `POINTER` attribute, then, in all other implicit tasks that belong to the
`parallel` region, the list item receives, as if by pointer assignment, the same association status of
the corresponding list item in the implicit task that is associated with the thread that executed the
structured block.

The order in which any final subroutines for different variables of a finalizable type are called is
unspecified.

\[
\text{Fortran}
\]

Note – The `copyprivate` clause is an alternative to using a shared variable for the value when
providing such a shared variable would be difficult (for example, in a recursion requiring a different
variable at each level).
Restrictions
Restrictions to the `copyprivate` clause are as follows:

• All list items that appear in the `copyprivate` clause must be either threadprivate or private in the enclosing context.

• A list item that appears in a `copyprivate` clause may not appear in a `private` or `firstprivate` clause on the `single` construct.

• A variable of class type (or array thereof) that appears in a `copyprivate` clause requires an accessible unambiguous copy assignment operator for the class type.

• A common block that appears in a `copyprivate` clause must be threadprivate.

• Pointers with the `INTENT(IN)` attribute may not appear in the `copyprivate` clause.

• Any list item with the `ALLOCATABLE` attribute must have the allocation status of allocated when the intrinsic assignment is performed.

• If a list item is a polymorphic variable with the `ALLOCATABLE` attribute, the behavior is unspecified.

Cross References
• `parallel` construct, see Section 2.6.
• `threadprivate` directive, see Section 2.21.2.
• `private` clause, see Section 2.21.4.3.

2.21.7 Data-Mapping Attribute Rules, Clauses, and Directives
This section describes how the data-mapping and data-sharing attributes of any variable referenced in a `target` region are determined. When specified, explicit data-sharing attribute, `map`, `is_device_ptr` or `has_device_addr` clauses on `target` directives determine these attributes. Otherwise, the first matching rule from the following implicit data-mapping rules applies for variables referenced in a `target` construct that are not declared in the construct and do not appear in one of the data-sharing attribute, `map, is_device_ptr` or `has_device_addr` clauses.

• If a variable appears in a `to` or `link` clause on a declare target directive that does not have a `device_type(nohost)` clause then it is treated as if it had appeared in a `map` clause with a `map-type` of `tofrom`. 
• If a list item appears in a `reduction`, `lastprivate` or `linear` clause on a combined target construct then it is treated as if it also appears in a `map` clause with a `map-type` of `tofrom`.

• If a list item appears in an `in_reduction` clause on a `target` construct then it is treated as if it also appears in a `map` clause with a `map-type` of `tofrom` and a `map-type-modifier` of `always`.

• If a `defaultmap` clause is present for the category of the variable and specifies an implicit behavior other than `default`, the data-mapping attribute is determined by that clause.

• If the `target` construct is within a class non-static member function, and a variable is an accessible data member of the object for which the non-static data member function is invoked, the variable is treated as if the `this[:1]` expression had appeared in a `map` clause with a `map-type` of `tofrom`. Additionally, if the variable is of type pointer or reference to pointer, it is also treated as if it had appeared in a `map` clause as a zero-length array section.

• If the `this` keyword is referenced inside a `target` construct within a class non-static member function, it is treated as if the `this[:1]` expression had appeared in a `map` clause with a `map-type` of `tofrom`.

- A variable that is of type pointer, but not a function pointer or (for C++) a pointer to a member function, is treated as if it is the base pointer of a zero-length array section that had appeared as a list item in a `map` clause.

- A variable that is of type reference to pointer, but not a function pointer or a reference to a pointer to a member function is treated as if it had appeared in a `map` clause as a zero-length array section.

• If a variable is not a scalar then it is treated as if it had appeared in a `map` clause with a `map-type` of `tofrom`.

- If a scalar variable has the `TARGET`, `ALLOCATABLE` or `POINTER` attribute then it is treated as if it had appeared in a `map` clause with a `map-type` of `tofrom`.

- If none of the above rules applies then a scalar variable is not mapped, but instead has an implicit data-sharing attribute of `firstprivate` (see Section 2.21.1.1).
2.21.7.1 map Clause

Summary
The map clause specifies how an original list item is mapped from the current task’s data environment to a corresponding list item in the device data environment of the device identified by the construct.

Syntax
The syntax of the map clause is as follows:

```plaintext
map (map-type: locator-list)
```

where map-type is one of the following:

- `to`
- `from`
- `tofrom`
- `alloc`
- `release`
- `delete`

and map-type-modifier is one of the following:

- `always`
- `close`
- `mapper(mapper-identifier)`
- `present`
- `iterator(iterators-definition)`

Description
The list items that appear in a map clause may include array sections and structure elements.

The map-type and map-type-modifier specify the effect of the map clause, as described below.

For a given construct, the effect of a map clause with the to, from, or tofrom map-type is ordered before the effect of a map clause with the alloc, release, or delete map-type. If a map clause with a present map-type-modifier appears in a map clause, then the effect of the clause is ordered before all other map clauses that do not have the present modifier.

If the mapper map-type-modifier is not present, the behavior is as if the mapper(default) modifier was specified. The map behavior of a list item in a map clause is modified by a visible user-defined mapper (see Section 2.21.7.4) if the mapper has the same mapper-identifier as the mapper-identifier in the mapper map-type-modifier and is specified for a type that matches the type of the list item. The effect of the mapper is to remove the list item from the map clause, if the
present modifier does not also appear, and to apply the clauses specified in the declared mapper to the construct on which the map clause appears. In the clauses applied by the mapper, references to var are replaced with references to the list item and the map-type is replaced with a final map type that is determined according to the rules of map-type decay (see Section 2.21.7.4).

A list item that is an array or array section of a type for which a user-defined mapper exists is mapped as if the map type decays to alloc, release, or delete, and then each array element is mapped with the original map type, as if by a separate construct, according to the mapper.

A list item in a map clause may reference iterators defined by an iterators-definition of an iterator modifier.

If a list item in a map clause is a variable of structure type then it is treated as if each structure element contained in the variable is a list item in the clause.

If a list item in a map clause is a structure element then all other structure elements of the containing structure variable form a structure sibling list. The map clause and the structure sibling list are associated with the same construct. If a corresponding list item of the structure sibling list item is present in the device data environment when the construct is encountered then:

- If the structure sibling list item does not appear in a map clause on the construct then:
  - If the construct is a target, target data, or target enter data construct then the structure sibling list item is treated as if it is a list item in a map clause on the construct with a map-type of alloc.
  - If the construct is target exit data construct, then the structure sibling list item is treated as if it is a list item in a map clause on the construct with a map-type of release.

- If the map clause in which the structure element appears as a list item has a map-type of delete and the structure sibling list item does not appear as a list item in a map clause on the construct with a map-type of delete then the structure sibling list item is treated as if it is a list item in a map clause on the construct with a map-type of delete.

Fortran

If a component of a derived type variable is a list item that results from the above rules for mapped structures and mapped structure elements, and it does not explicitly appear as another list item or as the base expression of another list item in a map clause on the construct, then:

- If it has the POINTER attribute, the map clause treats its association status as if it is undefined; and

- If it has the ALLOCATABLE attribute and an allocated allocation status, and it is present in the device data environment when the construct is encountered, the map clause may treat its allocation status as if it is unallocated if the corresponding component does not have allocated storage.
Given \textit{item}_1 is a list item in a \texttt{map} clause, and \textit{item}_2 is another list item in a \texttt{map} clause on the same construct, if \textit{item}_2 has a base pointer that is, or is part of, \textit{item}_1, then:

- If the construct is a \texttt{target}, \texttt{target data}, or \texttt{target enter data} construct, then, on entry to the corresponding region, the effect of the \texttt{map} clause on \textit{item}_1 is ordered to occur before the effect of the \texttt{map} clause on \textit{item}_2.

- If the construct is a \texttt{target}, \texttt{target data}, or \texttt{target exit data} construct, then, on exit from the corresponding region, the effect of the \texttt{map} clause on \textit{item}_2 is ordered to occur before the effect of the \texttt{map} clause on \textit{item}_1.

If a list item in a \texttt{map} clause is an associated pointer and the pointer is not the base pointer of another list item in a \texttt{map} clause on the same construct, then it is treated as if its pointer target is implicitly mapped in the same clause. For the purposes of the \texttt{map} clause, the mapped pointer target is treated as if its base pointer is the associated pointer.

If a list item in a \texttt{map} clause has a base pointer, a pointer variable is present in the device data environment that corresponds to the base pointer when the effect of the \texttt{map} clause occurs, and the corresponding pointer or the corresponding list item is created in the device data environment on entry to the construct, then:

1. The corresponding pointer variable is assigned an address such that the corresponding list item can be accessed through the pointer in a \texttt{target} region.

1. The corresponding pointer variable is associated with a pointer target that has the same rank and bounds as the pointer target of the original pointer, such that the corresponding list item can be accessed through the pointer in a \texttt{target} region.

2. The corresponding pointer variable becomes an attached pointer for the corresponding list item.

3. If the original base pointer and the corresponding attached pointer share storage, then the original list item and the corresponding list item must share storage.

If a \texttt{lambda} is mapped explicitly or implicitly, variables that are captured by the \texttt{lambda} behave as follows:

- The variables that are of pointer type are treated as if they had appeared in a \texttt{map} clause as zero-length array sections; and

- The variables that are of reference type are treated as if they had appeared in a \texttt{map} clause.
If a member variable is captured by a lambda in class scope, and the lambda is later mapped explicitly or implicitly with its full static type, the this pointer is treated as if it had appeared on a map clause.

The original and corresponding list items may share storage such that writes to either item by one task followed by a read or write of the other item by another task without intervening synchronization can result in data races. They are guaranteed to share storage if the map clause appears on a target construct that corresponds to an inactive target region, or if it appears on a target data, target enter data, or target exit data construct that applies to the device data environment of the host device.

If a map clause appears on a target, target data, target enter data or target exit data construct with a present map-type-modifier then on entry to the region if the corresponding list item does not appear in the device data environment then an error occurs and the program terminates.

If a map clause appears on a target, target data, or target enter data construct then on entry to the region the following sequence of steps occurs as if they are performed as a single atomic operation:

1. If a corresponding list item of the original list item is not present in the device data environment, then:
   a) A new list item with language-specific attributes is derived from the original list item and created in the device data environment;
   b) The new list item becomes the corresponding list item of the original list item in the device data environment;
   c) The corresponding list item has a reference count that is initialized to zero; and
   d) The value of the corresponding list item is undefined;

2. If the reference count of the corresponding list item was not already incremented because of the effect of a map clause on the construct then:
   a) The reference count is incremented by one;

3. If the reference count of the corresponding list item is one or the always map-type-modifier is present, and if the map-type is to or tofrom, then:
   a) For each part of the list item that is an attached pointer, that part of the corresponding list item will have the value that it had at the point immediately prior to the effect of the map clause; and
a) For each part of the list item that is an attached pointer, that part of the corresponding list item, if associated, will be associated with the same pointer target that it was associated with at the point immediately prior to the effect of the \texttt{map} clause.

b) For each part of the list item that is not an attached pointer, the value of that part of the original list item is assigned to that part of the corresponding list item.

\textbf{Note} – If the effect of the \texttt{map} clauses on a construct would assign the value of an original list item to a corresponding list item more than once, then an implementation is allowed to ignore additional assignments of the same value to the corresponding list item.

In all cases on entry to the region, concurrent reads or updates of any part of the corresponding list item must be synchronized with any update of the corresponding list item that occurs as a result of the \texttt{map} clause to avoid data races.

If the \texttt{map} clause appears on a \texttt{target}, \texttt{target data}, or \texttt{target exit data} construct and a corresponding list item of the original list item is not present in the device data environment on exit from the region then the list item is ignored. Alternatively, if the \texttt{map} clause appears on a \texttt{target}, \texttt{target data}, or \texttt{target exit data} construct and a corresponding list item of the original list item is present in the device data environment on exit from the region, then the following sequence of steps occurs as if performed as a single atomic operation:

1. If the \texttt{map-type} is not \texttt{delete} and the reference count of the corresponding list item is finite and was not already decremented because of the effect of a \texttt{map} clause on the construct then:
   a) The reference count of the corresponding list item is decremented by one;

2. If the \texttt{map-type} is \texttt{delete} and the reference count of the corresponding list item is finite then:
   a) The reference count of the corresponding list item is set to zero;

3. If the \texttt{map-type} is \texttt{from} or \texttt{tofrom} and if the reference count of the corresponding list item is zero or the \texttt{always map-type-modifier} is present then:
   a) For each part of the list item that is an attached pointer, that part of the original list item will have the value that it had at the point immediately prior to the effect of the \texttt{map} clause; and

   a) For each part of the list item that is an attached pointer, that part of the original list item, if associated, will be associated with the same pointer target with which it was associated at the point immediately prior to the effect of the \texttt{map} clause; and
b) For each part of the list item that is not an attached pointer, the value of that part of the corresponding list item is assigned to that part of the original list item; and

4. If the reference count of the corresponding list item is zero then the corresponding list item is removed from the device data environment.

\begin{itemize}
\item Note – If the effect of the map clauses on a construct would assign the value of a corresponding list item to an original list item more than once, then an implementation is allowed to ignore additional assignments of the same value to the original list item.
\end{itemize}

In all cases on exit from the region, concurrent reads or updates of any part of the original list item must be synchronized with any update of the original list item that occurs as a result of the map clause to avoid data races.

If a single contiguous part of the original storage of a list item with an implicit data-mapping attribute has corresponding storage in the device data environment prior to a task encountering the construct that is associated with the map clause, only that part of the original storage will have corresponding storage in the device data environment as a result of the map clause.

If a list item with an implicit data-mapping attribute does not have any corresponding storage in the device data environment prior to a task encountering the construct associated with the map clause, and one or more contiguous parts of the original storage are either list items or base pointers to list items that are explicitly mapped on the construct, only those parts of the original storage will have corresponding storage in the device data environment as a result of the map clauses on the construct.

\begin{itemize}
\item C / C++
\item Fortran
\end{itemize}

23 If a new list item is created then a new list item of the same type, with automatic storage duration, is allocated for the construct. The size and alignment of the new list item are determined by the static type of the variable. This allocation occurs if the region references the list item in any statement. Initialization and assignment of the new list item are through bitwise copy.

\begin{itemize}
\item C / C++
\item Fortran
\end{itemize}

26 If a new list item is created then a new list item of the same type, type parameter, and rank is allocated. The new list item inherits all default values for the type parameters from the original list item. The value of the new list item becomes that of the original list item in the map initialization and assignment.

If the allocation status of an original list item that has the ALLOCATABLE attribute is changed while a corresponding list item is present in the device data environment, the allocation status of the corresponding list item is unspecified until the list item is again mapped with an always modifier on entry to a target, target data or target enter data region.

\begin{itemize}
\item Fortran
\end{itemize}
The map-type determines how the new list item is initialized.

If a map-type is not specified, the map-type defaults to tofrom.

The close map-type-modifier is a hint to the runtime to allocate memory close to the target device.

### Execution Model Events

The target-map event occurs when a thread maps data to or from a target device.

The target-data-op-begin event occurs before a thread initiates a data operation on a target device.

The target-data-op-end event occurs after a thread initiates a data operation on a target device.

### Tool Callbacks

A thread dispatches a registered ompt_callback_target_map or ompt_callback_target_map_emi callback for each occurrence of a target-map event in that thread. The callback occurs in the context of the target task and has type signature ompt_callback_target_map_t or ompt_callback_target_map_emi_t, respectively.

A thread dispatches a registered ompt_callback_target_data_op_emi callback with ompt_scope_begin as its endpoint argument for each occurrence of a target-data-op-begin event in that thread. Similarly, a thread dispatches a registered ompt_callback_target_data_op_emi callback with ompt_scope_end as its endpoint argument for each occurrence of a target-data-op-end event in that thread. These callbacks have type signature ompt_callback_target_data_op_emi_t.

A thread dispatches a registered ompt_callback_target_data_op callback for each occurrence of a target-data-op-end event in that thread. The callback occurs in the context of the target task and has type signature ompt_callback_target_data_op_t.

### Restrictions

Restrictions to the map clause are as follows:

- Each of the map-type-modifier modifiers can appear at most once in the map clause.
- List items of the map clauses on the same construct must not share original storage unless they are the same lvalue expression or array section.
- If a list item is an array section, it must specify contiguous storage.
- If an expression that is used to form a list item in a map clause contains an iterator identifier, the list item instances that would result from different values of the iterator must not have the same containing array and must not have base pointers that share original storage.
• If multiple list items are explicitly mapped on the same construct and have the same containing array or have base pointers that share original storage, and if any of the list items do not have corresponding list items that are present in the device data environment prior to a task encountering the construct, then the list items must refer to the same array elements of either the containing array or the implicit array of the base pointers.

• If any part of the original storage of a list item with an explicit data-mapping attribute has corresponding storage in the device data environment prior to a task encountering the construct associated with the map clause, all of the original storage must have corresponding storage in the device data environment prior to the task encountering the construct.

• If an array appears as a list item in a map clause, multiple parts of the array have corresponding storage in the device data environment prior to a task encountering the construct associated with the map clause, and the corresponding storage for those parts was created by maps from more than one earlier construct, the behavior is unspecified.

• If a list item is an element of a structure, and a different element of the structure has a corresponding list item in the device data environment prior to a task encountering the construct associated with the map clause, then the list item must also have a corresponding list item in the device data environment prior to the task encountering the construct.

• A list item must have a mappable type.

• Threadprivate variables cannot appear in a map clause.

• If a mapper map-type-modifier appears in a map clause, the type on which the specified mapper operates must match the type of the list items in the clause.

• Memory spaces and memory allocators cannot appear as a list item in a map clause.

• If the type of a list item is a reference to a type $T$ then the reference in the device data environment is initialized to refer to the object in the device data environment that corresponds to the object referenced by the list item. If mapping occurs, it occurs as though the object were mapped through a pointer with an array section of type $T$ and length one.

• No type mapped through a reference can contain a reference to its own type, or any references to types that could produce a cycle of references.

• If a list item is a lambda, any pointers and references captured by the lambda must have the corresponding list item in the device data environment prior to the task encountering the construct.
- A list item cannot be a variable that is a member of a structure of a union type.
- A bit-field cannot appear in a `map` clause.
- A pointer that has a corresponding attached pointer must not be modified for the duration of the lifetime of the list item to which the corresponding pointer is attached in the device data environment.

- List items of the `map` clauses on the same construct must not share original storage unless they are the same variable or array section.
- If a list item of a `map` clause is an allocatable variable or is the subobject of an allocatable variable, the original allocatable variable may not be allocated, deallocated or reshaped while the corresponding allocatable variable has allocated storage.
- A pointer that has a corresponding attached pointer and is associated with a given pointer target must not become associated with a different pointer target for the duration of the lifetime of the list item to which the corresponding pointer is attached in the device data environment.
- If an array section is mapped and the size of the section is smaller than that of the whole array, the behavior of referencing the whole array in the `target` region is unspecified.
- A list item must not be a whole array of an assumed-size array.

Cross References
- Array sections, see Section 2.1.5.
- Iterators, see Section 2.1.6.
- `declare mapper` directive, see Section 2.21.7.4.
- `ompt_callback_target_data_op_t` or `ompt_callback_target_data_op_emi_t` callback type, see Section 4.5.2.25.
- `ompt_callback_target_map_t` or `ompt_callback_target_map_emi_t` callback type, see Section 4.5.2.27.
2.21.7.2 Pointer Initialization for Device Data Environments

This section describes how a pointer that is predetermined firstprivate for a target construct may
be assigned an initial value that is the address of an object that exists in a device data environment
and corresponds to a matching mapped list item.

All previously mapped list items that have corresponding storage in a given device data
environment constitute the set of currently mapped list items. If a currently mapped list item has a
base pointer, the base address of the currently mapped list item is the value of its base pointer.
Otherwise, the base address is determined by the following steps:

1. Let $X$ refer to the currently mapped list item.
2. If $X$ refers to an array section or array element, let $X$ refer to its base array.
3. If $X$ refers to a structure element, let $X$ refer to its containing structure and return to step 2.
4. The base address for the currently mapped list item is the address of $X$.

Additionally, each currently mapped list item has a starting address and an ending address. The
starting address is the address of the first storage location associated with the list item, and the
ending address is the address of the storage location that immediately follows the last storage
location associated with the list item.

The mapped address range of the currently mapped list item is the range of addresses that starts
from the starting address and ends with the ending address. The extended address range of the
currently mapped list item is the range of addresses that starts from the minimum of the starting
address and the base address and that ends with the maximum of the ending address and the base
address.

If the value of a given pointer is in the mapped address range of a currently mapped list item then
that currently mapped list item is a matching mapped list item. Otherwise, if the value of the
pointer is in the extended address range of a currently mapped list item then that currently mapped
list item is a matching mapped list item.

If multiple matching mapped list items are found and they all appear as part of the same containing
structure, the one that has the lowest starting address is treated as the sole matching mapped list
item. Otherwise, if multiple matching mapped list items are found then the behavior is unspecified.

If a matching mapped list item is found, the initial value that is assigned to the pointer is a device
address such that the corresponding list item in the device data environment can be accessed
through the pointer in a target region.

If a matching mapped list item is not found, the assigned initial value of the pointer is NULL unless
otherwise specified (see Section 2.5.1).
Cross References

- **requires** directive, see Section 2.5.1.
- **target** construct, see Section 2.14.5.
- **map** clause, see Section 2.21.7.1.

2.21.7.3 *defaultmap* Clause

Summary

The *defaultmap* clause explicitly determines the data-mapping attributes of variables that are referenced in a *target* construct for which the data-mapping attributes would otherwise be implicitly determined (see Section 2.21.7).

Syntax

The syntax of the *defaultmap* clause is as follows:

```
defaultmap(implicit-behavior[:variable-category])
```

Where *implicit-behavior* is one of:

- `alloc`
- `to`
- `from`
- `tofrom`
- `firstprivate`
- `none`
- `default`
- `present`

and *variable-category* is one of:

- `scalar`
- `aggregate`
- `pointer`
and variable-category is one of:

- scalar
- aggregate
- allocatable
- pointer

**Description**

The **defaultmap** clause sets the implicit data-mapping attribute for all variables referenced in the construct. If variable-category is specified, the effect of the **defaultmap** clause is as follows:

- If variable-category is **scalar**, all scalar variables of non-pointer type or all non-pointer non-allocatable scalar variables that have an implicitly determined data-mapping or data-sharing attribute will have a data-mapping or data-sharing attribute specified by implicit-behavior.

- If variable-category is **aggregate** or **allocatable**, all aggregate or allocatable variables that have an implicitly determined data-mapping or data-sharing attribute will have a data-mapping or data-sharing attribute specified by implicit-behavior.

- If variable-category is **pointer**, all variables of pointer type or with the **POINTER** attribute that have implicitly determined data-mapping or data-sharing attributes will have a data-mapping or data-sharing attribute specified by implicit-behavior.

If no variable-category is specified in the clause then implicit-behavior specifies the implicitly determined data-mapping or data-sharing attribute for all variables referenced in the construct. If implicit-behavior is **none**, each variable referenced in the construct that does not have a predetermined data-sharing attribute and does not appear in a to or link clause on a declare target directive must be listed in a data-mapping attribute clause, a data-sharing attribute clause (including a data-sharing attribute clause on a combined construct where target is one of the constituent constructs), an **is_device_ptr** clause or a **has_device_addr** clause. If implicit-behavior is **default**, then the clause has no effect for the variables in the category specified by variable-category. If implicit-behavior is **present**, each variable referenced in the construct in the category specified by variable-category is treated as if it had been listed in a map clause with the map-type of alloc and map-type-modifier of present.

### 2.21.7.4 declare mapper Directive

**Summary**

The **declare mapper** directive declares a user-defined mapper for a given type, and may define a mapper-identifier that can be used in a map clause. The **declare mapper** directive is a declarative directive.
Syntax

C / C++

The syntax of the `declare mapper` directive is as follows:

```c
#pragma omp declare mapper([mapper-identifier]:type var) \    
clause[,clause] ... new-line
```

Fortran

The syntax of the `declare mapper` directive is as follows:

```fortran
!$omp declare mapper([mapper-identifier]:type :: var) &    
clause[,clause] ...
```

where:

- `mapper-identifier` is a base language identifier or `default`
- `type` is a valid type in scope
- `var` is a valid base language identifier
- `clause` is `map([map-type-modifier[,map-type-modifier][,...]] map-type: list)`, where
  - `map-type` is one of the following:
    - `alloc`
    - `to`
    - `from`
    - `tofrom`
  - and where `map-type-modifier` is one of the following:
    - `always`
    - `close`

Description

User-defined mappers can be defined using the `declare mapper` directive. The `type` and an optional `mapper-identifier` uniquely identify the mapper for use in a `map` clause or motion clause later in the program. The visibility and accessibility of this declaration are the same as those of a variable declared at the same location in the program.

If `mapper-identifier` is not specified, the behavior is as if `mapper-identifier` is `default`.

The variable declared by `var` is available for use in all `map` clauses on the directive, and no part of the variable to be mapped is mapped by default.
The default mapper for all types $T$, designated by the predefined mapper-identifier \texttt{default}, is defined as if by the following directive, unless a user-defined mapper is specified for that type.

\begin{verbatim}
#pragma omp declare mapper(T v) map(tofrom: v)
\end{verbatim}

A \texttt{declare mapper} directive that uses the \texttt{default} mapper-identifier overrides the predefined default mapper for the given type, making it the default mapper for variables of that type.

The effect that a user-defined mapper has on either a \texttt{map} clause that maps a list item of the given type or a motion clause that invokes the mapper and updates a list item of the given type is to replace the map or update with a set of \texttt{map} clauses or updates derived from the \texttt{map} clauses specified by the mapper, as described in Section 2.21.7.1 and Section 2.14.6.

The final map types that a mapper applies for a \texttt{map} clause that maps a list item of the given type are determined according to the rules of map-type decay, defined according to Table 2.13.

Table 2.13 shows the final map type that is determined by the combination of two map types, where the rows represent the map type specified by the mapper and the columns represent the map type specified by a \texttt{map} clause that invokes the mapper. For a \texttt{target exit data} construct that invokes a mapper with a \texttt{map} clause that has the \texttt{from} map type, if a \texttt{map} clause in the mapper specifies an \texttt{alloc} or \texttt{to} map type then the result is a \texttt{release} map type.

A list item in a \texttt{map} clause that appears on a \texttt{declare mapper} directive may include array sections.

All \texttt{map} clauses that are introduced by a mapper are further subject to mappers that are in scope, except a \texttt{map} clause with list item \texttt{var} maps \texttt{var} without invoking a mapper.

\begin{table}[h]
\centering
\caption{Map-Type Decay of Map Type Combinations}
\begin{tabular}{|c|c|c|c|c|c|}
\hline
 & alloc & to & from & tofrom & release & delete \\
\hline
 alloc & alloc & alloc & alloc (release) & alloc & release & delete \\
 to & alloc & to & alloc (release) & to & release & delete \\
 from & alloc & alloc & from & from & release & delete \\
 tofrom & alloc & to & from & tofrom & release & delete \\
\hline
\end{tabular}
\end{table}
The `declare mapper` directive can also appear at locations in the program at which a static data member could be declared. In this case, the visibility and accessibility of the declaration are the same as those of a static data member declared at the same location in the program.

**Restrictions**

Restrictions to the `declare mapper` directive are as follows:

- No instance of type can be mapped as part of the mapper, either directly or indirectly through another type, except the instance `var` that is passed as the list item. If a set of `declare mapper` directives results in a cyclic definition then the behavior is unspecified.
- The type must not declare a new type.
- At least one `map` clause that maps `var` or at least one element of `var` is required.
- List items in `map` clauses on the `declare mapper` directive may only refer to the declared variable `var` and entities that could be referenced by a procedure defined at the same location.
- Each `map-type-modifier` can appear at most once on the `map` clause.
- Multiple `declare mapper` directives that specify the same `mapper-identifier` for the same type or for compatible types, according to the base language rules, may not appear in the same scope.
- `type` must be a `struct` or `union` type.
- `type` must be a `struct`, `union`, or `class` type.
- `type` must not be an intrinsic type or an abstract type.

**Cross References**

- `target update` construct, see Section 2.14.6.
- `map` clause, see Section 2.21.7.1.
2.22 Nesting of Regions

This section describes a set of restrictions on the nesting of regions. The restrictions on nesting are as follows:

- A loop region that binds to a parallel region or a worksharing region may not be closely nested inside a worksharing, loop, task, taskloop, critical, ordered, atomic, or masked region.
- A barrier region may not be closely nested inside a worksharing, loop, task, taskloop, critical, ordered, atomic, or masked region.
- A masked region may not be closely nested inside a worksharing, loop, atomic, task, or taskloop region.
- An ordered region that corresponds to an ordered construct without any clause or with the threads or depend clause may not be closely nested inside a critical, ordered, loop, atomic, task, or taskloop region.
- An ordered region that corresponds to an ordered construct without the simd clause specified must be closely nested inside a worksharing-loop region.
- An ordered region that corresponds to an ordered construct with the simd clause specified must be closely nested inside a simd or worksharing-loop SIMD region.
- An ordered region that corresponds to an ordered construct with both the simd and threads clauses must be closely nested inside a worksharing-loop SIMD region or closely nested inside a worksharing-loop and simd region.
- A critical region may not be nested (closely or otherwise) inside a critical region with the same name. This restriction is not sufficient to prevent deadlock.
- OpenMP constructs may not be encountered during execution of an atomic region.
- The only OpenMP constructs that can be encountered during execution of a simd (or worksharing-loop SIMD) region are the atomic construct, the loop construct without a defined binding region, the simd construct and the ordered construct with the simd clause.
- If a target update, target data, target enter data, or target exit data construct is encountered during execution of a target region, the behavior is unspecified.
- If a target construct is encountered during execution of a target region and a device clause in which the ancestor device-modifier appears is not present on the construct, the behavior is unspecified.
- A teams region must be strictly nested either within the implicit parallel region that surrounds the whole OpenMP program or within a target region. If a teams construct is nested within a target construct, that target construct must contain no statements, declarations or directives outside of the teams construct.
• **distribute, distribute simd**, distribute parallel worksharing-loop, distribute parallel worksharing-loop SIMD, **loop, parallel** regions, including any **parallel** regions arising from combined constructs, **omp_get_num_teams()** regions, and **omp_get_team_num()** regions are the only OpenMP regions that may be strictly nested inside the **teams** region.

• A **loop** region that binds to a **teams** region must be strictly nested inside a **teams** region.

• A **distribute** region must be strictly nested inside a **teams** region.

• If **construct-type-clause** is **taskgroup**, the **cancel** construct must be closely nested inside a **task** construct and the **cancel** region must be closely nested inside a **taskgroup** region. If **construct-type-clause** is **sections**, the **cancel** construct must be closely nested inside a **sections** or **section** construct. Otherwise, the **cancel** construct must be closely nested inside an OpenMP construct that matches the type specified in **construct-type-clause** of the **cancel** construct.

• A **cancellation point** construct for which **construct-type-clause** is **taskgroup** must be closely nested inside a **task** construct, and the **cancellation point** region must be closely nested inside a **taskgroup** region. A **cancellation point** construct for which **construct-type-clause** is **sections** must be closely nested inside a **sections** or **section** construct. Otherwise, a **cancellation point** construct must be closely nested inside an OpenMP construct that matches the type specified in **construct-type-clause**.

• The only constructs that may be encountered inside a region that corresponds to a construct with an **order** clause that specifies **concurrent** are the **loop** construct, the **parallel** construct, the **simd** construct, and combined constructs for which the first construct is a **parallel** construct.

• A region that corresponds to a construct with an **order** clause that specifies **concurrent** may not contain calls to procedures that contain OpenMP directives or calls to the OpenMP Runtime API.

• A **scope** region may not be closely nested inside a worksharing, **loop, task, taskloop, critical, ordered, atomic**, or **masked** region.
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3 Runtime Library Routines

This chapter describes the OpenMP API runtime library routines and queryable runtime states. In this chapter, *true* and *false* are used as generic terms to simplify the description of the routines.

- **C / C++**
  - *true* means a non-zero integer value and *false* means an integer value of zero.

- **Fortran**
  - *true* means a logical value of `.TRUE.` and *false* means a logical value of `.FALSE.`.

**Restrictions**

The following restrictions apply to all OpenMP runtime library routines:

- OpenMP runtime library routines may not be called from `PURE` or `ELEMENTAL` procedures.
- OpenMP runtime library routines may not be called in `DO CONCURRENT` constructs.

**3.1 Runtime Library Definitions**

For each base language, a compliant implementation must supply a set of definitions for the OpenMP API runtime library routines and the special data types of their parameters. The set of definitions must contain a declaration for each OpenMP API runtime library routine and variable and a definition of each required data type listed below. In addition, each set of definitions may specify other implementation specific values.

- **C / C++**
  - The library routines are external functions with “C” linkage.

Prototypes for the C/C++ runtime library routines described in this chapter shall be provided in a header file named `omp.h`. This file also defines the following:
• The type `omp_allocator_handle_t`, which must be an implementation-defined (for C++ possibly scoped) enum type with at least the `omp_null_allocator` enumerator with the value zero and an enumerator for each predefined memory allocator in Table 2.10;

• `omp_atv_default`, which is an instance of a type compatible with `omp_uintptr_t` with the value -1;

• The type `omp_control_tool_result_t`;

• The type `omp_control_tool_t`;

• The type `omp_depend_t`;

• The type `omp_event_handle_t`, which must be an implementation-defined (for C++ possibly scoped) enum type;

• The type `omp_intptr_t`, which is a signed integer type that is at least the size of a pointer on any device;

• The type `omp_interop_t`, which must be an implementation-defined integral or pointer type;

• The type `omp_interop_fr_t`, which must be an implementation-defined enum type with enumerators named `omp_ifr_name` where `name` is a foreign runtime name that is defined in the OpenMP Additional Definitions document;

• The type `omp_lock_hint_t` (deprecated);

• The type `omp_lock_t`;

• The type `omp_memspace_handle_t`, which must be an implementation-defined (for C++ possibly scoped) enum type with an enumerator for at least each predefined memory space in Table 2.8;

• The type `omp_nest_lock_t`;

• The type `omp_pause_resource_t`;

• The type `omp_proc_bind_t`;

• The type `omp_sched_t`;

• The type `omp_sync_hint_t`; and

• The type `omp_uintptr_t`, which is an unsigned integer type capable of holding a pointer on any device.
The OpenMP enumeration types provided in the `omp.h` header file shall not be scoped enumeration types unless explicitly allowed.

The `omp.h` header file also defines a class template that models the `Allocator` concept in the `omp::allocator` namespace for each predefined memory allocator in Table 2.10 for which the name includes neither the `omp_` prefix nor the `_alloc` suffix.

The OpenMP Fortran API runtime library routines are external procedures. The return values of these routines are of default kind, unless otherwise specified.

Interface declarations for the OpenMP Fortran runtime library routines described in this chapter shall be provided in the form of a Fortran `module` named `omp_lib` or a Fortran `include` file named `omp_lib.h`. Whether the `omp_lib.h` file provides derived-type definitions or those routines that require an explicit interface is implementation defined. Whether the `include` file or the `module` file (or both) is provided is also implementation defined.

These files also define the following:

- The default integer named constant `omp_allocator_handle_kind`;
- An integer named constant of kind `omp_allocator_handle_kind` for each predefined memory allocator in Table 2.10;
- The default integer named constant `omp_alloctrait_key_kind`;
- The default integer named constant `omp_alloctrait_val_kind`;
- The default integer named constant `omp_control_tool_kind`;
- The default integer named constant `omp_control_tool_result_kind`;
- The default integer named constant `omp_depend_kind`;
- The default integer named constant `omp_event_handle_kind`;
- The default integer named constant `omp_interop_kind`;
- The default integer named constant `omp_interop_fr_kind`;
- An integer named constant `omp_ifr_name` of kind `omp_interop_fr_kind` for each `name` that is a foreign runtime name that is defined in the OpenMP Additional Definitions document;
- The default integer named constant `omp_lock_hint_kind` (deprecated);
- The default integer named constant `omp_lock_kind`;
- The default integer named constant `omp_memspace_handle_kind`;
• An integer named constant of kind *omp_memspace_handle_kind* for each predefined memory space in Table 2.8;

• The default integer named constant *omp_nest_lock_kind*;

• The default integer named constant *omp_pause_resource_kind*;

• The default integer named constant *omp_proc_bind_kind*;

• The default integer named constant *omp_sched_kind*;

• The default integer named constant *omp_sync_hint_kind*; and

• The default integer named constant *openmp_version* with a value *yyyymm* where *yyy* and *mm* are the year and month designations of the version of the OpenMP Fortran API that the implementation supports; this value matches that of the C preprocessor macro *_OPENMP*, when a macro preprocessor is supported (see Section 2.2).

Whether any of the OpenMP runtime library routines that take an argument are extended with a generic interface so arguments of different *KIND* type can be accommodated is implementation defined.

---

### 3.2 Thread Team Routines

This section describes routines that affect and monitor thread teams in the current contention group.

#### 3.2.1 omp_set_num_threads

**Summary**

The *omp_set_num_threads* routine affects the number of threads to be used for subsequent parallel regions that do not specify a *num_threads* clause, by setting the value of the first element of the *nthreads-var ICV* of the current task.

**Format**

```c
void omp_set_num_threads(int num_threads);
```

```fortran
subroutine omp_set_num_threads(num_threads)
  integer num_threads
```

---
Constraints on Arguments
The value of the argument passed to this routine must evaluate to a positive integer, or else the behavior of this routine is implementation defined.

Binding
The binding task set for an `omp_set_num_threads` region is the generating task.

Effect
The effect of this routine is to set the value of the first element of the `nthreads-var` ICV of the current task to the value specified in the argument.

Cross References
- `nthreads-var` ICV, see Section 2.4.
- `parallel` construct and `num_threads` clause, see Section 2.6.
- Determining the number of threads for a `parallel` region, see Section 2.6.1.
- `omp_get_num_threads` routine, see Section 3.2.2.
- `omp_get_max_threads` routine, see Section 3.2.3.
- `OMP_NUM_THREADS` environment variable, see Section 6.2.

3.2.2omp_get_num_threads
Summary
The `omp_get_num_threads` routine returns the number of threads in the current team.

Format

C / C++

```c
int omp_get_num_threads(void);
```

Fortran

```fortran
integer function omp_get_num_threads()
```

Binding
The binding region for an `omp_get_num_threads` region is the innermost enclosing `parallel` region.

Effect
The `omp_get_num_threads` routine returns the number of threads in the team that is executing the `parallel` region to which the routine region binds. If called from the sequential part of a program, this routine returns 1.
Cross References
• \textit{nthreads-var} ICV, see Section 2.4.
• \texttt{parallel} construct and \texttt{num_threads} clause, see Section 2.6.
• Determining the number of threads for a \texttt{parallel} region, see Section 2.6.1.
• \texttt{omp_set_num_threads} routine, see Section 3.2.1.
• \texttt{OMP_NUM_THREADS} environment variable, see Section 6.2.

3.2.3 \texttt{omp_get_max_threads}

Summary
The \texttt{omp_get_max_threads} routine returns an upper bound on the number of threads that could be used to form a new team if a \texttt{parallel} construct without a \texttt{num_threads} clause were encountered after execution returns from this routine.

Format
\begin{verbatim}
int omp_get_max_threads(void);

integer function omp_get_max_threads()
\end{verbatim}

Binding
The binding task set for an \texttt{omp_get_max_threads} region is the generating task.

Effect
The value returned by \texttt{omp_get_max_threads} is the value of the first element of the \textit{nthreads-var} ICV of the current task. This value is also an upper bound on the number of threads that could be used to form a new team if a parallel region without a \texttt{num_threads} clause were encountered after execution returns from this routine.

Note – The return value of the \texttt{omp_get_max_threads} routine can be used to allocate sufficient storage dynamically for all threads in the team formed at the subsequent active \texttt{parallel} region.
Cross References

- *nthreads-var* ICV, see Section 2.4.
- `parallel` construct and `num_threads` clause, see Section 2.6.
- Determining the number of threads for a `parallel` region, see Section 2.6.1.
- `omp_set_num_threads` routine, see Section 3.2.1.
- `omp_get_num_threads` routine, see Section 3.2.2.
- `omp_get_thread_num` routine, see Section 3.2.4.
- `OMP_NUM_THREADS` environment variable, see Section 6.2.

3.2.4 `omp_get_thread_num`

Summary

The `omp_get_thread_num` routine returns the thread number, within the current team, of the calling thread.

Format

```c
int omp_get_thread_num(void);
```

```fortran
integer function omp_get_thread_num()
```

Binding

The binding thread set for an `omp_get_thread_num` region is the current team. The binding region for an `omp_get_thread_num` region is the innermost enclosing `parallel` region.

Effect

The `omp_get_thread_num` routine returns the thread number of the calling thread, within the team that is executing the `parallel` region to which the routine region binds. The thread number is an integer between 0 and one less than the value returned by `omp_get_num_threads`, inclusive. The thread number of the primary thread of the team is 0. The routine returns 0 if it is called from the sequential part of a program.

Note – The thread number may change during the execution of an untied task. The value returned by `omp_get_thread_num` is not generally useful during the execution of such a task region.
Cross References

- \textit{nthreads-var} ICV, see Section 2.4.
- \texttt{parallel} construct and \texttt{num_threads} clause, see Section 2.6.
- Determining the number of threads for a \texttt{parallel} region, see Section 2.6.1.
- \texttt{omp_set_num_threads} routine, see Section 3.2.1.
- \texttt{omp_get_num_threads} routine, see Section 3.2.2.
- \texttt{OMP_NUM_THREADS} environment variable, see Section 6.2.

3.2.5 \texttt{omp_in_parallel}

Summary

The \texttt{omp_in_parallel} routine returns \texttt{true} if the \texttt{active-levels-var} ICV is greater than zero; otherwise, it returns \texttt{false}.

Format

\begin{verbatim}
int omp_in_parallel(void);
logical function omp_in_parallel()
\end{verbatim}

Binding

The binding task set for an \texttt{omp_in_parallel} region is the generating task.

Effect

The effect of the \texttt{omp_in_parallel} routine is to return \texttt{true} if the current task is enclosed by an active \texttt{parallel} region, and the \texttt{parallel} region is enclosed by the outermost initial task region on the device; otherwise it returns \texttt{false}.

Cross References

- \textit{active-levels-var}, see Section 2.4.
- \texttt{parallel} construct, see Section 2.6.
- \texttt{omp_get_num_threads} routine, see Section 3.2.2.
- \texttt{omp_get_active_level} routine, see Section 3.2.20.
3.2.6 omp_set_dynamic

Summary
The `omp_set_dynamic` routine enables or disables dynamic adjustment of the number of threads available for the execution of subsequent `parallel` regions by setting the value of the `dyn-var` ICV.

Format

```
C / C++
void omp_set_dynamic(int dynamic_threads);
```

Fortran

```
subroutine omp_set_dynamic(dynamic_threads)
  logical dynamic_threads
```

Binding
The binding task set for an `omp_set_dynamic` region is the generating task.

Effect
For implementations that support dynamic adjustment of the number of threads, if the argument to `omp_set_dynamic` evaluates to `true`, dynamic adjustment is enabled for the current task; otherwise, dynamic adjustment is disabled for the current task. For implementations that do not support dynamic adjustment of the number of threads, this routine has no effect: the value of `dyn-var` remains `false`.

Cross References
- `dyn-var` ICV, see Section 2.4.
- Determining the number of threads for a `parallel` region, see Section 2.6.1.
- `omp_get_num_threads` routine, see Section 3.2.2.
- `omp_get_dynamic` routine, see Section 3.2.7.
- `OMP_DYNAMIC` environment variable, see Section 6.3.

3.2.7 omp_get_dynamic

Summary
The `omp_get_dynamic` routine returns the value of the `dyn-var` ICV, which determines whether dynamic adjustment of the number of threads is enabled or disabled.
**Format**

<table>
<thead>
<tr>
<th>C / C++</th>
<th>C / C++</th>
<th>Fortran</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>int omp_get_dynamic(void);</code></td>
<td><code>logical function omp_get_dynamic()</code></td>
<td></td>
</tr>
</tbody>
</table>

**Binding**

The binding task set for an `omp_get_dynamic` region is the generating task.

**Effect**

This routine returns `true` if dynamic adjustment of the number of threads is enabled for the current task; it returns `false`, otherwise. If an implementation does not support dynamic adjustment of the number of threads, then this routine always returns `false`.

**Cross References**

- `dyn-var` ICV, see Section 2.4.
- Determining the number of threads for a `parallel` region, see Section 2.6.1.
- `omp_set_dynamic` routine, see Section 3.2.6.
- `OMP_DYNAMIC` environment variable, see Section 6.3.

### 3.2.8 omp_get_cancellation

**Summary**

The `omp_get_cancellation` routine returns the value of the `cancel-var` ICV, which determines if cancellation is enabled or disabled.

**Format**

<table>
<thead>
<tr>
<th>C / C++</th>
<th>C / C++</th>
<th>Fortran</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>int omp_get_cancellation(void);</code></td>
<td><code>logical function omp_get_cancellation()</code></td>
<td></td>
</tr>
</tbody>
</table>

**Binding**

The binding task set for an `omp_get_cancellation` region is the whole program.
Effect
This routine returns true if cancellation is enabled. It returns false otherwise.

Cross References
• cancel-var ICV, see Section 2.4.1.
• cancel construct, see Section 2.20.1.
• OMP_CANCELLATION environment variable, see Section 6.11.

3.2.9 omp_set_nested (Deprecated)

Summary
The deprecated omp_set_nested routine enables or disables nested parallelism by setting the max-active-levels-var ICV.

Format

C / C++

```c
void omp_set_nested(int nested);
```

Fortran

```fortran
subroutine omp_set_nested(nested)
    logical nested
```

Binding
The binding task set for an omp_set_nested region is the generating task.

Effect
If the argument to omp_set_nested evaluates to true, the value of the max-active-levels-var ICV is set to the number of active levels of parallelism that the implementation supports; otherwise, if the value of max-active-levels-var is greater than 1 then it is set to 1. This routine has been deprecated.

Cross References
• max-active-levels-var ICV, see Section 2.4.
• Determining the number of threads for a parallel region, see Section 2.6.1.
• omp_get_nested routine, see Section 3.2.10.
• omp_set_max_active_levels routine, see Section 3.2.15.
• omp_get_max_active_levels routine, see Section 3.2.16.
• OMP_NESTED environment variable, see Section 6.9.
3.2.10 omp_get_nested (Deprecated)

**Summary**
The deprecated `omp_get_nested` routine returns whether nested parallelism is enabled or disabled, according to the value of the `max-active-levels-var` ICV.

**Format**

```c
int omp_get_nested(void);
```

**Binding**
The binding task set for an `omp_get_nested` region is the generating task.

**Effect**
This routine returns `true` if `max-active-levels-var` is greater than 1 and greater than `active-levels-var` for the current task; it returns `false`, otherwise. If an implementation does not support nested parallelism, this routine always returns `false`. This routine has been deprecated.

**Cross References**
- `max-active-levels-var` ICV, see Section 2.4.
- Determining the number of threads for a `parallel` region, see Section 2.6.1.
- `omp_set_nested` routine, see Section 3.2.9.
- `omp_set_max_active_levels` routine, see Section 3.2.15.
- `omp_get_max_active_levels` routine, see Section 3.2.16.
- `OMP_NESTED` environment variable, see Section 6.9.

3.2.11 omp_set_schedule

**Summary**
The `omp_set_schedule` routine affects the schedule that is applied when `runtime` is used as schedule kind, by setting the value of the `run-sched-var` ICV.
Format

C / C++

```c
void omp_set_schedule(omp_sched_t kind, int chunk_size);
```

Fortran

```fortran
subroutine omp_set_schedule(kind, chunk_size)
  integer (kind=omp_sched_kind) kind
  integer chunk_size
```

Constraints on Arguments

The first argument passed to this routine can be one of the valid OpenMP schedule kinds (except for `runtime`) or any implementation-specific schedule. The C/C++ header file (`omp.h`) and the Fortran include file (`omp_lib.h`) and/or Fortran module file (`omp_lib`) define the valid constants. The valid constants must include the following, which can be extended with implementation-specific values:

C / C++

```c
typedef enum omp_sched_t {
  // schedule kinds
  omp_sched_static = 0x1,
  omp_sched_dynamic = 0x2,
  omp_sched_guided = 0x3,
  omp_sched_auto = 0x4,

  // schedule modifier
  omp_sched_monotonic = 0x80000000u
} omp_sched_t;
```
Fortran

! schedule kinds
integer(kind=omp_sched_kind), &
  parameter :: omp_sched_static = &
           int(Z'1', kind=omp_sched_kind)
integer(kind=omp_sched_kind), &
  parameter :: omp_sched_dynamic = &
           int(Z'2', kind=omp_sched_kind)
integer(kind=omp_sched_kind), &
  parameter :: omp_sched_guided = &
           int(Z'3', kind=omp_sched_kind)
integer(kind=omp_sched_kind), &
  parameter :: omp_sched__auto = &
           int(Z'4', kind=omp_sched_kind)

! schedule modifier
integer(kind=omp_sched_kind), &
  parameter :: omp_sched_monotonic = &
           int(Z'80000000', kind=omp_sched_kind)

Binding
The binding task set for an omp_set_schedule region is the generating task.

Effect
The effect of this routine is to set the value of the run-sched-var ICV of the current task to the values specified in the two arguments. The schedule is set to the schedule kind that is specified by the first argument kind. It can be any of the standard schedule kinds or any other implementation-specific one. For the schedule kinds static, dynamic, and guided the chunk_size is set to the value of the second argument, or to the default chunk_size if the value of the second argument is less than 1; for the schedule kind auto the second argument has no meaning; for implementation-specific schedule kinds, the values and associated meanings of the second argument are implementation defined.

Each of the schedule kinds can be combined with the omp_sched_monotonic modifier by using the + or | operators in C/C++ or the + operator in Fortran. If the schedule kind is combined with the omp_sched_monotonic modifier, the schedule is modified as if the monotonic schedule modifier was specified. Otherwise, the schedule modifier is nonmonotonic.
Cross References

- run-sched-var ICV, see Section 2.4.
- Determining the schedule of a worksharing-loop, see Section 2.11.4.1.
- omp_get_schedule routine, see Section 3.2.12.
- OMP_SCHEDULE environment variable, see Section 6.1.

3.2.12 omp_get_schedule

Summary
The omp_get_schedule routine returns the schedule that is applied when the runtime schedule is used.

Format

C / C++
void omp_get_schedule(omp_sched_t *kind, int *chunk_size);

C / C++

Fortran

subroutine omp_get_schedule(kind, chunk_size)
integer (kind=omp_sched_kind) kind
integer chunk_size

Binding
The binding task set for an omp_get_schedule region is the generating task.

Effect
This routine returns the run-sched-var ICV in the task to which the routine binds. The first argument kind returns the schedule to be used. It can be any of the standard schedule kinds as defined in Section 3.2.11, or any implementation-specific schedule kind. The second argument chunk_size returns the chunk size to be used, or a value less than 1 if the default chunk size is to be used, if the returned schedule kind is static, dynamic, or guided. The value returned by the second argument is implementation defined for any other schedule kinds.

Cross References

- run-sched-var ICV, see Section 2.4.
- Determining the schedule of a worksharing-loop, see Section 2.11.4.1.
- omp_set_schedule routine, see Section 3.2.11.
- OMP_SCHEDULE environment variable, see Section 6.1.
3.2.13 omp_get_thread_limit

Summary
The omp_get_thread_limit routine returns the maximum number of OpenMP threads available to participate in the current contention group.

Format

\[
\begin{align*}
\text{C / C++} & : & \text{int} & \text{omp_get_thread_limit}(\text{void}); \\
\text{Fortran} & : & \text{integer function} & \text{omp_get_thread_limit}() \\
\end{align*}
\]

Binding
The binding task set for an omp_get_thread_limit region is the generating task.

Effect
The omp_get_thread_limit routine returns the value of the thread-limit-var ICV.

Cross References
- thread-limit-var ICV, see Section 2.4.
- omp_get_num_threads routine, see Section 3.2.2.
- OMP_NUM_THREADS environment variable, see Section 6.2.
- OMP_THREAD_LIMIT environment variable, see Section 6.10.

3.2.14 omp_get_supported_active_levels

Summary
The omp_get_supported_active_levels routine returns the number of active levels of parallelism supported by the implementation.

Format

\[
\begin{align*}
\text{C / C++} & : & \text{int} & \text{omp_get_supported_active_levels}(\text{void}); \\
\text{Fortran} & : & \text{integer function} & \text{omp_get_supported_active_levels}() \\
\end{align*}
\]
The binding task set for an `omp_get_supported_active_levels` region is the generating task.

The `omp_get_supported_active_levels` routine returns the number of active levels of parallelism supported by the implementation. The `max-active-levels-var` ICV may not have a value that is greater than this number. The value returned by the `omp_get_supported_active_levels` routine is implementation defined, but it must be greater than 0.

Cross References

- `max-active-levels-var` ICV, see Section 2.4.
- `omp_set_max_active_levels` routine, see Section 3.2.15.
- `omp_get_max_active_levels` routine, see Section 3.2.16.

### 3.2.15 omp_set_max_active_levels

Summary

The `omp_set_max_active_levels` routine limits the number of nested active parallel regions when a new nested parallel region is generated by the current task by setting the `max-active-levels-var` ICV.

Format

```
C / C++
void omp_set_max_active_levels(int max_levels);
```

```
Fortran
subroutine omp_set_max_active_levels(max_levels)
   integer max_levels
```

Constraints on Arguments

The value of the argument passed to this routine must evaluate to a non-negative integer, otherwise the behavior of this routine is implementation defined.

The binding task set for an `omp_set_max_active_levels` region is the generating task.
Effect
The effect of this routine is to set the value of the max-active-levels-var ICV to the value specified in the argument.

If the number of active levels requested exceeds the number of active levels of parallelism supported by the implementation, the value of the max-active-levels-var ICV will be set to the number of active levels supported by the implementation.

Cross References
• max-active-levels-var ICV, see Section 2.4.
• parallel construct, see Section 2.6.
• omp_get_supported_active_levels routine, see Section 3.2.14.
• omp_get_max_active_levels routine, see Section 3.2.16.
• OMP_MAX_ACTIVE_LEVELS environment variable, see Section 6.8.

3.2.16 omp_get_max_active_levels
Summary
The omp_get_max_active_levels routine returns the value of the max-active-levels-var ICV, which determines the maximum number of nested active parallel regions when the innermost parallel region is generated by the current task.

Format
```
C / C++
int omp_get_max_active_levels(void);
```

Binding
The binding task set for an omp_get_max_active_levels region is the generating task.

Effect
The omp_get_max_active_levels routine returns the value of the max-active-levels-var ICV. The current task may only generate an active parallel region if the returned value is greater than the value of the active-levels-var ICV.
Cross References

• max-active-levels-var ICV, see Section 2.4.

• parallel construct, see Section 2.6.

• omp_get_supported_active_levels routine, see Section 3.2.14.

• omp_set_max_active_levels routine, see Section 3.2.15.

• OMP_MAX_ACTIVE_LEVELS environment variable, see Section 6.8.

3.2.17 omp_get_level

Summary

The omp_get_level routine returns the value of the levels-var ICV.

Format

<table>
<thead>
<tr>
<th>int omp_get_level(void);</th>
<th>C / C++</th>
</tr>
</thead>
<tbody>
<tr>
<td>integer function omp_get_level()</td>
<td>Fortran</td>
</tr>
</tbody>
</table>

Binding

The binding task set for an omp_get_level region is the generating task.

Effect

The effect of the omp_get_level routine is to return the number of nested parallel regions (whether active or inactive) that enclose the current task such that all of the parallel regions are enclosed by the outermost initial task region on the current device.

Cross References

• levels-var ICV, see Section 2.4.

• parallel construct, see Section 2.6.

• omp_get_active_level routine, see Section 3.2.20.

• OMP_MAX_ACTIVE_LEVELS environment variable, see Section 6.8.
3.2.18 omp_get_ancestor_thread_num

Summary
The **omp_get_ancestor_thread_num** routine returns, for a given nested level of the current thread, the thread number of the ancestor of the current thread.

Format

<table>
<thead>
<tr>
<th>C / C++</th>
<th>Fortran</th>
</tr>
</thead>
<tbody>
<tr>
<td>int omp_get_ancestor_thread_num(int level);</td>
<td>integer function omp_get_ancestor_thread_num(level)</td>
</tr>
</tbody>
</table>

Binding
The binding thread set for an **omp_get_ancestor_thread_num** region is the encountering thread. The binding region for an **omp_get_ancestor_thread_num** region is the innermost enclosing parallel region.

Effect
The **omp_get_ancestor_thread_num** routine returns the thread number of the ancestor at a given nest level of the current thread or the thread number of the current thread. If the requested nest level is outside the range of 0 and the nest level of the current thread, as returned by the **omp_get_level** routine, the routine returns -1.

Note – When the **omp_get_ancestor_thread_num** routine is called with a value of level=0, the routine always returns 0. If level=omp_get_level(), the routine has the same effect as the **omp_get_thread_num** routine.

Cross References
- parallel construct, see Section 2.6.
- **omp_get_num_threads** routine, see Section 3.2.2.
- **omp_get_thread_num** routine, see Section 3.2.4.
- **omp_get_level** routine, see Section 3.2.17.
- **omp_get_team_size** routine, see Section 3.2.19.
3.2.19 omp_get_team_size

Summary
The `omp_get_team_size` routine returns, for a given nested level of the current thread, the size
of the thread team to which the ancestor or the current thread belongs.

Format
```
C / C++
int omp_get_team_size(int level);
```
```fortran
integer function omp_get_team_size(level)
integer level
```

Binding
The binding thread set for an `omp_get_team_size` region is the encountering thread. The
binding region for an `omp_get_team_size` region is the innermost enclosing `parallel`
region.

Effect
The `omp_get_team_size` routine returns the size of the thread team to which the ancestor or
the current thread belongs. If the requested nested level is outside the range of 0 and the nested
level of the current thread, as returned by the `omp_get_level` routine, the routine returns -1.
Inactive parallel regions are regarded like active parallel regions executed with one thread.

Note – When the `omp_get_team_size` routine is called with a value of `level`=0, the routine
always returns 1. If `level=omp_get_level()`, the routine has the same effect as the
`omp_get_num_threads` routine.

Cross References
• `omp_get_num_threads` routine, see Section 3.2.2.
• `omp_get_level` routine, see Section 3.2.17.
• `omp_get_ancestor_thread_num` routine, see Section 3.2.18.

3.2.20 omp_get_active_level

Summary
The `omp_get_active_level` routine returns the value of the `active-level-var` ICV.
The binding task set for the an `omp_get_active_level` region is the generating task.

The effect of the `omp_get_active_level` routine is to return the number of nested active parallel regions enclosing the current task such that all of the parallel regions are enclosed by the outermost initial task region on the current device.

Cross References

- `active-levels-var` ICV, see Section 2.4.
- `omp_set_max_active_levels` routine, see Section 3.2.15.
- `omp_get_max_active_levels` routine, see Section 3.2.16.
- `omp_get_level` routine, see Section 3.2.17.
- `OMP_MAX_ACTIVE_LEVELS` environment variable, see Section 6.8.

### 3.3 Thread Affinity Routines

This section describes routines that affect and access thread affinity policies that are in effect.

#### 3.3.1 omp_get_proc_bind

**Summary**

The `omp_get_proc_bind` routine returns the thread affinity policy to be used for the subsequent nested parallel regions that do not specify a `proc_bind` clause.
integer (kind=omp_proc_bind_kind) function omp_get_proc_bind()

**Constraints on Arguments**

The value returned by this routine must be one of the valid affinity policy kinds. The C/C++ header file (omp.h) and the Fortran include file (omp_lib.h) and/or Fortran module file (omp_lib) define the valid constants. The valid constants must include the following:

```c
typedef enum omp_proc_bind_t {
    omp_proc_bind_false = 0,
    omp_proc_bind_true = 1,
    omp_proc_bind_primary = 2,
    omp_proc_bind_master = omp_proc_bind_primary, // (deprecated)
    omp_proc_bind_close = 3,
    omp_proc_bind_spread = 4
} omp_proc_bind_t;
```

```fortran
integer (kind=omp_proc_bind_kind), parameter :: omp_proc_bind_false = 0
integer (kind=omp_proc_bind_kind), parameter :: omp_proc_bind_true = 1
integer (kind=omp_proc_bind_kind), parameter :: omp_proc_bind_primary = 2
integer (kind=omp_proc_bind_kind), parameter :: omp_proc_bind_master = omp_proc_bind_primary ! (deprecated)
integer (kind=omp_proc_bind_kind), parameter :: omp_proc_bind_close = 3
integer (kind=omp_proc_bind_kind), parameter :: omp_proc_bind_spread = 4
```

**Binding**

The binding task set for an **omp_get_proc_bind** region is the generating task.

**Effect**

The effect of this routine is to return the value of the first element of the *bind-var* ICV of the current task. See Section 2.6.2 for the rules that govern the thread affinity policy.
Cross References

• bind-var ICV, see Section 2.4.
• Controlling OpenMP thread affinity, see Section 2.6.2.
• omp_get_num_places routine, see Section 3.3.2.
• OMP_PROC_BIND environment variable, see Section 6.4.
• OMP_PLACES environment variable, see Section 6.5.

3.3.2 omp_get_num_places

Summary
The omp_get_num_places routine returns the number of places available to the execution environment in the place list.

Format

```
C / C++
int omp_get_num_places(void);
```

```
C / C++
Fortran
integer function omp_get_num_places()
```

Binding
The binding thread set for an omp_get_num_places region is all threads on a device. The effect of executing this routine is not related to any specific region corresponding to any construct or API routine.

Effect
The omp_get_num_places routine returns the number of places in the place list. This value is equivalent to the number of places in the place-partition-var ICV in the execution environment of the initial task.

Cross References

• place-partition-var ICV, see Section 2.4.
• Controlling OpenMP thread affinity, see Section 2.6.2.
• omp_get_place_num routine, see Section 3.3.5.
• OMP_PLACES environment variable, see Section 6.5.
3.3.3 omp_get_place_num_procs

Summary
The `omp_get_place_num_procs` routine returns the number of processors available to the
execution environment in the specified place.

Format

```
C / C++
int omp_get_place_num_procs(int place_num);
```

Fortran

```
integer function omp_get_place_num_procs(place_num)
integer place_num
```

Binding
The binding thread set for an `omp_get_place_num_procs` region is all threads on a device.
The effect of executing this routine is not related to any specific region corresponding to any
construct or API routine.

Effect
The `omp_get_place_num_procs` routine returns the number of processors associated with
the place numbered `place_num`. The routine returns zero when `place_num` is negative, or is greater
than or equal to the value returned by `omp_get_num_places()`.

Cross References
- `place-partition-var ICV`, see Section 2.4.
- Controlling OpenMP thread affinity, see Section 2.6.2.
- `omp_get_num_places` routine, see Section 3.3.2.
- `omp_get_place_proc_ids` routine, see Section 3.3.4.
- `OMP_PLACES` environment variable, see Section 6.5.

3.3.4 omp_get_place_proc_ids

Summary
The `omp_get_place_proc_ids` routine returns the numerical identifiers of the processors
available to the execution environment in the specified place.

Format

```
C / C++
void omp_get_place_proc_ids(int place_num, int *ids);
```

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subroutine omp_get_place_proc_ids(place_num, ids)
integer place_num
integer ids(*)

The binding thread set for an omp_get_place_proc_ids region is all threads on a device.
The effect of executing this routine is not related to any specific region corresponding to any
construct or API routine.

The omp_get_place_proc_ids routine returns the numerical identifiers of each processor
associated with the place numbered place_num. The numerical identifiers are non-negative and
their meaning is implementation defined. The numerical identifiers are returned in the array ids and
their order in the array is implementation defined. The array must be sufficiently large to contain
omp_get_place_num_procs(place_num) integers; otherwise, the behavior is unspecified.
The routine has no effect when place_num has a negative value or a value greater than or equal to
omp_get_num_places().

Cross References
• place-partition-var ICV, see Section 2.4.
• Controlling OpenMP thread affinity, see Section 2.6.2.
• omp_get_num_places routine, see Section 3.3.2.
• omp_get_place_num_procs routine, see Section 3.3.3.
• OMP PLACES environment variable, see Section 6.5.

3.3.5 omp_get_place_num

The omp_get_place_num routine returns the place number of the place to which the
encountering thread is bound.

Format
int omp_get_place_num(void);

integer function omp_get_place_num()
The binding thread set for an `omp_get_place_num` region is the encountering thread.

When the encountering thread is bound to a place, the `omp_get_place_num` routine returns the place number associated with the thread. The returned value is between 0 and one less than the value returned by `omp_get_num_places()`, inclusive. When the encountering thread is not bound to a place, the routine returns -1.

Cross References

- `place-partition-var` ICV, see Section 2.4.
- Controlling OpenMP thread affinity, see Section 2.6.2.
- `omp_get_num_places` routine, see Section 3.3.2.
- `OMP_PLACES` environment variable, see Section 6.5.

### 3.3.6 `omp_get_partition_num_places`

**Summary**
The `omp_get_partition_num_places` routine returns the number of places in the place partition of the innermost implicit task.

**Format**

```c/c++
int omp_get_partition_num_places(void);
```

```fortran
integer function omp_get_partition_num_places()
```

**Binding**
The binding task set for an `omp_get_partition_num_places` region is the encountering implicit task.

**Effect**
The `omp_get_partition_num_places` routine returns the number of places in the `place-partition-var` ICV.
Cross References

- *place-partition-var* ICV, see Section 2.4.
- Controlling OpenMP thread affinity, see Section 2.6.2.
- *omp_get_num_places* routine, see Section 3.3.2.
- *OMP_PLACES* environment variable, see Section 6.5.

3.3.7 omp_get_partition_place_nums

Summary

The *omp_get_partition_place_nums* routine returns the list of place numbers corresponding to the places in the *place-partition-var* ICV of the innermost implicit task.

Format

```
C / C++
void omp_get_partition_place_nums(int *place_nums);
```

```
C / C++
fortran
subroutine omp_get_partition_place_nums(place_nums)
integer place_nums(*)
```

Binding

The binding task set for an *omp_get_partition_place_nums* region is the encountering implicit task.

Effect

The *omp_get_partition_place_nums* routine returns the list of place numbers that correspond to the places in the *place-partition-var* ICV of the innermost implicit task. The array must be sufficiently large to contain *omp_get_partition_num_places()* integers; otherwise, the behavior is unspecified.

Cross References

- *place-partition-var* ICV, see Section 2.4.
- Controlling OpenMP thread affinity, see Section 2.6.2.
- *omp_get_partition_num_places* routine, see Section 3.3.6.
- *OMP_PLACES* environment variable, see Section 6.5.
3.3.8 omp_set_affinity_format

Summary
The omp_set_affinity_format routine sets the affinity format to be used on the device by setting the value of the affinity-format-var ICV.

Format

| C / C++ | void omp_set_affinity_format(const char *format); |
| Fortran | subroutine omp_set_affinity_format(format) |
| character(len=*) | : : format |

Binding
When called from a sequential part of the program, the binding thread set for an omp_set_affinity_format region is the encountering thread. When called from within any parallel or teams region, the binding thread set (and binding region, if required) for the omp_set_affinity_format region is implementation defined.

Effect
The effect of omp_set_affinity_format routine is to copy the character string specified by the format argument into the affinity-format-var ICV on the current device.

This routine has the described effect only when called from a sequential part of the program. When called from within a parallel or teams region, the effect of this routine is implementation defined.

Cross References
- Controlling OpenMP thread affinity, see Section 2.6.2.
- omp_get_affinity_format routine, see Section 3.3.9.
- omp_display_affinity routine, see Section 3.3.10.
- omp_capture_affinity routine, see Section 3.3.11.
- OMP_DISPLAY_AFFINITY environment variable, see Section 6.13.
- OMP_AFFINITY_FORMAT environment variable, see Section 6.14.
3.3.9 omp_get_affinity_format

Summary
The omp_get_affinity_format routine returns the value of the affinity-format-var ICV on the device.

Format

C / C++
size_t omp_get_affinity_format(char *buffer, size_t size);

C / C++
Fortran
integer function omp_get_affinity_format(buffer)
character(len=*) , intent(out) :: buffer

Binding
When called from a sequential part of the program, the binding thread set for an omp_get_affinity_format region is the encountering thread. When called from within any parallel or teams region, the binding thread set (and binding region, if required) for the omp_get_affinity_format region is implementation defined.

Effect

C / C++
The omp_get_affinity_format routine returns the number of characters in the affinity-format-var ICV on the current device, excluding the terminating null byte (’\0’) and if size is non-zero, writes the value of the affinity-format-var ICV on the current device to buffer followed by a null byte. If the return value is larger or equal to size, the affinity format specification is truncated, with the terminating null byte stored to buffer[size-1]. If size is zero, nothing is stored and buffer may be NULL.

Fortran
The omp_get_affinity_format routine returns the number of characters that are required to hold the affinity-format-var ICV on the current device and writes the value of the affinity-format-var ICV on the current device to buffer. If the return value is larger than len(buffer), the affinity format specification is truncated.

Fortran
If the buffer argument does not conform to the specified format then the result is implementation defined.
Cross References

• Controlling OpenMP thread affinity, see Section 2.6.2.

• omp_set_affinity_format routine, see Section 3.3.8.

• omp_display_affinity routine, see Section 3.3.10.

• omp_capture_affinity routine, see Section 3.3.11.

• OMP_DISPLAY_AFFINITY environment variable, see Section 6.13.

• OMP_AFFINITY_FORMAT environment variable, see Section 6.14.

3.3.10 omp_display_affinity

Summary

The omp_display_affinity routine prints the OpenMP thread affinity information using the format specification provided.

Format

C / C++

void omp_display_affinity(const char *format);

C / C++

Fortran


Fortran

subroutine omp_display_affinity(format)
character(len=*) , intent (in) :: format

Binding

The binding thread set for an omp_display_affinity region is the encountering thread.

Effect

The omp_display_affinity routine prints the thread affinity information of the current thread in the format specified by the format argument, followed by a new-line. If the format is NULL (for C/C++) or a zero-length string (for Fortran and C/C++), the value of the affinity-format-var ICV is used. If the format argument does not conform to the specified format then the result is implementation defined.
3.3.11 omp_capture_affinity

Summary
The `omp_capture_affinity` routine prints the OpenMP thread affinity information into a buffer using the format specification provided.

Format

```c
size_t omp_capture_affinity(
    char *buffer,
    size_t size,
    const char *format
);
```

```fortran
integer function omp_capture_affinity(buffer, format)
    character(len=*) intent(out), :: buffer
    character(len=*) intent(in), :: format
```

Binding
The binding thread set for an `omp_capture_affinity` region is the encountering thread.

Effect
The `omp_capture_affinity` routine returns the number of characters in the entire thread affinity information string excluding the terminating null byte (‘\0’) and if `size` is non-zero, writes the thread affinity information of the current thread in the format specified by the `format` argument into the character string `buffer` followed by a null byte. If the return value is larger or equal to `size`, the thread affinity information string is truncated, with the terminating null byte stored to `buffer[size-1]`. If `size` is zero, nothing is stored and `buffer` may be NULL. If the `format` is NULL or a zero-length string, the value of the `affinity-format-var` ICV is used.
The `omp_capture_affinity` routine returns the number of characters required to hold the entire thread affinity information string and prints the thread affinity information of the current thread into the character string `buffer` with the size of `len(buffer)` in the format specified by the `format` argument. If the `format` is a zero-length string, the value of the `affinity-format-var` ICV is used. If the return value is larger than `len(buffer)`, the thread affinity information string is truncated. If the `format` is a zero-length string, the value of the `affinity-format-var` ICV is used.

If the `format` argument does not conform to the specified format then the result is implementation defined.

**Cross References**
- Controlling OpenMP thread affinity, see Section 2.6.2.
- `omp_set_affinity_format` routine, see Section 3.3.8.
- `omp_get_affinity_format` routine, see Section 3.3.9.
- `omp_display_affinity` routine, see Section 3.3.10.
- `OMP_DISPLAY_AFFINITY` environment variable, see Section 6.13.
- `OMP_AFFINITY_FORMAT` environment variable, see Section 6.14.

### 3.4 Teams Region Routines

This section describes routines that affect and monitor the league of teams that may execute a teams region.

#### 3.4.1 `omp_get_num_teams`

**Summary**
The `omp_get_num_teams` routine returns the number of initial teams in the current teams region.

**Format**

```c
int omp_get_num_teams(void);
```

```fortran
integer function omp_get_num_teams()
```

If the `format` argument does not conform to the specified format then the result is implementation defined.
Binding
The binding task set for an `omp_get_num_teams` region is the generating task.

Effect
The effect of this routine is to return the number of initial teams in the current `teams` region. The routine returns 1 if it is called from outside of a `teams` region.

Cross References
- `teams` construct, see Section 2.7.
- `omp_get_team_num` routine, see Section 3.4.2.

3.4.2 `omp_get_team_num`

Summary
The `omp_get_team_num` routine returns the initial team number of the calling thread.

Format
```
<table>
<thead>
<tr>
<th>C / C++</th>
</tr>
</thead>
<tbody>
<tr>
<td>int omp_get_team_num(void);</td>
</tr>
<tr>
<td>----------</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Fortran</th>
</tr>
</thead>
<tbody>
<tr>
<td>integer function omp_get_team_num()</td>
</tr>
</tbody>
</table>
```

Binding
The binding task set for an `omp_get_team_num` region is the generating task.

Effect
The `omp_get_team_num` routine returns the initial team number of the calling thread. The initial team number is an integer between 0 and one less than the value returned by `omp_get_num_teams()`, inclusive. The routine returns 0 if it is called outside of a `teams` region.

Cross References
- `teams` construct, see Section 2.7.
- `omp_get_num_teams` routine, see Section 3.4.1.
3.4.3 omp_set_num_teams

Summary
The `omp_set_num_teams` routine affects the number of threads to be used for subsequent `teams` regions that do not specify a `num_teams` clause, by setting the value of the `nteams-var` ICV of the current task.

Format

```
C / C++
void omp_set_num_teams(int num_teams);
```

```
C / C++
Fortran

subroutine omp_set_num_teams(num_teams)
integer num_teams
```

Constraints on Arguments
The value of the argument passed to this routine must evaluate to a positive integer, or else the behavior of this routine is implementation defined.

Binding
The binding task set for an `omp_set_num_teams` region is the generating task.

Effect
The effect of this routine is to set the value of the `nteams-var` ICV of the current task to the value specified in the argument.

Restrictions
Restrictions to the `omp_set_num_teams` routine are as follows:
- The routine may not be called from within a parallel region that is not the implicit parallel region that surrounds the whole OpenMP program.

Cross References
- `nteams-var` ICV, see Section 2.4.
- `teams` construct and `num_teams` clause, see Section 2.7.
- `omp_get_num_teams` routine, see Section 3.4.1.
- `omp_get_max_teams` routine, see Section 3.4.4.
- `OMP_NUM_TEAMS` environment variable, see Section 6.23.
3.4.4 omp_get_max_teams

Summary
The omp_get_max_teams routine returns an upper bound on the number of teams that could be created by a teams construct without a num_teams clause that is encountered after execution returns from this routine.

Format

C / C++

```c
int omp_get_max_teams(void);
```

Fortran

```fortran
integer function omp_get_max_teams()
```

Binding
The binding task set for an omp_get_max_teams region is the generating task.

Effect
The value returned by omp_get_max_teams is the value of the nteams-var ICV of the current task. This value is also an upper bound on the number of teams that can be created by a teams construct without a num_teams clause that is encountered after execution returns from this routine.

Cross References
- nteams-var ICV, see Section 2.4.
- teams construct and num_teams clause, see Section 2.7.
- omp_get_num_teams routine, see Section 3.4.1.
- omp_set_num_teams routine, see Section 3.4.3.

3.4.5 omp_set_teams_thread_limit

Summary
The omp_set_teams_thread_limit routine defines the maximum number of OpenMP threads that can participate in each contention group created by a teams construct.

Format

C / C++

```c
void omp_set_teams_thread_limit(int thread_limit);
```

Fortran
subroutine omp_set_teams_thread_limit(thread_limit)
  integer thread_limit

Constraints on Arguments
The value of the argument passed to this routine must evaluate to a positive integer, or else the
behavior of this routine is implementation defined.

Binding
The binding task set for an omp_set_teams_thread_limit region is the generating task.

Effect
The omp_set_teams_thread_limit routine sets the value of the teams-thread-limit-var ICV to the value of the thread_limit argument.

If the value of thread_limit exceeds the number of OpenMP threads that an implementation
supports for each contention group created by a teams construct, the value of the
teams-thread-limit-var ICV will be set to the number that is supported by the implementation.

Restrictions
Restrictions to the omp_set_teams_thread_limit routine are as follows:

- The routine may not be called from within a parallel region other than the implicit parallel region
  that surrounds the whole OpenMP program.

Cross References
- teams_thread-limit-var ICV, see Section 2.4.
- teams construct and thread_limit clause, see Section 2.7.
- omp_get_teams_thread_limit routine, see Section 3.4.6.
- OMP_TEAMS_THREAD_LIMIT environment variable, see Section 6.24.

3.4.6 omp_get_teams_thread_limit

Summary
The omp_get_teams_thread_limit routine returns the maximum number of OpenMP
threads available to participate in each contention group created by a teams construct.

Format
C / C++

int omp_get_teams_thread_limit(void);
3.5 Tasking Routines

This section describes routines that pertain to OpenMP explicit tasks.

3.5.1 omp_get_max_task_priority

Summary
The `omp_get_max_task_priority` routine returns the maximum value that can be specified in the `priority` clause.

Format

```c
int omp_get_max_task_priority(void);
```

Binding
The binding thread set for an `omp_get_max_task_priority` region is all threads on the device. The effect of executing this routine is not related to any specific region that corresponds to any construct or API routine.
Effect
The `omp_get_max_task_priority` routine returns the value of the `max-task-priority-var` ICV, which determines the maximum value that can be specified in the `priority` clause.

Cross References
- `max-task-priority-var`, see Section 2.4.
- `task` construct, see Section 2.12.1.

3.5.2 omp_in_final

Summary
The `omp_in_final` routine returns `true` if the routine is executed in a final task region; otherwise, it returns `false`.

Format

```
C / C++
int omp_in_final(void);
```

`C / C++`  
`Fortran`

```
logical function omp_in_final()
```

Binding
The binding task set for an `omp_in_final` region is the generating task.

Effect
`omp_in_final` returns `true` if the enclosing task region is final. Otherwise, it returns `false`.

Cross References
- `task` construct, see Section 2.12.1.
3.6 Resource Relinquishing Routines

This section describes routines that relinquish resources used by the OpenMP runtime.

3.6.1 omp_pause_resource

Summary

The `omp_pause_resource` routine allows the runtime to relinquish resources used by OpenMP on the specified device.

Format

```c
int omp_pause_resource(
    omp_pause_resource_t kind,
    int device_num
);
```

```fortran
integer function omp_pause_resource(kind, device_num)
integer (kind=omp_pause_resource_kind) kind
integer device_num
```

Constraints on Arguments

The first argument passed to this routine can be one of the valid OpenMP pause kind, or any implementation specific pause kind. The C/C++ header file (`omp.h`) and the Fortran include file (`omp_lib.h`) and/or Fortran module file (`omp_lib`) define the valid constants. The valid constants must include the following, which can be extended with implementation-specific values:

Format

```c
typedef enum omp_pause_resource_t {
    omp_pause_soft = 1,
    omp_pause_hard = 2
} omp_pause_resource_t;
```

```fortran
integer (kind=omp_pause_resource_kind), parameter :: &
    omp_pause_soft = 1
integer (kind=omp_pause_resource_kind), parameter :: &
    omp_pause_hard = 2
```
The second argument passed to this routine indicates the device that will be paused. The `device_num` parameter must be greater than or equal to zero and less than or equal to the result of `omp_get_num_devices()`.

**Binding**
The binding task set for an `omp_pause_resource` region is the whole program.

**Effect**
The `omp_pause_resource` routine allows the runtime to relinquish resources used by OpenMP on the specified device.

If successful, the `omp_pause_hard` value results in a hard pause for which the OpenMP state is not guaranteed to persist across the `omp_pause_resource` call. A hard pause may relinquish any data allocated by OpenMP on a given device, including data allocated by memory routines for that device as well as data present on the device as a result of a declare target directive or `target data` construct. A hard pause may also relinquish any data associated with a `threadprivate` directive. When relinquished and when applicable, base language appropriate deallocation/finalization is performed. When relinquished and when applicable, mapped data on a device will not be copied back from the device to the host.

If successful, the `omp_pause_soft` value results in a soft pause for which the OpenMP state is guaranteed to persist across the call, with the exception of any data associated with a `threadprivate` directive, which may be relinquished across the call. When relinquished and when applicable, base language appropriate deallocation/finalization is performed.

---

**Note** – A hard pause may relinquish more resources, but may resume processing OpenMP regions more slowly. A soft pause allows OpenMP regions to restart more quickly, but may relinquish fewer resources. An OpenMP implementation will reclaim resources as needed for OpenMP regions encountered after the `omp_pause_resource` region. Since a hard pause may unmapped data on the specified device, appropriate data mapping is required before using data on the specified device after the `omp_pause_region` region.

The routine returns zero in case of success, and non-zero otherwise.

**Tool Callbacks**
If the tool is not allowed to interact with the specified device after encountering this call, then the runtime must call the tool finalizer for that device.

**Restrictions**
Restrictions to the `omp_pause_resource` routine are as follows:

- The `omp_pause_resource` region may not be nested in any explicit OpenMP region.
- The routine may only be called when all explicit tasks have finalized execution.
3.6.2 omp_pause_resource_all

Summary
The omp_pause_resource_all routine allows the runtime to relinquish resources used by OpenMP on all devices.

Format

\[
\begin{align*}
\text{C / C++} & : \quad \text{int} \ omp\_pause\_resource\_all(\text{omp\_pause\_resource\_t} \ kind); \\
\text{C / C++} & : \quad \text{integer function} \ omp\_pause\_resource\_all(\text{kind}) \\
\text{Fortran} & : \quad \text{integer (kind=omp\_pause\_resource\_kind)} \ kind
\end{align*}
\]

Binding
The binding task set for an omp_pause_resource_all region is the whole program.

Effect
The omp_pause_resource_all routine allows the runtime to relinquish resources used by OpenMP on all devices. It is equivalent to calling the omp_pause_resource routine once for each available device, including the host device.

The argument kind passed to this routine can be one of the valid OpenMP pause kind as defined in Section 3.6.1, or any implementation-specific pause kind.

Tool Callbacks
If the tool is not allowed to interact with a given device after encountering this call, then the runtime must call the tool finalizer for that device.
Restrictions
Restrictions to the `omp_pause_resource_all` routine are as follows:

- The `omp_pause_resource_all` region may not be nested in any explicit OpenMP region.
- The routine may only be called when all explicit tasks have finalized execution.

Cross References
- `target` construct, see Section 2.14.5.
- Declare target directive, see Section 2.14.7.
- To pause resources on a specific device only, see Section 3.6.1.

3.7 Device Information Routines
This section describes routines that pertain to the set of devices that are accessible to an OpenMP program.

3.7.1 omp_get_num_procs
Summary
The `omp_get_num_procs` routine returns the number of processors available to the device.

Format

```c
int omp_get_num_procs(void);
```

Effect
The `omp_get_num_procs` routine returns the number of processors that are available to the device at the time the routine is called. This value may change between the time that it is determined by the `omp_get_num_procs` routine and the time that it is read in the calling context due to system actions outside the control of the OpenMP implementation.
Cross References

• `omp_get_num_places` routine, see Section 3.3.2.
• `omp_get_place_num_procs` routine, see Section 3.3.3.
• `omp_get_place_proc_ids` routine, see Section 3.3.4.
• `omp_get_place_num` routine, see Section 3.3.5.

3.7.2 `omp_set_default_device`

Summary

The `omp_set_default_device` routine controls the default target device by assigning the value of the default-device-var ICV.

Format

```c
void omp_set_default_device(int device_num);
```

```fortran
subroutine omp_set_default_device(device_num)
  integer device_num
```

Binding

The binding task set for an `omp_set_default_device` region is the generating task.

Effect

The effect of this routine is to set the value of the default-device-var ICV of the current task to the value specified in the argument. When called from within a `target` region the effect of this routine is unspecified.

Cross References

• default-device-var, see Section 2.4.
• target construct, see Section 2.14.5.
• `omp_get_default_device`, see Section 3.7.3.
• OMP_DEFAULT_DEVICE environment variable, see Section 6.15.

3.7.3 `omp_get_default_device`

Summary

The `omp_get_default_device` routine returns the default target device.
3.7.4 omp_get_default_device

Summary
The `omp_get_default_device` routine returns the value of the `default-device-var` ICV of the current task. When called from within a `target` region the effect of this routine is unspecified.

Cross References
- `default-device-var`, see Section 2.4.
- `target` construct, see Section 2.14.5.
- `omp_set_default_device`, see Section 3.7.2.
- `OMP_DEFAULT_DEVICE` environment variable, see Section 6.15.

3.7.4 omp_get_num_devices

Summary
The `omp_get_num_devices` routine returns the number of non-host devices available for offloading code or data.

Binding
The binding task set for an `omp_get_num_devices` region is the generating task.
The `omp_get_num_devices` routine returns the number of available non-host devices onto which code or data may be offloaded. When called from within a `target` region the effect of this routine is unspecified.

**Cross References**
- `target` construct, see Section 2.14.5.
- `omp_get_default_device`, see Section 3.7.3.
- `omp_get_device_num`, see Section 3.7.5.

### 3.7.5 `omp_get_device_num`

**Summary**
The `omp_get_device_num` routine returns the device number of the device on which the calling thread is executing.

**Format**

```
C / C++
int omp_get_device_num(void);
```

```
Fortran
integer function omp_get_device_num()
```

**Binding**
The binding task set for an `omp_get_device_num` region is the generating task.

**Effect**
The `omp_get_device_num` routine returns the device number of the device on which the calling thread is executing. When called on the host device, it will return the same value as the `omp_get_initial_device` routine.

**Cross References**
- `target` construct, see Section 2.14.5.
- `omp_get_default_device`, see Section 3.7.3.
- `omp_get_num_devices`, see Section 3.7.4.
- `omp_get_initial_device` routine, see Section 3.7.7.
3.7.6 omp_is_initial_device

Summary
The `omp_is_initial_device` routine returns `true` if the current task is executing on the host device; otherwise, it returns `false`.

Format

```c
int omp_is_initial_device(void);
```

```fortran
logical function omp_is_initial_device()
```

Binding
The binding task set for an `omp_is_initial_device` region is the generating task.

Effect
The effect of this routine is to return `true` if the current task is executing on the host device; otherwise, it returns `false`.

Cross References
- `omp_get_initial_device` routine, see Section 3.7.7.
- Device memory routines, see Section 3.8.

3.7.7 omp_get_initial_device

Summary
The `omp_get_initial_device` routine returns a device number that represents the host device.

Format

```c
int omp_get_initial_device(void);
```

```fortran
integer function omp_get_initial_device()
```

Binding
The binding task set for an `omp_get_initial_device` region is the generating task.
Effect
The effect of this routine is to return the device number of the host device. The value of the device
number is the value returned by the `omp_get_num_devices` routine. When called from within
a `target` region the effect of this routine is unspecified.

Cross References
- `target` construct, see Section 2.14.5.
- `omp_is_initial_device` routine, see Section 3.7.6.
- Device memory routines, see Section 3.8.

3.8 Device Memory Routines

This section describes routines that support allocation of memory and management of pointers in
the data environments of target devices.

3.8.1 `omp_target_alloc`

Summary
The `omp_target_alloc` routine allocates memory in a device data environment and returns a
device pointer to that memory.

Format

<table>
<thead>
<tr>
<th>C / C++</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>void* omp_target_alloc(size_t size, int device_num);</code></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>C / C++</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>type(c_ptr) function omp_target_alloc(size, device_num) bind(c)</code></td>
</tr>
<tr>
<td>use, intrinsic :: iso_c_binding, only : c_ptr, c_size_t, c_int</td>
</tr>
<tr>
<td>integer(c_size_t), value :: size</td>
</tr>
<tr>
<td>integer(c_int), value :: device_num</td>
</tr>
</tbody>
</table>

Constraints on Arguments
The `device_num` argument must be greater than or equal to zero and less than or equal to the result
of `omp_get_num_devices()`.

Binding
The binding task set for an `omp_target_alloc` region is the generating task, which is the `target`
task generated by the call to the `omp_target_alloc` routine.
Effect
The `omp_target_alloc` routine returns a device pointer that references the device address of a storage location of size bytes. The storage location is dynamically allocated in the device data environment of the device specified by `device_num`.

The `omp_target_alloc` routine executes as if part of a target task that is generated by the call to the routine and that is an included task.

The `omp_target_alloc` routine returns NULL (or, `C_NULL_PTR`, for Fortran) if it cannot dynamically allocate the memory in the device data environment.

The device pointer returned by `omp_target_alloc` can be used in an `is_device_ptr` clause, Section 2.14.5.

<table>
<thead>
<tr>
<th>Fortran</th>
</tr>
</thead>
<tbody>
<tr>
<td>The <code>omp_target_alloc</code> routine requires an explicit interface and so might not be provided in <code>omp_lib.h</code>.</td>
</tr>
</tbody>
</table>

 Execution Model Events
The `target-data-allocation-begin` event occurs before a thread initiates a data allocation on a target device.

The `target-data-allocation-end` event occurs after a thread initiates a data allocation on a target device.

Tool Callbacks
A thread dispatches a registered `ompt_callback_target_data_op_emi` callback with `ompt_scope_begin` as its endpoint argument for each occurrence of a `target-data-allocation-begin` event in that thread. Similarly, a thread dispatches a registered `ompt_callback_target_data_op_emi` callback with `ompt_scope_end` as its endpoint argument for each occurrence of a `target-data-allocation-end` event in that thread. These callbacks have type signature `ompt_callback_target_data_op_emi_t`.

A thread dispatches a registered `ompt_callback_target_data_op` callback for each occurrence of a `target-data-allocation-begin` event in that thread. The callback occurs in the context of the target task and has type signature `ompt_callback_target_data_op_t`.

Restrictions
Restrictions to the `omp_target_alloc` routine are as follows.

- Freeing the storage returned by `omp_target_alloc` with any routine other than `omp_target_free` results in unspecified behavior.

- When called from within a `target` region the effect is unspecified.
• Unless the `unified_address` clause appears on a `requires` directive in the compilation unit, pointer arithmetic is not supported on the device pointer returned by `omp_target_alloc`.

Cross References

• `target` construct, see Section 2.14.5.
• `omp_get_num_devices` routine, see Section 3.7.4.
• `omp_target_free` routine, see Section 3.8.2.
• `ompt_callback_target_data_op_t` or `ompt_callback_target_data_op_emi_t` callback type, see Section 4.5.2.25.

3.8.2 omp_target_free

Summary

The `omp_target_free` routine frees the device memory allocated by the `omp_target_alloc` routine.

Format

```c++
void omp_target_free(void *device_ptr, int device_num);
```

```fortran
subroutine omp_target_free(device_ptr, device_num) bind(c)
use, intrinsic :: iso_c_binding, only : c_ptr, c_int
   type(c_ptr), value :: device_ptr
   integer(c_int), value :: device_num
```

Constraints on Arguments

A program that calls `omp_target_free` with a non-null pointer that does not have a value returned from `omp_target_alloc` is non-conforming. The `device_num` argument must be greater than or equal to zero and less than or equal to the result of `omp_get_num_devices()`.

Binding

The binding task set for an `omp_target_free` region is the generating task, which is the `target` task generated by the call to the `omp_target_free` routine.
Effect

The `omp_target_free` routine frees the memory in the device data environment associated with `device_ptr`. If `device_ptr` is `NULL` (or `C_NULL_PTR`, for Fortran), the operation is ignored.

The `omp_target_free` routine executes as if part of a target task that is generated by the call to the routine and that is an included task.

Synchronization must be inserted to ensure that all accesses to `device_ptr` are completed before the call to `omp_target_free`.

---

**Fortran**

The `omp_target_free` routine requires an explicit interface and so might not be provided in `omp_lib.h`.

---

Execution Model Events

The `target-data-free-begin` event occurs before a thread initiates a data free on a target device.

The `target-data-free-end` event occurs after a thread initiates a data free on a target device.

Tool Callbacks

A thread dispatches a registered `ompt_callback_target_data_op_eml` callback with `ompt_scope_begin` as its endpoint argument for each occurrence of a `target-data-free-begin` event in that thread. Similarly, a thread dispatches a registered `ompt_callback_target_data_op_eml` callback with `ompt_scope_end` as its endpoint argument for each occurrence of a `target-data-free-end` event in that thread. These callbacks have type signature `ompt_callback_target_data_op_eml_t`.

A thread dispatches a registered `ompt_callback_target_data_op` callback for each occurrence of a `target-data-free-begin` event in that thread. The callback occurs in the context of the target task and has type signature `ompt_callback_target_data_op_t`.

Restrictions

Restrictions to the `omp_target_free` routine are as follows.

- When called from within a `target` region the effect is unspecified.

Cross References

- `target` construct, see Section 2.14.5.
- `omp_get_num_devices` routine, see Section 3.7.4.
- `omp_target_alloc` routine, see Section 3.8.1.
- `ompt_callback_target_data_op_t` or `ompt_callback_target_data_op_eml_t` callback type, see Section 4.5.2.25.
3.8.3 omp_target_is_present

Summary
The **omp_target_is_present** routine tests whether a host pointer refers to storage that is
mapped to a given device.

Format

```
C / C++
int omp_target_is_present(const void *ptr, int device_num);
```

```
C / C++
Fortran
integer(c_int) function omp_target_is_present(ptr, device_num) &
   bind(c)
use, intrinsic :: iso_c_binding, only : c_ptr, c_int
type(c_ptr), value :: ptr
integer(c_int), value :: device_num
```

Constraints on Arguments
The value of *ptr* must be a valid host pointer or **NULL** (or **C_NULL_PTR**, for Fortran). The
*device_num* argument must be greater than or equal to zero and less than or equal to the result of
**omp_get_num_devices**().

Binding
The binding task set for an **omp_target_is_present** region is the encountering task.

Effect
The **omp_target_is_present** routine returns **true** if *device_num* refers to the host device or
if *ptr* refers to storage that has corresponding storage in the device data environment of device
*device_num*. Otherwise, the routine returns **false**.

Restrictions
Restrictions to the **omp_target_is_present** routine are as follows.

- When called from within a **target** region the effect is unspecified.

Cross References
- **target** construct, see Section 2.14.5.
- **map** clause, see Section 2.21.7.1.
- **omp_get_num_devices** routine, see Section 3.7.4.
3.8.4 omp_target_is_accessible

Summary
The **omp_target_is_accessible** routine tests whether host memory is accessible from a given device.

Format

**C / C++**

```c
int omp_target_is_accessible( const void *ptr, size_t size, int device_num);
```

**C / C++**

```c
integer(c_int) function omp_target_is_accessible( &ptr, size, device_num) bind(c)
```

**Fortran**

```fortran
use, intrinsic :: iso_c_binding, only : c_ptr, c_size_t, c_int
```

```fortran
type(c_ptr), value :: ptr
```

```fortran
integer(c_size_t), value :: size
```

```fortran
integer(c_int), value :: device_num
```

Constraints on Arguments

The value of *ptr* must be a valid host pointer or **NULL** (or **C_NULL_PTR**, for Fortran). The *device_num* argument must be greater than or equal to zero and less than or equal to the result of **omp_get_num_devices**().

Binding

The binding task set for an **omp_target_is_accessible** region is the encountering task.

Effect

This routine returns **true** if the storage of *size* bytes starting at the address given by *ptr* is accessible from device *device_num*. Otherwise, it returns **false**.

Restrictions

Restrictions to the **omp_target_is_accessible** routine are as follows.

- When called from within a **target** region the effect is unspecified.

Cross References

- **target** construct, see Section 2.14.5.
- **omp_get_num_devices** routine, see Section 3.7.4.
3.8.5 omp_target_memcpy

Summary
The omp_target_memcpy routine copies memory between any combination of host and device
pointers.

Format

C / C++

```c
int omp_target_memcpy(
    void *dst,
    const void *src,
    size_t length,
    size_t dst_offset,
    size_t src_offset,
    int dst_device_num,
    int src_device_num
);
```

C / C++

Fortran

```fortran
integer(c_int) function omp_target_memcpy(dst, src, length, &
    dst_offset, src_offset, dst_device_num, src_device_num) bind(c)
use, intrinsic :: iso_c_binding, only : c_ptr, c_int, c_size_t
    type(c_ptr), value :: dst, src
    integer(c_size_t), value :: length, dst_offset, src_offset
    integer(c_int), value :: dst_device_num, src_device_num
```

Constraints on Arguments
Each device pointer specified must be valid for the device on the same side of the copy. The
dst_device_num and src_device_num arguments must be greater than or equal to zero and less than
or equal to the result of omp_get_num_devices().

Binding
The binding task set for an omp_target_memcpy region is the generating task, which is the
target task generated by the call to the omp_target_memcpy routine.

Effect
This routine copies length bytes of memory at offset src_offset from src in the device data
environment of device src_device_num to dst starting at offset dst_offset in the device data
environment of device dst_device_num.

The omp_target_memcpy routine executes as if part of a target task that is generated by the call
to the routine and that is an included task.

The return value is zero on success and non-zero on failure. This routine contains a task scheduling
point.
Fortran

The `omp_target_memcpy` routine requires an explicit interface and so might not be provided in `omp_lib.h`.

Execution Model Events

The `target-data-op-begin` event occurs before a thread initiates a data transfer.

The `target-data-op-end` event occurs after a thread initiated a data transfer.

Tool Callbacks

A thread dispatches a registered `ompt_callback_target_data_op_emi` callback with `ompt_scope_begin` as its endpoint argument for each occurrence of a `target-data-op-begin` event in that thread. Similarly, a thread dispatches a registered `ompt_callback_target_data_op_emi` callback with `ompt_scope_end` as its endpoint argument for each occurrence of a `target-data-op-end` event in that thread. These callbacks have type signature `ompt_callback_target_data_op_emi_t`.

A thread dispatches a registered `ompt_callback_target_data_op` callback for each occurrence of a `target-data-op-end` event in that thread. The callback occurs in the context of the target task and has type signature `ompt_callback_target_data_op_t`.

Restrictions

Restrictions to the `omp_target_memcpy` routine are as follows.

- When called from within a `target` region the effect is unspecified.

Cross References

- `target` construct, see Section 2.14.5.
- `omp_get_num_devices` routine, see Section 3.7.4.
- `ompt_callback_target_data_op_t` or `ompt_callback_target_data_op_emi_t` callback type, see Section 4.5.2.25.

3.8.6 `omp_target_memcpy_rect`

Summary

The `omp_target_memcpy_rect` routine copies a rectangular subvolume from a multi-dimensional array to another multi-dimensional array. The `omp_target_memcpy_rect` routine performs a copy between any combination of host and device pointers.
```
int omp_target_memcpy_rect(
  void *dst,
  const void *src,
  size_t element_size,
  int num_dims,
  const size_t *volume,
  const size_t *dst_offsets,
  const size_t *src_offsets,
  const size_t *dst_dimensions,
  const size_t *src_dimensions,
  int dst_device_num,
  int src_device_num
);
```

```fortran
integer(c_int) function omp_target_memcpy_rect(dst, src, element_size, &
  num_dims, volume, dst_offsets, src_offsets, dst_dimensions, src_dimensions, &
  dst_device_num, src_device_num) bind(c)
use, intrinsic :: iso_c_binding, only : c_ptr, c_int, c_size_t
type(c_ptr), value :: dst, src
integer(c_size_t), value :: element_size
integer(c_int), value :: num_dims, dst_device_num, src_device_num
integer(c_size_t), intent(in) :: volume(*), dst_offsets(*), &
  src_offsets(*), dst_dimensions(*), src_dimensions(*)
```

**Constraints on Arguments**

Each device pointer specified must be valid for the device on the same side of the copy. The `dst_device_num` and `src_device_num` arguments must be greater than or equal to zero and less than or equal to the result of `omp_get_num_devices()`.

The length of the offset and dimension arrays must be at least the value of `num_dims`. The value of `num_dims` must be between 1 and the implementation-defined limit, which must be at least three.

Because the interface binds directly to a C language function the function assumes C memory ordering.

**Binding**

The binding task set for an `omp_target_memcpy_rect` region is the generating task, which is the target task generated by the call to the `omp_target_memcpy_rect` routine.
Effect
This routine copies a rectangular subvolume of src, in the device data environment of device src_device_num, to dst, in the device data environment of device dst_device_num. The volume is specified in terms of the size of an element, number of dimensions, and constant arrays of length num_dims. The maximum number of dimensions supported is at least three; support for higher dimensionality is implementation defined. The volume array specifies the length, in number of elements, to copy in each dimension from src to dst. The dst_offsets (src_offsets) parameter specifies the number of elements from the origin of dst (src) in elements. The dst_dimensions (src_dimensions) parameter specifies the length of each dimension of dst (src).

The omp_target_memcpy_rect routine executes as if part of a target task that is generated by the call to the routine and that is an included task.

The routine returns zero if successful. Otherwise, it returns a non-zero value. The routine contains a task scheduling point.

An application can determine the inclusive number of dimensions supported by an implementation by passing NULL pointers for both dst and src. The routine returns the number of dimensions supported by the implementation for the specified device numbers. No copy operation is performed.

Fortran
The omp_target_memcpy_rect routine requires an explicit interface and so might not be provided in omp_lib.h.

Execution Model Events
The target-data-op-begin event occurs before a thread initiates a data transfer.

The target-data-op-end event occurs after a thread initiated a data transfer.

Tool Callbacks
A thread dispatches a registered ompt_callback_target_data_op_emi callback with ompt_scope_begin as its endpoint argument for each occurrence of a target-data-op-begin event in that thread. Similarly, a thread dispatches a registered ompt_callback_target_data_op_emi callback with ompt_scope_end as its endpoint argument for each occurrence of a target-data-op-end event in that thread. These callbacks have type signature ompt_callback_target_data_op_emi_t.

A thread dispatches a registered ompt_callback_target_data_op callback for each occurrence of a target-data-op-end event in that thread. The callback occurs in the context of the target task and has type signature ompt_callback_target_data_op_t.

Restrictions
Restrictions to the omp_target_memcpy_rect routine are as follows.

• When called from within a target region the effect is unspecified.
Cross References
• target construct, see Section 2.14.5.
• omp_get_num_devices routine, see Section 3.7.4.
• ompt_callback_target_data_op_t or ompt_callback_target_data_op_emi_t callback type, see Section 4.5.2.25.

3.8.7 omp_target_memcpyAsync

Summary
The omp_target_memcpyAsync routine asynchronously performs a copy between any combination of host and device pointers.

Format

```
int omp_target_memcpyAsync(
    void *dst,
    const void *src,
    size_t length,
    size_t dst_offset,
    size_t src_offset,
    int dst_device_num,
    int src_device_num,
    int depobj_count,
    omp_depend_t *depobj_list
);
```

Constraints on Arguments
Each device pointer specified must be valid for the device on the same side of the copy. The dst_device_num and src_device_num arguments must be greater than or equal to zero and less than or equal to the result of omp_get_num_devices().
Binding
The binding task set for an \texttt{omp\_target\_memcpy\_async} region is the generating task, which
is the \textit{target task} generated by the call to the \texttt{omp\_target\_memcpy\_async} routine.

Effect
This routine performs an asynchronous memory copy where \textit{length} bytes of memory at offset
\texttt{src\_offset} from \texttt{src} in the device data environment of device \texttt{src\_device\_num} are copied to \texttt{dst}
starting at offset \texttt{dst\_offset} in the device data environment of device \texttt{dst\_device\_num}.

The \texttt{omp\_target\_memcpy\_async} routine executes as if part of a target task that is generated
by the call to the routine and for which execution may be deferred.

Task dependences are expressed with zero or more \texttt{omp\_depend\_t} objects. The dependences are
specified by passing the number of \texttt{omp\_depend\_t} objects followed by an array of
\texttt{omp\_depend\_t} objects. The generated target task is not a dependent task if the program passes
in a count of zero for \texttt{depobj\_count}. \texttt{depobj\_list} is ignored if the value of \texttt{depobj\_count} is zero.

The routine returns zero if successful. Otherwise, it returns a non-zero value. The routine contains
a task scheduling point.

\begin{Verbatim}
Fortran
\end{Verbatim}

The \texttt{omp\_target\_memcpy\_async} routine requires an explicit interface and so might not be
provided in \texttt{omp\_lib.h}.

\begin{Verbatim}
Fortran
\end{Verbatim}

Execution Model Events
The \textit{target-data-op-begin} event occurs before a thread initiates a data transfer.

The \textit{target-data-op-end} event occurs after a thread initiated a data transfer.

Tool Callbacks
A thread dispatches a registered \texttt{ompt\_callback\_target\_data\_op\_emi} callback with
\texttt{ompt\_scope\_begin} as its endpoint argument for each occurrence of a \textit{target-data-op-begin}
event in that thread. Similarly, a thread dispatches a registered
\texttt{ompt\_callback\_target\_data\_op\_emi} callback with \texttt{ompt\_scope\_end} as its endpoint
argument for each occurrence of a \textit{target-data-op-end} event in that thread. These callbacks have
type signature \texttt{ompt\_callback\_target\_data\_op\_emi\_t}.

A thread dispatches a registered \texttt{ompt\_callback\_target\_data\_op} callback for each
occurrence of a \textit{target-data-op-end} event in that thread. The callback occurs in the context of the
target task and has type signature \texttt{ompt\_callback\_target\_data\_op\_t}.

Restrictions
Restrictions to the \texttt{omp\_target\_memcpy\_async} routine are as follows.

- When called from within a \texttt{target} region the effect is unspecified.
Cross References

- target construct, see Section 2.14.5.
- Depend objects, see Section 2.19.10.
- omp_get_num_devices routine, see Section 3.7.4.
- omp_callback_target_data_op_t or ompt_callback_target_data_op_emi_t callback type, see Section 4.5.2.25.

3.8.8 omp_target_memcpy_rect_async

Summary

The omp_target_memcpy_rect_async routine asynchronously performs a copy between any combination of host and device pointers.

Format

```c
int omp_target_memcpy_rect_async(
    void *dst,
    const void *src,
    size_t element_size,
    int num_dims,
    const size_t *volume,
    const size_t *dstOffsets,
    const size_t *srcOffsets,
    const size_t *dstDimensions,
    const size_t *srcDimensions,
    int dstDeviceNum,
    int srcDeviceNum,
    int depobj_count,
    omp_depend_t *depobj_list
);```

```fortran
integer(c_int) function omp_target_memcpy_rect_async(dst, src, &
    element_size, num_dims, volume, dst_offsets, src_offsets, &
    dst_dimensions, src_dimensions, dst_device_num, src_device_num, &
    depobj_count, depobj_list) bind(c)
use, intrinsic :: iso_c_binding, only : c_ptr, c_int, c_size_t
  type(c_ptr), value :: dst, src
  integer(c_size_t), value :: element_size
  integer(c_int), value :: num_dims, dst_device_num, src_device_num, &
    depobj_count
```
integer(c_size_t), intent(in) :: volume(*), dst_offsets(*), &
src_offsets(*), dst_dimensions(*), src_dimensions(*)
integer(omp_depobj_kind), optional :: depobj_list(*)

Constraints on Arguments
Each device pointer specified must be valid for the device on the same side of the copy. The
dst_device_num and src_device_num arguments must be greater than or equal to zero and less than
or equal to the result of omp_get_num_devices().

The length of the offset and dimension arrays must be at least the value of num_dims. The value of
num_dims must be between 1 and the implementation-defined limit, which must be at least three.

Because the interface binds directly to a C language function the function assumes C memory
ordering.

Binding
The binding task set for an omp_target_memcpy_rect_async region is the generating task,
which is the target task generated by the call to the omp_target_memcpy_rect_async
routine.

Effect
This routine copies a rectangular subvolume of src, in the device data environment of device
src_device_num, to dst, in the device data environment of device dst_device_num. The volume is
specified in terms of the size of an element, number of dimensions, and constant arrays of length
num_dims. The maximum number of dimensions supported is at least three; support for higher
dimensionality is implementation defined. The volume array specifies the length, in number of
elements, to copy in each dimension from src to dst. The dst_offsets (src_offsets) parameter
specifies the number of elements from the origin of dst (src) in elements. The dst_dimensions
(src_dimensions) parameter specifies the length of each dimension of dst (src).

The omp_target_memcpy_rect_async routine executes as if part of a target task that is
generated by the call to the routine and for which execution may be deferred.

Task dependences are expressed with zero or more omp_depend_t objects. The dependences are
specified by passing the number of omp_depend_t objects followed by an array of
omp_depend_t objects. The generated target task is not a dependent task if the program passes
in a count of zero for depobj_count. depobj_list is ignored if the value of depobj_count is zero.

The routine returns zero if successful. Otherwise, it returns a non-zero value. The routine contains
a task scheduling point.

An application can determine the number of inclusive dimensions supported by an implementation
by passing NULL pointers (or C_NULL_PTR, for Fortran) for both dst and src. The routine returns
the number of dimensions supported by the implementation for the specified device numbers. No copy operation is performed.

```fortran
Fortran
```

The `omp_target_memcpy_rect_async` routine requires an explicit interface and so might not be provided in `omp_lib.h`.

```
Fortran
```

**Execution Model Events**

The `target-data-op-begin` event occurs before a thread initiates a data transfer.

The `target-data-op-end` event occurs after a thread initiated a data transfer.

**Tool Callbacks**

A thread dispatches a registered `ompt_callback_target_data_op_emi` callback with `ompt_scope_begin` as its endpoint argument for each occurrence of a `target-data-op-begin` event in that thread. Similarly, a thread dispatches a registered `ompt_callback_target_data_op_emi` callback with `ompt_scope_end` as its endpoint argument for each occurrence of a `target-data-op-end` event in that thread. These callbacks have type signature `ompt_callback_target_data_op_emi_t`.

A thread dispatches a registered `ompt_callback_target_data_op` callback for each occurrence of a `target-data-op-end` event in that thread. The callback occurs in the context of the target task and has type signature `ompt_callback_target_data_op_t`.

**Restrictions**

Restrictions to the `omp_target_memcpy_rect_async` routine are as follows.

- When called from within a `target` region the effect is unspecified.

**Cross References**

- `target` construct, see Section 2.14.5.
- `Depend objects`, see Section 2.19.10.
- `omp_get_num_devices` routine, see Section 3.7.4.
- `ompt_callback_target_data_op_t` or `ompt_callback_target_data_op_emi_t` callback type, see Section 4.5.2.25.

**3.8.9 omp_target_associate_ptr**

**Summary**

The `omp_target_associate_ptr` routine maps a device pointer, which may be returned from `omp_target_alloc` or implementation-defined runtime routines, to a host pointer.
Format

C / C++

```c
int omp_target_associate_ptr(
    const void *host_ptr,
    const void *device_ptr,
    size_t size,
    size_t device_offset,
    int device_num
);
```

C / C++

```fortran
integer(c_int) function omp_target_associate_ptr(host_ptr, &
    device_ptr, size, device_offset, device_num) bind(c)
    use, intrinsic :: iso_c_binding, only : c_ptr, c_size_t, c_int
    type(c_ptr), value :: host_ptr, device_ptr
    integer(c_size_t), value :: size, device_offset
    integer(c_int), value :: device_num
```

Constraints on Arguments

The value of `device_ptr` value must be a valid pointer to device memory for the device denoted by the value of `device_num`. The `device_num` argument must be greater than or equal to zero and less than or equal to the result of `omp_get_num_devices()`.

Binding

The binding task set for an `omp_target_associate_ptr` region is the generating task, which is the target task generated by the call to the `omp_target_associate_ptr` routine.

Effect

The `omp_target_associate_ptr` routine associates a device pointer in the device data environment of device `device_num` with a host pointer such that when the host pointer appears in a subsequent `map` clause, the associated device pointer is used as the target for data motion associated with that host pointer. The `device_offset` parameter specifies the offset into `device_ptr` that is used as the base address for the device side of the mapping. The reference count of the resulting mapping will be infinite. After being successfully associated, the buffer to which the device pointer points is invalidated and accessing data directly through the device pointer results in unspecified behavior. The pointer can be retrieved for other uses by using the `omp_target_disassociate_ptr` routine to disassociate it.

The `omp_target_associate_ptr` routine executes as if part of a target task that is generated by the call to the routine and that is an included task.

The routine returns zero if successful. Otherwise it returns a non-zero value.
Only one device buffer can be associated with a given host pointer value and device number pair. Attempting to associate a second buffer will return non-zero. Associating the same pair of pointers on the same device with the same offset has no effect and returns zero. Associating pointers that share underlying storage will result in unspecified behavior. The `omp_target_is_present` function can be used to test whether a given host pointer has a corresponding variable in the device data environment.

The `omp_target_associate_ptr` routine requires an explicit interface and so might not be provided in `omp_lib.h`.

**Execution Model Events**

The `target-data-associate` event occurs before a thread initiates a device pointer association on a target device.

**Tool Callbacks**

A thread dispatches a registered `ompt_callback_target_data_op` callback, or a registered `ompt_callback_target_data_op_emi` callback with `ompt_scope_beginend` as its endpoint argument for each occurrence of a `target-data-associate` event in that thread. These callbacks have type signature `ompt_callback_target_data_op_t` or `ompt_callback_target_data_op_emi_t`, respectively.

**Restrictions**

Restrictions to the `omp_target_associate_ptr` routine are as follows.

- When called from within a `target` region the effect is unspecified.

**Cross References**

- `target` construct, see Section 2.14.5.
- `map` clause, see Section 2.21.7.1.
- `omp_get_num_devices` routine, see Section 3.7.4.
- `omp_target_alloc` routine, see Section 3.8.1.
- `omp_target_is_present` routine, see Section 3.8.3.
- `omp_target_disassociate_ptr` routine, see Section 3.8.10.
- `omp_get_mapped_ptr` routine, see Section 3.8.11.
- `ompt_callback_target_data_op_t` or `ompt_callback_target_data_op_emi_t` callback type, see Section 4.5.2.25.
3.8.10 omp_target_disassociate_ptr

Summary
The `omp_target_disassociate_ptr` removes the associated pointer for a given device from a host pointer.

Format

```c
C / C++

int omp_target_disassociate_ptr(const void *ptr, int device_num);
```

```fortran
C / C++

integer(c_int) function omp_target_disassociate_ptr(ptr, &device_num) bind(c)
use, intrinsic :: iso_c_binding, only : c_ptr, c_int
type(c_ptr), value :: ptr
integer(c_int), value :: device_num
```

Constraints on Arguments
The `device_num` argument must be greater than or equal to zero and less than or equal to the result of `omp_get_num_devices()`.

Binding
The binding task set for an `omp_target_disassociate_ptr` region is the generating task, which is the `target task` generated by the call to the `omp_target_disassociate_ptr` routine.

Effect
The `omp_target_disassociate_ptr` removes the associated device data on device `device_num` from the presence table for host pointer `ptr`. A call to this routine on a pointer that is not `NULL` (or `C_NULL_PTR`, for Fortran) and does not have associated data on the given device results in unspecified behavior. The reference count of the mapping is reduced to zero, regardless of its current value.

The `omp_target_disassociate_ptr` routine executes as if part of a target task that is generated by the call to the routine and that is an included task.

The routine returns zero if successful. Otherwise it returns a non-zero value.

After a call to `omp_target_disassociate_ptr`, the contents of the device buffer are invalidated.

Fortran

The `omp_target_disassociate_ptr` routine requires an explicit interface and so might not be provided in `omp_lib.h`.

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Execution Model Events
The target-data-disassociate event occurs before a thread initiates a device pointer disassociation on a target device.

Tool Callbacks
A thread dispatches a registered ompt_callback_target_data_op callback, or a registered ompt_callback_target_data_op_emi callback with ompt_scope_beginend as its endpoint argument for each occurrence of a target-data-disassociate event in that thread. These callbacks have type signature ompt_callback_target_data_op_t or ompt_callback_target_data_op_emi_t, respectively.

Restrictions
Restrictions to the omp_target_disassociate_ptr routine are as follows.

- When called from within a target region the effect is unspecified.

Cross References
- target construct, see Section 2.14.5.
- omp_get_num_devices routine, see Section 3.7.4.
- omp_target_associate_ptr routine, see Section 3.8.9.
- ompt_callback_target_data_op_t or ompt_callback_target_data_op_emi_t callback type, see Section 4.5.2.25.

3.8.11 omp_get_mapped_ptr

Summary
The omp_get_mapped_ptr routine returns the device pointer that is associated with a host pointer for a given device.

Format
C / C++

void * omp_get_mapped_ptr(const void *ptr, int device_num);

C / C++

Fortran

type(c_ptr) function omp_get_mapped_ptr(ptr, &device_num) bind(c)
use, intrinsic :: iso_c_binding, only : c_ptr, c_int
type(c_ptr), value :: ptr
integer(c_int), value :: device_num

Fortran
Constraints on Arguments

The `device_num` argument must be greater than or equal to zero and less than or equal to the result of `omp_get_num_devices()`.

Binding

The binding task set for an `omp_get_mapped_ptr` region is the encountering task.

Effect

The `omp_get_mapped_ptr` routine returns the associated device pointer on device `device_num`. A call to this routine for a pointer that is not `NULL` (or `C_NULL_PTR`, for Fortran) and does not have an associated pointer on the given device results in a `NULL` pointer. The routine returns `NULL` (or `C_NULL_PTR`, for Fortran) if unsuccessful. Otherwise it returns the device pointer, which is `ptr` if `device_num` is the value returned by `omp_get_initial_device()`.

Fortran

The `omp_get_mapped_ptr` routine requires an explicit interface and so might not be provided in `omp_lib.h`.

Execution Model Events

No events are associated with this routine.

Restrictions

Restrictions to the `omp_get_mapped_ptr` routine are as follows.

- When called from within a `target` region the effect is unspecified.

Cross References

- `omp_get_num_devices` routine, see Section 3.7.4.
- `omp_get_initial_device` routine, see Section 3.7.7.
3.9 Lock Routines

The OpenMP runtime library includes a set of general-purpose lock routines that can be used for synchronization. These general-purpose lock routines operate on OpenMP locks that are represented by OpenMP lock variables. OpenMP lock variables must be accessed only through the routines described in this section; programs that otherwise access OpenMP lock variables are non-conforming.

An OpenMP lock can be in one of the following states: uninitialized; unlocked; or locked. If a lock is in the unlocked state, a task can set the lock, which changes its state to locked. The task that sets the lock is then said to own the lock. A task that owns a lock can unset that lock, returning it to the unlocked state. A program in which a task unsets a lock that is owned by another task is non-conforming.

Two types of locks are supported: simple locks and nestable locks. A nestable lock can be set multiple times by the same task before being unset; a simple lock cannot be set if it is already owned by the task trying to set it. Simple lock variables are associated with simple locks and can only be passed to simple lock routines. Nestable lock variables are associated with nestable locks and can only be passed to nestable lock routines.

Each type of lock can also have a synchronization hint that contains information about the intended usage of the lock by the application code. The effect of the hint is implementation defined. An OpenMP implementation can use this hint to select a usage-specific lock, but hints do not change the mutual exclusion semantics of locks. A conforming implementation can safely ignore the hint.

Constraints on the state and ownership of the lock accessed by each of the lock routines are described with the routine. If these constraints are not met, the behavior of the routine is unspecified.

The OpenMP lock routines access a lock variable such that they always read and update the most current value of the lock variable. An OpenMP program does not need to include explicit flush directives to ensure that the lock variable’s value is consistent among different tasks.

Binding

The binding thread set for all lock routine regions is all threads in the contention group. As a consequence, for each OpenMP lock, the lock routine effects relate to all tasks that call the routines, without regard to which teams in the contention group the threads that are executing the tasks belong.

Simple Lock Routines

The type `omp_lock_t` represents a simple lock. For the following routines, a simple lock variable must be of `omp_lock_t` type. All simple lock routines require an argument that is a pointer to a variable of type `omp_lock_t`. 
For the following routines, a simple lock variable must be an integer variable of kind=omp_lock_kind.

The simple lock routines are as follows:

- The omp_init_lock routine initializes a simple lock;
- The omp_init_lock_with_hint routine initializes a simple lock and attaches a hint to it;
- The omp_destroy_lock routine uninitializes a simple lock;
- The omp_set_lock routine waits until a simple lock is available and then sets it;
- The omp_unset_lock routine unsets a simple lock; and
- The omp_test_lock routine tests a simple lock and sets it if it is available.

Nestable Lock Routines

The type omp_nest_lock_t represents a nestable lock. For the following routines, a nestable lock variable must be of omp_nest_lock_t type. All nestable lock routines require an argument that is a pointer to a variable of type omp_nest_lock_t.

For the following routines, a nestable lock variable must be an integer variable of kind=omp_nest_lock_kind.

The nestable lock routines are as follows:

- The omp_init_nest_lock routine initializes a nestable lock;
- The omp_init_nest_lock_with_hint routine initializes a nestable lock and attaches a hint to it;
- The omp_destroy_nest_lock routine uninitializes a nestable lock;
- The omp_set_nest_lock routine waits until a nestable lock is available and then sets it;
- The omp_unset_nest_lock routine unsets a nestable lock; and
- The omp_test_nest_lock routine tests a nestable lock and sets it if it is available.

Restrictions

Restrictions to OpenMP lock routines are as follows:

- The use of the same OpenMP lock in different contention groups results in unspecified behavior.
3.9.1 **omp_init_lock** and **omp_init_nest_lock**

**Summary**
These routines initialize an OpenMP lock without a hint.

**Format**

```c
void omp_init_lock(omp_lock_t *lock);
void omp_init_nest_lock(omp_nest_lock_t *lock);
```

```fortran
subroutine omp_init_lock(svar)
    integer (kind=omp_lock_kind) svar
subroutine omp_init_nest_lock(nvar)
    integer (kind=omp_nest_lock_kind) nvar
```

**Constraints on Arguments**
A program that accesses a lock that is not in the uninitialized state through either routine is non-conforming.

**Effect**
The effect of these routines is to initialize the lock to the unlocked state; that is, no task owns the lock. In addition, the nesting count for a nestable lock is set to zero.

**Execution Model Events**
The `lock-init` event occurs in a thread that executes an `omp_init_lock` region after initialization of the lock, but before it finishes the region. The `nest-lock-init` event occurs in a thread that executes an `omp_init_nest_lock` region after initialization of the lock, but before it finishes the region.

**Tool Callbacks**
A thread dispatches a registered `ompt_callback_lock_init` callback with `omp_sync_hint_none` as the `hint` argument and `ompt_mutex_lock` as the `kind` argument for each occurrence of a `lock-init` event in that thread. Similarly, a thread dispatches a registered `ompt_callback_lock_init` callback with `omp_sync_hint_none` as the `hint` argument and `ompt_mutex_nest_lock` as the `kind` argument for each occurrence of a `nest-lock-init` event in that thread. These callbacks have the type signature `ompt_callback_mutex_acquire_t` and occur in the task that encounters the routine.

**Cross References**
- `ompt_callback_mutex_acquire_t`, see Section 4.5.2.14.
3.9.2 omp_init_lock_with_hint and omp_init_nest_lock_with_hint

Summary
These routines initialize an OpenMP lock with a hint. The effect of the hint is implementation-defined. The OpenMP implementation can ignore the hint without changing program semantics.

Format

```c
void omp_init_lock_with_hint(
    omp_lock_t *lock,
    omp_sync_hint_t hint
);

void omp_init_nest_lock_with_hint(
    omp_nest_lock_t *lock,
    omp_sync_hint_t hint
);
```

```fortran
subroutine omp_init_lock_with_hint(svar, hint)
    integer (kind=omp_lock_kind) svar
    integer (kind=omp_sync_hint_kind) hint
  end subroutine

subroutine omp_init_nest_lock_with_hint(nvar, hint)
    integer (kind=omp_nest_lock_kind) nvar
    integer (kind=omp_sync_hint_kind) hint
  end subroutine
```

Constraints on Arguments
A program that accesses a lock that is not in the uninitialized state through either routine is non-conforming.

The second argument passed to these routines (hint) is a hint as described in Section 2.19.12.

Effect
The effect of these routines is to initialize the lock to the unlocked state and, optionally, to choose a specific lock implementation based on the hint. After initialization no task owns the lock. In addition, the nesting count for a nestable lock is set to zero.
Execution Model Events
The `lock-init-with-hint` event occurs in a thread that executes an `omp_init_lock_with_hint` region after initialization of the lock, but before it finishes the region. The `nest-lock-init-with-hint` event occurs in a thread that executes an `omp_init_nest_lock` region after initialization of the lock, but before it finishes the region.

Tool Callbacks
A thread dispatches a registered `ompt_callback_lock_init` callback with the same value for its `hint` argument as the `hint` argument of the call to `omp_init_lock_with_hint` and `ompt_mutex_lock` as the `kind` argument for each occurrence of a `lock-init-with-hint` event in that thread. Similarly, a thread dispatches a registered `ompt_callback_lock_init` callback with the same value for its `hint` argument as the `hint` argument of the call to `omp_init_nest_lock_with_hint` and `ompt_mutex_nest_lock` as the `kind` argument for each occurrence of a `nest-lock-init-with-hint` event in that thread. These callbacks have the type signature `ompt_callback_mutex_acquire_t` and occur in the task that encounters the routine.

Cross References
- Synchronization Hints, see Section 2.19.12.
- `ompt_callback_mutex_acquire_t`, see Section 4.5.2.14.

3.9.3 `omp_destroy_lock` and `omp_destroy_nest_lock` Summary
These routines ensure that the OpenMP lock is uninitialized.

Format

```
C / C++

void omp_destroy_lock(omp_lock_t *lock);

void omp_destroy_nest_lock(omp_nest_lock_t *lock);

C / C++

Fortran

subroutine omp_destroy_lock(svar)
integer (kind=omp_lock_kind) svar

subroutine omp_destroy_nest_lock(nvar)
integer (kind=omp_nest_lock_kind) nvar
```

Fortran
Constraints on Arguments
A program that accesses a lock that is not in the unlocked state through either routine is non-conforming.

Effect
The effect of these routines is to change the state of the lock to uninitialized.

Execution Model Events
The lock-destroy event occurs in a thread that executes an `omp_destroy_lock` region before it finishes the region. The nest-lock-destroy event occurs in a thread that executes an `omp_destroy_nest_lock` region before it finishes the region.

Tool Callbacks
A thread dispatches a registered `ompt_callback_lock_destroy` callback with `ompt_mutex_lock` as the kind argument for each occurrence of a lock-destroy event in that thread. Similarly, a thread dispatches a registered `ompt_callback_lock_destroy` callback with `ompt_mutex_nest_lock` as the kind argument for each occurrence of a nest-lock-destroy event in that thread. These callbacks have the type signature `ompt_callback_mutex_t` and occur in the task that encounters the routine.

Cross References
- `ompt_callback_mutex_t`, see Section 4.5.2.15.

3.9.4 `omp_set_lock` and `omp_set_nest_lock`

Summary
These routines provide a means of setting an OpenMP lock. The calling task region behaves as if it was suspended until the lock can be set by this task.

Format

<table>
<thead>
<tr>
<th>C / C++</th>
<th>Fortran</th>
</tr>
</thead>
<tbody>
<tr>
<td>void omp_set_lock(omp_lock_t *lock);</td>
<td>subroutine omp_set_lock(svar)</td>
</tr>
<tr>
<td>void omp_set_nest_lock(omp_nest_lock_t *lock);</td>
<td></td>
</tr>
<tr>
<td>subroutine omp_set_lock(svar)</td>
<td>integer (kind=omp_lock_kind) svar</td>
</tr>
<tr>
<td>subroutine omp_set_nest_lock(nvar)</td>
<td>integer (kind=omp_nest_lock_kind) nvar</td>
</tr>
</tbody>
</table>
Constraints on Arguments

A program that accesses a lock that is in the uninitialized state through either routine is non-conforming. A simple lock accessed by `omp_set_lock` that is in the locked state must not be owned by the task that contains the call or deadlock will result.

Effect

Each of these routines has an effect equivalent to suspension of the task that is executing the routine until the specified lock is available.

Note – The semantics of these routines is specified as if they serialize execution of the region guarded by the lock. However, implementations may implement them in other ways provided that the isolation properties are respected so that the actual execution delivers a result that could arise from some serialization.

A simple lock is available if it is unlocked. Ownership of the lock is granted to the task that executes the routine.

A nestable lock is available if it is unlocked or if it is already owned by the task that executes the routine. The task that executes the routine is granted, or retains, ownership of the lock, and the nesting count for the lock is incremented.

Execution Model Events

The lock-acquire event occurs in a thread that executes an `omp_set_lock` region before the associated lock is requested. The nest-lock-acquire event occurs in a thread that executes an `omp_set_nest_lock` region before the associated lock is requested.

The lock-acquired event occurs in a thread that executes an `omp_set_lock` region after it acquires the associated lock but before it finishes the region. The nest-lock-acquired event occurs in a thread that executes an `omp_set_nest_lock` region if the thread did not already own the lock, after it acquires the associated lock but before it finishes the region.

The nest-lock-owned event occurs in a thread when it already owns the lock and executes an `omp_set_nest_lock` region. The event occurs after the nesting count is incremented but before the thread finishes the region.

Tool Callbacks

A thread dispatches a registered `ompt_callback_mutex_acquire` callback for each occurrence of a lock-acquire or nest-lock-acquire event in that thread. This callback has the type signature `ompt_callback_mutex_acquire_t`.

A thread dispatches a registered `ompt_callback_mutex_acquired` callback for each occurrence of a lock-acquired or nest-lock-acquired event in that thread. This callback has the type signature `ompt_callback_mutex_t`.
A thread dispatches a registered `ompt_callback_nest_lock` callback with `ompt_scope_begin` as its `endpoint` argument for each occurrence of a `nest-lock-owned` event in that thread. This callback has the type signature `ompt_callback_nest_lock_t`.

The above callbacks occur in the task that encounters the lock function. The `kind` argument of these callbacks is `ompt_mutex_lock` when the events arise from an `omp_set_lock` region while it is `ompt_mutex_nest_lock` when the events arise from an `omp_set_nest_lock` region.

**Cross References**
- `ompt_callback_mutex_acquire_t`, see Section 4.5.2.14.
- `ompt_callback_mutex_t`, see Section 4.5.2.15.
- `ompt_callback_nest_lock_t`, see Section 4.5.2.16.

### 3.9.5 `omp_unset_lock` and `omp_unset_nest_lock`

**Summary**
These routines provide the means of unsetting an OpenMP lock.

**Format**

```c
C / C++
void omp_unset_lock(omp_lock_t *lock);
void omp_unset_nest_lock(omp_nest_lock_t *lock);
```

```fortran
C / C++
subroutine omp_unset_lock(svar)
integer (kind=omp_lock_kind) svar
```

```fortran
subroutine omp_unset_nest_lock(nvar)
integer (kind=omp_nest_lock_kind) nvar
```

**Constraints on Arguments**
A program that accesses a lock that is not in the locked state or that is not owned by the task that contains the call through either routine is non-conforming.

**Effect**
For a simple lock, the `omp_unset_lock` routine causes the lock to become unlocked.

For a nestable lock, the `omp_unset_nest_lock` routine decrements the nesting count, and causes the lock to become unlocked if the resulting nesting count is zero.

For either routine, if the lock becomes unlocked, and if one or more task regions were effectively suspended because the lock was unavailable, the effect is that one task is chosen and given ownership of the lock.
Execution Model Events

The lock-release event occurs in a thread that executes an `omp_unset_lock` region after it releases the associated lock but before it finishes the region. The nest-lock-release event occurs in a thread that executes an `omp_unset_nest_lock` region after it releases the associated lock but before it finishes the region.

The nest-lock-held event occurs in a thread that executes an `omp_unset_nest_lock` region before it finishes the region when the thread still owns the lock after the nesting count is decremented.

Tool Callbacks

A thread dispatches a registered `ompt_callback_mutex_released` callback with `ompt_mutex_lock` as the kind argument for each occurrence of a lock-release event in that thread. Similarly, a thread dispatches a registered `ompt_callback_mutex_released` callback with `ompt_mutex_nest_lock` as the kind argument for each occurrence of a nest-lock-release event in that thread. These callbacks have the type signature `ompt_callback_mutex_t` and occur in the task that encounters the routine.

A thread dispatches a registered `ompt_callback_nest_lock` callback with `ompt_scope_end` as its endpoint argument for each occurrence of a nest-lock-held event in that thread. This callback has the type signature `ompt_callback_nest_lock_t`.

Cross References

- `ompt_callback_mutex_t`, see Section 4.5.2.15.
- `ompt_callback_nest_lock_t`, see Section 4.5.2.16.

3.9.6 `omp_test_lock` and `omp_test_nest_lock`

Summary

These routines attempt to set an OpenMP lock but do not suspend execution of the task that executes the routine.

Format

```c
int omp_test_lock(omp_lock_t *lock);
int omp_test_nest_lock(omp_nest_lock_t *lock);
```

```fortran
logical function omp_test_lock(svar)
integer (kind=omp_lock_kind) svar

integer function omp_test_nest_lock(nvar)
integer (kind=omp_nest_lock_kind) nvar
```
Constraints on Arguments
A program that accesses a lock that is in the uninitialized state through either routine is non-conforming. The behavior is unspecified if a simple lock accessed by \texttt{omp_test_lock} is in the locked state and is owned by the task that contains the call.

Effect
These routines attempt to set a lock in the same manner as \texttt{omp_set_lock} and \texttt{omp_set_nest_lock}, except that they do not suspend execution of the task that executes the routine.

For a simple lock, the \texttt{omp_test_lock} routine returns \textit{true} if the lock is successfully set; otherwise, it returns \textit{false}.

For a nestable lock, the \texttt{omp_test_nest_lock} routine returns the new nesting count if the lock is successfully set; otherwise, it returns zero.

Execution Model Events
The \textit{lock-test} event occurs in a thread that executes an \texttt{omp_test_lock} region before the associated lock is tested. The \textit{nest-lock-test} event occurs in a thread that executes an \texttt{omp_test_nest_lock} region before the associated lock is tested.

The \textit{lock-test-acquired} event occurs in a thread that executes an \texttt{omp_test_lock} region before it finishes the region if the associated lock was acquired. The \textit{nest-lock-test-acquired} event occurs in a thread that executes an \texttt{omp_test_nest_lock} region before it finishes the region if the associated lock was acquired and the thread did not already own the lock.

The \textit{nest-lock-owned} event occurs in a thread that executes an \texttt{omp_test_nest_lock} region before it finishes the region after the nesting count is incremented if the thread already owned the lock.

Tool Callbacks
A thread dispatches a registered \texttt{ompt_callback_mutex_acquire} callback for each occurrence of a \textit{lock-test} or \textit{nest-lock-test} event in that thread. This callback has the type signature \texttt{ompt_callback_mutex_acquire_t}.

A thread dispatches a registered \texttt{ompt_callback_mutex_acquired} callback for each occurrence of a \textit{lock-test-acquired} or \textit{nest-lock-test-acquired} event in that thread. This callback has the type signature \texttt{ompt_callback_mutex_t}.

A thread dispatches a registered \texttt{ompt_callback_nest_lock} callback with \texttt{ompt_scope_begin} as its \texttt{endpoint} argument for each occurrence of a \textit{nest-lock-owned} event in that thread. This callback has the type signature \texttt{ompt_callback_nest_lock_t}.

The above callbacks occur in the task that encounters the lock function. The \texttt{kind} argument of these callbacks is \texttt{ompt_mutex_test_lock} when the events arise from an \texttt{omp_test_lock} region while it is \texttt{ompt_mutex_test_nest_lock} when the events arise from an \texttt{omp_test_nest_lock} region.
Cross References

- ompt_callback_mutex_acquire_t, see Section 4.5.2.14.
- ompt_callback_mutex_t, see Section 4.5.2.15.
- ompt_callback_nest_lock_t, see Section 4.5.2.16.

3.10 Timing Routines

This section describes routines that support a portable wall clock timer.

3.10.1 omp_get_wtime

Summary

The `omp_get_wtime` routine returns elapsed wall clock time in seconds.

Format

```c
double omp_get_wtime(void);
```

Effect

The `omp_get_wtime` routine returns a value equal to the elapsed wall clock time in seconds since some `time-in-the-past`. The actual `time-in-the-past` is arbitrary, but it is guaranteed not to change during the execution of the application program. The time returned is a _per-thread time_, so it is not required to be globally consistent across all threads that participate in an application.

3.10.2 omp_get_wtick

Summary

The `omp_get_wtick` routine returns the precision of the timer used by `omp_get_wtime`.

Format

```c
double omp_get_wtick(void);
```

Binding

The binding thread set for an `omp_get_wtime` region is the encountering thread. The routine’s return value is not guaranteed to be consistent across any set of threads.
Fortran

double precision function omp_get_wtick()

Binding
The binding thread set for an `omp_get_wtick` region is the encountering thread. The routine’s return value is not guaranteed to be consistent across any set of threads.

Effect
The `omp_get_wtick` routine returns a value equal to the number of seconds between successive clock ticks of the timer used by `omp_get_wtime`.

3.11 Event Routine

This section describes a routine that supports OpenMP event objects.

Binding
The binding thread set for all event routine regions is the encountering thread.

3.11.1 `omp_fulfill_event`

Summary
This routine fulfills and destroys an OpenMP event.

Format

```c
void omp_fulfill_event(omp_event_handle_t event);
```

```fortran
subroutine omp_fulfill_event(event)
  integer (kind=omp_event_handle_kind) event
```

Constraints on Arguments
A program that calls this routine on an event that was already fulfilled is non-conforming. A program that calls this routine with an event handle that was not created by the `detach` clause is non-conforming.

Effect
The effect of this routine is to fulfill the event associated with the event handle argument. The effect of fulfilling the event will depend on how the event was created. The event is destroyed and cannot be accessed after calling this routine, and the event handle becomes unassociated with any event.
Execution Model Events
The task-fulfill event occurs in a thread that executes an `omp_fulfill_event` region before the event is fulfilled if the OpenMP event object was created by a `detach` clause on a task.

Tool Callbacks
A thread dispatches a registered `ompt_callback_task_schedule` callback with `NULL` as its `next_task_data` argument while the argument `prior_task_data` binds to the detached task for each occurrence of a task-fulfill event. If the task-fulfill event occurs before the detached task finished the execution of the associated structured-block, the callback has `ompt_task_early_fulfill` as its `prior_task_status` argument; otherwise the callback has `ompt_task_late_fulfill` as its `prior_task_status` argument. This callback has type signature `ompt_callback_task_schedule_t`.

Cross References
- `detach` clause, see Section 2.12.1.
- `ompt_callback_task_schedule_t`, see Section 4.5.2.10.

3.12 Interoperability Routines
The interoperability routines provide mechanisms to inspect the properties associated with an `omp_interop_t` object. Such objects may be initialized, destroyed or otherwise used by an `interop` construct. Additionally, an `omp_interop_t` object can be initialized to `omp_interop_none`, which is defined to be zero. An `omp_interop_t` object may only be accessed or modified through OpenMP directives and API routines.

An `omp_interop_t` object can be copied without affecting, or copying, the underlying state. Destruction of an `omp_interop_t` object destroys the state to which all copies of the object refer.

OpenMP reserves all negative values for properties, as listed in Table 3.1; implementation-defined properties may use zero and positive values. The special property, `omp_ipr_first`, will always have the lowest property value which may change in future versions of this specification. Valid values and types for the properties that Table 3.1 lists are specified in the *OpenMP Additional Definitions* document or are implementation defined unless otherwise specified.

Table 3.2 lists the return codes used by routines that take an `int* ret_code` argument.

Binding
The binding task set for all interoperability routine regions is the generating task.
### TABLE 3.1: Required Values of the `omp_interop_property_t` enum Type

<table>
<thead>
<tr>
<th>enum name</th>
<th>contexts</th>
<th>name</th>
<th>property</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>omp_ipr_fr_id = -1</code></td>
<td>all</td>
<td><code>fr_id</code></td>
<td>An <code>intptr_t</code> value that represents the foreign runtime id of context</td>
</tr>
<tr>
<td><code>omp_ipr_fr_name = -2</code></td>
<td>all</td>
<td><code>fr_name</code></td>
<td>C string value that represents the foreign runtime name of context</td>
</tr>
<tr>
<td><code>omp_ipr_vendor = -3</code></td>
<td>all</td>
<td><code>vendor</code></td>
<td>An <code>intptr_t</code> that represents the vendor of context</td>
</tr>
<tr>
<td><code>omp_ipr_vendor_name = -4</code></td>
<td>all</td>
<td><code>vendor_name</code></td>
<td>C string value that represents the vendor of context</td>
</tr>
<tr>
<td><code>omp_ipr_device_num = -5</code></td>
<td>all</td>
<td><code>device_num</code></td>
<td>The OpenMP device ID for the device in the range 0 to ( \text{omp_get_num_devices}() ) inclusive</td>
</tr>
<tr>
<td><code>omp_ipr_platform = -6</code></td>
<td>target</td>
<td><code>platform</code></td>
<td>A foreign platform handle usually spanning multiple devices</td>
</tr>
<tr>
<td><code>omp_ipr_device = -7</code></td>
<td>target</td>
<td><code>device</code></td>
<td>A foreign device handle</td>
</tr>
<tr>
<td><code>omp_ipr_device_context = -8</code></td>
<td>target</td>
<td><code>device_context</code></td>
<td>A handle to an instance of a foreign device context</td>
</tr>
<tr>
<td><code>omp_ipr_targetsync = -9</code></td>
<td>target</td>
<td><code>targetsync</code></td>
<td>A handle to a synchronization object of a foreign execution context</td>
</tr>
<tr>
<td><code>omp_ipr_first = -9</code></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
TABLE 3.2: Required Values for the `omp_interop_rc_t` enum Type

<table>
<thead>
<tr>
<th>enum name</th>
<th>description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>omp_irc_no_value</code> = 1</td>
<td>Parameters valid, no meaningful value available</td>
</tr>
<tr>
<td><code>omp_irc_success</code> = 0</td>
<td>Successful, value is usable</td>
</tr>
<tr>
<td><code>omp_irc_empty</code> = -1</td>
<td>The object provided is equal to <code>omp_interop_none</code></td>
</tr>
<tr>
<td><code>omp_irc_out_of_range</code> = -2</td>
<td>Property ID is out of range, see Table 3.1</td>
</tr>
<tr>
<td><code>omp_irc_type_int</code> = -3</td>
<td>Property type is int; use <code>omp_get_interop_int</code></td>
</tr>
<tr>
<td><code>omp_irc_type_ptr</code> = -4</td>
<td>Property type is pointer; use <code>omp_get_interop_ptr</code></td>
</tr>
<tr>
<td><code>omp_irc_type_str</code> = -5</td>
<td>Property type is string; use <code>omp_get_interop_str</code></td>
</tr>
<tr>
<td><code>omp_irc_other</code> = -6</td>
<td>Other error; use <code>omp_get_interop_rc_desc</code></td>
</tr>
</tbody>
</table>

3.12.1 `omp_get_num_interop_properties`

Summary
The `omp_get_num_interop_properties` routine retrieves the number of implementation-defined properties available for an `omp_interop_t` object.

Format
```c
int omp_get_num_interop_properties(const omp_interop_t interop);
```

Effect
The `omp_get_num_interop_properties` routine returns the number of implementation-defined properties available for `interop`. The total number of properties available for `interop` is the returned value minus `omp_ipr_first`.

Cross References
- `interop` construct, see Section 2.15.1.

3.12.2 `omp_get_interop_int`

Summary
The `omp_get_interop_int` routine retrieves an integer property from an `omp_interop_t` object.

Format
```c
omp_intptr_t omp_get_interop_int(const omp_interop_t interop, omp_interop_property_t property_id, int *ret_code);
```
C/C++ (cont.)

Effect
The `omp_get_interop_int` routine returns the requested integer property, if available, and zero if an error occurs or no value is available.

If the `interop` is `omp_interop_none`, an empty error occurs.

If the `property_id` is smaller than `omp_ipr_first` or not smaller than `omp_get_num_interop_properties(interop)`, an out of range error occurs.

If the requested property value is not convertible into an integer value, a type error occurs.

If a non-null pointer is passed to `ret_code`, an `omp_interop_rc_t` value that indicates the return code is stored in the object to which `ret_code` points. If an error occurred, the stored value will be negative and it will match the error as defined in Table 3.2. On success, zero will be stored. If no error occurred but no meaningful value can be returned, `omp irc_no_value`, which is one, will be stored.

Restrictions
Restrictions to the `omp_get_interop_int` routine are as follows:

- The behavior of the routine is unspecified if an invalid `omp_interop_t` object is provided.

Cross References
- `interop` construct, see Section 2.15.1.
- `omp_get_num_interop_properties` routine, see Section 3.12.1.

3.12.3 `omp_get_interop_ptr`

Summary
The `omp_get_interop_ptr` routine retrieves a pointer property from an `omp_interop_t` object.

Format
```c
void* omp_get_interop_ptr(const omp_interop_t interop,
                          omp_interop_property_t property_id,
                          int *ret_code);
```

Effect
The `omp_get_interop_ptr` routine returns the requested pointer property, if available, and `NULL` if an error occurs or no value is available.

If the `interop` is `omp_interop_none`, an empty error occurs.

If the `property_id` is smaller than `omp_ipr_first` or not smaller than `omp_get_num_interop_properties(interop)`, an out of range error occurs.

If the requested property value is not convertible into a pointer value, a type error occurs.
If a non-null pointer is passed to `ret_code`, an `omp_interop_rc_t` value that indicates the return code is stored in the object to which the `ret_code` points. If an error occurred, the stored value will be negative and it will match the error as defined in Table 3.2. On success, zero will be stored. If no error occurred but no meaningful value can be returned, `omp_irc_no_value`, which is one, will be stored.

**Restrictions**

Restrictions to the `omp_get_interop_ptr` routine are as follows:

- The behavior of the routine is unspecified if an invalid `omp_interop_t` object is provided.
- Memory referenced by the pointer returned by the `omp_get_interop_ptr` routine is managed by the OpenMP implementation and should not be freed or modified.

**Cross References**

- `interop` construct, see Section 2.15.1.
- `omp_get_num_interop_properties` routine, see Section 3.12.1.

### 3.12.4 `omp_get_interop_str`

**Summary**

The `omp_get_interop_str` routine retrieves a string property from an `omp_interop_t` object.

**Format**

```c
const char* omp_get_interop_str(const omp_interop_t interop,
                                omp_interop_property_t property_id,
                                int *ret_code);
```

**Effect**

The `omp_get_interop_str` routine returns the requested string property as a C string, if available, and `NULL` if an error occurs or no value is available.

If the `interop` is `omp_interop_none`, an empty error occurs.

If the `property_id` is smaller than `omp_ipr_first` or not smaller than `omp_get_num_interop_properties(interop)`, an out of range error occurs.

If the requested property value is not convertible into a string value, a type error occurs.

If a non-null pointer is passed to `ret_code`, an `omp_interop_rc_t` value that indicates the return code is stored in the object to which the `ret_code` points. If an error occurred, the stored value will be negative and it will match the error as defined in Table 3.2. On success, zero will be stored. If no error occurred but no meaningful value can be returned, `omp_irc_no_value`, which is one, will be stored.
Restrictions
Restrictions to the `omp_get_interop_str` routine are as follows:

- The behavior of the routine is unspecified if an invalid `omp_interop_t` object is provided.
- Memory referenced by the pointer returned by the `omp_get_interop_str` routine is managed by the OpenMP implementation and should not be freed or modified.

Cross References
- `interop` construct, see Section 2.15.1.
- `omp_get_num_interop_properties` routine, see Section 3.12.1.

### 3.12.5 `omp_get_interop_name`

Summary
The `omp_get_interop_name` routine retrieves a property name from an `omp_interop_t` object.

Format
```
const char* omp_get_interop_name(const omp_interop_t interop,
                                 omp_interop_property_t property_id)
```

Effect
The `omp_get_interop_name` routine returns the name of the property identified by `property_id` as a C string.

Property names for non-implementation defined properties are listed in Table 3.1.

If the `property_id` is smaller than `omp_ipr_first` or not smaller than `omp_get_num_interop_properties(interop)`, NULL is returned.

Restrictions
Restrictions to the `omp_get_interop_name` routine are as follows:

- The behavior of the routine is unspecified if an invalid object is provided.
- Memory referenced by the pointer returned by the `omp_get_interop_name` routine is managed by the OpenMP implementation and should not be freed or modified.

Cross References
- `interop` construct, see Section 2.15.1.
- `omp_get_num_interop_properties` routine, see Section 3.12.1.
3.12.6 omp_get_interop_type_desc

Summary
The `omp_get_interop_type_desc` routine retrieves a description of the type of a property associated with an `omp_interop_t` object.

Format
```
const char* omp_get_interop_type_desc(const omp_interop_t interop,
                                      omp_interop_property_t property_id) ;
```

Effect
The `omp_get_interop_type_desc` routine returns a C string that describes the type of the property identified by `property_id` in human-readable form. That may contain a valid C type declaration possibly followed by a description or name of the type.

If `interop` has the value `omp_interop_none`, `NULL` is returned.

If the `property_id` is smaller than `omp_ipr_first` or not smaller than `omp_get_num_interop_properties(interop)`, `NULL` is returned.

Restrictions
Restrictions to the `omp_get_interop_type_desc` routine are as follows:

- The behavior of the routine is unspecified if an invalid object is provided.
- Memory referenced by the pointer returned from the `omp_get_interop_type_desc` routine is managed by the OpenMP implementation and should not be freed or modified.

Cross References
- `interop` construct, see Section 2.15.1.
- `omp_get_num_interop_properties` routine, see Section 3.12.1.

3.12.7 omp_get_interop_rc_desc

Summary
The `omp_get_interop_rc_desc` routine retrieves a description of the return code associated with an `omp_interop_t` object.

Format
```
const char* omp_get_interop_rc_desc(const omp_interop_t interop,
                                     omp_interop_rc_t ret_code) ;
```

Effect
The `omp_get_interop_rc_desc` routine returns a C string that describes the return code `ret_code` in human-readable form.
Restrictions
Restrictions to the `omp_get_interop_rc_desc` routine are as follows:

- The behavior of the routine is unspecified if an invalid object is provided or if `ret_code` was not last written by an interoperability routine invoked with the `omp_interop_t` object `interop`.
- Memory referenced by the pointer returned by the `omp_get_interop_rc_desc` routine is managed by the OpenMP implementation and should not be freed or modified.

Cross References
- `interop` construct, see Section 2.15.1.
- `omp_get_num_interop_properties` routine, see Section 3.12.1.

3.13 Memory Management Routines

This section describes routines that support memory management on the current device.

Instances of memory management types must be accessed only through the routines described in this section; programs that otherwise access instances of these types are non-conforming.

3.13.1 Memory Management Types

The following type definitions are used by the memory management routines:

```c
typedef enum omp_alloctrait_key_t {
    omp_atk_sync_hint = 1,
    omp_atk_alignment = 2,
    omp_atk_access = 3,
    omp_atk_pool_size = 4,
    omp_atk_fallback = 5,
    omp_atk_fb_data = 6,
    omp_atk_pinned = 7,
    omp_atk_partition = 8
} omp_alloctrait_key_t;

typedef enum omp_alloctrait_value_t {
    omp_atv_false = 0,
    omp_atv_true = 1,
    omp_atv_contended = 3,
    omp_atv_uncontended = 4,
    omp_atv_serialized = 5,
} omp_alloctrait_value_t;
```
omp_atv_sequential = omp_atv_serialized, // (deprecated)
omp_atv_private = 6,
omp_atv_all = 7,
omp_atv_thread = 8,
omp_atv_pteam = 9,
omp_atv_cgroup = 10,
omp_atv_default_mem_fb = 11,
omp_atv_null_fb = 12,
omp_atv_abort_fb = 13,
omp_atv_allocator_fb = 14,
omp_atv_environment = 15,
omp_atv_nearest = 16,
omp_atv_blocked = 17,
omp_atv_interleaved = 18
}
omp_alloctrait_value_t;

typedef struct omp_alloctrait_t {
omp_alloctrait_key_t key;
omp_uintptr_t value;
} omp_alloctrait_t;

integer(kind=omp_alloctrait_key_kind), &
parameter :: omp_atk_sync_hint = 1
integer(kind=omp_alloctrait_key_kind), &
parameter :: omp_atk_alignment = 2
integer(kind=omp_alloctrait_key_kind), &
parameter :: omp_atk_access = 3
integer(kind=omp_alloctrait_key_kind), &
parameter :: omp_atk_pool_size = 4
integer(kind=omp_alloctrait_key_kind), &
parameter :: omp_atk_fallback = 5
integer(kind=omp_alloctrait_key_kind), &
parameter :: omp_atk_fb_data = 6
integer(kind=omp_alloctrait_key_kind), &
parameter :: omp_atk_pinned = 7
integer(kind=omp_alloctrait_key_kind), &
parameter :: omp_atk_partition = 8
integer(kind=omp_alloctrait_val_kind), &
parameter :: omp_atv_default = -1
integer(kind=omp_alloctrait_val_kind), &
Fortran (cont.)

```fortran
parameter :: omp_atv_false = 0
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_true = 1
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_contended = 3
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_uncontended = 4
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_serialized = 5
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_sequential = &
omp_atv_serialized ! (deprecated)
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_private = 6
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_all = 7
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_thread = 8
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_pteam = 9
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_cgroup = 10
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_default_mem_fb = 11
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_null_fb = 12
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_abort_fb = 13
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_allocator_fb = 14
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_environment = 15
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_nearest = 16
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_blocked = 17
integer(kind=omp_alloctrait_val_kind), &
  parameter :: omp_atv_interleaved = 18

! omp_alloctrait might not be provided in omp_lib.h.
type omp_alloctrait
  integer(kind=omp_alloctrait_key_kind) key
  integer(kind=omp_alloctrait_val_kind) value
```

CHAPTER 3. RUNTIME LIBRARY ROUTINES
3.13.2 omp_init_allocator

Summary
The `omp_init_allocator` routine initializes an allocator and associates it with a memory space.

Format

```c
omp_allocator_handle_t omp_init_allocator ( 
   omp_memspace_handle_t memspace, 
   int ntraits, 
   const omp_alloctrait_t traits[] 
);
```

### C / C++

```fortran
integer(kind=omp_allocator_handle_kind) & 
function omp_init_allocator ( memspace, ntraits, traits ) 
integer(kind=omp_memspace_handle_kind),intent(in) :: memspace 
integer,intent(in) :: ntraits 
type(omp_alloctrait),intent(in) :: traits(*)
```

Constraints on Arguments
The `memspace` argument must be one of the predefined memory spaces defined in Table 2.8.

If the `ntraits` argument is greater than zero then the `traits` argument must specify at least that many traits. If it specifies fewer than `ntraits` traits the behavior is unspecified.

Binding
The binding thread set for an `omp_init_allocator` region is all threads on a device. The effect of executing this routine is not related to any specific region that corresponds to any construct or API routine.
Effect
The `omp_init_allocator` routine creates a new allocator that is associated with the `memspace` memory space and returns a handle to it. All allocations through the created allocator will behave according to the allocator traits specified in the `traits` argument. The number of traits in the `traits` argument is specified by the `ntraits` argument. Specifying the same allocator trait more than once results in unspecified behavior. The routine returns a handle for the created allocator. If the special `omp_atv_default` value is used for a given trait, then its value will be the default value specified in Table 2.9 for that given trait.

If `memspace` is `omp_default_mem_space` and the `traits` argument is an empty set this routine will always return a handle to an allocator. Otherwise if an allocator based on the requirements cannot be created then the special `omp_null_allocator` handle is returned.

Restrictions
The restrictions to the `omp_init_allocator` routine are as follows:

- The use of an allocator returned by this routine on a device other than the one on which it was created results in unspecified behavior.
- Unless a `requires` directive with the `dynamicAllocators` clause is present in the same compilation unit, using this routine in a `target` region results in unspecified behavior.

Cross References
- Memory Spaces, see Section 2.13.1.
- Memory Allocators, see Section 2.13.2.

3.13.3 `omp_destroy_allocator`

Summary
The `omp_destroy_allocator` routine releases all resources used by the allocator handle.

Format

```
C / C++
void omp_destroy_allocator (omp_allocator_handle_t allocator);
```

```
FORTRAN
subroutine omp_destroy_allocator (allocator)
    integer(kind=omp_allocator_handle_kind),intent(in) :: allocator
```

Constraints on Arguments
The `allocator` argument must not represent a predefined memory allocator.
**Binding**
The binding thread set for an `omp_destroy_allocator` region is all threads on a device. The effect of executing this routine is not related to any specific region that corresponds to any construct or API routine.

**Effect**
The `omp_destroy_allocator` routine releases all resources used to implement the `allocator` handle.

If `allocator` is `omp_null_allocator` then this routine will have no effect.

**Restrictions**
The restrictions to the `omp_destroy_allocator` routine are as follows:

- Accessing any memory allocated by the `allocator` after this call results in unspecified behavior.
- Unless a `requires` directive with the `dynamic_allocators` clause is present in the same compilation unit, using this routine in a `target` region results in unspecified behavior.

**Cross References**
- Memory Allocators, see Section 2.13.2.

### 3.13.4 `omp_set_default_allocator`

**Summary**
The `omp_set_default_allocator` routine sets the default memory allocator to be used by allocation calls, `allocate` directives and `allocate` clauses that do not specify an allocator.

**Format**

```c
default_allocator (omp_allocator_handle_t allocator);
```

```fortran
subroutine omp_set_default_allocator ( allocator )
integer(kind=omp_allocator_handle_kind),intent(in) :: allocator
```

**Constraints on Arguments**
The `allocator` argument must be a valid memory allocator handle.

**Binding**
The binding task set for an `omp_set_default_allocator` region is the binding implicit task.
Effect
The effect of this routine is to set the value of the `def-allocator-var` ICV of the binding implicit task to the value specified in the `allocator` argument.

Cross References
- `def-allocator-var` ICV, see Section 2.4.
- Memory Allocators, see Section 2.13.2.
- `omp_alloc` routine, see Section 3.13.6.

### 3.13.5 omp_get_default_allocator

Summary
The `omp_get_default_allocator` routine returns a handle to the memory allocator to be used by allocation calls, `allocate` directives and `allocate` clauses that do not specify an allocator.

Format

C / C++

```
omp_allocator_handle_t omp_get_default_allocator (void);
```

Fortran

```
integer(kind=omp_allocator_handle_kind)&
function omp_get_default_allocator ()
```

Binding
The binding task set for an `omp_get_default_allocator` region is the binding implicit task.

Effect
The effect of this routine is to return the value of the `def-allocator-var` ICV of the binding implicit task.

Cross References
- `def-allocator-var` ICV, see Section 2.4.
- Memory Allocators, see Section 2.13.2.
- `omp_alloc` routine, see Section 3.13.6.
3.13.6 **omp_alloc** and **omp_aligned_alloc**

**Summary**

The `omp_alloc` and `omp_aligned_alloc` routines request a memory allocation from a memory allocator.

**Format**

```c
void *omp_alloc(size_t size, omp_allocator_handle_t allocator);

void *omp_aligned_alloc(
    size_t alignment,
    size_t size,
    omp_allocator_handle_t allocator);
```

```c++
void *omp_alloc(
    size_t size,
    omp_allocator_handle_t allocator=omp_null_allocator
);

void *omp_aligned_alloc(
    size_t alignment,
    size_t size,
    omp_allocator_handle_t allocator=omp_null_allocator
);
```

```fortran
use, intrinsic :: iso_c_binding, only : c_ptr, c_size_t
integer(c_size_t), value :: size
integer(omp_allocator_handle_kind), value :: allocator

type(c_ptr) function omp_alloc(size, allocator) bind(c)

use, intrinsic :: iso_c_binding, only : c_ptr, c_size_t
integer(c_size_t), value :: allocation, size
integer(omp_allocator_handle_kind), value :: allocator
```

`OpenMP API – Version 5.1 November 2020`
Constraints on Arguments

Unless `dynamic_allocators` appears on a `requires` directive in the same compilation unit, `omp_alloc` and `omp_aligned_alloc` invocations that appear in `target` regions must not pass `omp_null_allocator` as the `allocator` argument, which must be a constant expression that evaluates to one of the predefined memory allocator values.

The `alignment` argument to `omp_aligned_alloc` must be a power of two and the `size` argument must be a multiple of `alignment`.

Binding

The binding task set for an `omp_alloc` or `omp_aligned_alloc` region is the generating task.

Effect

The `omp_alloc` and `omp_aligned_alloc` routines request a memory allocation of `size` bytes from the specified memory allocator. If the `allocator` argument is `omp_null_allocator` the memory allocator used by the routines will be the one specified by the `def-allocator-var` ICV of the binding implicit task. Upon success they return a pointer to the allocated memory. Otherwise, the behavior that the `fallback` trait of the allocator specifies will be followed.

If `size` is 0, `omp_alloc` and `omp_aligned_alloc` will return `NULL` (or, `C_NULL_PTR`, for Fortran).

Memory allocated by `omp_alloc` will be byte-aligned to at least the maximum of the alignment required by `malloc` and the `alignment` trait of the allocator.

Memory allocated by `omp_aligned_alloc` will be byte-aligned to at least the maximum of the alignment required by `malloc`, the `alignment` trait of the allocator and the `alignment` argument value.

Cross References

- Memory allocators, see Section 2.13.2.

3.13.7 `omp_free`

Summary

The `omp_free` routine deallocates previously allocated memory.

Format

C

```c
void omp_free (void *ptr, omp_allocator_handle_t allocator);
```
void omp_free(
    void *ptr,
    omp_allocator_handle_t allocator=omp_null_allocator
);

subroutine omp_free(ptr, allocator) bind(c)  
use, intrinsic :: iso_c_binding, only : c_ptr  
type(c_ptr), value :: ptr  
integer(omp_allocator_handle_kind), value :: allocator

The binding task set for an *omp_free* region is the generating task.

**Effect**
The *omp_free* routine deallocates the memory to which *ptr* points. The *ptr* argument must have been returned by an OpenMP allocation routine. If the *allocator* argument is specified it must be the memory allocator to which the allocation request was made. If the *allocator* argument is *omp_null_allocator* the implementation will determine that value automatically.

If *ptr* is *NULL* (or, *C_NULL_PTR*, for Fortran), no operation is performed.

The *omp_free* routine requires an explicit interface and so might not be provided in *omp_lib.h*.

**Restrictions**
The restrictions to the *omp_free* routine are as follows:

- Using *omp_free* on memory that was already deallocated or that was allocated by an allocator that has already been destroyed with *omp_destroy_allocator* results in unspecified behavior.

**Cross References**
- Memory allocators, see Section 2.13.2.
3.13.8 **omp_calloc and omp_aligned_calloc**

**Summary**

The `omp_calloc` and `omp_aligned_calloc` routines request a zero initialized memory allocation from a memory allocator.

**Format**

```c
void *omp_calloc(
    size_t nmemb,
    size_t size,
    omp_allocator_handle_t allocator
);
void *omp_aligned_calloc(
    size_t alignment,
    size_t nmemb,
    size_t size,
    omp_allocator_handle_t allocator
);
```

```c++
void *omp_calloc(
    size_t nmemb,
    size_t size,
    omp_allocator_handle_t allocator=omp_null_allocator
);
void *omp_aligned_calloc(
    size_t alignment,
    size_t nmemb,
    size_t size,
    omp_allocator_handle_t allocator=omp_null_allocator
);
```
**Constraints on Arguments**

Unless `dynamic_allocators` appears on a `requires` directive in the same compilation unit, `omp_calloc` and `omp_aligned_calloc` invocations that appear in `target` regions must not pass `omp_null_allocator` as the `allocator` argument, which must be a constant expression that evaluates to one of the predefined memory allocator values.

The `alignment` argument to `omp_aligned_calloc` must be a power of two and the `size` argument must be a multiple of `alignment`.

**Binding**

The binding task set for an `omp_calloc` or `omp_aligned_calloc` region is the generating task.

**Effect**

The `omp_calloc` and `omp_aligned_calloc` routines request a memory allocation from the specified memory allocator for an array of `nmemb` elements each of which has a size of `size` bytes. If the `allocator` argument is `omp_null_allocator` the memory allocator used by the routines will be the one specified by the `def-allocator-var` ICV of the binding implicit task. Upon success they return a pointer to the allocated memory. Otherwise, the behavior that the `fallback` trait of the allocator specifies will be followed. Any memory allocated by these routines will be set to zero before returning.

If either `nmemb` or `size` is 0, `omp_calloc` will return `NULL` (or, `C_NULL_PTR`, for Fortran).

Memory allocated by `omp_calloc` will be byte-aligned to at least the maximum of the alignment required by `malloc` and the `alignment` trait of the allocator.

Memory allocated by `omp_aligned_calloc` will be byte-aligned to at least the maximum of the alignment required by `malloc`, the `alignment` trait of the allocator and the `alignment` argument value.
The `omp_calloc` and `omp_aligned_calloc` routines require an explicit interface and so might not be provided in `omp_lib.h`.

Cross References

- Memory allocators, see Section 2.13.2.

### 3.13.9 `omp_realloc`

**Summary**

The `omp_realloc` routine deallocates previously allocated memory and requests a memory allocation from a memory allocator.

**Format**

```c
void *omp_realloc(
    void *ptr,
    size_t size,
    omp_allocator_handle_t allocator,
    omp_allocator_handle_t free_allocator
);
```

```c++
void *omp_realloc(
    void *ptr,
    size_t size,
    omp_allocator_handle_t allocator=omp_null_allocator,
    omp_allocator_handle_t free_allocator=omp_null_allocator
);
```

```fortran
type(c_ptr) &
function omp_realloc(ptr, size, allocator, free_allocator) bind(c)
use, intrinsic :: iso_c_binding, only : c_ptr, c_size_t
type(c_ptr), value :: ptr
integer(c_size_t), value :: size
integer(omp_allocator_handle_kind), value :: allocator, free_allocator
```

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Constraints on Arguments

Unless a `dynamic_allocators` clause appears on a `requires` directive in the same compilation unit, `omp_realloc` invocations that appear in `target` regions must not pass `omp_null_allocator` as the `allocator` or `free_allocator` argument, which must be constant expressions that evaluate to one of the predefined memory allocator values.

Binding

The binding task set for an `omp_realloc` region is the generating task.

Effect

The `omp_realloc` routine deallocates the memory to which `ptr` points and requests a new memory allocation of `size` bytes from the specified memory allocator. If the `free_allocator` argument is specified, it must be the memory allocator to which the previous allocation request was made. If the `free_allocator` argument is `omp_null_allocator` the implementation will determine that value automatically. If the `allocator` argument is `omp_null_allocator` the behavior is as if the memory allocator that allocated the memory to which `ptr` argument points is passed to the `allocator` argument. Upon success it returns a (possibly moved) pointer to the allocated memory and the contents of the new object shall be the same as that of the old object prior to deallocation, up to the minimum size of old allocated size and `size`. Any bytes in the new object beyond the old allocated size will have unspecified values. If the allocation failed, the behavior that the `fallback` trait of the `allocator` specifies will be followed.

If `ptr` is `NULL` (or, `C_NULL_PTR`, for Fortran), `omp_realloc` will behave the same as `omp_alloc` with the same `size` and `allocator` arguments.

If `size` is 0, `omp_realloc` will return `NULL` (or, `C_NULL_PTR`, for Fortran) and the old allocation will be deallocated.

If `size` is not 0, the old allocation will be deallocated if and only if the function returns a non-`NULL` value (or, a non-`C_NULL_PTR` value, for Fortran).

Memory allocated by `omp_realloc` will be byte-aligned to at least the maximum of the alignment required by `malloc` and the `alignment` trait of the allocator.

```
                    Fortran
```

The `omp_realloc` routine requires an explicit interface and so might not be provided in `omp_lib.h`.

```
                    Fortran
```

Restrictions

The restrictions to the `omp_realloc` routine are as follows:

- The `ptr` argument must have been returned by an OpenMP allocation routine.

- Using `omp_realloc` on memory that was already deallocated or that was allocated by an allocator that has already been destroyed with `omp_destroy_allocator` results in unspecified behavior.
Cross References

- Memory allocators, see Section 2.13.2.

3.14 Tool Control Routine

Summary

The `omp_control_tool` routine enables a program to pass commands to an active tool.

Format

C / C++

```
int omp_control_tool(int command, int modifier, void *arg);
```

C / C++

```
integer function omp_control_tool(command, modifier)
integer (kind=omp_control_tool_kind) command
integer modifier
```

Constraints on Arguments

The following enumeration type defines four standard commands. Table 3.3 describes the actions that these commands request from a tool.

C / C++

```
typedef enum omp_control_tool_t {
    omp_control_tool_start = 1,
    omp_control_tool_pause = 2,
    omp_control_tool_flush = 3,
    omp_control_tool_end = 4
} omp_control_tool_t;
```

C / C++

```
integer (kind=omp_control_tool_kind), &
    parameter :: omp_control_tool_start = 1
integer (kind=omp_control_tool_kind), &
    parameter :: omp_control_tool_pause = 2
integer (kind=omp_control_tool_kind), &
    parameter :: omp_control_tool_flush = 3
integer (kind=omp_control_tool_kind), &
    parameter :: omp_control_tool_end = 4
```
Tool-specific values for command must be greater or equal to 64. Tools must ignore command values that they are not explicitly designed to handle. Other values accepted by a tool for command, and any values for modifier and arg are tool-defined.

**TABLE 3.3:** Standard Tool Control Commands

<table>
<thead>
<tr>
<th>Command</th>
<th>Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>omp_control_tool_start</td>
<td>Start or restart monitoring if it is off. If monitoring is already on, this command is idempotent. If monitoring has already been turned off permanently, this command will have no effect.</td>
</tr>
<tr>
<td>omp_control_tool_pause</td>
<td>Temporarily turn monitoring off. If monitoring is already off, it is idempotent.</td>
</tr>
<tr>
<td>omp_control_tool_flush</td>
<td>Flush any data buffered by a tool. This command may be applied whether monitoring is on or off.</td>
</tr>
<tr>
<td>omp_control_tool_end</td>
<td>Turn monitoring off permanently; the tool finalizes itself and flushes all output.</td>
</tr>
</tbody>
</table>

**Binding**

The binding task set for an omp_control_tool region is the generating task.

**Effect**

An OpenMP program may use omp_control_tool to pass commands to a tool. An application can use omp_control_tool to request that a tool starts or restarts data collection when a code region of interest is encountered, that a tool pauses data collection when leaving the region of interest, that a tool flushes any data that it has collected so far, or that a tool ends data collection. Additionally, omp_control_tool can be used to pass tool-specific commands to a particular tool.

The following types correspond to return values from omp_control_tool:

```c
typedef enum omp_control_tool_result_t {
    omp_control_tool_notool = -2,
    omp_control_tool_nocallback = -1,
    omp_control_tool_success = 0,
    omp_control_tool_ignored = 1
} omp_control_tool_result_t;
```
If the OMPT interface state is inactive, the OpenMP implementation returns `omp_control_tool_notool`. If the OMPT interface state is active, but no callback is registered for the `tool-control` event, the OpenMP implementation returns `omp_control_tool_nocallback`. An OpenMP implementation may return other implementation-defined negative values strictly smaller than -64; an application may assume that any negative return value indicates that a tool has not received the command. A return value of `omp_control_tool_success` indicates that the tool has performed the specified command. A return value of `omp_control_tool_ignored` indicates that the tool has ignored the specified command. A tool may return other positive values strictly greater than 64 that are tool-defined.

### Execution Model Events

The `tool-control` event occurs in the thread that encounters a call to `omp_control_tool` at a point inside its corresponding OpenMP region.

### Tool Callbacks

A thread dispatches a registered `ompt_callback_control_tool` callback for each occurrence of a `tool-control` event. The callback executes in the context of the call that occurs in the user program and has type signature `ompt_callback_control_tool_t`. The callback may return any non-negative value, which will be returned to the application by the OpenMP implementation as the return value of the `omp_control_tool` call that triggered the callback.

Arguments passed to the callback are those passed by the user to `omp_control_tool`. If the call is made in Fortran, the tool will be passed `NULL` as the third argument to the callback. If any of the four standard commands is presented to a tool, the tool will ignore the modifier and arg argument values.

### Restrictions

Restrictions on access to the state of an OpenMP first-party tool are as follows:

- An application may access the tool state modified by an OMPT callback only by using `omp_control_tool`.  

```fortran
integer (kind=omp_control_tool_result_kind), &
parameter :: omp_control_tool_notool = -2
integer (kind=omp_control_tool_result_kind), &
parameter :: omp_control_tool_nocallback = -1
integer (kind=omp_control_tool_result_kind), &
parameter :: omp_control_tool_success = 0
integer (kind=omp_control_tool_result_kind), &
parameter :: omp_control_tool_ignored = 1
```
Cross References

- OMPT Interface, see Chapter 4
- `ompt_callback_control_tool_t`, see Section 4.5.2.29.

3.15 Environment Display Routine

Summary

The `omp_display_env` routine displays the OpenMP version number and the initial values of ICVs associated with the environment variables described in Chapter 6.

Format

```c
void omp_display_env(int verbose);
```

```fortran
subroutine omp_display_env( verbose )
logical,intent(in) :: verbose
```

Binding

The binding thread set for an `omp_display_env` region is the encountering thread.

Effect

Each time the `omp_display_env` routine is invoked, the runtime system prints the OpenMP version number and the initial values of the ICVs associated with the environment variables described in Chapter 6. The displayed values are the values of the ICVs after they have been modified according to the environment variable settings and before the execution of any OpenMP construct or API routine.

The display begins with "OPENMP DISPLAY ENVIRONMENT BEGIN", followed by the `__OPENMP` version macro (or the `openmp_version` named constant for Fortran) and ICV values, in the format `NAME '=' VALUE`. `NAME` corresponds to the macro or environment variable name, optionally prepended with a bracketed `DEVICE`. `VALUE` corresponds to the value of the macro or ICV associated with this environment variable. Values are enclosed in single quotes. `DEVICE` corresponds to the device on which the value of the ICV is applied. The display is terminated with "OPENMP DISPLAY ENVIRONMENT END".

For the `OMP_NESTED` environment variable, the printed value is `true` if the `max-active-levels-var` ICV is initialized to a value greater than 1; otherwise the printed value is `false`. The `OMP_NESTED` environment variable has been deprecated.
If the `verbose` argument is set to 0 (or `false` in Fortran), the runtime displays the OpenMP version number defined by the `_OPENMP` version macro (or the `omp_version` named constant for Fortran) value and the initial ICV values for the environment variables listed in Chapter 6. If the `verbose` argument is set to 1 (or `true` for Fortran), the runtime may also display the values of vendor-specific ICVs that may be modified by vendor-specific environment variables.

Example output:

```
OPENMP DISPLAY ENVIRONMENT BEGIN
  _OPENMP='202011'
  [host] OMP_SCHEDULE='GUIDED,4'
  [host] OMP_NUM_THREADS='4,3,2'
  [device] OMP_NUM_THREADS='2'
  [host,device] OMP_DYNAMIC='TRUE'
  [host] OMP_PLACES='{0:4},{4:4},{8:4},{12:4}'
... OPENMP DISPLAY ENVIRONMENT END
```

Cross References

- `OMP_DISPLAY_ENV` environment variable, see Section 6.12.
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This chapter describes OMPT, which is an interface for first-party tools. First-party tools are linked or loaded directly into the OpenMP program. OMPT defines mechanisms to initialize a tool, to examine OpenMP state associated with an OpenMP thread, to interpret the call stack of an OpenMP thread, to receive notification about OpenMP events, to trace activity on OpenMP target devices, to assess implementation-dependent details of an OpenMP implementation (such as supported states and mutual exclusion implementations), and to control a tool from an OpenMP application.

### 4.1 OMPT Interfaces Definitions

A compliant implementation must supply a set of definitions for the OMPT runtime entry points, OMPT callback signatures, and the special data types of their parameters and return values. These definitions, which are listed throughout this chapter, and their associated declarations shall be provided in a header file named `omp-tools.h`. In addition, the set of definitions may specify other implementation-specific values.

The `ompt_start_tool` function is an external function with C linkage.

### 4.2 Activating a First-Party Tool

To activate a tool, an OpenMP implementation first determines whether the tool should be initialized. If so, the OpenMP implementation invokes the initializer of the tool, which enables the tool to prepare to monitor execution on the host. The tool may then also arrange to monitor computation that executes on target devices. This section explains how the tool and an OpenMP implementation interact to accomplish these tasks.

#### 4.2.1 ompt_start_tool

**Summary**

In order to use the OMPT interface provided by an OpenMP implementation, a tool must implement the `ompt_start_tool` function, through which the OpenMP implementation initializes the tool.
Format

```c
ompt_start_tool_result_t *ompt_start_tool(
    unsigned int omp_version,
    const char *runtime_version
);
```

Description

For a tool to use the OMPT interface that an OpenMP implementation provides, the tool must define a globally-visible implementation of the function `ompt_start_tool`. The tool indicates that it will use the OMPT interface that an OpenMP implementation provides by returning a non-null pointer to an `ompt_start_tool_result_t` structure from the `ompt_start_tool` implementation that it provides. The `ompt_start_tool_result_t` structure contains pointers to tool initialization and finalization callbacks as well as a tool data word that an OpenMP implementation must pass by reference to these callbacks. A tool may return `NULL` from `ompt_start_tool` to indicate that it will not use the OMPT interface in a particular execution.

A tool may use the `omp_version` argument to determine if it is compatible with the OMPT interface that the OpenMP implementation provides.

Description of Arguments

The argument `omp_version` is the value of the `_OPENMP` version macro associated with the OpenMP API implementation. This value identifies the OpenMP API version that an OpenMP implementation supports, which specifies the version of the OMPT interface that it supports.

The argument `runtime_version` is a version string that unambiguously identifies the OpenMP implementation.

Constraints on Arguments

The argument `runtime_version` must be an immutable string that is defined for the lifetime of a program execution.

Effect

If a tool returns a non-null pointer to an `ompt_start_tool_result_t` structure, an OpenMP implementation will call the tool initializer specified by the `initialize` field in this structure before beginning execution of any OpenMP construct or completing execution of any environment routine invocation; the OpenMP implementation will call the tool finalizer specified by the `finalize` field in this structure when the OpenMP implementation shuts down.

Cross References

- `ompt_start_tool_result_t`, see Section 4.4.1.
4.2.2 Determining Whether a First-Party Tool Should be Initialized

An OpenMP implementation examines the `tool-var` ICV as one of its first initialization steps. If the value of `tool-var` is disabled, the initialization continues without a check for the presence of a tool and the functionality of the OMPT interface will be unavailable as the program executes. In this case, the OMPT interface state remains inactive.

Otherwise, the OMPT interface state changes to pending and the OpenMP implementation activates any first-party tool that it finds. A tool can provide a definition of `ompt_start_tool` to an OpenMP implementation in three ways:

- By statically-linking its definition of `ompt_start_tool` into an OpenMP application;
- By introducing a dynamically-linked library that includes its definition of `ompt_start_tool` into the application’s address space; or
- By providing, in the `tool-libraries-var` ICV, the name of a dynamically-linked library that is appropriate for the architecture and operating system used by the application and that includes a definition of `ompt_start_tool`. 

![First-Party Tool Activation Flow Chart](image.png)
If the value of tool-var is enabled, the OpenMP implementation must check if a tool has provided an implementation of ompt_start_tool. The OpenMP implementation first checks if a tool-provided implementation of ompt_start_tool is available in the address space, either statically-linked into the application or in a dynamically-linked library loaded in the address space. If multiple implementations of ompt_start_tool are available, the OpenMP implementation will use the first tool-provided implementation of ompt_start_tool that it finds.

If the implementation does not find a tool-provided implementation of ompt_start_tool in the address space, it consults the tool-libraries-var ICV, which contains a (possibly empty) list of dynamically-linked libraries. As described in detail in Section 6.19, the libraries in tool-libraries-var are then searched for the first usable implementation of ompt_start_tool that one of the libraries in the list provides.

If the implementation finds a tool-provided definition of ompt_start_tool, it invokes that method; if a NULL pointer is returned, the OMPT interface state remains pending and the implementation continues to look for implementations of ompt_start_tool; otherwise a non-null pointer to an ompt_start_tool_result_t structure is returned, the OMPT interface state changes to active and the OpenMP implementation makes the OMPT interface available as the program executes. In this case, as the OpenMP implementation completes its initialization, it initializes the OMPT interface.

If no tool can be found, the OMPT interface state changes to inactive.

Cross References

• tool-libraries-var ICV, see Section 2.4.

• tool-var ICV, see Section 2.4.

• ompt_start_tool function, see Section 4.2.1.

• ompt_start_tool_result_t type, see Section 4.4.1.

4.2.3 Initializing a First-Party Tool

To initialize the OMPT interface, the OpenMP implementation invokes the tool initializer that is specified in the ompt_start_tool_result_t structure that is indicated by the non-null pointer that ompt_start_tool returns. The initializer is invoked prior to the occurrence of any OpenMP event.

A tool initializer, described in Section 4.5.1.1, uses the function specified in its lookup argument to look up pointers to OMPT interface runtime entry points that the OpenMP implementation provides; this process is described in Section 4.2.3.1. Typically, a tool initializer obtains a pointer to the ompt_set_callback runtime entry point with type signature ompt_set_callback_t and then uses this runtime entry point to register tool callbacks for OpenMP events, as described in Section 4.2.4.
A tool initializer may use the `ompt Enumerate States` runtime entry point, which has type signature `ompt Enumerate States_t`, to determine the thread states that an OpenMP implementation employs. Similarly, it may use the `ompt Enumerate Mutex Impls` runtime entry point, which has type signature `ompt Enumerate Mutex Impls_t`, to determine the mutual exclusion implementations that the OpenMP implementation employs.

If a tool initializer returns a non-zero value, the OMPT interface state remains `active` for the execution; otherwise, the OMPT interface state changes to `inactive`.

**Cross References**
- `ompt Start Tool` function, see Section 4.2.1.
- `ompt Start Tool Result_t` type, see Section 4.4.1.
- `ompt Initialize_t` type, see Section 4.5.1.1.
- `ompt Callback Thread Begin_t` type, see Section 4.5.2.1.
- `ompt Enumerate States_t` type, see Section 4.6.1.1.
- `ompt Enumerate Mutex Impls_t` type, see Section 4.6.1.2.
- `ompt Set Callback_t` type, see Section 4.6.1.3.
- `ompt Function Lookup_t` type, see Section 4.6.3.

### 4.2.3.1 Binding Entry Points in the OMPT Callback Interface

Functions that an OpenMP implementation provides to support the OMPT interface are not defined as global function symbols. Instead, they are defined as runtime entry points that a tool can only identify through the `lookup` function that is provided as an argument with type signature `ompt Function Lookup_t` to the tool initializer. A tool can use this function to obtain a pointer to each of the runtime entry points that an OpenMP implementation provides to support the OMPT interface. Once a tool has obtained a `lookup` function, it may employ it at any point in the future.

For each runtime entry point in the OMPT interface for the host device, Table 4.1 provides the string name by which it is known and its associated type signature. Implementations can provide additional implementation-specific names and corresponding entry points. Any names that begin with `ompt_` are reserved names.

During initialization, a tool should look up each runtime entry point in the OMPT interface by name and bind a pointer maintained by the tool that can later be used to invoke the entry point. The entry points described in Table 4.1 enable a tool to assess the thread states and mutual exclusion implementations that an OpenMP implementation supports to register tool callbacks, to inspect registered callbacks, to introspect OpenMP state associated with threads, and to use tracing to monitor computations that execute on target devices.
Detailed information about each runtime entry point listed in Table 4.1 is included as part of the description of its type signature.

**Cross References**
- `ompt Enumerate States T` type, see Section 4.6.1.1.
- `ompt Enumerate Mutex Impls T` type, see Section 4.6.1.2.
- `ompt Set Callback T` type, see Section 4.6.1.3.
- `ompt Get Callback T` type, see Section 4.6.1.4.
- `ompt Get Thread Data T` type, see Section 4.6.1.5.
- `ompt Get Num Procs T` type, see Section 4.6.1.6.
- `ompt Get Num Places T` type, see Section 4.6.1.7.
- `ompt Get Place Proc Ids T` type, see Section 4.6.1.8.
- `ompt Get Place Num T` type, see Section 4.6.1.9.
- `ompt Get Partition Place_nums T` type, see Section 4.6.1.10.
- `ompt Get Proc Id T` type, see Section 4.6.1.11.
- `ompt Get State T` type, see Section 4.6.1.12.
- `ompt Get Parallel Info T` type, see Section 4.6.1.13.
- `ompt Get Task Info T` type, see Section 4.6.1.14.
- `ompt Get Task Memory T` type, see Section 4.6.1.15.
- `ompt Get Target Info T` type, see Section 4.6.1.16.
- `ompt Get Num Devices T` type, see Section 4.6.1.17.
- `ompt Get Unique Id T` type, see Section 4.6.1.18.
- `ompt Finalize Tool T` type, see Section 4.6.1.19.
- `ompt Function Lookup T` type, see Section 4.6.3.

### 4.2.4 Monitoring Activity on the Host with OMPT

To monitor the execution of an OpenMP program on the host device, a tool initializer must register to receive notification of events that occur as an OpenMP program executes. A tool can use the `ompt_set_callback` runtime entry point to register callbacks for OpenMP events. The return codes for `ompt_set_callback` use the `ompt_set_result_t` enumeration type. If the `ompt_set_callback` runtime entry point is called outside a tool initializer, registration of supported callbacks may fail with a return value of `ompt_set_error`. 
**Table 4.1:** OMPT Callback Interface Runtime Entry Point Names and Their Type Signatures

<table>
<thead>
<tr>
<th>Entry Point String Name</th>
<th>Type signature</th>
</tr>
</thead>
<tbody>
<tr>
<td>“omptEnumerateStates”</td>
<td>omptEnumerateStates_t</td>
</tr>
<tr>
<td>“omptEnumerateMutexImpls”</td>
<td>omptEnumerateMutexImpls_t</td>
</tr>
<tr>
<td>“omptSetCallback”</td>
<td>omptSetCallback_t</td>
</tr>
<tr>
<td>“omptGetCallback”</td>
<td>omptGetCallback_t</td>
</tr>
<tr>
<td>“omptGetThreadData”</td>
<td>omptGetThreadData_t</td>
</tr>
<tr>
<td>“omptGetNumPlaces”</td>
<td>omptGetNumPlaces_t</td>
</tr>
<tr>
<td>“omptGetPlaceProcIds”</td>
<td>omptGetPlaceProcIds_t</td>
</tr>
<tr>
<td>“omptGetPlaceNum”</td>
<td>omptGetPlaceNum_t</td>
</tr>
<tr>
<td>“omptGetPartitionPlaceNums”</td>
<td>omptGetPartitionPlaceNums_t</td>
</tr>
<tr>
<td>“omptGetProcId”</td>
<td>omptGetProcId_t</td>
</tr>
<tr>
<td>“omptGetState”</td>
<td>omptGetState_t</td>
</tr>
<tr>
<td>“omptGetParallelInfo”</td>
<td>omptGetParallelInfo_t</td>
</tr>
<tr>
<td>“omptGetTaskInfo”</td>
<td>omptGetTaskInfo_t</td>
</tr>
<tr>
<td>“omptGetTaskMemory”</td>
<td>omptGetTaskMemory_t</td>
</tr>
<tr>
<td>“omptGetNumDevices”</td>
<td>omptGetNumDevices_t</td>
</tr>
<tr>
<td>“omptGetNumProcs”</td>
<td>omptGetNumProcs_t</td>
</tr>
<tr>
<td>“omptGetTargetInfo”</td>
<td>omptGetTargetInfo_t</td>
</tr>
<tr>
<td>“omptGetUniqueId”</td>
<td>omptGetUniqueId_t</td>
</tr>
<tr>
<td>“omptFinalizeTool”</td>
<td>omptFinalizeTool_t</td>
</tr>
</tbody>
</table>

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All callbacks registered with `ompt_set_callback` or returned by `ompt_get_callback` use
the dummy type signature `ompt_callback_t`.

For callbacks listed in Table 4.2, `ompt_set_always` is the only registration return code that is
allowed. An OpenMP implementation must guarantee that the callback will be invoked every time
that a runtime event that is associated with it occurs. Support for such callbacks is required in a
minimal implementation of the OMPT interface.

For callbacks listed in Table 4.3, the `ompt_set_callback` runtime entry may return any
non-error code. Whether an OpenMP implementation invokes a registered callback never,
sometimes, or always is implementation defined. If registration for a callback allows a return code
of `omp_set_never`, support for invoking such a callback may not be present in a minimal
implementation of the OMPT interface. The return code from registering a callback indicates the
implementation-defined level of support for the callback.

Two techniques reduce the size of the OMPT interface. First, in cases where events are naturally
paired, for example, the beginning and end of a region, and the arguments needed by the callback at
each endpoint are identical, a tool registers a single callback for the pair of events, with
`ompt_scope_begin` or `ompt_scope_end` provided as an argument to identify for which
endpoint the callback is invoked. Second, when a class of events is amenable to uniform treatment,
OMPT provides a single callback for that class of events, for example, an
`ompt_callback_sync_region_wait` callback is used for multiple kinds of synchronization
regions, such as barrier, taskwait, and taskgroup regions. Some events, for example,
`ompt_callback_sync_region_wait`, use both techniques.

Cross References
- `ompt_set_result_t` type, see Section 4.4.4.2.
- `ompt_set_callback_t` type, see Section 4.6.1.3.
- `ompt_get_callback_t` type, see Section 4.6.1.4.

4.2.5 Tracing Activity on Target Devices with OMPT

A target device may or may not initialize a full OpenMP runtime system. Unless it does,
monitoring activity on a device using a tool interface based on callbacks may not be possible. To
accommodate such cases, the OMPT interface defines a monitoring interface for tracing activity on
target devices. Tracing activity on a target device involves the following steps:

- To prepare to trace activity on a target device, a tool must register for an
  `ompt_callback_device_initialize` callback. A tool may also register for an
  `ompt_callback_device_load` callback to be notified when code is loaded onto a target
device or an `ompt_callback_device_unload` callback to be notified when code is
  unloaded from a target device. A tool may also optionally register an
  `ompt_callback_device_finalize` callback.
<table>
<thead>
<tr>
<th>Callback name</th>
</tr>
</thead>
<tbody>
<tr>
<td>ompt_callback_thread_begin</td>
</tr>
<tr>
<td>ompt_callback_thread_end</td>
</tr>
<tr>
<td>ompt_callback_parallel_begin</td>
</tr>
<tr>
<td>ompt_callback_parallel_end</td>
</tr>
<tr>
<td>ompt_callback_task_create</td>
</tr>
<tr>
<td>ompt_callback_task_schedule</td>
</tr>
<tr>
<td>ompt_callback_implicit_task</td>
</tr>
<tr>
<td>ompt_callback_target</td>
</tr>
<tr>
<td>ompt_callback_target_emi</td>
</tr>
<tr>
<td>ompt_callback_target_data_op</td>
</tr>
<tr>
<td>ompt_callback_target_data_op_emi</td>
</tr>
<tr>
<td>ompt_callback_target_submit</td>
</tr>
<tr>
<td>ompt_callback_target_submit_emi</td>
</tr>
<tr>
<td>ompt_callback_control_tool</td>
</tr>
<tr>
<td>ompt_callback_device_initialize</td>
</tr>
<tr>
<td>ompt_callback_device_finalize</td>
</tr>
<tr>
<td>ompt_callback_device_load</td>
</tr>
<tr>
<td>ompt_callback_device_unload</td>
</tr>
<tr>
<td>Callback name</td>
</tr>
<tr>
<td>--------------------------------------------------</td>
</tr>
<tr>
<td>ompt_callback_sync_region_wait</td>
</tr>
<tr>
<td>ompt_callback_mutex_released</td>
</tr>
<tr>
<td>ompt_callback_dependencies</td>
</tr>
<tr>
<td>ompt_callback_task_dependence</td>
</tr>
<tr>
<td>ompt_callback_work</td>
</tr>
<tr>
<td>ompt_callback_master // (deprecated)</td>
</tr>
<tr>
<td>ompt_callback_masked</td>
</tr>
<tr>
<td>ompt_callback_target_map</td>
</tr>
<tr>
<td>ompt_callback_target_map_emi</td>
</tr>
<tr>
<td>ompt_callback_sync_region</td>
</tr>
<tr>
<td>ompt_callback_reduction</td>
</tr>
<tr>
<td>ompt_callback_lock_init</td>
</tr>
<tr>
<td>ompt_callback_lock_destroy</td>
</tr>
<tr>
<td>ompt_callback_mutex_acquire</td>
</tr>
<tr>
<td>ompt_callback_mutex_acquired</td>
</tr>
<tr>
<td>ompt_callback_nest_lock</td>
</tr>
<tr>
<td>ompt_callback_flush</td>
</tr>
<tr>
<td>ompt_callback_cancel</td>
</tr>
<tr>
<td>ompt_callback_dispatch</td>
</tr>
</tbody>
</table>
When an OpenMP implementation initializes a target device, the OpenMP implementation dispatches the device initialization callback of the tool on the host device. If the OpenMP implementation or target device does not support tracing, the OpenMP implementation passes NULL to the device initializer of the tool for its lookup argument; otherwise, the OpenMP implementation passes a pointer to a device-specific runtime entry point with type signature ompt_function_lookup_t to the device initializer of the tool.

If a non-null lookup pointer is provided to the device initializer of the tool, the tool may use it to determine the runtime entry points in the tracing interface that are available for the device and may bind the returned function pointers to tool variables. Table 4.4 indicates the names of runtime entry points that may be available for a device; an implementations may provide additional implementation-defined names and corresponding entry points. The driver for the device provides the runtime entry points that enable a tool to control the trace collection interface of the device. The native trace format that the interface uses may be device specific and the available kinds of trace records are implementation defined. Some devices may allow a tool to collect traces of records in a standard format known as OMPT trace records. Each OMPT trace record serves as a substitute for an OMPT callback that cannot be made on the device. The fields in each trace record type are defined in the description of the callback that the record represents. If this type of record is provided then the lookup function returns values for the runtime entry points ompt_set_trace_ompt and ompt_get_record_ompt, which support collecting and decoding OMPT traces. If the native tracing format for a device is the OMPT format then tracing can be controlled using the runtime entry points for native or OMPT tracing.

The tool uses the ompt_set_trace_native and/or the ompt_set_trace_ompt runtime entry point to specify what types of events or activities to monitor on the device. The return codes for ompt_set_trace_ompt and ompt_set_trace_native use the ompt_set_result_t enumeration type. If the ompt_set_trace_native or the ompt_set_trace_ompt runtime entry point is called outside a device initializer, registration of supported callbacks may fail with a return code of ompt_set_error.

The tool initiates tracing on the device by invoking ompt_start_trace. Arguments to ompt_start_trace include two tool callbacks through which the OpenMP implementation can manage traces associated with the device. One callback allocates a buffer in which the device can deposit trace events. The second callback processes a buffer of trace events from the device.

If the device requires a trace buffer, the OpenMP implementation invokes the tool-supplied callback function on the host device to request a new buffer.

The OpenMP implementation monitors the execution of OpenMP constructs on the device and records a trace of events or activities into a trace buffer. If possible, device trace records are marked with a host_op_id—an identifier that associates device activities with the target operation that the host initiated to cause these activities. To correlate activities on the host with activities on a device, a tool can register a ompt_callback_target_submit_emi callback. Before and after the host initiates creation of an initial task on a device associated with a structured block for a target construct, the OpenMP implementation dispatches the ompt_callback_target_submit_emi callback on the host in the thread that is executing...
TABLE 4.4: OMPT Tracing Interface Runtime Entry Point Names and Their Type Signatures

<table>
<thead>
<tr>
<th>Entry Point String Name</th>
<th>Type Signature</th>
</tr>
</thead>
<tbody>
<tr>
<td>&quot;ompt_get_device_num_procs&quot;</td>
<td>ompt_get_device_num_procs_t</td>
</tr>
<tr>
<td>&quot;ompt_get_device_time&quot;</td>
<td>ompt_get_device_time_t</td>
</tr>
<tr>
<td>&quot;ompt_translate_time&quot;</td>
<td>ompt_translate_time_t</td>
</tr>
<tr>
<td>&quot;ompt_set_trace_ompt&quot;</td>
<td>ompt_set_trace_ompt_t</td>
</tr>
<tr>
<td>&quot;ompt_set_trace_native&quot;</td>
<td>ompt_set_trace_native_t</td>
</tr>
<tr>
<td>&quot;ompt_start_trace&quot;</td>
<td>ompt_start_trace_t</td>
</tr>
<tr>
<td>&quot;ompt_pause_trace&quot;</td>
<td>ompt_pause_trace_t</td>
</tr>
<tr>
<td>&quot;ompt_flush_trace&quot;</td>
<td>ompt_flush_trace_t</td>
</tr>
<tr>
<td>&quot;ompt_stop_trace&quot;</td>
<td>ompt_stop_trace_t</td>
</tr>
<tr>
<td>&quot;ompt_advance_buffer_cursor&quot;</td>
<td>ompt_advance_buffer_cursor_t</td>
</tr>
<tr>
<td>&quot;ompt_get_record_type&quot;</td>
<td>ompt_get_record_type_t</td>
</tr>
<tr>
<td>&quot;ompt_get_record_ompt&quot;</td>
<td>ompt_get_record_ompt_t</td>
</tr>
<tr>
<td>&quot;ompt_get_record_native&quot;</td>
<td>ompt_get_record_native_t</td>
</tr>
<tr>
<td>&quot;ompt_get_record_abstract&quot;</td>
<td>ompt_get_record_abstract_t</td>
</tr>
</tbody>
</table>
the task that encounters the target construct. This callback provides the tool with a pair of identifiers: one that identifies the target region and a second that uniquely identifies the initial task associated with that region. These identifiers help the tool correlate activities on the target device with their target region.

- When appropriate, for example, when a trace buffer fills or needs to be flushed, the OpenMP implementation invokes the tool-supplied buffer completion callback to process a non-empty sequence of records in a trace buffer that is associated with the device.

- The tool-supplied buffer completion callback may return immediately, ignoring records in the trace buffer, or it may iterate through them using the ompt_advance_buffer_cursor entry point to inspect each record. A tool may use the ompt_get_record_type runtime entry point to inspect the type of the record at the current cursor position. Three runtime entry points (ompt_get_record_ompt, ompt_get_record_native, and ompt_get_record_abstract) allow tools to inspect the contents of some or all records in a trace buffer. The ompt_get_record_native runtime entry point uses the native trace format of the device. The ompt_get_record_abstract runtime entry point decodes the contents of a native trace record and summarizes them as an ompt_record_abstract_t record. The ompt_get_record_ompt runtime entry point can only be used to retrieve records in OMPT format.

- Once tracing has been started on a device, a tool may pause or resume tracing on the device at any time by invoking ompt_pause_trace with an appropriate flag value as an argument.

- A tool may invoke the ompt_flush_trace runtime entry point for a device at any time between device initialization and finalization to cause the device to flush pending trace records.

- At any time, a tool may use the ompt_start_trace runtime entry point to start tracing or the ompt_stop_trace runtime entry point to stop tracing on a device. When tracing is stopped on a device, the OpenMP implementation eventually gathers all trace records already collected on the device and presents them to the tool using the buffer completion callback.

- An OpenMP implementation can be shut down while device tracing is in progress.

- When an OpenMP implementation is shut down, it finalizes each device. Device finalization occurs in three steps. First, the OpenMP implementation halts any tracing in progress for the device. Second, the OpenMP implementation flushes all trace records collected for the device and uses the buffer completion callback associated with that device to present them to the tool. Finally, the OpenMP implementation dispatches any ompt_callback_device_finalize callback registered for the device.

**Restrictions**

Restrictions on tracing activity on devices are as follows:

- Implementation-defined names must not start with the prefix ompt__, which is reserved for the OpenMP specification.
4.3 Finalizing a First-Party Tool

If the OMPT interface state is active, the tool finalizer, which has type signature
ompt_finalize_t and is specified by the finalize field in the
ompt_start_tool_result_t structure returned from the ompt_start_tool function, is
called when the OpenMP implementation shuts down.

Cross References
• ompt_finalize_t callback type, see Section 4.5.1.2
4.4 OMPT Data Types

The C/C++ header file (`omp-tools.h`) provides the definitions of the types that are specified throughout this subsection.

4.4.1 Tool Initialization and Finalization

**Summary**

A tool’s implementation of `ompt_start_tool` returns a pointer to an `ompt_start_tool_result_t` structure, which contains pointers to the tool’s initialization and finalization callbacks as well as an `ompt_data_t` object for use by the tool.

**Format**

```c
typedef struct ompt_start_tool_result_t {
    ompt_initialize_t initialize;
    ompt_finalize_t finalize;
    ompt_data_t tool_data;
} ompt_start_tool_result_t;
```

**Restrictions**

Restrictions to the `ompt_start_tool_result_t` type are as follows:

- The `initialize` and `finalize` callback pointer values in an `ompt_start_tool_result_t` structure that `ompt_start_tool` returns must be non-null.

**Cross References**

- `ompt_start_tool` function, see Section 4.2.1.
- `ompt_data_t` type, see Section 4.4.4.4.
- `ompt_initialize_t` callback type, see Section 4.5.1.1.
- `ompt_finalize_t` callback type, see Section 4.5.1.2.

4.4.2 Callbacks

**Summary**

The `ompt_callbacks_t` enumeration type indicates the integer codes used to identify OpenMP callbacks when registering or querying them.
```c
typedef enum ompt_callbacks_t {
    ompt_callback_thread_begin = 1,
    ompt_callback_thread_end = 2,
    ompt_callback_parallel_begin = 3,
    ompt_callback_parallel_end = 4,
    ompt_callback_task_create = 5,
    ompt_callback_task_schedule = 6,
    ompt_callback_implicit_task = 7,
    ompt_callback_target = 8,
    ompt_callback_target_data_op = 9,
    ompt_callback_target_submit = 10,
    ompt_callback_control_tool = 11,
    ompt_callback_device_initialize = 12,
    ompt_callback_device_finalize = 13,
    ompt_callback_device_load = 14,
    ompt_callback_device_unload = 15,
    ompt_callback_sync_region_wait = 16,
    ompt_callback_mutex_released = 17,
    ompt_callback_dependences = 18,
    ompt_callback_task_dependence = 19,
    ompt_callback_work = 20,
    ompt_callback_masked = 21,
    ompt_callback_master /*(deprecated)*/ = ompt_callback_masked,
    ompt_callback_target_map = 22,
    ompt_callback_sync_region = 23,
    ompt_callback_lock_init = 24,
    ompt_callback_lock_destroy = 25,
    ompt_callback_mutex_acquire = 26,
    ompt_callback_mutex_acquired = 27,
    ompt_callback_nest_lock = 28,
    ompt_callback_flush = 29,
    ompt_callback_cancel = 30,
    ompt_callback_reduction = 31,
    ompt_callback_dispatch = 32,
    ompt_callback_target_emi = 33,
    ompt_callback_target_data_op_emi = 34,
    ompt_callback_target_submit_emi = 35,
    ompt_callback_target_map_emi = 36,
    ompt_callback_error = 37
} ompt_callbacks_t;
```
4.4.3 Tracing
OpenMP provides type definitions that support tracing with OMPT.

4.4.3.1 Record Type
Summary
The `ompt_record_t` enumeration type indicates the integer codes used to identify OpenMP trace record formats.

Format
```c
typedef enum ompt_record_t {
    ompt_record_ompt = 1,
    ompt_record_native = 2,
    ompt_record_invalid = 3
} ompt_record_t;
```

4.4.3.2 Native Record Kind
Summary
The `ompt_record_native_t` enumeration type indicates the integer codes used to identify OpenMP native trace record contents.

Format
```c
typedef enum ompt_record_native_t {
    ompt_record_native_info = 1,
    ompt_record_native_event = 2
} ompt_record_native_t;
```

4.4.3.3 Native Record Abstract Type
Summary
The `ompt_record_abstract_t` type provides an abstract trace record format that is used to summarize native device trace records.
Format

```c
typedef struct ompt_record_abstract_t {
    ompt_record_native_t rclass;
    const char *type;
    ompt_device_time_t start_time;
    ompt_device_time_t end_time;
    ompt_hwid_t hwid;
} ompt_record_abstract_t;
```

Description

An `ompt_record_abstract_t` record contains information that a tool can use to process a native record that it may not fully understand. The `rclass` field indicates that the record is informational or that it represents an event; this information can help a tool determine how to present the record. The record `type` field points to a statically-allocated, immutable character string that provides a meaningful name that a tool can use to describe the event to a user. The `start_time` and `end_time` fields are used to place an event in time. The times are relative to the device clock. If an event does not have an associated `start_time (end_time)`, the value of the `start_time (end_time)` field is `ompt_time_none`. The hardware identifier field, `hwid`, indicates the location on the device where the event occurred. A `hwid` may represent a hardware abstraction such as a core or a hardware thread identifier. The meaning of a `hwid` value for a device is implementation defined. If no hardware abstraction is associated with the record then the value of `hwid` is `ompt_hwid_none`.

4.4.3.4 Record Type

Summary

The `ompt_record_ompt_t` type provides a standard complete trace record format.

Format

```c
typedef struct ompt_record_ompt_t {
    ompt_callbacks_t type;
    ompt_device_time_t time;
    ompt_id_t thread_id;
    ompt_id_t target_id;
    union {
        ompt_record_thread_begin_t thread_begin;
        ompt_record_parallel_begin_t parallel_begin;
        ompt_record_parallel_end_t parallel_end;
        ompt_record_work_t work;
        ompt_record_dispatch_t dispatch;
        ompt_record_task_create_t task_create;
    }
} ompt_record_ompt_t;
```
ompt_record_dependences_t  dependences;
ompt_record_task_dependence_t  task_dependence;
ompt_record_task_schedule_t  task_schedule;
ompt_record_implicit_task_t  implicit_task;
ompt_record_masked_t  masked;
ompt_record_sync_region_t  sync_region;
ompt_record_mutex_acquire_t  mutex_acquire;
ompt_record_mutex_t  mutex;
ompt_record_nest_lock_t  nest_lock;
ompt_record_flush_t  flush;
ompt_record_cancel_t  cancel;
ompt_record_target_t  target;
ompt_record_target_data_op_t  target_data_op;
ompt_record_target_map_t  target_map;
ompt_record_target_kernel_t  target_kernel;
ompt_record_control_tool_t  control_tool;
ompt_record_error_t  error;
}  record;
}  ompt_record_ompt_t;

C / C++]

Description
The field type specifies the type of record provided by this structure. According to the type, event specific information is stored in the matching record entry.

Restrictions
Restrictions to the ompt_record_ompt_t type are as follows:

- If type is set to ompt_callback_thread_end_t then the value of record is undefined.

4.4.4 Miscellaneous Type Definitions
This section describes miscellaneous types and enumerations used by the tool interface.

4.4.4.1 ompt_callback_t
Summary
Pointers to tool callback functions with different type signatures are passed to the ompt_set_callback runtime entry point and returned by the ompt_get_callback runtime entry point. For convenience, these runtime entry points expect all type signatures to be cast to a dummy type ompt_callback_t.
typedef void (*ompt_callback_t) (void);

4.4.4.2 ompt_set_result_t

Summary
The ompt_set_result_t enumeration type corresponds to values that the
ompt_set_callback, ompt_set_trace_ompt and ompt_set_trace_native
runtime entry points return.

typedef enum ompt_set_result_t {
    ompt_set_error = 0,
    ompt_set_never = 1,
    ompt_set_impossible = 2,
    ompt_set_sometimes = 3,
    ompt_set_sometimes_paired = 4,
    ompt_set_always = 5
} ompt_set_result_t;

Description
Values of ompt_set_result_t, may indicate several possible outcomes. The
omp_set_error value indicates that the associated call failed. Otherwise, the value indicates
when an event may occur and, when appropriate, dispatching a callback event leads to the
invocation of the callback. The ompt_set_never value indicates that the event will never occur
or that the callback will never be invoked at runtime. The ompt_set_impossible value
indicates that the event may occur but that tracing of it is not possible. The
ompt_set_sometimes value indicates that the event may occur and, for an
implementation-defined subset of associated event occurrences, will be traced or the callback will
be invoked at runtime. The ompt_set_sometimes_paired value indicates the same result as
ompt_set_sometimes and, in addition, that a callback with an endpoint value of
ompt_scope_begin will be invoked if and only if the same callback with an endpoint value of
ompt_scope_end will also be invoked sometime in the future. The ompt_set_always value
indicates that, whenever an associated event occurs, it will be traced or the callback will be invoked.
Cross References

- Monitoring activity on the host with OMPT, see Section 4.2.4.
- Tracing activity on target devices with OMPT, see Section 4.2.5.
- `ompt_set_callback` runtime entry point, see Section 4.6.1.3.
- `ompt_set_trace_ompt` runtime entry point, see Section 4.6.2.4.
- `ompt_set_trace_native` runtime entry point, see Section 4.6.2.5.

4.4.4.3 ompt_id_t

Summary

The `ompt_id_t` type is used to provide various identifiers to tools.

Format

```
C / C++
typedef uint64_t ompt_id_t;
```

Description

When tracing asynchronous activity on devices, identifiers enable tools to correlate target regions and operations that the host initiates with associated activities on a target device. In addition, OMPT provides identifiers to refer to parallel regions and tasks that execute on a device. These various identifiers are of type `ompt_id_t`.

`ompt_id_none` is defined as an instance of type `ompt_id_t` with the value 0.

Restrictions

Restrictions to the `ompt_id_t` type are as follows:

- Identifiers created on each device must be unique from the time an OpenMP implementation is initialized until it is shut down. Identifiers for each target region and target data operation instance that the host device initiates must be unique over time on the host. Identifiers for parallel and task region instances that execute on a device must be unique over time within that device.
4.4.4.4 ompt_data_t

Summary
The ompt_data_t type represents data associated with threads and with parallel and task regions.

Format
```c
typedef union ompt_data_t {
    uint64_t value;
    void *ptr;
} ompt_data_t;
```

Description
The ompt_data_t type represents data that is reserved for tool use and that is related to a thread or to a parallel or task region. When an OpenMP implementation creates a thread or an instance of a parallel, teams, task, or target region, it initializes the associated ompt_data_t object with the value ompt_data_none, which is an instance of the type with the data and pointer fields equal to 0.

4.4.4.5 ompt_device_t

Summary
The ompt_device_t opaque object type represents a device.

Format
```c
typedef void ompt_device_t;
```

4.4.4.6 ompt_device_time_t

Summary
The ompt_device_time_t type represents raw device time values.

Format
```c
typedef uint64_t ompt_device_time_t;
```
The `ompt_device_time_t` opaque object type represents raw device time values. `ompt_time_none` refers to an unknown or unspecified time and is defined as an instance of type `ompt_device_time_t` with the value 0.

### 4.4.4.7 ompt_buffer_t

**Summary**
The `ompt_buffer_t` opaque object type is a handle for a target buffer.

**Format**
```c
typedef void ompt_buffer_t;
```

### 4.4.4.8 ompt_buffer_cursor_t

**Summary**
The `ompt_buffer_cursor_t` opaque type is a handle for a position in a target buffer.

**Format**
```c
typedef uint64_t ompt_buffer_cursor_t;
```

### 4.4.4.9 ompt_dependence_t

**Summary**
The `ompt_dependence_t` type represents a task dependence.

**Format**
```c
typedef struct ompt_dependence_t {
    ompt_data_t variable;
    ompt_dependence_type_t dependence_type;
} ompt_dependence_t;
```
Description
The `ompt_dependence_t` type is a structure that holds information about a depend clause. For task dependences, the `variable` field points to the storage location of the dependence. For `doacross` dependences, the `variable` field contains the value of a vector element that describes the dependence. The `dependence_type` field indicates the type of the dependence.

Cross References
• `ompt_dependence_type_t` type, see Section 4.4.4.23.

4.4.4.10 `ompt_thread_t`

Summary
The `ompt_thread_t` enumeration type defines the valid thread type values.

Format
```c
typedef enum ompt_thread_t {
    ompt_thread_initial = 1,
    ompt_thread_worker = 2,
    ompt_thread_other = 3,
    ompt_thread_unknown = 4
} ompt_thread_t;
```

Description
Any `initial thread` has thread type `ompt_thread_initial`. All `OpenMP threads` that are not initial threads have thread type `ompt_thread_worker`. A thread that an OpenMP implementation uses but that does not execute user code has thread type `ompt_thread_other`. Any thread that is created outside an OpenMP implementation and that is not an `initial thread` has thread type `ompt_thread_unknown`.

4.4.4.11 `ompt_scope_endpoint_t`

Summary
The `ompt_scope_endpoint_t` enumeration type defines valid scope endpoint values.

Format
```c
typedef enum ompt_scope_endpoint_t {
    ompt_scope_begin = 1,
    ompt_scope_end = 2,
    ompt_scope_beginend = 3
} ompt_scope_endpoint_t;
```
4.4.4.12 ompt_dispatch_t

Summary
The ompt_dispatch_t enumeration type defines the valid dispatch kind values.

Format

```c
typedef enum ompt_dispatch_t {
    ompt_dispatch_iteration = 1,
    ompt_dispatch_section = 2
} ompt_dispatch_t;
```

4.4.4.13 ompt_sync_region_t

Summary
The ompt_sync_region_t enumeration type defines the valid synchronization region kind values.

Format

```c
typedef enum ompt_sync_region_t {
    ompt_sync_region_barrier = 1, // deprecated
    ompt_sync_region_barrier_implicit = 2, // deprecated
    ompt_sync_region_barrier_explicit = 3,
    ompt_sync_region_barrier_implementation = 4,
    ompt_sync_region_taskwait = 5,
    ompt_sync_region_taskgroup = 6,
    ompt_sync_region_reduction = 7,
    ompt_sync_region_barrier_implicit_workshare = 8,
    ompt_sync_region_barrier_implicit_parallel = 9,
    ompt_sync_region_barrier_teams = 10
} ompt_sync_region_t;
```
4.4.14 ompt_target_data_op_t

**Summary**
The `ompt_target_data_op_t` enumeration type defines the valid target data operation values.

**Format**
```c
typedef enum ompt_target_data_op_t {
    ompt_target_data_alloc = 1,
    ompt_target_data_transfer_to_device = 2,
    ompt_target_data_transfer_from_device = 3,
    ompt_target_data_delete = 4,
    ompt_target_data_associate = 5,
    ompt_target_data_disassociate = 6,
    ompt_target_data_alloc_async = 17,
    ompt_target_data_transfer_to_device_async = 18,
    ompt_target_data_transfer_from_device_async = 19,
    ompt_target_data_delete_async = 20
} ompt_target_data_op_t;
```

4.4.15 ompt_work_t

**Summary**
The `ompt_work_t` enumeration type defines the valid work type values.

**Format**
```c
typedef enum ompt_work_t {
    ompt_work_loop = 1,
    ompt_work_sections = 2,
    ompt_work_single_executor = 3,
    ompt_work_single_other = 4,
    ompt_work_workshare = 5,
    ompt_work_distribute = 6,
    ompt_work_taskloop = 7,
    ompt_work_scope = 8
} ompt_work_t;
```
4.4.4.16 ompt_mutex_t

Summary
The ompt_mutex_t enumeration type defines the valid mutex kind values.

Format

```c
typedef enum ompt_mutex_t {
    ompt_mutex_lock = 1,
    ompt_mutex_test_lock = 2,
    ompt_mutex_nest_lock = 3,
    ompt_mutex_test_nest_lock = 4,
    ompt_mutex_critical = 5,
    ompt_mutex_atomic = 6,
    ompt_mutex_ordered = 7
} ompt_mutex_t;
```

4.4.4.17 ompt_native_mon_flag_t

Summary
The ompt_native_mon_flag_t enumeration type defines the valid native monitoring flag values.

Format

```c
typedef enum ompt_native_mon_flag_t {
    ompt_native_data_motion_explicit = 0x01,
    ompt_native_data_motion_implicit = 0x02,
    ompt_native_kernel_invocation = 0x04,
    ompt_native_kernel_execution = 0x08,
    ompt_native_driver = 0x10,
    ompt_native_runtime = 0x20,
    ompt_native_overhead = 0x40,
    ompt_native_idleness = 0x80
} ompt_native_mon_flag_t;
```
4.4.4.18 ompt_task_flag_t

Summary
The ompt_task_flag_t enumeration type defines valid task types.

Format

```c
typedef enum ompt_task_flag_t {
    ompt_task_initial = 0x00000001,
    ompt_task_implicit = 0x00000002,
    ompt_task_explicit = 0x00000004,
    ompt_task_target = 0x00000008,
    ompt_task_taskwait = 0x00000010,
    ompt_task_undeferred = 0x08000000,
    ompt_task_untied = 0x10000000,
    ompt_task_final = 0x20000000,
    ompt_task_mergeable = 0x40000000,
    ompt_task_merged = 0x80000000
} ompt_task_flag_t;
```

Description
The ompt_task_flag_t enumeration type defines valid task type values. The least significant byte provides information about the general classification of the task. The other bits represent properties of the task.

4.4.4.19 ompt_task_status_t

Summary
The ompt_task_status_t enumeration type indicates the reason that a task was switched when it reached a task scheduling point.

Format

```c
typedef enum ompt_task_status_t {
    ompt_task_complete = 1,
    ompt_task_yield = 2,
    ompt_task_cancel = 3,
    ompt_task_detach = 4,
    ompt_task_early_fulfill = 5,
    ompt_task_late_fulfill = 6,
    ompt_task_switch = 7,
    ompt_taskwait_complete = 8
} ompt_task_status_t;
```
Description
The value `ompt_task_complete` of the `ompt_task_status_t` type indicates that the task that encountered the task scheduling point completed execution of the associated structured block and an associated `allow-completion` event was fulfilled. The value `ompt_task_yield` indicates that the task encountered a `taskyield` construct. The value `ompt_task_cancel` indicates that the task was canceled when it encountered an active cancellation point. The value `ompt_task_detach` indicates that a task for which the `detach` clause was specified completed execution of the associated structured block and is waiting for an `allow-completion` event to be fulfilled. The value `ompt_task_early_fulfill` indicates that the `allow-completion` event of the task was fulfilled before the task completed execution of the associated structured block. The value `ompt_task_late_fulfill` indicates that the `allow-completion` event of the task was fulfilled after the task completed execution of the associated structured block. The value `ompt_taskwait_complete` indicates completion of the dependent task that results from a `taskwait` construct with one or more `depend` clauses. The value `ompt_task_switch` is used for all other cases that a task was switched.

4.4.4.20 ompt_target_t

Summary
The `ompt_target_t` enumeration type defines the valid target type values.

Format

```c
typedef enum ompt_target_t {
    ompt_target           = 1,
    ompt_target_enter_data = 2,
    ompt_target_exit_data = 3,
    ompt_target_update    = 4,
    ompt_target_nowait    = 9,
    ompt_target_enter_data_nowait = 10,
    ompt_target_exit_data_nowait = 11,
    ompt_target_update_nowait = 12
} ompt_target_t;
```
4.4.4.21 ompt_parallel_flag_t

Summary
The ompt_parallel_flag_t enumeration type defines valid invoker values.

Format

C / C++

typedef enum ompt_parallel_flag_t {
  ompt_parallel_invoker_program = 0x00000001,
  ompt_parallel_invoker_runtime = 0x00000002,
  ompt_parallel_league = 0x40000000,
  ompt_parallel_team = 0x80000000
} ompt_parallel_flag_t;

Description
The ompt_parallel_flag_t enumeration type defines valid invoker values, which indicate how an outlined function is invoked.

The value ompt_parallel_invoker_program indicates that the outlined function associated with implicit tasks for the region is invoked directly by the application on the primary thread for a parallel region.

The value ompt_parallel_invoker_runtime indicates that the outlined function associated with implicit tasks for the region is invoked by the runtime on the primary thread for a parallel region.

The value ompt_parallel_league indicates that the callback is invoked due to the creation of a league of teams by a teams construct.

The value ompt_parallel_team indicates that the callback is invoked due to the creation of a team of threads by a parallel construct.
4.4.4.22 ompt_target_map_flag_t

Summary
The ompt_target_map_flag_t enumeration type defines the valid target map flag values.

Format
```c
typedef enum ompt_target_map_flag_t {
    ompt_target_map_flag_to = 0x01,
    ompt_target_map_flag_from = 0x02,
    ompt_target_map_flag_alloc = 0x04,
    ompt_target_map_flag_release = 0x08,
    ompt_target_map_flag_delete = 0x10,
    ompt_target_map_flag_implicit = 0x20
} ompt_target_map_flag_t;
```

4.4.4.23 ompt_dependence_type_t

Summary
The ompt_dependence_type_t enumeration type defines the valid task dependence type values.

Format
```c
typedef enum ompt_dependence_type_t {
    ompt_dependence_type_in = 1,
    ompt_dependence_type_out = 2,
    ompt_dependence_type_inout = 3,
    ompt_dependence_type_mutexinoutset = 4,
    ompt_dependence_type_source = 5,
    ompt_dependence_type_sink = 6,
    ompt_dependence_type_inoutset = 7
} ompt_dependence_type_t;
```
4.4.4.24 ompt_severity_t

Summary
The ompt_severity_t enumeration type defines the valid severity values.

Format
```c
typedef enum ompt_severity_t {
    ompt_warning = 1,
    ompt_fatal = 2
} ompt_severity_t;
```

4.4.4.25 ompt_cancel_flag_t

Summary
The ompt_cancel_flag_t enumeration type defines the valid cancel flag values.

Format
```c
typedef enum ompt_cancel_flag_t {
    ompt_cancel_parallel = 0x01,
    ompt_cancel_sections = 0x02,
    ompt_cancel_loop = 0x04,
    ompt_cancel_taskgroup = 0x08,
    ompt_cancel_activated = 0x10,
    ompt_cancel_detected = 0x20,
    ompt_cancel_discarded_task = 0x40
} ompt_cancel_flag_t;
```

4.4.4.26 ompt_hwid_t

Summary
The ompt_hwid_t opaque type is a handle for a hardware identifier for a target device.

Format
```c
typedef uint64_t ompt_hwid_t;
```
Description

The `ompt_hwid_t` opaque type is a handle for a hardware identifier for a target device.

`ompt_hwid_none` is an instance of the type that refers to an unknown or unspecified hardware identifier and that has the value 0. If no `hwid` is associated with an `ompt_record_abstract_t` then the value of `hwid` is `ompt_hwid_none`.

Cross References

- `ompt_record_abstract_t` type, see Section 4.4.3.3.

4.4.4.27 `ompt_state_t`

Summary

If the OMPT interface is in the active state then an OpenMP implementation must maintain thread state information for each thread. The thread state maintained is an approximation of the instantaneous state of a thread.

Format

```
 typedef enum ompt_state_t {
   ompt_state_work_serial = 0x000,
   ompt_state_work_parallel = 0x001,
   ompt_state_work_reduction = 0x002,
   ompt_state_wait_barrier = 0x010, // deprecated
   ompt_state_wait_barrier_implicit_parallel = 0x011,
   ompt_state_wait_barrier_implicit_workshare = 0x012,
   ompt_state_wait_barrier_implicit = 0x013, // deprecated
   ompt_state_wait_barrier_explicit = 0x014,
   ompt_state_wait_barrier_implementation = 0x015,
   ompt_state_wait_barrier_teams = 0x016,
   ompt_state_wait_taskwait = 0x020,
   ompt_state_wait_taskgroup = 0x021,
   ompt_state_wait_mutex = 0x040,
   ompt_state_wait_lock = 0x041,
   ompt_state_wait_critical = 0x042,
   ompt_state_wait_atomic = 0x043,
   ompt_state_wait_ordered = 0x044,
}CHAPTER 4. OMPT INTERFACE 503
```
A tool can query the OpenMP state of a thread at any time. If a tool queries the state of a thread that is not associated with OpenMP then the implementation reports the state as `ompt_state_undefined`.

The value `ompt_state_work_serial` indicates that the thread is executing code outside all parallel regions.

The value `ompt_state_work_parallel` indicates that the thread is executing code within the scope of a parallel region.

The value `ompt_state_work_reduction` indicates that the thread is combining partial reduction results from threads in its team. An OpenMP implementation may never report a thread in this state; a thread that is combining partial reduction results may have its state reported as `ompt_state_work_parallel` or `ompt_state_overhead`.

The value `ompt_state_wait_barrier_implicit_parallel` indicates that the thread is waiting at the implicit barrier at the end of a parallel region.

The value `ompt_state_wait_barrier_implicit_workshare` indicates that the thread is waiting at an implicit barrier at the end of a worksharing construct.

The value `ompt_state_wait_barrier_explicit` indicates that the thread is waiting in an explicit barrier region.

The value `ompt_state_wait_barrier_implementation` indicates that the thread is waiting in a barrier not required by the OpenMP standard but introduced by an OpenMP implementation.

The value `ompt_state_wait_barrier_teams` indicates that the thread is waiting at a barrier at the end of a teams region.

The value `ompt_state_wait_taskwait` indicates that the thread is waiting at a taskwait construct.

The value `ompt_state_wait_taskgroup` indicates that the thread is waiting at the end of a taskgroup construct.
The value `ompt_state_wait_mutex` indicates that the thread is waiting for a mutex of an unspecified type.

The value `ompt_state_wait_lock` indicates that the thread is waiting for a lock or nestable lock.

The value `ompt_state_wait_critical` indicates that the thread is waiting to enter a critical region.

The value `ompt_state_wait_atomic` indicates that the thread is waiting to enter an atomic region.

The value `ompt_state_wait_ordered` indicates that the thread is waiting to enter an ordered region.

The value `ompt_state_wait_target` indicates that the thread is waiting for a target region to complete.

The value `ompt_state_wait_target_map` indicates that the thread is waiting for a target data mapping operation to complete. An implementation may report `ompt_state_wait_target` for target data constructs.

The value `ompt_state_wait_target_update` indicates that the thread is waiting for a target update operation to complete. An implementation may report `ompt_state_wait_target` for target update constructs.

The value `ompt_state_idle` indicates that the thread is idle, that is, it is not part of an OpenMP team.

The value `ompt_state_overhead` indicates that the thread is in the overhead state at any point while executing within the OpenMP runtime, except while waiting at a synchronization point.

The value `ompt_state_undefined` indicates that the native thread is not created by the OpenMP implementation.

### 4.4.4.28 ompt_frame_t

**Summary**

The `ompt_frame_t` type describes procedure frame information for an OpenMP task.

**Format**

```c
typedef struct ompt_frame_t {
  ompt_data_t exit_frame;
  ompt_data_t enter_frame;
  int exit_frame_flags;
  int enter_frame_flags;
} ompt_frame_t;
```
Description
Each `ompt_frame_t` object is associated with the task to which the procedure frames belong. Each non-merged initial, implicit, explicit, or target task with one or more frames on the stack of a native thread has an associated `ompt_frame_t` object.

The `exit_frame` field of an `ompt_frame_t` object contains information to identify the first procedure frame executing the task region. The `exit_frame` for the `ompt_frame_t` object associated with the initial task that is not nested inside any OpenMP construct is `NULL`.

The `enter_frame` field of an `ompt_frame_t` object contains information to identify the latest still active procedure frame executing the task region before entering the OpenMP runtime implementation or before executing a different task. If a task with frames on the stack has not been suspended, the value of `enter_frame` for the `ompt_frame_t` object associated with the task may contain `NULL`.

For `exit_frame`, the `exit_frame_flags` and, for `enter_frame`, the `enter_frame_flags` field indicates that the provided frame information points to a runtime or an application frame address. The same fields also specify the kind of information that is provided to identify the frame. These fields are a disjunction of values in the `ompt_frame_flag_t` enumeration type.

The lifetime of an `ompt_frame_t` object begins when a task is created and ends when the task is destroyed. Tools should not assume that a frame structure remains at a constant location in memory throughout the lifetime of the task. A pointer to an `ompt_frame_t` object is passed to some callbacks; a pointer to the `ompt_frame_t` object of a task can also be retrieved by a tool at any time, including in a signal handler, by invoking the `ompt_get_task_info` runtime entry point (described in Section 4.6.1.14). A pointer to an `ompt_frame_t` object that a tool retrieved is valid as long as the tool does not pass back control to the OpenMP implementation.

Note – A monitoring tool that uses asynchronous sampling can observe values of `exit_frame` and `enter_frame` at inconvenient times. Tools must be prepared to handle `ompt_frame_t` objects observed just prior to when their field values will be set or cleared.

4.4.4.29 `ompt_frame_flag_t`
Summary
The `ompt_frame_flag_t` enumeration type defines valid frame information flags.
Format

```c
typedef enum ompt_frame_flag_t {
    ompt_frame_runtime  = 0x00,
    ompt_frame_application = 0x01,
    ompt_frame_cfa = 0x10,
    ompt_frame_framepointer = 0x20,
    ompt_frame_stackaddress = 0x30
} ompt_frame_flag_t;
```

Description

The value `ompt_frame_runtime` of the `ompt_frame_flag_t` type indicates that a frame address is a procedure frame in the OpenMP runtime implementation. The value `ompt_frame_application` of the `ompt_frame_flag_t` type indicates that a frame address is a procedure frame in the OpenMP application.

Higher order bits indicate the kind of provided information that is unique for the particular frame pointer. The value `ompt_frame_cfa` indicates that a frame address specifies a canonical frame address. The value `ompt_frame_framepointer` indicates that a frame address provides the value of the frame pointer register. The value `ompt_frame_stackaddress` indicates that a frame address specifies a pointer address that is contained in the current stack frame.

4.4.4.30 ompt_wait_id_t

Summary

The `ompt_wait_id_t` type describes wait identifiers for an OpenMP thread.

Format

```c
typedef uint64_t ompt_wait_id_t;
```

Description

Each thread maintains a `wait identifier` of type `ompt_wait_id_t`. When a task that a thread executes is waiting for mutual exclusion, the wait identifier of the thread indicates the reason that the thread is waiting. A wait identifier may represent a critical section `name`, a lock, a program variable accessed in an atomic region, or a synchronization object that is internal to an OpenMP implementation. When a thread is not in a wait state then the value of the wait identifier of the thread is undefined.

`ompt_wait_id_none` is defined as an instance of type `ompt_wait_id_t` with the value 0.
4.5 OMPT Tool Callback Signatures and Trace Records

The C/C++ header file (`omp-tools.h`) provides the definitions of the types that are specified throughout this subsection. Restrictions to the OpenMP tool callbacks are as follows:

Restrictions
- Tool callbacks may not use OpenMP directives or call any runtime library routines described in Section 3.
- Tool callbacks must exit by either returning to the caller or aborting.

4.5.1 Initialization and Finalization Callback Signature

4.5.1.1 `ompt_initialize_t`

Summary
A callback with type signature `ompt_initialize_t` initializes use of the OMPT interface.

Format
```
C / C++

typedef int (*ompt_initialize_t) (ompt_function_lookup_t lookup,
                                   int initial_device_num,
                                   ompt_data_t *tool_data);
```

Description
To use the OMPT interface, an implementation of `ompt_start_tool` must return a non-null pointer to an `ompt_start_tool_result_t` structure that contains a pointer to a tool initializer function with type signature `ompt_initialize_t`. An OpenMP implementation will call the initializer after fully initializing itself but before beginning execution of any OpenMP construct or runtime library routine.

The initializer returns a non-zero value if it succeeds; otherwise the OMPT interface state changes to `inactive` as described in Section 4.2.3.

Description of Arguments
The `lookup` argument is a callback to an OpenMP runtime routine that must be used to obtain a pointer to each runtime entry point in the OMPT interface. The `initial_device_num` argument provides the value of `omp_get_initial_device()`. The `tool_data` argument is a pointer to the `tool_data` field in the `ompt_start_tool_result_t` structure that `ompt_start_tool` returned.
Cross References

• `omp_get_initial_device` routine, see Section 3.7.7.

• `ompt_start_tool` function, see Section 4.2.1.

• `ompt_start_tool_result_t` type, see Section 4.4.1.

• `ompt_data_t` type, see Section 4.4.4.

• `ompt_function_lookup_t` type, see Section 4.6.3.

4.5.1.2 `ompt_finalize_t`

Summary

A tool implements a finalizer with the type signature `ompt_finalize_t` to finalize its use of the OMPT interface.

Format

```
C / C++

typedef void (*ompt_finalize_t) (ompt_data_t *tool_data);
```

Description

To use the OMPT interface, an implementation of `ompt_start_tool` must return a non-null pointer to an `ompt_start_tool_result_t` structure that contains a non-null pointer to a tool finalizer with type signature `ompt_finalize_t`. An OpenMP implementation must call the tool finalizer after the last OMPT event as the OpenMP implementation shuts down.

Description of Arguments

The `tool_data` argument is a pointer to the `tool_data` field in the `ompt_start_tool_result_t` structure returned by `ompt_start_tool`.

Cross References

• `ompt_start_tool` function, see Section 4.2.1.

• `ompt_start_tool_result_t` type, see Section 4.4.1.

• `ompt_data_t` type, see Section 4.4.4.
4.5.2 Event Callback Signatures and Trace Records

This section describes the signatures of tool callback functions that an OMPT tool may register and that are called during runtime of an OpenMP program. An implementation may also provide a trace of events per device. Along with the callbacks, the following defines standard trace records. For the trace records, tool data arguments are replaced by an ID, which must be initialized by the OpenMP implementation. Each of parallel_id, task_id, and thread_id must be unique per target region. Tool implementations of callbacks are not required to be async signal safe.

Cross References
- ompt_id_t type, see Section 4.4.4.3.
- ompt_data_t type, see Section 4.4.4.4.

4.5.2.1 ompt_callback_thread_begin_t

Summary
The ompt_callback_thread_begin_t type is used for callbacks that are dispatched when native threads are created.

Format
```c
typedef void (*ompt_callback_thread_begin_t) (ompt_thread_t thread_type,
                                          ompt_data_t *thread_data);
```

Trace Record
```c
typedef struct ompt_record_thread_begin_t {
  ompt_thread_t thread_type;
} ompt_record_thread_begin_t;
```

Description of Arguments
The thread_type argument indicates the type of the new thread: initial, worker, or other. The binding of the thread_data argument is the new thread.
Cross References
• parallel construct, see Section 2.6.
• teams construct, see Section 2.7.
• Initial task, see Section 2.12.5.
• ompt_data_t type, see Section 4.4.4.4.
• ompt_thread_t type, see Section 4.4.10.

4.5.2.2 ompt_callback_thread_end_t

Summary
The ompt_callback_thread_end_t type is used for callbacks that are dispatched when native threads are destroyed.

Format

C / C++

typedef void (*ompt_callback_thread_end_t) (ompt_data_t *thread_data);

C / C++

Description of Arguments
The binding of the thread_data argument is the thread that will be destroyed.

Cross References
• parallel construct, see Section 2.6.
• teams construct, see Section 2.7.
• Initial task, see Section 2.12.5.
• ompt_record_ompt_t type, see Section 4.4.3.4.
• ompt_data_t type, see Section 4.4.4.4.

4.5.2.3 ompt_callback_parallel_begin_t

Summary
The ompt_callback_parallel_begin_t type is used for callbacks that are dispatched when a parallel or teams region starts.
typedef void (*ompt_callback_parallel_begin_t) (ompt_data_t *encountering_task_data, const ompt_frame_t *encountering_task_frame, ompt_data_t *parallel_data, unsigned int requested_parallelism, int flags, const void *codeptr_ra);

typedef struct ompt_record_parallel_begin_t {
    ompt_id_t encountering_task_id;
    ompt_id_t parallel_id;
    unsigned int requested_parallelism;
    int flags;
    const void *codeptr_ra;
} ompt_record_parallel_begin_t;

Description of Arguments

The binding of the `encountering_task_data` argument is the encountering task.

The `encountering_task_frame` argument points to the frame object that is associated with the encountering task. Accessing the frame object after the callback returned can cause a data race.

The binding of the `parallel_data` argument is the parallel or teams region that is beginning.

The `requested_parallelism` argument indicates the number of threads or teams that the user requested.

The `flags` argument indicates whether the code for the region is inlined into the application or invoked by the runtime and also whether the region is a parallel or teams region. Valid values for `flags` are a disjunction of elements in the enum `ompt_parallel_flag_t`.

The `codeptr_ra` argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature `ompt_callback_parallel_begin_t` then `codeptr_ra` contains the return address of the call to that runtime routine. If the implementation of the region is inlined then `codeptr_ra` contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, `codeptr_ra` may be `NULL`.
Cross References

- **parallel** construct, see Section 2.6.
- **teams** construct, see Section 2.7.
- ompt_data_t type, see Section 4.4.4.
- ompt_parallel_flag_t type, see Section 4.4.21.
- ompt_frame_t type, see Section 4.4.28.

4.5.2.4 ompt_callback_parallel_end_t

Summary
The ompt_callback_parallel_end_t type is used for callbacks that are dispatched when a parallel or teams region ends.

Format

```c
typedef void (*ompt_callback_parallel_end_t) ( 
    ompt_data_t *parallel_data, 
    ompt_data_t *encountering_task_data, 
    int flags, 
    const void *codeptr Ra
);
```

Trace Record

```c
typedef struct ompt_record_parallel_end_t {
    ompt_id_t parallel_id; 
    ompt_id_t encountering_task_id; 
    int flags; 
    const void *codeptr Ra; 
} ompt_record_parallel_end_t;
```

Description of Arguments
The binding of the parallel_data argument is the parallel or teams region that is ending.

The binding of the encountering_task_data argument is the encountering task.

The flags argument indicates whether the execution of the region is inlined into the application or invoked by the runtime and also whether it is a parallel or teams region. Values for flags are a disjunction of elements in the enum ompt_parallel_flag_t.
The `codeptr_ra` argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature `ompt_callback_parallel_end_t` then `codeptr_ra` contains the return address of the call to that runtime routine. If the implementation of the region is inlined then `codeptr_ra` contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, `codeptr_ra` may be `NULL`.

Cross References
- `parallel` construct, see Section 2.6.
- `teams` construct, see Section 2.7.
- `ompt_data_t` type, see Section 4.4.4.
- `ompt_parallel_flag_t` type, see Section 4.4.21.

4.5.2.5 `ompt_callback_work_t`

Summary
The `ompt_callback_work_t` type is used for callbacks that are dispatched when worksharing regions, loop-related regions, `taskloop` regions and `scope` regions begin and end.

Format

```c
typedef void (*ompt_callback_work_t) (ompt_work_t wstype,
                 ompt_scope_endpoint_t endpoint,
                 ompt_data_t *parallel_data,
                 ompt_data_t *task_data,
                 uint64_t count,
                 const void *codeptr_ra);
```

Trace Record

```c
typedef struct ompt_record_work_t {
    ompt_work_t wstype;
    ompt_scope_endpoint_t endpoint;
    ompt_id_t parallel_id;
    ompt_id_t task_id;
    uint64_t count;
    const void *codeptr_ra;
} ompt_record_work_t;
```
Description of Arguments

The wstype argument indicates the kind of region.

The endpoint argument indicates that the callback signals the beginning of a scope or the end of a scope.

The binding of the parallel_data argument is the current parallel region.

The binding of the task_data argument is the current task.

The count argument is a measure of the quantity of work involved in the construct. For a worksharing-loop or taskloop construct, count represents the number of iterations in the iteration space, which may be the result of collapsing several associated loops. For a sections construct, count represents the number of sections. For a workshare construct, count represents the units of work, as defined by the workshare construct. For a single or scope construct, count is always 1. When the endpoint argument signals the end of a scope, a count value of 0 indicates that the actual count value is not available.

The codeptr_ra argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature ompt_callback_work_t then codeptr_ra contains the return address of the call to that runtime routine. If the implementation of the region is inlined then codeptr_ra contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, codeptr_ra may be NULL.

Cross References

• Worksharing constructs, see Section 2.10.
• Loop-related directives, see Section 2.11.
• Worksharing-Loop construct, see Section 2.11.4.
• taskloop construct, see Section 2.12.2.
• ompt_data_t type, see Section 4.4.4.4.
• ompt_scope_endpoint_t type, see Section 4.4.4.11.
• ompt_work_t type, see Section 4.4.4.15.

4.5.2.6 ompt_callback_dispatch_t

Summary

The ompt_callback_dispatch_t type is used for callbacks that are dispatched when a thread begins to execute a section or loop iteration.
Format

C / C++

```c
typedef void (*ompt_callback_dispatch_t) (
    ompt_data_t *parallel_data,
    ompt_data_t *task_data,
    ompt_dispatch_t kind,
    ompt_data_t instance
);
```

Trace Record

C / C++

```c
typedef struct ompt_record_dispatch_t {
    ompt_id_t parallel_id;
    ompt_id_t task_id;
    ompt_dispatch_t kind;
    ompt_data_t instance;
} ompt_record_dispatch_t;
```

Description of Arguments

The binding of the `parallel_data` argument is the current parallel region.

The binding of the `task_data` argument is the implicit task that executes the structured block of the parallel region.

The `kind` argument indicates whether a loop iteration or a section is being dispatched.

For a loop iteration, the `instance.value` argument contains the logical iteration number. For a structured block in the `sections` construct, `instance.ptr` contains a code address that identifies the structured block. In cases where a runtime routine implements the structured block associated with this callback, `instance.ptr` contains the return address of the call to the runtime routine. In cases where the implementation of the structured block is inlined, `instance.ptr` contains the return address of the invocation of this callback.

Cross References

- `sections` and `section` constructs, see Section 2.10.1.
- Worksharing-loop construct, see Section 2.11.4.
- `taskloop` construct, see Section 2.12.2.
- `ompt_data_t` type, see Section 4.4.4.4.
- `ompt_dispatch_t` type, see Section 4.4.4.12.
4.5.2.7 ompt_callback_task_create_t

Summary
The `ompt_callback_task_create_t` type is used for callbacks that are dispatched when task regions are generated.

Format

```c/c++
typedef void (*ompt_callback_task_create_t) (  
    ompt_data_t *encountering_task_data,  
    const ompt_frame_t *encountering_task_frame,  
    ompt_data_t *new_task_data,  
    int flags,  
    int has_dependences,  
    const void *codeptr_ra  
);```

Trace Record

```c/c++
typedef struct ompt_record_task_create_t {  
    ompt_id_t encountering_task_id;  
    ompt_id_t new_task_id;  
    int flags;  
    int has_dependences;  
    const void *codeptr_ra;  
} ompt_record_task_create_t;```

Description of Arguments

The binding of the `encountering_task_data` argument is the encountering task.

The `encountering_task_frame` argument points to the frame object associated with the encountering task. Accessing the frame object after the callback returned can cause a data race.

The binding of the `new_task_data` argument is the generated task.

The `flags` argument indicates the kind of task (explicit or target) that is generated. Values for `flags` are a disjunction of elements in the `ompt_task_flag_t` enumeration type.

The `has_dependences` argument is `true` if the generated task has dependences and `false` otherwise.

The `codeptr_ra` argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature `ompt_callback_task_create_t` then `codeptr_ra` contains the return address of the call to that runtime routine. If the implementation of the region is inlined then `codeptr_ra` contains the
return address of the invocation of the callback. If attribution to source code is impossible or 
inappropriate, `codeptr_ra` may be NULL.

Cross References

- **task** construct, see Section 2.12.1.
- Initial task, see Section 2.12.5.
- `ompt_data_t` type, see Section 4.4.4.
- `ompt_task_flag_t` type, see Section 4.4.18.
- `ompt_frame_t` type, see Section 4.4.28.

4.5.2.8 `ompt_callback_dependences_t`

Summary

The `ompt_callback_dependences_t` type is used for callbacks that are related to 
dependences and that are dispatched when new tasks are generated and when `ordered` constructs 
are encountered.

Format

```c
typedef void (*ompt_callback_dependences_t) ( 
    ompt_data_t *task_data, 
    const ompt_dependence_t *deps, 
    int ndeps 
);
```

Trace Record

```c
typedef struct ompt_record_dependences_t { 
    ompt_id_t task_id; 
    ompt_dependence_t dep; 
    int ndeps; 
} ompt_record_dependences_t;
```

Description of Arguments

The binding of the `task_data` argument is the generated task for a depend clause on a task construct, 
the target task for a depend clause on a target construct respectively depend object in an 
asynchronous runtime routine, or the encountering implicit task for a depend clause of the ordered 
construct.
The *deps* argument lists dependences of the new task or the dependence vector of the ordered construct. Dependences denoted with dependency objects are described in terms of their dependency semantics.

The *ndeps* argument specifies the length of the list passed by the *deps* argument. The memory for *deps* is owned by the caller; the tool cannot rely on the data after the callback returns.

The performance monitor interface for tracing activity on target devices provides one record per dependence.

**Cross References**
- *ordered* construct, see Section 2.19.9.
- *depend* clause, see Section 2.19.11.
- *ompt_data_t* type, see Section 4.4.4.4.
- *ompt_dependence_t* type, see Section 4.4.4.9.

### 4.5.2.9 ompt_callback_task_dependence_t

**Summary**
The *ompt_callback_task_dependence_t* type is used for callbacks that are dispatched when unfulfilled task dependences are encountered.

**Format**
```
C / C++
typedef void (*ompt_callback_task_dependence_t) (  
    ompt_data_t *src_task_data,  
    ompt_data_t *sink_task_data  
);  
```

**Trace Record**
```
C / C++
typedef struct ompt_record_task_dependence_t {  
    ompt_id_t src_task_id;  
    ompt_id_t sink_task_id;  
} ompt_record_task_dependence_t;  
```

**Description of Arguments**
The binding of the *src_task_data* argument is a running task with an outgoing dependence.

The binding of the *sink_task_data* argument is a task with an unsatisfied incoming dependence.
Cross References

- **depend clause**, see Section 2.19.11.
- **ompt_data_t type**, see Section 4.4.4.4.

4.5.2.10 *ompt_callback_task_schedule_t*

Summary

The *ompt_callback_task_schedule_t* type is used for callbacks that are dispatched when task scheduling decisions are made.

Format

```c
typedef void (*ompt_callback_task_schedule_t) (
    ompt_data_t *prior_task_data,
    ompt_task_status_t prior_task_status,
    ompt_data_t *next_task_data
);```

Trace Record

```c
typedef struct ompt_record_task_schedule_t {
    ompt_id_t prior_task_id;
    ompt_task_status_t prior_task_status;
    ompt_id_t next_task_id;
} ompt_record_task_schedule_t;
```

Description of Arguments

- The *prior_task_status* argument indicates the status of the task that arrived at a task scheduling point.
- The binding of the *prior_task_data* argument is the task that arrived at the scheduling point.
- The binding of the *next_task_data* argument is the task that is resumed at the scheduling point. This argument is **NULL** if the callback is dispatched for a *task-fulfill* event or if the callback signals completion of a taskwait construct.

Cross References

- Task scheduling, see Section 2.12.6.
- **ompt_data_t type**, see Section 4.4.4.4.
- **ompt_task_status_t type**, see Section 4.4.4.19.
4.5.2.11 ompt_callback_implicit_task_t

Summary

The ompt_callback_implicit_task_t type is used for callbacks that are dispatched when initial tasks and implicit tasks are generated and completed.

Format

```c
typedef void (*ompt_callback_implicit_task_t) (  
  ompt_scope_endpoint_t endpoint,  
  ompt_data_t *parallel_data,  
  ompt_data_t *task_data,  
  unsigned int actual_parallelism,  
  unsigned int index,  
  int flags  
);
```

Trace Record

```c
typedef struct ompt_record_implicit_task_t {  
  ompt_scope_endpoint_t endpoint;  
  ompt_id_t parallel_id;  
  ompt_id_t task_id;  
  unsigned int actual_parallelism;  
  unsigned int index;  
  int flags;  
} ompt_record_implicit_task_t;
```

Description of Arguments

The endpoint argument indicates that the callback signals the beginning of a scope or the end of a scope.

The binding of the parallel_data argument is the current parallel or teams region. For the implicit-task-end and the initial-task-end events, this argument is NULL.

The binding of the task_data argument is the implicit task that executes the structured block of the parallel or teams region.

The actual_parallelism argument indicates the number of threads in the parallel region or the number of teams in the teams region. For initial tasks, that are not closely nested in a teams construct, this argument is 1. For the implicit-task-end and the initial-task-end events, this argument is 0.
The index argument indicates the thread number or team number of the calling thread, within the team or league that is executing the parallel or teams region to which the implicit task region binds. For initial tasks, that are not created by a teams construct, this argument is 1.

The flags argument indicates the kind of task (initial or implicit).

Cross References
- parallel construct, see Section 2.6.
- teams construct, see Section 2.7.
- ompt_data_t type, see Section 4.4.4.
- ompt_scope_endpoint_t enumeration type, see Section 4.4.11.

4.5.2.12 ompt_callback_masked_t

Summary
The ompt_callback_masked_t type is used for callbacks that are dispatched when masked regions start and end.

Format

```c
typedef void (*ompt_callback_masked_t) (  
    ompt_scope_endpoint_t endpoint,  
    ompt_data_t *parallel_data,  
    ompt_data_t *task_data,  
    const void *codeptr_ra  
  );
```

Trace Record

```c
typedef struct ompt_record_masked_t {  
    ompt_scope_endpoint_t endpoint;  
    ompt_id_t parallel_id;  
    ompt_id_t task_id;  
    const void *codeptr_ra;  
  } ompt_record_masked_t;
```
Description of Arguments

The *endpoint* argument indicates that the callback signals the beginning of a scope or the end of a scope.

The binding of the *parallel_data* argument is the current parallel region.

The binding of the *task_data* argument is the encountering task.

The *codeptr_ra* argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature `ompt_callback_masked_t` then *codeptr_ra* contains the return address of the call to that runtime routine. If the implementation of the region is inlined then *codeptr_ra* contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, *codeptr_ra* may be *NULL*.

Cross References

- *masked* construct, see Section 2.8.
- *ompt_data_t* type, see Section 4.4.4.4.
- *ompt_scope_endpoint_t* type, see Section 4.4.11.

4.5.2.13 `ompt_callback_sync_region_t`

Summary

The `ompt_callback_sync_region_t` type is used for callbacks that are dispatched when barrier regions, *taskwait* regions, and *taskgroup* regions begin and end and when waiting begins and ends for them as well as for when reductions are performed.

Format

```c
typedef void (*ompt_callback_sync_region_t) (ompt_sync_region_t kind,
                                           ompt_scope_endpoint_t endpoint,
                                           ompt_data_t *parallel_data,
                                           ompt_data_t *task_data,
                                           const void *codeptr_ra);
```
typedef struct ompt_record_sync_region_t {
    ompt_sync_region_t kind;
    ompt_scope_endpoint_t endpoint;
    ompt_id_t parallel_id;
    ompt_id_t task_id;
    const void *codeptr_ra;
} ompt_record_sync_region_t;

Description of Arguments
The `kind` argument indicates the kind of synchronization.

The `endpoint` argument indicates that the callback signals the beginning of a scope or the end of a scope.

The binding of the `parallel_data` argument is the current parallel region. For the `barrier-end` event at the end of a parallel region this argument is `NULL`.

The binding of the `task_data` argument is the current task.

The `codeptr_ra` argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature `ompt_callback_sync_region_t` then `codeptr_ra` contains the return address of the call to that runtime routine. If the implementation of the region is inlined then `codeptr_ra` contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, `codeptr_ra` may be `NULL`.

Cross References
• `barrier` construct, see Section 2.19.2.
• Implicit barriers, see Section 2.19.3.
• `taskwait` construct, see Section 2.19.5.
• `taskgroup` construct, see Section 2.19.6.
• Properties common to all reduction clauses, see Section 2.21.5.1.
• `ompt_data_t` type, see Section 4.4.4.4.
• `ompt_scope_endpoint_t` type, see Section 4.4.4.11.
• `ompt_sync_region_t` type, see Section 4.4.4.13.
4.5.2.14 ompt_callback_mutex_acquire_t

Summary
The ompt_callback_mutex_acquire_t type is used for callbacks that are dispatched when
locks are initialized, acquired and tested and when critical regions, atomic regions, and
ordered regions are begun.

Format

```c
typedef void (*ompt_callback_mutex_acquire_t) (
    ompt_mutex_t kind,
    unsigned int hint,
    unsigned int impl,
    ompt_wait_id_t wait_id,
    const void *codeptr_ra
);
```

Trace Record

```c
typedef struct ompt_record_mutex_acquire_t {
    ompt_mutex_t kind;
    unsigned int hint;
    unsigned int impl;
    ompt_wait_id_t wait_id;
    const void *codeptr_ra;
} ompt_record_mutex_acquire_t;
```

Description of Arguments
The kind argument indicates the kind of mutual exclusion event.

The hint argument indicates the hint that was provided when initializing an implementation of
mutual exclusion. If no hint is available when a thread initiates acquisition of mutual exclusion, the
runtime may supply omp_sync_hint_none as the value for hint.

The impl argument indicates the mechanism chosen by the runtime to implement the mutual
exclusion.

The wait_id argument indicates the object being awaited.

The codeptr_ra argument relates the implementation of an OpenMP region to its source code. If a
runtime routine implements the region associated with a callback that has type signature
ompt_callback_mutex_acquire_t then codeptr_ra contains the return address of the call
to that runtime routine. If the implementation of the region is inlined then codeptr_ra contains the
return address of the invocation of the callback. If attribution to source code is impossible or
inappropriate, codeptr_ra may be NULL.
Cross References
• critical construct, see Section 2.19.1.
• atomic construct, see Section 2.19.7.
• ordered construct, see Section 2.19.9.
• omp_init_lock and omp_init_nest_lock routines, see Section 3.9.1.
• ompt_mutex_t type, see Section 4.4.4.16.
• ompt_wait_id_t type, see Section 4.4.4.30.

4.5.2.15 ompt_callback_mutex_t

Summary
The ompt_callback_mutex_t type is used for callbacks that indicate important synchronization events.

Format
C / C++

```c
typedef void (*ompt_callback_mutex_t) (  
  ompt_mutex_t kind,  
  ompt_wait_id_t wait_id,  
  const void *codeptr_ra  
);  
```

Trace Record
C / C++

```c
typedef struct ompt_record_mutex_t {  
  ompt_mutex_t kind;  
  ompt_wait_id_t wait_id;  
  const void *codeptr_ra;  
} ompt_record_mutex_t;  
```

Description of Arguments

The kind argument indicates the kind of mutual exclusion event.

The wait_id argument indicates the object being awaited.

The codeptr_ra argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature ompt_callback_mutex_t then codeptr_ra contains the return address of the call to that runtime routine. If the implementation of the region is inlined then codeptr_ra contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, codeptr_ra may be NULL.
Cross References

- critical construct, see Section 2.19.1.
- atomic construct, see Section 2.19.7.
- ordered construct, see Section 2.19.9.
- omp_destroy_lock and omp_destroy_nest_lock routines, see Section 3.9.3.
- omp_set_lock and omp_set_nest_lock routines, see Section 3.9.4.
- omp_unset_lock and omp_unset_nest_lock routines, see Section 3.9.5.
- omp_test_lock and omp_test_nest_lock routines, see Section 3.9.6.
- ompt_mutex_t type, see Section 4.4.4.16.
- ompt_wait_id_t type, see Section 4.4.4.30.

4.5.2.16 ompt_callback_nest_lock_t

Summary
The ompt_callback_nest_lock_t type is used for callbacks that indicate that a thread that owns a nested lock has performed an action related to the lock but has not relinquished ownership of it.

Format

```c
typedef void (*ompt_callback_nest_lock_t) (  
    ompt_scope_endpoint_t endpoint,  
    ompt_wait_id_t wait_id,  
    const void *codeptr_ra  
);  
```

Trace Record

```c
typedef struct ompt_record_nest_lock_t {
    ompt_scope_endpoint_t endpoint;  
    ompt_wait_id_t wait_id;  
    const void *codeptr_ra;  
} ompt_record_nest_lock_t;  
```

Description of Arguments

The endpoint argument indicates that the callback signals the beginning of a scope or the end of a scope.

The wait_id argument indicates the object being awaited.
The `codeptr_ra` argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature `ompt_callback_nest_lock_t` then `codeptr_ra` contains the return address of the call to that runtime routine. If the implementation of the region is inlined then `codeptr_ra` contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, `codeptr_ra` may be `NULL`.

Cross References
- `omp_set_nest_lock` routine, see Section 3.9.4.
- `omp_unset_nest_lock` routine, see Section 3.9.5.
- `omp_test_nest_lock` routine, see Section 3.9.6.
- `ompt_scope_endpoint_t` type, see Section 4.4.4.11.
- `ompt_wait_id_t` type, see Section 4.4.4.30.

### 4.5.2.17 ompt_callback_flush_t

**Summary**
The `ompt_callback_flush_t` type is used for callbacks that are dispatched when `flush` constructs are encountered.

**Format**

```c
typedef void (*ompt_callback_flush_t) (ompt_data_t *thread_data, const void *codeptr_ra);
```

**Trace Record**

```c
typedef struct ompt_record_flush_t {
    const void *codeptr_ra;
} ompt_record_flush_t;
```

**Description of Arguments**

The binding of the `thread_data` argument is the executing thread.

The `codeptr_ra` argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature `ompt_callback_flush_t` then `codeptr_ra` contains the return address of the call to that runtime routine. If the implementation of the region is inlined then `codeptr_ra` contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, `codeptr_ra` may be `NULL`. 
Cross References
• flush construct, see Section 2.19.8.
• ompt_data_t type, see Section 4.4.4.4.

4.5.2.18 ompt_callback_cancel_t

Summary
The ompt_callback_cancel_t type is used for callbacks that are dispatched for cancellation, cancel and discarded-task events.

Format

```c
typedef void (*ompt_callback_cancel_t) (ompt_data_t *task_data,
                                           int flags,
                                           const void *codeptr_ra);
```

Trace Record

```c
typedef struct ompt_record_cancel_t {
  ompt_id_t task_id;
  int flags;
  const void *codeptr_ra;
} ompt_record_cancel_t;
```

Description of Arguments
The binding of the task_data argument is the task that encounters a cancel construct, a cancellation point construct, or a construct defined as having an implicit cancellation point.

The flags argument, defined by the ompt_cancel_flag_t enumeration type, indicates whether cancellation is activated by the current task, or detected as being activated by another task. The construct that is being canceled is also described in the flags argument. When several constructs are detected as being concurrently canceled, each corresponding bit in the argument will be set.

The codeptr_ra argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature ompt_callback_cancel_t then codeptr_ra contains the return address of the call to that runtime routine. If the implementation of the region is inlined then codeptr_ra contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, codeptr_ra may be NULL.
Cross References

• `omp_cancel_flag_t` enumeration type, see Section 4.4.25.

4.5.2.19 `ompt_callback_device_initialize_t`

Summary

The `ompt_callback_device_initialize_t` type is used for callbacks that initialize device tracing interfaces.

Format

```c
C / C++

typedef void (*ompt_callback_device_initialize_t) (  
    int device_num,  
    const char *type,  
    ompt_device_t *device,  
    ompt_function_lookup_t lookup,  
    const char *documentation
);
```

Description

Registration of a callback with type signature `ompt_callback_device_initialize_t` for the `ompt_callback_device_initialize` event enables asynchronous collection of a trace for a device. The OpenMP implementation invokes this callback after OpenMP is initialized for the device but before execution of any OpenMP construct is started on the device.

Description of Arguments

The `device_num` argument identifies the logical device that is being initialized.

The `type` argument is a character string that indicates the type of the device. A device type string is a semicolon-separated character string that includes at a minimum the vendor and model name of the device. These names may be followed by a semicolon-separated sequence of properties that describe the hardware or software of the device.

The `device` argument is a pointer to an opaque object that represents the target device instance. Functions in the device tracing interface use this pointer to identify the device that is being addressed.

The `lookup` argument points to a runtime callback that a tool must use to obtain pointers to runtime entry points in the device’s OMPT tracing interface. If a device does not support tracing then `lookup` is `NULL`.

The `documentation` argument is a string that describes how to use any device-specific runtime entry points that can be obtained through the `lookup` argument. This documentation string may be a pointer to external documentation, or it may be inline descriptions that include names and type
signatures for any device-specific interfaces that are available through the *lookup* argument along with descriptions of how to use these interface functions to control monitoring and analysis of device traces.

**Constraints on Arguments**
The *type* and *documentation* arguments must be immutable strings that are defined for the lifetime of program execution.

**Effect**
A device initializer must fulfill several duties. First, the *type* argument should be used to determine if any special knowledge about the hardware and/or software of a device is employed. Second, the *lookup* argument should be used to look up pointers to runtime entry points in the OMPT tracing interface for the device. Finally, these runtime entry points should be used to set up tracing for the device.

Initialization of tracing for a target device is described in Section 4.2.5.

**Cross References**
- *ompt_function_lookup_t* type, see Section 4.6.3.

### 4.5.2.20 ompt_callback_device_finalize_t

**Summary**
The *ompt_callback_device_initialize_t* type is used for callbacks that finalize device tracing interfaces.

**Format**

```c
C / C++
typedef void (*ompt_callback_device_finalize_t) (int device_num);
```

**Description of Arguments**
The *device_num* argument identifies the logical device that is being finalized.

**Description**
A registered callback with type signature *ompt_callback_device_finalize_t* is dispatched for a device immediately prior to finalizing the device. Prior to dispatching a finalization callback for a device on which tracing is active, the OpenMP implementation stops tracing on the device and synchronously flushes all trace records for the device that have not yet been reported. These trace records are flushed through one or more buffer completion callbacks with type signature *ompt_callback_buffer_complete_t* as needed prior to the dispatch of the callback with type signature *ompt_callback_device_finalize_t*. 
Cross References

- `ompt_callback_buffer_complete_t` callback type, see Section 4.5.2.24.

### 4.5.2.21 ompt_callback_device_load_t

#### Summary

The `ompt_callback_device_load_t` type is used for callbacks that the OpenMP runtime invokes to indicate that it has just loaded code onto the specified device.

#### Format

```c
typedef void (*ompt_callback_device_load_t) (  
    int device_num,  
    const char *filename,  
    int64_t offset_in_file,  
    void *vma_in_file,  
    size_t bytes,  
    void *host_addr,  
    void *device_addr,  
    uint64_t module_id

);
```

#### Description of Arguments

The `device_num` argument specifies the device.

The `filename` argument indicates the name of a file in which the device code can be found. A NULL `filename` indicates that the code is not available in a file in the file system.

The `offset_in_file` argument indicates an offset into `filename` at which the code can be found. A value of -1 indicates that no offset is provided.

`ompt_addr_none` is defined as a pointer with the value ~0.

The `vma_in_file` argument indicates a virtual address in `filename` at which the code can be found. A value of `ompt_addr_none` indicates that a virtual address in the file is not available.

The `bytes` argument indicates the size of the device code object in bytes.

The `host_addr` argument indicates the address at which a copy of the device code is available in host memory. A value of `ompt_addr_none` indicates that a host code address is not available.

The `device_addr` argument indicates the address at which the device code has been loaded in device memory. A value of `ompt_addr_none` indicates that a device code address is not available.

The `module_id` argument is an identifier that is associated with the device code object.
Cross References
• Device directives, see Section 2.14.

4.5.2.22 ompt_callback_device_unload_t

Summary
The `ompt_callback_device_unload_t` type is used for callbacks that the OpenMP runtime invokes to indicate that it is about to unload code from the specified device.

Format

```c
typedef void (*ompt_callback_device_unload_t) (  
    int device_num,  
    uint64_t module_id
);
```

Description of Arguments
The `device_num` argument specifies the device.
The `module_id` argument is an identifier that is associated with the device code object.

Cross References
• Device directives, see Section 2.14.

4.5.2.23 ompt_callback_buffer_request_t

Summary
The `ompt_callback_buffer_request_t` type is used for callbacks that are dispatched when a buffer to store event records for a device is requested.

Format

```c
typedef void (*ompt_callback_buffer_request_t) (  
    int device_num,  
    ompt_buffer_t **buffer,  
    size_t *bytes
);
```

Description
A callback with type signature `ompt_callback_buffer_request_t` requests a buffer to store trace records for the specified device. A buffer request callback may set `*bytes` to 0 if it does not provide a buffer. If a callback sets `*bytes` to 0, further recording of events for the device is disabled until the next invocation of `ompt_start_trace`. This action causes the device to drop future trace records until recording is restarted.
Description of Arguments

The `device_num` argument specifies the device.

The `*buffer` argument points to a buffer where device events may be recorded. The `*bytes` argument indicates the length of that buffer.

Cross References

• ompt_buffer_t type, see Section 4.4.4.7.

4.5.2.24 ompt_callback_buffer_complete_t

Summary

The ompt_callback_buffer_complete_t type is used for callbacks that are dispatched when devices will not record any more trace records in an event buffer and all records written to the buffer are valid.

Format

```
C / C++
```

```
typedef void (*ompt_callback_buffer_complete_t) (
    int device_num,
    ompt_buffer_t *buffer,
    size_t bytes,
    ompt_buffer_cursor_t begin,
    int buffer_owned
);
```

Description

A callback with type signature ompt_callback_buffer_complete_t provides a buffer that contains trace records for the specified device. Typically, a tool will iterate through the records in the buffer and process them.

The OpenMP implementation makes these callbacks on a thread that is not an OpenMP primary or worker thread.

The callee may not delete the buffer if the buffer_owned argument is 0.

The buffer completion callback is not required to be async signal safe.

Description of Arguments

The `device_num` argument indicates the device for which the buffer contains events.

The `buffer` argument is the address of a buffer that was previously allocated by a buffer request callback.

The `bytes` argument indicates the full size of the buffer.
The `begin` argument is an opaque cursor that indicates the position of the beginning of the first record in the buffer.

The `buffer_owned` argument is 1 if the data to which the buffer points can be deleted by the callback and 0 otherwise. If multiple devices accumulate trace events into a single buffer, this callback may be invoked with a pointer to one or more trace records in a shared buffer with `buffer_owned = 0`. In this case, the callback may not delete the buffer.

**Cross References**
- `ompt_buffer_t` type, see Section 4.4.4.7.
- `ompt_buffer_cursor_t` type, see Section 4.4.4.8.

### 4.5.2.25 `ompt_callback_target_data_op_emi_t` and `ompt_callback_target_data_op_t`

**Summary**
The `ompt_callback_target_data_op_emi_t` and `ompt_callback_target_data_op_t` types are used for callbacks that are dispatched when a thread maps data to a device.

**Format**

```c
typedef void (*ompt_callback_target_data_op_emi_t) (  
    ompt_scope_endpoint_t endpoint,  
    ompt_data_t *target_task_data,  
    ompt_data_t *target_data,  
    ompt_id_t *host_op_id,  
    ompt_target_data_op_t optype,  
    void *src_addr,  
    int src_device_num,  
    void *dest_addr,  
    int dest_device_num,  
    size_t bytes,  
    const void *codeptr_ra
);
```

```c
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```
typedef void (*ompt_callback_target_data_op_t) (  
    ompt_id_t target_id,  
    ompt_id_t host_op_id,  
    ompt_target_data_op_t optype,  
    void *src_addr,  
    int src_device_num,  
    void *dest_addr,  
    int dest_device_num,  
    size_t bytes,  
    const void *codeptr_ra  
);
**Description of Arguments**

The *endpoint* argument indicates that the callback signals the beginning or end of a scope.

The binding of the *target_task_data* argument is the target task region.

The binding of the *target_data* argument is the target region.

The *host_op_id* argument points to a tool controlled integer value, which identifies a data operation on a target device.

The *optype* argument indicates the kind of data operation.

The *src_addr* argument indicates the data address before the operation, where applicable.

The *src_device_num* argument indicates the source device number for the data operation, where applicable.

The *dest_addr* argument indicates the data address after the operation.

The *dest_device_num* argument indicates the destination device number for the data operation.

Whether in some operations *src_addr* or *dest_addr* may point to an intermediate buffer is implementation defined.

The *bytes* argument indicates the size of data.

The *codeptr_ra* argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature

```c
ompt_callback_target_data_op_emi_t
```

or

```c
ompt_callback_target_data_op_t
```

then *codeptr_ra* contains the return address of the call to that runtime routine. If the implementation of the region is inlined then *codeptr_ra* contains the return address of the invocation of the callback.

If attribution to source code is impossible or inappropriate, *codeptr_ra* may be **NULL**.

**Restrictions**

Restrictions to the *ompt_callback_target_data_op_emi* and

```c
ompt_callback_target_data_op
```

callbacks are as follows:

- These callbacks must not be registered at the same time.

**Cross References**

- *map* clause, see Section 2.21.7.1.
- *ompt_id_t* type, see Section 4.4.4.3.
- *ompt_data_t* type, see Section 4.4.4.4.
- *ompt_scope_endpoint_t* type, see Section 4.4.11.
- *ompt_target_data_op_t* type, see Section 4.4.14.
4.5.2.26 ompt_callback_target_emi_t and ompt_callback_target_t

Summary
The ompt_callback_target_emi_t and ompt_callback_target_t types are used for callbacks that are dispatched when a thread begins to execute a device construct.

Format

```c
typedef void (*ompt_callback_target_emi_t) (
    ompt_target_t kind,
    ompt_scope_endpoint_t endpoint,
    int device_num,
    ompt_data_t *task_data,
    ompt_data_t *target_task_data,
    ompt_data_t *target_data,
    const void *codeptr_ra
);
```

```c
typedef void (*ompt_callback_target_t) (
    ompt_target_t kind,
    ompt_scope_endpoint_t endpoint,
    int device_num,
    ompt_data_t *task_data,
    ompt_id_t target_id,
    const void *codeptr_ra
);
```

Trace Record

```c
typedef struct ompt_record_target_t {
    ompt_target_t kind;
    ompt_scope_endpoint_t endpoint;
    int device_num;
    ompt_id_t task_id;
    ompt_id_t target_id;
    const void *codeptr_ra;
} ompt_record_target_t;
```
Description of Arguments

The *kind* argument indicates the kind of target region.

The *endpoint* argument indicates that the callback signals the beginning of a scope or the end of a scope.

The *device_num* argument indicates the device number of the device that will execute the target region.

The binding of the *task_data* argument is the generating task.

The binding of the *target_task_data* argument is the target region.

The binding of the *target_data* argument is the target region.

The *codeptr_ra* argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature `ompt_callback_target_emi_t` or `ompt_callback_target_t` then `codeptr_ra` contains the return address of the call to that runtime routine. If the implementation of the region is inlined then `codeptr_ra` contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, `codeptr_ra` may be `NULL`.

Restrictions

Restrictions to the *ompt_callback_target_emi* and *ompt_callback_target* callbacks are as follows:

- These callbacks must not be registered at the same time.

Cross References

- *target data* construct, see Section 2.14.2.
- *target enter data* construct, see Section 2.14.3.
- *target exit data* construct, see Section 2.14.4.
- *target* construct, see Section 2.14.5.
- *target update* construct, see Section 2.14.6.
- *ompt_id_t* type, see Section 4.4.4.3.
- *ompt_data_t* type, see Section 4.4.4.4.
- *ompt_scope_endpoint_t* type, see Section 4.4.4.11.
- *ompt_target_t* type, see Section 4.4.4.20.
4.5.2.27 `ompt_callback_target_map_emi_t` and `ompt_callback_target_map_t`

Summary
The `ompt_callback_target_map_emi_t` and `ompt_callback_target_map_t` types are used for callbacks that are dispatched to indicate data mapping relationships.

Format

```c
typedef void (*ompt_callback_target_map_emi_t) (  
    ompt_data_t *target_data,  
    unsigned int nitems,  
    void **host_addr,  
    void **device_addr,  
    size_t *bytes,  
    unsigned int *mapping_flags,  
    const void *codeptr_ra  
);
```

```c
typedef void (*ompt_callback_target_map_t) (  
    ompt_id_t target_id,  
    unsigned int nitems,  
    void **host_addr,  
    void **device_addr,  
    size_t *bytes,  
    unsigned int *mapping_flags,  
    const void *codeptr_ra  
);
```

Trace Record

```c
typedef struct ompt_record_target_map_t  
{  
    ompt_id_t target_id;  
    unsigned int nitems;  
    void **host_addr;  
    void **device_addr;  
    size_t *bytes;  
    unsigned int *mapping_flags;  
    const void *codeptr_ra;  
} ompt_record_target_map_t;
```
Description
An instance of a \texttt{target, target data, target enter data, or target exit data} construct may contain one or more \texttt{map} clauses. An OpenMP implementation may report the set of mappings associated with \texttt{map} clauses for a construct with a single \texttt{ompt_callback_target_map_emi} or \texttt{ompt_callback_target_map} callback to report the effect of all mappings or multiple \texttt{ompt_callback_target_map_emi} or \texttt{ompt_callback_target_map} callbacks with each reporting a subset of the mappings. Furthermore, an OpenMP implementation may omit mappings that it determines are unnecessary. If an OpenMP implementation issues multiple \texttt{ompt_callback_target_map_emi} or \texttt{ompt_callback_target_map} callbacks, these callbacks may be interleaved with \texttt{ompt_callback_target_data_op_emi} or \texttt{ompt_callback_target_data_op} callbacks used to report data operations associated with the mappings.

Description of Arguments
The binding of the \texttt{target_data} argument is the target region.

The \texttt{nitems} argument indicates the number of data mappings that this callback reports.

The \texttt{host_addr} argument indicates an array of host data addresses.

The \texttt{device_addr} argument indicates an array of device data addresses.

The \texttt{bytes} argument indicates an array of sizes of data.

The \texttt{mapping_flags} argument indicates the kind of data mapping. Flags for a mapping include one or more values specified by the \texttt{ompt_target_map_flag_t} type.

The \texttt{codeptr_ra} argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature \texttt{ompt_callback_target_map_t} or \texttt{ompt_callback_target_map_emi_t} then \texttt{codeptr_ra} contains the return address of the call to that runtime routine. If the implementation of the region is inlined then \texttt{codeptr_ra} contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, \texttt{codeptr_ra} may be \texttt{NULL}.

Restrictions
Restrictions to the \texttt{ompt_callback_target_data_map_emi} and \texttt{ompt_callback_target_data_map} callbacks are as follows:

\begin{itemize}
  \item These callbacks must not be registered at the same time.
\end{itemize}

Cross References
\begin{itemize}
  \item target data construct, see Section 2.14.2.
  \item target enter data construct, see Section 2.14.3.
  \item target exit data construct, see Section 2.14.4.
  \item target construct, see Section 2.14.5.
\end{itemize}
• ompt_id_t type, see Section 4.4.3.
• ompt_data_t type, see Section 4.4.4.
• ompt_target_map_flag_t type, see Section 4.4.22.

• ompt_callback_target_data_op_emi_t or
ompt_callback_target_data_op_t callback type, see Section 4.5.2.25.

4.5.2.28 ompt_callback_target_submit_emi_t and
ompt_callback_target_submit_t

Summary
The ompt_callback_target_submit_emi_t and
ompt_callback_target_submit_t types are used for callbacks that are dispatched before
and after the host initiates creation of an initial task on a device.

Format

```c
typedef void (*ompt_callback_target_submit_emi_t) (  
    ompt_scope_endpoint_t endpoint,  
    ompt_data_t *target_data,  
    ompt_id_t *host_op_id,  
    unsigned int requested_num_teams  
) ;
```

```c
typedef void (*ompt_callback_target_submit_t) (  
    ompt_id_t target_id,  
    ompt_id_t host_op_id,  
    unsigned int requested_num_teams  
) ;
```

Trace Record

```c
typedef struct ompt_record_target_kernel_t {
    ompt_id_t host_op_id;
    unsigned int requested_num_teams;
    unsigned int granted_num_teams;
    ompt_device_time_t end_time;
} ompt_record_target_kernel_t;
```
Description
A thread dispatches a registered `ompt_callback_target_submit_emi` or
`ompt_callback_target_submit` callback on the host before and after a target task initiates
creation of an initial task on a device.

Description of Arguments
The `endpoint` argument indicates that the callback signals the beginning or end of a scope.

The binding of the `target_data` argument is the target region.

The `host_op_id` argument points to a tool controlled integer value, which identifies an initial task
on a target device.

The `requested_num_teams` argument is the number of teams that the host requested to execute the
kernel. The actual number of teams that execute the kernel may be smaller and generally will not be
known until the kernel begins to execute on the device.

If `ompt_set_trace_ompt` has configured the device to trace kernel execution then the device
will log a `ompt_record_target_kernel_t` record in a trace. The fields in the record are as
follows:

- The `host_op_id` field contains a tool-controlled identifier that can be used to correlate a
  `ompt_record_target_kernel_t` record with its associated
  `ompt_callback_target_submit_emi` or `ompt_callback_target_submit`
callback on the host;

- The `requested_num_teams` field contains the number of teams that the host requested to execute
  the kernel;

- The `granted_num_teams` field contains the number of teams that the device actually used to
  execute the kernel;

- The time when the initial task began execution on the device is recorded in the `time` field of an
  enclosing `ompt_record_t` structure; and

- The time when the initial task completed execution on the device is recorded in the `end_time`
  field.

Restrictions
Restrictions to the `ompt_callback_target_submit_emi` and
`ompt_callback_target_submit` callbacks are as follows:

- These callbacks must not be registered at the same time.
Cross References

- **target** construct, see Section 2.14.5.
- **ompt_id_t** type, see Section 4.4.3.
- **ompt_data_t** type, see Section 4.4.4.
- **ompt_scope_endpoint_t** type, see Section 4.4.11.

4.5.2.29 ompt_callback_control_tool_t

Summary

The `ompt_callback_control_tool_t` type is used for callbacks that dispatch *tool-control* events.

Format

```c
typedef int (*ompt_callback_control_tool_t) (  
    uint64_t command,  
    uint64_t modifier,  
    void *arg,  
    const void *codeptr_ra
);
```

Trace Record

```c
typedef struct ompt_record_control_tool_t {
    uint64_t command;  
    uint64_t modifier;  
    const void *codeptr_ra;
} ompt_record_control_tool_t;
```

Description

Callbacks with type signature `ompt_callback_control_tool_t` may return any non-negative value, which will be returned to the application as the return value of the `omp_control_tool` call that triggered the callback.

Description of Arguments

The `command` argument passes a command from an application to a tool. Standard values for `command` are defined by `omp_control_tool_t` in Section 3.14.

The `modifier` argument passes a command modifier from an application to a tool.

The `command` and `modifier` arguments may have tool-specific values. Tools must ignore `command` values that they are not designed to handle.
The arg argument is a void pointer that enables a tool and an application to exchange arbitrary state. The arg argument may be NULL.

The codeptr_ra argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature ompt_callback_control_tool_t then codeptr_ra contains the return address of the call to that runtime routine. If the implementation of the region is inlined then codeptr_ra contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, codeptr_ra may be NULL.

Constraints on Arguments
Tool-specific values for command must be ≥ 64.

Cross References
• Tool control routine and types, see Section 3.14.

4.5.2.30 ompt_callback_error_t
Summary
The ompt_callback_error_t type is used for callbacks that dispatch runtime-error events.

Format

```c
typedef void (*ompt_callback_error_t) (ompt_severity_t severity, const char *message, size_t length, const void *codeptr_ra);
```

Description
A thread dispatches a registered ompt_callback_error_t callback when an error directive is encountered for which the at(execution) clause is specified.
Description of Arguments

The severity argument passes the specified severity level.

The message argument passes the string from the message clause.

The length argument provides the length of the string.

The codeptr_ra argument relates the implementation of an OpenMP region to its source code. If a runtime routine implements the region associated with a callback that has type signature ompt_callback_error_t then codeptr_ra contains the return address of the call to that runtime routine. If the implementation of the region is inlined then codeptr_ra contains the return address of the invocation of the callback. If attribution to source code is impossible or inappropriate, codeptr_ra may be NULL.

Cross References

• error directive, see Section 2.5.4.
• ompt_severity_t enumeration type, see Section 4.4.24.

4.6 OMPT Runtime Entry Points for Tools

OMPT supports two principal sets of runtime entry points for tools. One set of runtime entry points enables a tool to register callbacks for OpenMP events and to inspect the state of an OpenMP thread while executing in a tool callback or a signal handler. The second set of runtime entry points enables a tool to trace activities on a device. When directed by the tracing interface, an OpenMP implementation will trace activities on a device, collect buffers of trace records, and invoke callbacks on the host to process these records. OMPT runtime entry points should not be global symbols since tools cannot rely on the visibility of such symbols.

OMPT also supports runtime entry points for two classes of lookup routines. The first class of lookup routines contains a single member: a routine that returns runtime entry points in the OMPT callback interface. The second class of lookup routines includes a unique lookup routine for each kind of device that can return runtime entry points in a device’s OMPT tracing interface.

The omp-tools.h C/C++ header file provides the definitions of the types that are specified throughout this subsection.

Binding

The binding thread set for each of the entry points in this section is the encountering thread unless otherwise specified. The binding task set is the task executing on the encountering thread.

Restrictions

Restrictions on OMPT runtime entry points are as follows:

• OMPT runtime entry points must not be called from a signal handler on a native thread before a native-thread-begin or after a native-thread-end event.

• OMPT device runtime entry points must not be called after a device-finalize event for that device.
4.6.1 Entry Points in the OMPT Callback Interface

Entry points in the OMPT callback interface enable a tool to register callbacks for OpenMP events and to inspect the state of an OpenMP thread while executing in a tool callback or a signal handler. Pointers to these runtime entry points are obtained through the lookup function that is provided through the OMPT initializer.

4.6.1.1 omptEnumerateStates_t

Summary

The omptEnumerateStates_t type is the type signature of the omptEnumerateStates runtime entry point, which enumerates the thread states that an OpenMP implementation supports.

Format

```c
typedef int (*omptEnumerateStates_t) ( 
    int current_state, 
    int *next_state, 
    const char **next_state_name
 );
```

Description

An OpenMP implementation may support only a subset of the states defined by the omptState_t enumeration type. An OpenMP implementation may also support implementation-specific states. The omptEnumerateStates runtime entry point, which has type signature omptEnumerateStates_t, enables a tool to enumerate the supported thread states.

When a supported thread state is passed as current_state, the runtime entry point assigns the next thread state in the enumeration to the variable passed by reference in next_state and assigns the name associated with that state to the character pointer passed by reference in next_state_name.

Whenever one or more states are left in the enumeration, the omptEnumerateStates runtime entry point returns 1. When the last state in the enumeration is passed as current_state, omptEnumerateStates returns 0, which indicates that the enumeration is complete.

Description of Arguments

The current_state argument must be a thread state that the OpenMP implementation supports. To begin enumerating the supported states, a tool should pass omptStateUndefined as current_state. Subsequent invocations of omptEnumerateStates should pass the value assigned to the variable that was passed by reference in next_state to the previous call.
The value `ompt_state_undefined` is reserved to indicate an invalid thread state. `ompt_state_undefined` is defined as an integer with the value 0.

The `next_state` argument is a pointer to an integer in which `ompt_enumerate_states` returns the value of the next state in the enumeration.

The `next_state_name` argument is a pointer to a character string pointer through which `ompt_enumerate_states` returns a string that describes the next state.

**Constraints on Arguments**

Any string returned through the `next_state_name` argument must be immutable and defined for the lifetime of program execution.

**Cross References**

- `ompt_state_t` type, see Section 4.4.27.

### 4.6.1.2 `ompt_enumerate_mutex_impls_t`

**Summary**

The `ompt_enumerate_mutex_impls_t` type is the type signature of the `ompt_enumerate_mutex_impls` runtime entry point, which enumerates the kinds of mutual exclusion implementations that an OpenMP implementation employs.

**Format**

```c
typedef int (*ompt_enumerate_mutex_impls_t) (  
  int current_impl,  
  int *next_impl,  
  const char **next_impl_name
);
```

**Description**

Mutual exclusion for locks, `critical` sections, and `atomic` regions may be implemented in several ways. The `ompt_enumerate_mutex_impls` runtime entry point, which has type signature `ompt_enumerate_mutex_impls_t`, enables a tool to enumerate the supported mutual exclusion implementations.

When a supported mutex implementation is passed as `current_impl`, the runtime entry point assigns the next mutex implementation in the enumeration to the variable passed by reference in `next_impl` and assigns the name associated with that mutex implementation to the character pointer passed by reference in `next_impl_name`.

Whenever one or more mutex implementations are left in the enumeration, the `ompt_enumerate_mutex_impls` runtime entry point returns 1. When the last mutex
implementation in the enumeration is passed as current_impl, the runtime entry point returns 0, which indicates that the enumeration is complete.

**Description of Arguments**
The current_impl argument must be a mutex implementation that an OpenMP implementation supports. To begin enumerating the supported mutex implementations, a tool should pass ompt_mutex_impl_none as current_impl. Subsequent invocations of ompt Enumerate Mutex Impl should pass the value assigned to the variable that was passed in next_impl to the previous call.

The value ompt_mutex_impl_none is reserved to indicate an invalid mutex implementation. ompt_mutex_impl_none is defined as an integer with the value 0.

The next_impl argument is a pointer to an integer in which ompt Enumerate Mutex Impl returns the value of the next mutex implementation in the enumeration.

The next_impl_name argument is a pointer to a character string pointer in which ompt Enumerate Mutex Impl returns a string that describes the next mutex implementation.

**Constraints on Arguments**
Any string returned through the next_impl_name argument must be immutable and defined for the lifetime of a program execution.

**Cross References**

- ompt_mutex_t type, see Section 4.4.4.16.

**4.6.1.3 ompt_set_callback_t**

**Summary**
The ompt_set_callback_t type is the type signature of the ompt_set_callback runtime entry point, which registers a pointer to a tool callback that an OpenMP implementation invokes when a host OpenMP event occurs.

**Format**

```
C / C++
typedef ompt_set_result_t (*ompt_set_callback_t) (ompt_callbacks_t event,
                                                  ompt_callback_t callback);
```
Description
OpenMP implementations can use callbacks to indicate the occurrence of events during the execution of an OpenMP program. The `ompt_set_callback` runtime entry point, which has type signature `ompt_set_callback_t`, registers a callback for an OpenMP event on the current device. The return value of `ompt_set_callback` indicates the outcome of registering the callback.

Description of Arguments
The `event` argument indicates the event for which the callback is being registered.

The `callback` argument is a tool callback function. If `callback` is `NULL` then callbacks associated with `event` are disabled. If callbacks are successfully disabled then `ompt_set_always` is returned.

Constraints on Arguments
When a tool registers a callback for an event, the type signature for the callback must match the type signature appropriate for the event.

Restrictions
Restrictions on the `ompt_set_callback` runtime entry point are as follows:

- The entry point must not return `ompt_set_impossible`.

Cross References
- Monitoring activity on the host with OMPT, see Section 4.2.4.
- `ompt_callbacks_t` enumeration type, see Section 4.4.2.
- `ompt_callback_t` type, see Section 4.4.1.
- `ompt_set_result_t` type, see Section 4.4.2.
- `ompt_get_callback_t` host callback type signature, see Section 4.6.1.4.

4.6.1.4 `ompt_get_callback_t`

Summary
The `ompt_get_callback_t` type is the type signature of the `ompt_get_callback` runtime entry point, which retrieves a pointer to a registered tool callback routine (if any) that an OpenMP implementation invokes when a host OpenMP event occurs.

Format

```c
typedef int (*ompt_get_callback_t) (ompt_callbacks_t event, ompt_callback_t *callback);
```

C / C++
Description
The `ompt_get_callback` runtime entry point, which has type signature `ompt_get_callback_t`, retrieves a pointer to the tool callback that an OpenMP implementation may invoke when a host OpenMP event occurs. If a non-null tool callback is registered for the specified event, the pointer to the tool callback is assigned to the variable passed by reference in `callback` and `ompt_get_callback` returns 1; otherwise, it returns 0. If `ompt_get_callback` returns 0, the value of the variable passed by reference as `callback` is undefined.

Description of Arguments
The `event` argument indicates the event for which the callback would be invoked.

The `callback` argument returns a pointer to the callback associated with `event`.

Constraints on Arguments
The `callback` argument cannot be NULL and must point to valid storage.

Cross References
- `ompt_callbacks_t` enumeration type, see Section 4.4.2.
- `ompt_callback_t` type, see Section 4.4.4.1.
- `ompt_set_callback_t` type signature, see Section 4.6.1.3.

4.6.1.5 `ompt_get_thread_data_t`

Summary
The `ompt_get_thread_data_t` type is the type signature of the `ompt_get_thread_data` runtime entry point, which returns the address of the thread data object for the current thread.

Format
```
C / C++
typedef ompt_data_t *(*ompt_get_thread_data_t) (void);
```

Description
Each OpenMP thread can have an associated thread data object of type `ompt_data_t`. The `ompt_get_thread_data` runtime entry point, which has type signature `ompt_get_thread_data_t`, retrieves a pointer to the thread data object, if any, that is associated with the current thread. A tool may use a pointer to an OpenMP thread’s data object that `ompt_get_thread_data` retrieves to inspect or to modify the value of the data object. When an OpenMP thread is created, its data object is initialized with value `ompt_data_none`.

This runtime entry point is `async signal safe`.

Cross References
- `ompt_data_t` type, see Section 4.4.4.4.
4.6.1.6 ompt_get_num_procs_t

Summary
The ompt_get_num_procs_t type is the type signature of the ompt_get_num_procs runtime entry point, which returns the number of processors currently available to the execution environment on the host device.

Format
```c
typedef int (*ompt_get_num_procs_t) (void);
```

Binding
The binding thread set is all threads on the host device.

Description
The ompt_get_num_procs runtime entry point, which has type signature ompt_get_num_procs_t, returns the number of processors that are available on the host device at the time the routine is called. This value may change between the time that it is determined and the time that it is read in the calling context due to system actions outside the control of the OpenMP implementation.

This runtime entry point is async signal safe.

4.6.1.7 ompt_get_num_places_t

Summary
The ompt_get_num_places_t type is the type signature of the ompt_get_num_places runtime entry point, which returns the number of places currently available to the execution environment in the place list.

Format
```c
typedef int (*ompt_get_num_places_t) (void);
```

Binding
The binding thread set is all threads on a device.

Description
The ompt_get_num_places runtime entry point, which has type signature ompt_get_num_places_t, returns the number of places in the place list. This value is equivalent to the number of places in the place-partition-var ICV in the execution environment of the initial task.

This runtime entry point is async signal safe.
Cross References

- `place-partition-var ICV`, see Section 2.4.
- `OMP_PLACES` environment variable, see Section 6.5.

4.6.1.8 ompt_get_place_proc_ids_t

Summary
The `ompt_get_place_proc_ids_t` type is the type signature of the `ompt_get_num_place_procs_ids` runtime entry point, which returns the numerical identifiers of the processors that are available to the execution environment in the specified place.

Format

```c
typedef int (*ompt_get_place_proc_ids_t) (  
    int place_num,  
    int ids_size,  
    int *ids  
);  
```

Binding
The binding thread set is all threads on a device.

Description
The `ompt_get_place_proc_ids` runtime entry point, which has type signature `ompt_get_place_proc_ids_t`, returns the numerical identifiers of each processor that is associated with the specified place. These numerical identifiers are non-negative, and their meaning is implementation defined.

Description of Arguments
The `place_num` argument specifies the place that is being queried.

The `ids` argument is an array in which the routine can return a vector of processor identifiers in the specified place.

The `ids_size` argument indicates the size of the result array that is specified by `ids`.

Effect
If the `ids` array of size `ids_size` is large enough to contain all identifiers then they are returned in `ids` and their order in the array is implementation defined. Otherwise, if the `ids` array is too small, the values in `ids` when the function returns are unspecified. The routine always returns the number of numerical identifiers of the processors that are available to the execution environment in the specified place.
4.6.1.9 ompt_get_place_num_t

Summary
The ompt_get_place_num_t type is the type signature of the ompt_get_place_num runtime entry point, which returns the place number of the place to which the current thread is bound.

Format
```c/c++
typedef int (*ompt_get_place_num_t) (void);
```

Description
When the current thread is bound to a place, ompt_get_place_num returns the place number associated with the thread. The returned value is between 0 and one less than the value returned by ompt_get_num_places, inclusive. When the current thread is not bound to a place, the routine returns -1.

This runtime entry point is async signal safe.

4.6.1.10 ompt_get_partition_place_nums_t

Summary
The ompt_get_partition_place_nums_t type is the type signature of the ompt_get_partition_place_nums runtime entry point, which returns a list of place numbers that correspond to the places in the place-partition-var ICV of the innermost implicit task.

Format
```c/c++
typedef int (*ompt_get_partition_place_nums_t) (int place_nums_size,
                                       int *place_nums);
```

Description
The ompt_get_partition_place_nums runtime entry point, which has type signature ompt_get_partition_place_nums_t, returns a list of place numbers that correspond to the places in the place-partition-var ICV of the innermost implicit task.

This runtime entry point is async signal safe.
Description of Arguments

The `place_nums` argument is an array in which the routine can return a vector of place identifiers.

The `place_nums_size` argument indicates the size of the result array that the `place_nums` argument specifies.

Effect

If the `place_nums` array of size `place_nums_size` is large enough to contain all identifiers then they are returned in `place_nums` and their order in the array is implementation defined. Otherwise, if the `place_nums` array is too small, the values in `place_nums` when the function returns are unspecified.

The routine always returns the number of places in the `place-partition-var ICV` of the innermost implicit task.

Cross References

- `place-partition-var ICV`, see Section 2.4.
- `OMP_PLACES` environment variable, see Section 6.5.

4.6.1.11 ompt_get_proc_id_t

Summary

The `ompt_get_proc_id_t` type is the type signature of the `ompt_get_proc_id` runtime entry point, which returns the numerical identifier of the processor of the current thread.

Format

```
C / C++
typedef int (*ompt_get_proc_id_t) (void);
```

Description

The `ompt_get_proc_id` runtime entry point, which has type signature `ompt_get_proc_id_t`, returns the numerical identifier of the processor of the current thread. A defined numerical identifier is non-negative, and its meaning is implementation defined. A negative number indicates a failure to retrieve the numerical identifier.

This runtime entry point is `async signal safe`.

4.6.1.12 ompt_get_state_t

Summary

The `ompt_get_state_t` type is the type signature of the `ompt_get_state` runtime entry point, which returns the state and the wait identifier of the current thread.
typedef int (*ompt_get_state_t)(
    ompt_wait_id_t *wait_id
);

Description
Each OpenMP thread has an associated state and a wait identifier. If a thread’s state indicates that
the thread is waiting for mutual exclusion then its wait identifier contains an opaque handle that
indicates the data object upon which the thread is waiting. The `ompt_get_state` runtime entry
point, which has type signature `ompt_get_state_t`, retrieves the state and wait identifier of the
current thread. The returned value may be any one of the states predefined by `ompt_state_t` or
a value that represents an implementation-specific state. The tool may obtain a string representation
for each state with the `ompt_enumerate_states` function.

If the returned state indicates that the thread is waiting for a lock, nest lock, `critical` region,
atomic region, or `ordered` region then the value of the thread’s wait identifier is assigned to a
non-null wait identifier passed as the `wait_id` argument.

This runtime entry point is async signal safe.

Description of Arguments
The `wait_id` argument is a pointer to an opaque handle that is available to receive the value of the
wait identifier of the thread. If `wait_id` is not `NULL` then the entry point assigns the value of the
wait identifier of the thread to the object to which `wait_id` points. If the returned state is not one of
the specified wait states then the value of opaque object to which `wait_id` points is undefined after
the call.

Constraints on Arguments
The argument passed to the entry point must be a reference to a variable of the specified type or
`NULL`.

Cross References
- `ompt_state_t` type, see Section 4.4.4.27.
- `ompt_wait_id_t` type, see Section 4.4.4.30.
- `ompt_enumerate_states_t` type, see Section 4.6.1.1.

4.6.1.13 ompt_get_parallel_info_t
Summary
The `ompt_get_parallel_info_t` type is the type signature of the
`ompt_get_parallel_info` runtime entry point, which returns information about the parallel
region, if any, at the specified ancestor level for the current execution context.
C / C++

```c
typedef int (*ompt_get_parallel_info_t) (
    int ancestor_level,
    ompt_data_t **parallel_data,
    int *team_size
);
```

Description

During execution, an OpenMP program may employ nested parallel regions. The `ompt_get_parallel_info` runtime entry point, which has type signature `ompt_get_parallel_info_t`, retrieves information about the current parallel region and any enclosing parallel regions for the current execution context. The entry point returns 2 if a parallel region exists at the specified ancestor level and the information is available, 1 if a parallel region exists at the specified ancestor level but the information is currently unavailable, and 0 otherwise.

A tool may use the pointer to the data object of a parallel region that it obtains from this runtime entry point to inspect or to modify the value of the data object. When a parallel region is created, its data object will be initialized with the value `ompt_data_none`. This runtime entry point is async signal safe.

Between a `parallel-begin` event and an `implicit-task-begin` event, a call to `ompt_get_parallel_info(0,...)` may return information about the outer parallel team, the new parallel team or an inconsistent state.

If a thread is in the state `ompt_state_wait_barrier_implicit_parallel` then a call to `ompt_get_parallel_info` may return a pointer to a copy of the specified parallel region’s `parallel_data` rather than a pointer to the data word for the region itself. This convention enables the primary thread for a parallel region to free storage for the region immediately after the region ends, yet avoid having some other thread in the team that is executing the region potentially reference the `parallel_data` object for the region after it has been freed.

Description of Arguments

The `ancestor_level` argument specifies the parallel region of interest by its ancestor level. Ancestor level 0 refers to the innermost parallel region; information about enclosing parallel regions may be obtained using larger values for `ancestor_level`.

The `parallel_data` argument returns the parallel data if the argument is not `NULL`.

The `team_size` argument returns the team size if the argument is not `NULL`. 
Effect

If the runtime entry point returns 0 or 1, no argument is modified. Otherwise, `ompt_get_parallel_info` has the following effects:

- If a non-null value was passed for `parallel_data`, the value returned in `parallel_data` is a pointer to a data word that is associated with the parallel region at the specified level; and

- If a non-null value was passed for `team_size`, the value returned in the integer to which `team_size` point is the number of threads in the team that is associated with the parallel region.

Constraints on Arguments

While argument `ancestor_level` is passed by value, all other arguments to the entry point must be pointers to variables of the specified types or `NULL`.

Cross References

- `ompt_data_t` type, see Section 4.4.4.4.

4.6.1.14 `ompt_get_task_info_t`

Summary

The `ompt_get_task_info_t` type is the type signature of the `ompt_get_task_info` runtime entry point, which returns information about the task, if any, at the specified ancestor level in the current execution context.

Format

```c
typedef int (*ompt_get_task_info_t) (  
    int ancestor_level,  
    int *flags,  
    ompt_data_t **task_data,  
    ompt_frame_t **task_frame,  
    ompt_data_t **parallel_data,  
    int *thread_num  
);```

Description

During execution, an OpenMP thread may be executing an OpenMP task. Additionally, the stack of the thread may contain procedure frames that are associated with suspended OpenMP tasks or OpenMP runtime system routines. To obtain information about any task on the stack of the current thread, a tool uses the `ompt_get_task_info` runtime entry point, which has type signature `ompt_get_task_info_t`.

Ancestor level 0 refers to the active task; information about other tasks with associated frames present on the stack in the current execution context may be queried at higher ancestor levels.
The `ompt_get_task_info` runtime entry point returns 2 if a task region exists at the specified ancestor level and the information is available, 1 if a task region exists at the specified ancestor level but the information is currently unavailable, and 0 otherwise.

If a task exists at the specified ancestor level and the information is available then information is returned in the variables passed by reference to the entry point. If no task region exists at the specified ancestor level or the information is unavailable then the values of variables passed by reference to the entry point are undefined when `ompt_get_task_info` returns.

A tool may use a pointer to a data object for a task or parallel region that it obtains from `ompt_get_task_info` to inspect or to modify the value of the data object. When either a parallel region or a task region is created, its data object will be initialized with the value `ompt_data_none`.

This runtime entry point is *async signal safe*.

**Description of Arguments**

The `ancestor_level` argument specifies the task region of interest by its ancestor level. Ancestor level 0 refers to the active task; information about ancestor tasks found in the current execution context may be queried at higher ancestor levels.

The `flags` argument returns the task type if the argument is not `NULL`.

The `task_data` argument returns the task data if the argument is not `NULL`.

The `task_frame` argument returns the task frame pointer if the argument is not `NULL`.

The `parallel_data` argument returns the parallel data if the argument is not `NULL`.

The `thread_num` argument returns the thread number if the argument is not `NULL`.

**Effect**

If the runtime entry point returns 0 or 1, no argument is modified. Otherwise, `ompt_get_task_info` has the following effects:

- If a non-null value was passed for `flags` then the value returned in the integer to which `flags` points represents the type of the task at the specified level; possible task types include initial, implicit, explicit, and target tasks;

- If a non-null value was passed for `task_data` then the value that is returned in the object to which it points is a pointer to a data word that is associated with the task at the specified level;

- If a non-null value was passed for `task_frame` then the value that is returned in the object to which `task_frame` points is a pointer to the `ompt_frame_t` structure that is associated with the task at the specified level;

- If a non-null value was passed for `parallel_data` then the value that is returned in the object to which `parallel_data` points is a pointer to a data word that is associated with the parallel region that contains the task at the specified level or, if the task at the specified level is an initial task, `NULL`; and
If a non-null value was passed for `thread_num`, then the value that is returned in the object to which `thread_num` points indicates the number of the thread in the parallel region that is executing the task at the specified level.

**Constraints on Arguments**

While argument `ancestor_level` is passed by value, all other arguments to `ompt_get_task_info` must be pointers to variables of the specified types or NULL.

**Cross References**

- `ompt_data_t` type, see Section 4.4.4.4.
- `ompt_task_flag_t` type, see Section 4.4.4.18.
- `ompt_frame_t` type, see Section 4.4.4.28.

### 4.6.1.15 ompt_get_task_memory_t

**Summary**

The `ompt_get_task_memory_t` type is the type signature of the `ompt_get_task_memory` runtime entry point, which returns information about memory ranges that are associated with the task.

**Format**

```c/c++
typedef int (*ompt_get_task_memory_t)(
    void **addr,
    size_t *size,
    int block
);
```

**Description**

During execution, an OpenMP thread may be executing an OpenMP task. The OpenMP implementation must preserve the data environment from the creation of the task for the execution of the task. The `ompt_get_task_memory` runtime entry point, which has type signature `ompt_get_task_memory_t`, provides information about the memory ranges used to store the data environment for the current task.

Multiple memory ranges may be used to store these data. The `block` argument supports iteration over these memory ranges.

The `ompt_get_task_memory` runtime entry point returns 1 if more memory ranges are available, and 0 otherwise. If no memory is used for a task, `size` is set to 0. In this case, `addr` is unspecified.

This runtime entry point is async signal safe.
Description of Arguments
The `addr` argument is a pointer to a void pointer return value to provide the start address of a memory block.

The `size` argument is a pointer to a size type return value to provide the size of the memory block.

The `block` argument is an integer value to specify the memory block of interest.

4.6.1.16 ompt_get_target_info_t

Summary
The `ompt_get_target_info_t` type is the type signature of the `ompt_get_target_info` runtime entry point, which returns identifiers that specify a thread’s current target region and target operation ID, if any.

Format

```c
typedef int (*ompt_get_target_info_t) ( 
    uint64_t *device_num, 
    ompt_id_t *target_id, 
    ompt_id_t *host_op_id
);
```

Description
The `ompt_get_target_info` entry point, which has type signature `ompt_get_target_info_t`, returns 1 if the current thread is in a target region and 0 otherwise. If the entry point returns 0 then the values of the variables passed by reference as its arguments are undefined.

If the current thread is in a target region then `ompt_get_target_info` returns information about the current device, active target region, and active host operation, if any.

This runtime entry point is `async signal safe`.

Description of Arguments
The `device_num` argument returns the device number if the current thread is in a target region.

The `target_id` argument returns the target region identifier if the current thread is in a target region.

If the current thread is in the process of initiating an operation on a target device (for example, copying data to or from an accelerator or launching a kernel), then `host_op_id` returns the identifier for the operation; otherwise, `host_op_id` returns `ompt_id_none`. 
Constraints on Arguments
Arguments passed to the entry point must be valid references to variables of the specified types.

Cross References
• ompt_id_t type, see Section 4.4.3.

4.6.1.17 ompt_get_num_devices_t
Summary
The ompt_get_num_devices_t type is the type signature of the ompt_get_num_devices runtime entry point, which returns the number of available devices.

Format

```c
typedef int (*ompt_get_num_devices_t) (void);
```

Description
The ompt_get_num_devices runtime entry point, which has type signature ompt_get_num_devices_t, returns the number of devices available to an OpenMP program. This runtime entry point is async signal safe.

4.6.1.18 ompt_get_unique_id_t
Summary
The ompt_get_unique_id_t type is the type signature of the ompt_get_unique_id runtime entry point, which returns a unique number.

Format

```c
typedef uint64_t (*ompt_get_unique_id_t) (void);
```

Description
The ompt_get_unique_id runtime entry point, which has type signature ompt_get_unique_id_t, returns a number that is unique for the duration of an OpenMP program. Successive invocations may not result in consecutive or even increasing numbers. This runtime entry point is async signal safe.
4.6.1.19 ompt_finalize_tool_t

Summary
The ompt_finalize_tool_t type is the type signature of the ompt_finalize_tool runtime entry point, which enables a tool to finalize itself.

Format
C / C++

```c
typedef void (*ompt_finalize_tool_t) (void);
```

Description
A tool may detect that the execution of an OpenMP program is ending before the OpenMP implementation does. To facilitate clean termination of the tool, the tool may invoke the ompt_finalize_tool runtime entry point, which has type signature ompt_finalize_tool_t. Upon completion of ompt_finalize_tool, no OMPT callbacks are dispatched.

Effect
The ompt_finalize_tool routine detaches the tool from the runtime, unregisters all callbacks and invalidates all OMPT entry points passed to the tool in the lookup-function. Upon completion of ompt_finalize_tool, no further callbacks will be issued on any thread.

Before the callbacks are unregistered, the OpenMP runtime should attempt to dispatch all outstanding registered callbacks as well as the callbacks that would be encountered during shutdown of the runtime, if possible in the current execution context.

4.6.2 Entry Points in the OMPT Device Tracing Interface

The runtime entry points with type signatures of the types that are specified in this section enable a tool to trace activities on a device.

4.6.2.1 ompt_get_device_num_procs_t

Summary
The ompt_get_device_num_procs_t type is the type signature of the ompt_get_device_num_procs runtime entry point, which returns the number of processors currently available to the execution environment on the specified device.

Format
C / C++

```c
typedef int (*ompt_get_device_num_procs_t) (ompt_device_t *device);
```
Description
The `omp_get_device_num_procs` runtime entry point, which has type signature `omp_get_device_num_procs_t`, returns the number of processors that are available on the device at the time the routine is called. This value may change between the time that it is determined and the time that it is read in the calling context due to system actions outside the control of the OpenMP implementation.

Description of Arguments
The `device` argument is a pointer to an opaque object that represents the target device instance. The pointer to the device instance object is used by functions in the device tracing interface to identify the device being addressed.

Cross References
• `omp_device_t` type, see Section 4.4.5.

4.6.2.2 `omp_get_device_time_t`

Summary
The `omp_get_device_time_t` type is the type signature of the `omp_get_device_time` runtime entry point, which returns the current time on the specified device.

Format
```
C / C++
typedef ompt_device_time_t (*ompt_get_device_time_t) (ompt_device_t *device);
```

Description
Host and target devices are typically distinct and run independently. If host and target devices are different hardware components, they may use different clock generators. For this reason, a common time base for ordering host-side and device-side events may not be available.

The `omp_get_device_time` runtime entry point, which has type signature `omp_get_device_time_t`, returns the current time on the specified device. A tool can use this information to align time stamps from different devices.

Description of Arguments
The `device` argument is a pointer to an opaque object that represents the target device instance. The pointer to the device instance object is used by functions in the device tracing interface to identify the device being addressed.
Cross References
• ompt_device_t type, see Section 4.4.5.
• ompt_device_time_t type, see Section 4.4.6.

4.6.2.3 ompt_translate_time_t

Summary
The ompt_translate_time_t type is the type signature of the ompt_translate_time runtime entry point, which translates a time value that is obtained from the specified device to a corresponding time value on the host device.

Format

C / C++

```c
typedef double (*ompt_translate_time_t) (ompt_device_t *device,
ompt_device_time_t time);
```

Description
The ompt_translate_time runtime entry point, which has type signature ompt_translate_time_t, translates a time value obtained from the specified device to a corresponding time value on the host device. The returned value for the host time has the same meaning as the value returned from omp_get_wtime.

Note – The accuracy of time translations may degrade, if they are not performed promptly after a device time value is received and if either the host or device vary their clock speeds. Prompt translation of device times to host times is recommended.

Description of Arguments
The device argument is a pointer to an opaque object that represents the target device instance. The pointer to the device instance object is used by functions in the device tracing interface to identify the device being addressed.

The time argument is a time from the specified device.

Cross References
• omp_get_wtime routine, see Section 3.10.1.
• ompt_device_t type, see Section 4.4.5.
• ompt_device_time_t type, see Section 4.4.6.
4.6.2.4 ompt_set_trace_ompt_t

Summary
The ompt_set_trace_ompt_t type is the type signature of the ompt_set_trace_ompt runtime entry point, which enables or disables the recording of trace records for one or more types of OMPT events.

Format

```c
typedef ompt_set_result_t (*ompt_set_trace_ompt_t) (
    ompt_device_t *device,
    unsigned int enable,
    unsigned int etype
);
```

Description of Arguments

The device argument points to an opaque object that represents the target device instance. Functions in the device tracing interface use this pointer to identify the device that is being addressed.

The etype argument indicates the events to which the invocation of ompt_set_trace_ompt applies. If the value of etype is 0 then the invocation applies to all events. If etype is positive then it applies to the event in ompt_callbacks_t that matches that value.

The enable argument indicates whether tracing should be enabled or disabled for the event or events that the etype argument specifies. A positive value for enable indicates that recording should be enabled; a value of 0 for enable indicates that recording should be disabled.

Restrictions

Restrictions on the ompt_set_trace_ompt runtime entry point are as follows:

- The entry point must not return ompt_set_sometimes_paired.

Cross References

- Tracing activity on target devices with OMPT, see Section 4.2.5.
- ompt_callbacks_t type, see Section 4.4.2.
- ompt_set_result_t type, see Section 4.4.4.2.
- ompt_device_t type, see Section 4.4.4.5.
### 4.6.2.5 ompt_set_trace_native_t

**Summary**
The `ompt_set_trace_native_t` type is the type signature of the `ompt_set_trace_native` runtime entry point, which enables or disables the recording of native trace records for a device.

**Format**
```c
C / C++
typedef ompt_set_result_t (*ompt_set_trace_native_t) (
    ompt_device_t *device,
    int enable,
    int flags
);
```

**Description**
This interface is designed for use by a tool that cannot directly use native control functions for the device. If a tool can directly use the native control functions then it can invoke native control functions directly using pointers that the `lookup` function associated with the device provides and that are described in the `documentation` string that is provided to the device initializer callback.

**Description of Arguments**
The `device` argument points to an opaque object that represents the target device instance. Functions in the device tracing interface use this pointer to identify the device that is being addressed.

The `enable` argument indicates whether this invocation should enable or disable recording of events.

The `flags` argument specifies the kinds of native device monitoring to enable or to disable. Each kind of monitoring is specified by a flag bit. Flags can be composed by using logical `OR` to combine enumeration values from type `ompt_native_mon_flag_t`.

To start, to pause, to flush, or to stop tracing for a specific target device associated with `device`, a tool invokes the `ompt_start_trace`, `ompt_pause_trace`, `ompt_flush_trace`, or `ompt_stop_trace` runtime entry point for the device.

**Restrictions**
Restrictions on the `ompt_set_trace_native` runtime entry point are as follows:

- The entry point must not return `ompt_set_sometimes_paired`.

**Cross References**
- Tracing activity on target devices with OMPT, see Section 4.2.5.
- `ompt_set_result_t` type, see Section 4.4.4.2.
- `ompt_device_t` type, see Section 4.4.4.5.
4.6.2.6 ompt_start_trace_t

Summary
The ompt_start_trace_t type is the type signature of the ompt_start_trace runtime entry point, which starts tracing of activity on a specific device.

Format

```c
typedef int (*ompt_start_trace_t) (ompt_device_t *device,
                                  ompt_callback_buffer_request_t request,
                                  ompt_callback_buffer_complete_t complete);
```

Description
A device’s ompt_start_trace runtime entry point, which has type signature ompt_start_trace_t, initiates tracing on the device. Under normal operating conditions, every event buffer provided to a device by a tool callback is returned to the tool before the OpenMP runtime shuts down. If an exceptional condition terminates execution of an OpenMP program, the OpenMP runtime may not return buffers provided to the device.

An invocation of ompt_start_trace returns 1 if the command succeeds and 0 otherwise.

Description of Arguments
The device argument points to an opaque object that represents the target device instance. Functions in the device tracing interface use this pointer to identify the device that is being addressed.

The request argument specifies a tool callback that supplies a buffer in which a device can deposit events.

The complete argument specifies a tool callback that is invoked by the OpenMP implementation to empty a buffer that contains event records.

Cross References
- ompt_device_t type, see Section 4.4.4.5.
- ompt_callback_buffer_request_t callback type, see Section 4.5.2.23.
- ompt_callback_buffer_complete_t callback type, see Section 4.5.2.24.

4.6.2.7 ompt_pause_trace_t

Summary
The ompt_pause_trace_t type is the type signature of the ompt_pause_trace runtime entry point, which pauses or restarts activity tracing on a specific device.
**Format**

```c
typedef int (*ompt_pause_trace_t) (
    ompt_device_t *device,
    int begin_pause
);
```

**Description**

A device’s `ompt_pause_trace` runtime entry point, which has type signature `ompt_pause_trace_t`, pauses or resumes tracing on a device. An invocation of `ompt_pause_trace` returns 1 if the command succeeds and 0 otherwise. Redundant pause or resume commands are idempotent and will return the same value as the prior command.

**Description of Arguments**

The `device` argument points to an opaque object that represents the target device instance. Functions in the device tracing interface use this pointer to identify the device that is being addressed.

The `begin_pause` argument indicates whether to pause or to resume tracing. To resume tracing, zero should be supplied for `begin_pause`; To pause tracing, any other value should be supplied.

**Cross References**

- `ompt_device_t` type, see Section 4.4.4.5.

**4.6.2.8 ompt_flush_trace_t**

**Summary**

The `ompt_flush_trace_t` type is the type signature of the `ompt_flush_trace` runtime entry point, which causes all pending trace records for the specified device to be delivered.

```c
typedef int (*ompt_flush_trace_t) (
    ompt_device_t *device
);
```

**Description**

A device’s `ompt_flush_trace` runtime entry point, which has type signature `ompt_flush_trace_t`, causes the OpenMP implementation to issue a sequence of zero or more buffer completion callbacks to deliver all trace records that have been collected prior to the flush. An invocation of `ompt_flush_trace` returns 1 if the command succeeds and 0 otherwise.
Description of Arguments
The device argument points to an opaque object that represents the target device instance. Functions in the device tracing interface use this pointer to identify the device that is being addressed.

Cross References
• ompt_device_t type, see Section 4.4.4.5.

4.6.2.9 ompt_stop_trace_t
Summary
The ompt_stop_trace_t type is the type signature of the ompt_stop_trace runtime entry point, which stops tracing for a device.

Format
```c
typedef int (*ompt_stop_trace_t) (ompt_device_t *device);
```

Description
A device’s ompt_stop_trace runtime entry point, which has type signature ompt_stop_trace_t, halts tracing on the device and requests that any pending trace records are flushed. An invocation of ompt_stop_trace returns 1 if the command succeeds and 0 otherwise.

Description of Arguments
The device argument points to an opaque object that represents the target device instance. Functions in the device tracing interface use this pointer to identify the device that is being addressed.

Cross References
• ompt_device_t type, see Section 4.4.4.5.

4.6.2.10 ompt_advance_buffer_cursor_t
Summary
The ompt_advance_buffer_cursor_t type is the type signature of the ompt_advance_buffer_cursor runtime entry point, which advances a trace buffer cursor to the next record.
Format

```c
typedef int (*ompt_advance_buffer_cursor_t) (  
    ompt_device_t *device,  
    ompt_buffer_t *buffer,  
    size_t size,  
    ompt_buffer_cursor_t current,  
    ompt_buffer_cursor_t *next
);
```

Description

A device’s `ompt_advance_buffer_cursor` runtime entry point, which has type signature `ompt_advance_buffer_cursor_t`, advances a trace buffer pointer to the next trace record. An invocation of `ompt_advance_buffer_cursor` returns `true` if the advance is successful and the next position in the buffer is valid.

Description of Arguments

The `device` argument points to an opaque object that represents the target device instance. Functions in the device tracing interface use this pointer to identify the device that is being addressed.

The `buffer` argument indicates a trace buffer that is associated with the cursors.

The argument `size` indicates the size of `buffer` in bytes.

The `current` argument is an opaque buffer cursor.

The `next` argument returns the next value of an opaque buffer cursor.

Cross References

- `ompt_device_t` type, see Section 4.4.4.5.
- `ompt_buffer_cursor_t` type, see Section 4.4.4.8.

### 4.6.2.11 ompt_get_record_type_t

Summary

The `ompt_get_record_type_t` type is the type signature of the `ompt_get_record_type` runtime entry point, which inspects the type of a trace record.

Format

```c
typedef ompt_record_t (*ompt_get_record_type_t) (  
    ompt_buffer_t *buffer,  
    ompt_buffer_cursor_t current
);
```
Description
Trace records for a device may be in one of two forms: native record format, which may be device-specific, or OMPT record format, in which each trace record corresponds to an OpenMP event and most fields in the record structure are the arguments that would be passed to the OMPT callback for the event.

A device’s `ompt_get_record_type` runtime entry point, which has type signature `ompt_get_record_type_t`, inspects the type of a trace record and indicates whether the record at the current position in the trace buffer is an OMPT record, a native record, or an invalid record. An invalid record type is returned if the cursor is out of bounds.

Description of Arguments
The `buffer` argument indicates a trace buffer.

The `current` argument is an opaque buffer cursor.

Cross References
- `ompt_record_t` type, see Section 4.4.3.1.
- `ompt_buffer_t` type, see Section 4.4.4.7.
- `ompt_buffer_cursor_t` type, see Section 4.4.4.8.

4.6.2.12 `ompt_get_record_ompt_t`

Summary
The `ompt_get_record_ompt_t` type is the type signature of the `ompt_get_record_ompt` runtime entry point, which obtains a pointer to an OMPT trace record from a trace buffer associated with a device.

Format
```c
C / C++
typedef ompt_record_ompt_t *(*ompt_get_record_ompt_t) (  
    ompt_buffer_t *buffer,  
    ompt_buffer_cursor_t current  
);  
```

Description
A device’s `ompt_get_record_ompt` runtime entry point, which has type signature `ompt_get_record_ompt_t`, returns a pointer that may point to a record in the trace buffer, or it may point to a record in thread local storage in which the information extracted from a record was assembled. The information available for an event depends upon its type.

The return value of the `ompt_record_ompt_t` type includes a field of a union type that can represent information for any OMPT event record type. Another call to the runtime entry point may overwrite the contents of the fields in a record returned by a prior invocation.
Description of Arguments
The `buffer` argument indicates a trace buffer.

The `current` argument is an opaque buffer cursor.

Cross References
- `ompt_record_ompt_t` type, see Section 4.4.3.4.
- `ompt_device_t` type, see Section 4.4.4.5.
- `ompt_buffer_cursor_t` type, see Section 4.4.4.8.

4.6.2.13 `ompt_get_record_native_t`

Summary
The `ompt_get_record_native_t` type is the type signature of the `ompt_get_record_native` runtime entry point, which obtains a pointer to a native trace record from a trace buffer associated with a device.

Format

```c/
c++
typedef void *(*ompt_get_record_native_t) (  
    ompt_buffer_t *buffer,  
    ompt_buffer_cursor_t current,  
    ompt_id_t *host_op_id
  ) ;
```

Description
A device’s `ompt_get_record_native` runtime entry point, which has type signature `ompt_get_record_native_t`, returns a pointer that may point into the specified trace buffer, or into thread local storage in which the information extracted from a trace record was assembled. The information available for a native event depends upon its type. If the function returns a non-null result, it will also set the object to which `host_op_id` points to a host-side identifier for the operation that is associated with the record. A subsequent call to `ompt_get_record_native` may overwrite the contents of the fields in a record returned by a prior invocation.

Description of Arguments
The `buffer` argument indicates a trace buffer.

The `current` argument is an opaque buffer cursor.

The `host_op_id` argument is a pointer to an identifier that is returned by the function. The entry point sets the identifier to which `host_op_id` points to the value of a host-side identifier for an operation on a target device that was created when the operation was initiated by the host.
Cross References

• ompt_id_t type, see Section 4.4.3.
• ompt_buffer_t type, see Section 4.4.7.
• ompt_buffer_cursor_t type, see Section 4.4.8.

4.6.2.14 ompt_get_record_abstract_t

Summary
The ompt_get_record_abstract_t type is the type signature of the
ompt_get_record_abstract runtime entry point, which summarizes the context of a native
(device-specific) trace record.

Format
```
C / C++
```
```
typedef ompt_record_abstract_t *(*ompt_get_record_abstract_t) (void *native_record);
```

Description
An OpenMP implementation may execute on a device that logs trace records in a native
(device-specific) format that a tool cannot interpret directly. The
ompt_get_record_abstract runtime entry point of a device, which has type signature
ompt_get_record_abstract_t, translates a native trace record into a standard form.

Description of Arguments
The native_record argument is a pointer to a native trace record.

Cross References
• ompt_record_abstract_t type, see Section 4.4.3.3.

4.6.3 Lookup Entry Points: ompt_function_lookup_t

Summary
The ompt_function_lookup_t type is the type signature of the lookup runtime entry points
that provide pointers to runtime entry points that are part of the OMPT interface.

Format
```
C / C++
```
```
typedef void (*ompt_interface_fn_t) (void);
```
```
typedef ompt_interface_fn_t (*ompt_function_lookup_t) (const char *interface_function_name);
```

}
Description
An OpenMP implementation provides a pointer to a lookup routine that provides pointers to OMPT runtime entry points. When the implementation invokes a tool initializer to configure the OMPT callback interface, it provides a lookup function that provides pointers to runtime entry points that implement routines that are part of the OMPT callback interface. Alternatively, when it invokes a tool initializer to configure the OMPT tracing interface for a device, it provides a lookup function that provides pointers to runtime entry points that implement tracing control routines appropriate for that device.

If the provided function name is unknown to the OpenMP implementation, the function returns NULL. In a compliant implementation, the lookup function provided by the tool initializer for the OMPT callback interface returns a valid function pointer for any OMPT runtime entry point name listed in Table 4.1.

A compliant implementation of a lookup function passed to a tool’s ompt_device_initialize callback must provide non-NULL function pointers for all strings in Table 4.4, except for ompt_set_trace_ompt and ompt_get_record_ompt, as described in Section 4.2.5.

Description of Arguments
The interface_function_name argument is a C string that represents the name of a runtime entry point.

Cross References
• Tool initializer for a device’s OMPT tracing interface, see Section 4.2.5.
• Tool initializer for the OMPT callback interface, see Section 4.5.1.1.
• Entry points in the OMPT callback interface, see Table 4.1 for a list and Section 4.6.1 for detailed definitions.
• Entry points in the OMPT tracing interface, see Table 4.4 for a list and Section 4.6.2 for detailed definitions.
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5 OMPD Interface

This chapter describes OMPD, which is an interface for third-party tools. Third-party tools exist in
separate processes from the OpenMP program. To provide OMPD support, an OpenMP
implementation must provide an OMPD library that the third-party tool can load. An OpenMP
implementation does not need to maintain any extra information to support OMPD inquiries from
third-party tools unless it is explicitly instructed to do so.

OMPD allows third-party tools such as debuggers to inspect the OpenMP state of a live program or
core file in an implementation-agnostic manner. That is, a third-party tool that uses OMPD should
work with any conforming OpenMP implementation. An OpenMP implementer provides a library
for OMPD that a third-party tool can dynamically load. The third-party tool can use the interface
exported by the OMPD library to inspect the OpenMP state of a program. In order to satisfy
requests from the third-party tool, the OMPD library may need to read data from the OpenMP
program, or to find the addresses of symbols in it. The OMPD library provides this functionality
through a callback interface that the third-party tool must instantiate for the OMPD library.

To use OMPD, the third-party tool loads the OMPD library. The OMPD library exports the API
that is defined throughout this section, and the third-party tool uses the API to determine OpenMP
information about the OpenMP program. The OMPD library must look up the symbols and read
data out of the program. It does not perform these operations directly but instead directs the third-
party tool to perform them by using the callback interface that the third-party tool exports.

The OMPD design insulates third-party tools from the internal structure of the OpenMP runtime,
while the OMPD library is insulated from the details of how to access the OpenMP program. This
decoupled design allows for flexibility in how the OpenMP program and third-party tool are
deployed, so that, for example, the third-party tool and the OpenMP program are not required to
execute on the same machine.

Generally, the third-party tool does not interact directly with the OpenMP runtime but instead
interacts with the runtime through the OMPD library. However, a few cases require the third-party
tool to access the OpenMP runtime directly. These cases fall into two broad categories. The first is
during initialization where the third-party tool must look up symbols and read variables in the
OpenMP runtime in order to identify the OMPD library that it should use, which is discussed in
Section 5.2.2 and Section 5.2.3. The second category relates to arranging for the third-party tool to
be notified when certain events occur during the execution of the OpenMP program. For this
purpose, the OpenMP implementation must define certain symbols in the runtime code, as is
discussed in Section 5.6. Each of these symbols corresponds to an event type. The OpenMP
runtime must ensure that control passes through the appropriate named location when events occur.
If the third-party tool requires notification of an event, it can plant a breakpoint at the matching
location. The location can, but may not, be a function. It can, for example, simply be a label. However, the names of the locations must have external C linkage.

### 5.1 OMPD Interfaces Definitions

A compliant implementation must supply a set of definitions for the OMPD runtime entry points, OMPD third-party tool callback signatures, third-party tool interface functions and the special data types of their parameters and return values. These definitions, which are listed throughout this chapter, and their associated declarations shall be provided in a header file named `omp-tools.h`. In addition, the set of definitions may specify other implementation-specific values.

The `ompd_dll_locations` variable, all OMPD third-party tool interface functions, and all OMPD runtime entry points are external symbols with C linkage.

### 5.2 Activating a Third-Party Tool

The third-party tool and the OpenMP program exist as separate processes. Thus, coordination is required between the OpenMP runtime and the third-party tool for OMPD.

#### 5.2.1 Enabling Runtime Support for OMPD

In order to support third-party tools, the OpenMP runtime may need to collect and to store information that it may not otherwise maintain. The OpenMP runtime collects whatever information is necessary to support OMPD if the environment variable `OMP_DEBUG` is set to `enabled`.

**Cross References**

- Activating a first-party tool, see Section 4.2.
- `OMP_DEBUG` environment variable, see Section 6.21.

#### 5.2.2 ompd_dll_locations

**Summary**

The `ompd_dll_locations` global variable points to the locations of OMPD libraries that are compatible with the OpenMP implementation.

**Format**

```c
extern const char **ompd_dll_locations;
```
Description

An OpenMP runtime may have more than one OMPD library. The third-party tool must be able to locate the right library to use for the OpenMP program that it is examining. The OpenMP runtime system must provide a public variable `ompd_dll_locations`, which is an `argv`-style vector of filename string pointers that provides the names of any compatible OMPD libraries. This variable must have C linkage. The third-party tool uses the name of the variable verbatim and, in particular, does not apply any name mangling before performing the look up.

The architecture on which the third-party tool and, thus, the OMPD library execute does not have to match the architecture on which the OpenMP program that is being examined executes. The third-party tool must interpret the contents of `ompd_dll_locations` to find a suitable OMPD library that matches its own architectural characteristics. On platforms that support different architectures (for example, 32-bit vs 64-bit), OpenMP implementations are encouraged to provide an OMPD library for each supported architecture that can handle OpenMP programs that run on any supported architecture. Thus, for example, a 32-bit debugger that uses OMPD should be able to debug a 64-bit OpenMP program by loading a 32-bit OMPD implementation that can manage a 64-bit OpenMP runtime.

The `ompd_dll_locations` variable points to a NULL-terminated vector of zero or more NULL-terminated pathname strings that do not have any filename conventions. This vector must be fully initialized before `ompd_dll_locations` is set to a non-null value. Thus, if a third-party tool, such as a debugger, stops execution of the OpenMP program at any point at which `ompd_dll_locations` is non-null, the vector of strings to which it points shall be valid and complete.

Cross References

• `ompd_dll_locations_valid` global variable, see Section 5.2.3.

5.2.3 `ompd_dll_locations_valid`

Summary

The OpenMP runtime notifies third-party tools that `ompd_dll_locations` is valid by allowing execution to pass through a location that the symbol `ompd_dll_locations_valid` identifies.

Format

```c
void ompd_dll_locations_valid(void);
```
Since `ompd_dll_locations` may not be a static variable, it may require runtime initialization. The OpenMP runtime notifies third-party tools that `ompd_dll_locations` is valid by having execution pass through a location that the symbol `ompd_dll_locations_valid` identifies. If `ompd_dll_locations` is NULL, a third-party tool can place a breakpoint at `ompd_dll_locations_valid` to be notified that `ompd_dll_locations` is initialized. In practice, the symbol `ompd_dll_locations_valid` may not be a function; instead, it may be a labeled machine instruction through which execution passes once the vector is valid.

### 5.3 OMPD Data Types

This section defines OMPD data types.

#### 5.3.1 Size Type

**Summary**

The `ompd_size_t` type specifies the number of bytes in opaque data objects that are passed across the OMPD API.

**Format**

```c
typedef uint64_t ompd_size_t;
```

#### 5.3.2 Wait ID Type

**Summary**

A variable of `ompd_wait_id_t` type identifies the object on which a thread waits.

**Format**

```c
typedef uint64_t ompd_wait_id_t;
```

**Description**

The values and meaning of `ompd_wait_id_t` is the same as defined for the `ompt_wait_id_t` type.

**Cross References**

- `ompt_wait_id_t` type, see Section 4.4.4.30.
5.3.3 Basic Value Types

Summary
These definitions represent word, address, and segment value types.

Format

```
C / C++
typedef uint64_t ompd_addr_t;
typedef int64_t ompd_word_t;
typedef uint64_t ompd_seg_t;
```

Description
The `ompd_addr_t` type represents an address in an OpenMP process with an unsigned integer type. The `ompd_word_t` type represents a data word from the OpenMP runtime with a signed integer type. The `ompd_seg_t` type represents a segment value with an unsigned integer type.

5.3.4 Address Type

Summary
The `ompd_address_t` type is used to specify device addresses.

Format

```
C / C++
typedef struct ompd_address_t {
    ompd_seg_t segment;
    ompd_addr_t address;
} ompd_address_t;
```

Description
The `ompd_address_t` type is a structure that OMPD uses to specify device addresses, which may or may not be segmented. For non-segmented architectures, `ompd_segment_none` is used in the `segment` field of `ompd_address_t`; it is an instance of the `ompd_seg_t` type that has the value 0.
5.3.5 Frame Information Type

Summary
The ompd_frame_info_t type is used to specify frame information.

Format
<typename struct ompd_frame_info_t {
  ompd_address_t frame_address;
  ompd_word_t frame_flag;
} ompd_frame_info_t;

Description
The ompd_frame_info_t type is a structure that OMPD uses to specify frame information. The frame_address field of ompd_frame_info_t identifies a frame. The frame_flag field of ompd_frame_info_t indicates what type of information is provided in frame_address. The values and meaning is the same as defined for the ompt_frame_flag_t enumeration type.

Cross References
- ompt_frame_t type, see Section 4.4.4.28.

5.3.6 System Device Identifiers

Summary
The ompd_device_t type provides information about OpenMP devices.

Format
<typename struct uint64_t ompd_device_t;

Description
OpenMP runtimes may utilize different underlying devices, each represented by a device identifier. The device identifiers can vary in size and format and, thus, are not explicitly represented in the OMPD interface. Instead, a device identifier is passed across the interface via its ompd_device_t kind, its size in bytes and a pointer to where it is stored. The OMPD library and the third-party tool use the ompd_device_t kind to interpret the format of the device identifier that is referenced by the pointer argument. Each different device identifier kind is represented by a unique unsigned 64-bit integer value.

Recommended values of ompd_device_t kinds are defined in the ompd-types.h header file, which is available on http://www.openmp.org/.
5.3.7 Native Thread Identifiers

Summary

The `ompd_thread_id_t` type provides information about native threads.

Format

```c
typedef uint64_t ompd_thread_id_t;
```

Description

OpenMP runtimes may use different native thread implementations. Native thread identifiers for these implementations can vary in size and format and, thus, are not explicitly represented in the OMPD interface. Instead, a native thread identifier is passed across the interface via its `ompd_thread_id_t` kind, its size in bytes and a pointer to where it is stored. The OMPD library and the third-party tool use the `ompd_thread_id_t` kind to interpret the format of the native thread identifier that is referenced by the pointer argument. Each different native thread identifier kind is represented by a unique unsigned 64-bit integer value.

Recommended values of `ompd_thread_id_t` kinds, and formats for some corresponding native thread identifiers, are defined in the `ompd-types.h` header file, which is available on http://www.openmp.org/.

5.3.8 OMPD Handle Types

Summary

The OMPD library defines handles for referring to address spaces, threads, parallel regions and tasks. The internal structure of the handles are opaque to the third-party tool.

Format

```c
typedef struct _ompd_aspace_handle ompd_address_space_handle_t;
typedef struct _ompd_thread_handle ompd_thread_handle_t;
typedef struct _ompd_parallel_handle ompd_parallel_handle_t;
typedef struct _ompd_task_handle ompd_task_handle_t;
```

Description

OMPD uses handles for address spaces (`ompd_address_space_handle_t`), threads (`ompd_thread_handle_t`), parallel regions (`ompd_parallel_handle_t`), and tasks (`ompd_task_handle_t`). Each operation of the OMPD interface that applies to a particular address space, thread, parallel region or task must explicitly specify a corresponding handle. A handle for an entity is constant while the entity itself is alive. Handles are defined by the OMPD library and are opaque to the third-party tool.
Defining externally visible type names in this way introduces type safety to the interface, and helps
to catch instances where incorrect handles are passed by the third-party tool to the OMPD library.
The structures do not need to be defined; instead, the OMPD library must cast incoming (pointers
to) handles to the appropriate internal, private types.

### 5.3.9 OMPD Scope Types

#### Summary

The `ompd_scope_t` type identifies OMPD scopes.

#### Format

```c
typedef enum ompd_scope_t {
    ompd_scope_global = 1,
    ompd_scope_address_space = 2,
    ompd_scope_thread = 3,
    ompd_scope_parallel = 4,
    ompd_scope_implicit_task = 5,
    ompd_scope_task = 6
} ompd_scope_t;
```

#### Description

The `ompd_scope_t` type identifies OpenMP scopes, including those related to parallel regions
and tasks. When used in an OMPD interface function call, the scope type and the OMPD handle
must match according to Table 5.1.

<table>
<thead>
<tr>
<th>Scope types</th>
<th>Handles</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>ompd_scope_global</code></td>
<td>Address space handle for the host device</td>
</tr>
<tr>
<td><code>ompd_scope_address_space</code></td>
<td>Any address space handle</td>
</tr>
<tr>
<td><code>ompd_scope_thread</code></td>
<td>Any thread handle</td>
</tr>
<tr>
<td><code>ompd_scope_parallel</code></td>
<td>Any parallel region handle</td>
</tr>
<tr>
<td><code>ompd_scope_implicit_task</code></td>
<td>Task handle for an implicit task</td>
</tr>
<tr>
<td><code>ompd_scope_task</code></td>
<td>Any task handle</td>
</tr>
</tbody>
</table>
5.3.10 ICV ID Type

Summary
The `ompd_icv_id_t` type identifies an OpenMP implementation ICV.

Format

```c
typedef uint64_t ompd_icv_id_t;
```

The `ompd_icv_id_t` type identifies OpenMP implementation ICVs. `ompd_icv_undefined` is an instance of this type with the value 0.

5.3.11 Tool Context Types

Summary
A third-party tool defines contexts to identify abstractions uniquely. The internal structure of these contexts are opaque to the OMPD library.

Format

```c
typedef struct _ompd_aspace_cont ompd_address_space_context_t;
typedef struct _ompd_thread_cont ompd_thread_context_t;
```

Description
A third-party tool uniquely defines an address space context to identify the address space for the process that it is monitoring. Similarly, it uniquely defines a thread context to identify a native thread of the process that it is monitoring. These contexts are opaque to the OMPD library.

5.3.12 Return Code Types

Summary
The `ompd_rc_t` type is the return code type of an OMPD operation.

Format

```c
typedef enum ompd_rc_t {
    ompd_rc_ok = 0,
    ompd_rc_unavailable = 1,
    ompd_rc_stale_handle = 2,
    ompd_rc_bad_input = 3,
    ompd_rc_error = 4,
    ompd_rc_unsupported = 5,
    ompd_rc_needs_state_tracking = 6,
} C / C++
```
C / C++

Description
The ompd_rc_t type is used for the return codes of OMPD operations. The return code types and their semantics are defined as follows:

- `ompd_rc_ok` is returned when the operation is successful;
- `ompd_rc_unavailable` is returned when information is not available for the specified context;
- `ompd_rc_stale_handle` is returned when the specified handle is no longer valid;
- `ompd_rc_bad_input` is returned when the input parameters (other than handle) are invalid;
- `ompd_rc_error` is returned when a fatal error occurred;
- `ompd_rc_unsupported` is returned when the requested operation is not supported;
- `ompd_rc_needs_state_tracking` is returned when the state tracking operation failed because state tracking is not currently enabled;
- `ompd_rc_device_read_error` is returned when a read operation failed on the device;
- `ompd_rc_device_write_error` is returned when a write operation failed on the device;
- `ompd_rc_incompatible` is returned when this OMPD library is incompatible with the OpenMP program or is not capable of handling it;
- `ompd_rc_nomem` is returned when a memory allocation fails;
- `ompd_rc_incomplete` is returned when the information provided on return is incomplete, while the arguments are still set to valid values; and
- `ompd_rc_callback_error` is returned when the callback interface or any one of the required callback routines provided by the third-party tool is invalid.

5.3.13 Primitive Type Sizes
Summary
The ompd_device_type_sizes_t type provides the size of primitive types in the OpenMP architecture address space.
Format

```c
typedef struct ompd_device_type_sizes_t {
    uint8_t sizeof_char;
    uint8_t sizeof_short;
    uint8_t sizeof_int;
    uint8_t sizeof_long;
    uint8_t sizeof_long_long;
    uint8_t sizeof_pointer;
} ompd_device_type_sizes_t;
```

Description

The `ompd_device_type_sizes_t` type is used in operations through which the OMPD library can interrogate the third-party tool about the size of primitive types for the target architecture of the OpenMP runtime, as returned by the `sizeof` operator. The fields of `ompd_device_type_sizes_t` give the sizes of the eponymous basic types used by the OpenMP runtime. As the third-party tool and the OMPD library, by definition, execute on the same architecture, the size of the fields can be given as `uint8_t`.

Cross References

- `ompd_callback_sizeof_fn_t` type, see Section 5.4.2.2.

5.4 OMPD Third-Party Tool Callback Interface

For the OMPD library to provide information about the internal state of the OpenMP runtime system in an OpenMP process or core file, it must have a means to extract information from the OpenMP process that the third-party tool is examining. The OpenMP process on which the third-party tool is operating may be either a “live” process or a core file, and a thread may be either a “live” thread in an OpenMP process or a thread in a core file. To enable the OMPD library to extract state information from an OpenMP process or core file, the third-party tool must supply the OMPD library with callback functions to inquire about the size of primitive types in the device of the OpenMP process, to look up the addresses of symbols, and to read and to write memory in the device. The OMPD library uses these callbacks to implement its interface operations. The OMPD library only invokes the callback functions in direct response to calls made by the third-party tool to the OMPD library.

Description of Return Codes

All of the OMPD callback functions must return the following return codes or function-specific return codes:

- `ompd_rc_ok` on success; or
- `ompd_rc_stale_handle` if an invalid context argument is provided.
5.4.1 Memory Management of OMPD Library

The OMPD library must not access the heap manager directly. Instead, if it needs heap memory it must use the memory allocation and deallocation callback functions that are described in this section, `ompd_callback_memory_alloc_fn_t` (see Section 5.4.1.1) and `ompd_callback_memory_free_fn_t` (see Section 5.4.1.2), which are provided by the third-party tool to obtain and to release heap memory. This mechanism ensures that the library does not interfere with any custom memory management scheme that the third-party tool may use.

If the OMPD library is implemented in C++ then memory management operators, like `new` and `delete` and their variants, *must all* be overloaded and implemented in terms of the callbacks that the third-party tool provides. The OMPD library must be implemented in a manner such that any of its definitions of `new` or `delete` do not interfere with any that the third-party tool defines.

In some cases, the OMPD library must allocate memory to return results to the third-party tool. The third-party tool then owns this memory and has the responsibility to release it. Thus, the OMPD library and the third-party tool must use the same memory manager.

The OMPD library creates OMPD handles, which are opaque to the third-party tool and may have a complex internal structure. The third-party tool cannot determine if the handle pointers that the API returns correspond to discrete heap allocations. Thus, the third-party tool must not simply deallocate a handle by passing an address that it receives from the OMPD library to its own memory manager. Instead, the OMPD API includes functions that the third-party tool must use when it no longer needs a handle.

A third-party tool creates contexts and passes them to the OMPD library. The OMPD library does not release contexts; instead the third-party tool releases them after it releases any handles that may reference the contexts.

5.4.1.1 ompd_callback_memory_alloc_fn_t

Summary

The `ompd_callback_memory_alloc_fn_t` type is the type signature of the callback routine that the third-party tool provides to the OMPD library to allocate memory.

Format

```c
typedef ompd_rc_t (*ompd_callback_memory_alloc_fn_t)(
    ompd_size_t nbytes,
    void **ptr
);```

Description

The `ompd_callback_memory_alloc_fn_t` type is the type signature of the memory allocation callback routine that the third-party tool provides. The OMPD library may call the `ompd_callback_memory_alloc_fn_t` callback function to allocate memory.
Description of Arguments
The *nbytes* argument is the size in bytes of the block of memory to allocate.

The address of the newly allocated block of memory is returned in the location to which the *ptr* argument points. The newly allocated block is suitably aligned for any type of variable and is not guaranteed to be set to zero.

Description of Return Codes
Routines that use the `ompd_callback_memory_alloc_fn_t` type may return the general return codes listed at the beginning of Section 5.4.

Cross References
- `ompd_size_t` type, see Section 5.3.1.
- `ompd_rc_t` type, see Section 5.3.12.

5.4.1.2 `ompd_callback_memory_free_fn_t`
Summary
The `ompd_callback_memory_free_fn_t` type is the type signature of the callback routine that the third-party tool provides to the OMPD library to deallocate memory.

Format

```c
typedef ompd_rc_t (*ompd_callback_memory_free_fn_t) (void *ptr);
```

Description
The `ompd_callback_memory_free_fn_t` type is the type signature of the memory deallocation callback routine that the third-party tool provides. The OMPD library may call the `ompd_callback_memory_free_fn_t` callback function to deallocate memory that was obtained from a prior call to the `ompd_callback_memory_alloc_fn_t` callback function.

Description of Arguments
The *ptr* argument is the address of the block to be deallocated.

Description of Return Codes
Routines that use the `ompd_callback_memory_free_fn_t` type may return the general return codes listed at the beginning of Section 5.4.
Cross References
• `ompd_rc_t` type, see Section 5.3.12.
• `ompd_callback_memory_alloc_fn_t` type, see Section 5.4.1.1.
• `ompd_callbacks_t` type, see Section 5.4.6.

5.4.2 Context Management and Navigation

Summary
The third-party tool provides the OMPD library with callbacks to manage and to navigate context relationships.

5.4.2.1 `ompd_callback_get_thread_context_for_thread_id_fn_t`

Summary
The `ompd_callback_get_thread_context_for_thread_id_fn_t` is the type signature of the callback routine that the third-party tool provides to the OMPD library to map a native thread identifier to a third-party tool thread context.

Format
```c
typedef ompd_rc_t
(*ompd_callback_get_thread_context_for_thread_id_fn_t) (
  ompd_address_space_context_t *address_space_context,
  ompd_thread_id_t kind,
  ompd_size_t sizeof_thread_id,
  const void *thread_id,
  ompd_thread_context_t **thread_context
);```

Description
The `ompd_callback_get_thread_context_for_thread_id_fn_t` is the type signature of the context mapping callback routine that the third-party tool provides. This callback maps a native thread identifier to a third-party tool thread context. The native thread identifier is within the address space that `address_space_context` identifies. The OMPD library can use the thread context, for example, to access thread local storage.

Description of Arguments
The `address_space_context` argument is an opaque handle that the third-party tool provides to reference an address space. The `kind`, `sizeof_thread_id`, and `thread_id` arguments represent a native thread identifier. On return, the `thread_context` argument provides an opaque handle that maps a native thread identifier to a third-party tool thread context.
Description of Return Codes
In addition to the general return codes listed at the beginning of Section 5.4, routines that use the 
ompd_callback_get_thread_context_for_thread_id_fn_t type may also return 
the following return codes:

- **ompd_rc_bad_input** if a different value in sizeof_thread_id is expected for the native thread 
  identifier kind given by kind; or

- **ompd_rc_unsupported** if the native thread identifier kind is not supported.

Restrictions
Restrictions on routines that use 
ompd_callback_get_thread_context_for_thread_id_fn_t are as follows:

- The provided thread_context must be valid until the OMPD library returns from the OMPD 
  third-party tool interface routine.

Cross References
- ompd_size_t type, see Section 5.3.1.
- ompd_thread_id_t type, see Section 5.3.7.
- ompd_address_space_context_t type, see Section 5.3.11.
- ompd_thread_context_t type, see Section 5.3.11.
- ompd_rc_t type, see Section 5.3.12.

5.4.2.2 ompd_callback_sizeof_fn_t

Summary
The ompd_callback_sizeof_fn_t type is the type signature of the callback routine that the 
third-party tool provides to the OMPD library to determine the sizes of the primitive types in an 
address space.

Format

```c
typedef ompd_rc_t (*ompd_callback_sizeof_fn_t) ( 
    ompd_address_space_context_t *address_space_context, 
    ompd_device_type_sizes_t *sizes 
); 
```

Description
The ompd_callback_sizeof_fn_t is the type signature of the type-size query callback 
routine that the third-party tool provides. This callback provides the sizes of the basic primitive 
types for a given address space.
Description of Arguments
The callback returns the sizes of the basic primitive types used by the address space context that the `address_space_context` argument specifies in the location to which the `sizes` argument points.

Description of Return Codes
Routines that use the `ompd_callback_sizeof_fn_t` type may return the general return codes listed at the beginning of Section 5.4.

Cross References
- `ompd_address_space_context_t` type, see Section 5.3.11.
- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_device_type_sizes_t` type, see Section 5.3.13.
- `ompd_callbacks_t` type, see Section 5.4.6.

5.4.3 Accessing Memory in the OpenMP Program or Runtime
The OMPD library cannot directly read from or write to memory of the OpenMP program. Instead, the OMPD library must use callbacks that the third-party tool provides so that the third-party tool performs the operation.

5.4.3.1 `ompd_callback_symbol_addr_fn_t`

Summary
The `ompd_callback_symbol_addr_fn_t` type is the type signature of the callback that the third-party tool provides to look up the addresses of symbols in an OpenMP program.

Format
```c
typedef ompd_rc_t (*ompd_callback_symbol_addr_fn_t) (
    ompd_address_space_context_t *address_space_context,
    const_thread_context_t *thread_context,
    const char *symbol_name,
    ompd_address_t *symbol_addr,
    const char *file_name
);```

Description
The `ompd_callback_symbol_addr_fn_t` is the type signature of the symbol-address query callback routine that the third-party tool provides. This callback looks up addresses of symbols within a specified address space.
Description of Arguments

This callback looks up the symbol provided in the `symbol_name` argument.

The `address_space_context` argument is the third-party tool’s representation of the address space of the process, core file, or device.

The `thread_context` argument is NULL for global memory accesses. If `thread_context` is not NULL, `thread_context` gives the thread-specific context for the symbol lookup for the purpose of calculating thread local storage addresses. In this case, the thread to which `thread_context` refers must be associated with either the process or the device that corresponds to the `address_space_context` argument.

The third-party tool uses the `symbol_name` argument that the OMPD library supplies verbatim. In particular, no name mangling, demangling or other transformations are performed prior to the lookup. The `symbol_name` parameter must correspond to a statically allocated symbol within the specified address space. The symbol can correspond to any type of object, such as a variable, thread local storage variable, function, or untyped label. The symbol can have a local, global, or weak binding.

The `file_name` argument is an optional input parameter that indicates the name of the shared library in which the symbol is defined, and it is intended to help the third-party tool disambiguate symbols that are defined multiple times across the executable or shared library files. The shared library name may not be an exact match for the name seen by the third-party tool. If `file_name` is NULL then the third-party tool first tries to find the symbol in the executable file, and, if the symbol is not found, the third-party tool tries to find the symbol in the shared libraries in the order in which the shared libraries are loaded into the address space. If `file_name` is non-null then the third-party tool first tries to find the symbol in the libraries that match the name in the `file_name` argument, and, if the symbol is not found, the third-party tool then uses the same procedure as when `file_name` is NULL.

The callback does not support finding either symbols that are dynamically allocated on the call stack or statically allocated symbols that are defined within the scope of a function or subroutine.

The callback returns the address of the symbol in the location to which `symbol_addr` points.

Description of Return Codes

In addition to the general return codes listed at the beginning of Section 5.4, routines that use the `ompd_callback_symbol_addr_fn_t` type may also return the following return codes:

- `ompd_rc_error` if the requested symbol is not found; or
- `ompd_rc_bad_input` if no symbol name is provided.

Restrictions

Restrictions on routines that use the `ompd_callback_symbol_addr_fn_t` type are as follows:

- The `address_space_context` argument must be non-null.
- The symbol that the `symbol_name` argument specifies must be defined.
Cross References

- ompd_address_t type, see Section 5.3.4.
- ompd_address_space_context_t type, see Section 5.3.11.
- ompd_thread_context_t type, see Section 5.3.11.
- ompd_thread_context_t type, see Section 5.3.11.
- ompd_thread_context_t type, see Section 5.3.12.
- ompd_callbacks_t type, see Section 5.4.6.

5.4.3.2 ompd_callback_memory_read_fn_t

Summary

The ompd_callback_memory_read_fn_t type is the type signature of the callback that the third-party tool provides to read data (read_memory) or a string (read_string) from an OpenMP program.

Format

```c
typedef ompd_rc_t (*ompd_callback_memory_read_fn_t) (
    ompd_address_space_context_t *address_space_context,
    ompd_thread_context_t *thread_context,
    const ompd_address_t *addr,
    ompd_size_t nbytes,
    void *buffer
);
```

Description

The ompd_callback_memory_read_fn_t is the type signature of the read callback routines that the third-party tool provides.

The read_memory callback copies a block of data from addr within the address space given by address_space_context to the third-party tool buffer.

The read_string callback copies a string to which addr points, including the terminating null byte (’\0’), to the third-party tool buffer. At most nbytes bytes are copied. If a null byte is not among the first nbytes bytes, the string placed in buffer is not null-terminated.

Description of Arguments

The address from which the data are to be read in the OpenMP program that address_space_context specifies is given by addr. The nbytes argument is the number of bytes to be transferred. The thread_context argument is optional for global memory access, and in that case should be NULL. If it is non-null, thread_context identifies the thread-specific context for the memory access for the purpose of accessing thread local storage.
The data are returned through buffer, which is allocated and owned by the OMPD library. The contents of the buffer are unstructured, raw bytes. The OMPD library must arrange for any transformations such as byte-swapping that may be necessary (see Section 5.4.4) to interpret the data.

**Description of Return Codes**

In addition to the general return codes listed at the beginning of Section 5.4, routines that use the `ompd_callback_memory_read_fn_t` type may also return the following return codes:

- `ompd_rc_incomplete` if no terminating null byte is found while reading `nbytes` using the `read_string` callback; or
- `ompd_rc_error` if unallocated memory is reached while reading `nbytes` using either the `read_memory` or `read_string` callback.

**Cross References**

- `ompd_size_t` type, see Section 5.3.1.
- `ompd_address_t` type, see Section 5.3.4.
- `ompd_address_space_context_t` type, see Section 5.3.11.
- `ompd_thread_context_t` type, see Section 5.3.11.
- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_callback_device_host_fn_t` type, see Section 5.4.4.
- `ompd_callbacks_t` type, see Section 5.4.6.

### 5.4.3.3 ompd_callback_memory_write_fn_t

**Summary**

The `ompd_callback_memory_write_fn_t` type is the type signature of the callback that the third-party tool provides to write data to an OpenMP program.

**Format**

```c
typedef ompd_rc_t (*ompd_callback_memory_write_fn_t) (  
    ompd_address_space_context_t *address_space_context,  
    ompd_thread_context_t *thread_context,  
    const ompd_address_t *addr,  
    ompd_size_t nbytes,  
    const void *buffer  
);  
```
Description
The `ompd_callback_memory_write_fn_t` is the type signature of the write callback routine that the third-party tool provides. The OMPD library may call this callback to have the third-party tool write a block of data to a location within an address space from a provided buffer.

Description of Arguments
The address to which the data are to be written in the OpenMP program that `address_space_context` specifies is given by `addr`. The `nbytes` argument is the number of bytes to be transferred. The `thread_context` argument is optional for global memory access, and in that case should be NULL. If it is non-null then `thread_context` identifies the thread-specific context for the memory access for the purpose of accessing thread local storage.

The data to be written are passed through `buffer`, which is allocated and owned by the OMPD library. The contents of the buffer are unstructured, raw bytes. The OMPD library must arrange for any transformations such as byte-swapping that may be necessary (see Section 5.4.4) to render the data into a form that is compatible with the OpenMP runtime.

Description of Return Codes
Routines that use the `ompd_callback_memory_write_fn_t` type may return the general return codes listed at the beginning of Section 5.4.

Cross References
- `ompd_size_t` type, see Section 5.3.1.
- `ompd_address_t` type, see Section 5.3.4.
- `ompd_address_space_context_t` type, see Section 5.3.11.
- `ompd_thread_context_t` type, see Section 5.3.11.
- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_callback_device_host_fn_t` type, see Section 5.4.4.
- `ompd_callbacks_t` type, see Section 5.4.6.

5.4.4 Data Format Conversion:

`ompd_callback_device_host_fn_t`

Summary
The `ompd_callback_device_host_fn_t` type is the type signature of the callback that the third-party tool provides to convert data between the formats that the third-party tool and the OMPD library use and that the OpenMP program uses.
Format

```
typedef ompd_rc_t (*ompd_callback_device_host_fn_t) (
    ompd_address_space_context_t *address_space_context,
    const void *input,
    ompd_size_t unit_size,
    ompd_size_t count,
    void *output
);
```
5.4.5 ompd_callback_print_string_fn_t

Summary
The ompd_callback_print_string_fn_t type is the type signature of the callback that the third-party tool provides so that the OMPD library can emit output.

Format
```c
typedef ompd_rc_t (*ompd_callback_print_string_fn_t) (const char *string, int category);
```

Description
The OMPD library may call the ompd_callback_print_string_fn_t callback function to emit output, such as logging or debug information. The third-party tool may set the ompd_callback_print_string_fn_t callback function to NULL to prevent the OMPD library from emitting output. The OMPD library may not write to file descriptors that it did not open.

Description of Arguments
The string argument is the null-terminated string to be printed. No conversion or formatting is performed on the string.

The category argument is the implementation-defined category of the string to be printed.

Description of Return Codes
Routines that use the ompd_callback_print_string_fn_t type may return the general return codes listed at the beginning of Section 5.4.

Cross References
- ompd_rc_t type, see Section 5.3.12.
- ompd_callbacks_t type, see Section 5.4.6.

5.4.6 The Callback Interface

Summary
All OMPD library interactions with the OpenMP program must be through a set of callbacks that the third-party tool provides. These callbacks must also be used for allocating or releasing resources, such as memory, that the OMPD library needs.
typedef struct ompd_callbacks_t {
    ompd_callback_memory_alloc_fn_t alloc_memory;
    ompd_callback_memory_free_fn_t free_memory;
    ompd_callback_print_string_fn_t print_string;
    ompd_callback_sizeof_fn_t sizeof_type;
    ompd_callback_symbol_addr_fn_t symbol_addr_lookup;
    ompd_callback_memory_read_fn_t read_memory;
    ompd_callback_memory_write_fn_t write_memory;
    ompd_callback_memory_read_fn_t read_string;
    ompd_callback_device_host_fn_t device_to_host;
    ompd_callback_device_host_fn_t host_to_device;
    ompd_callback_get_thread_context_for_thread_id_fn_t
        get_thread_context_for_thread_id;
} ompd_callbacks_t;

Description
The set of callbacks that the OMPD library must use is collected in the *ompd_callbacks_t*
structure. An instance of this type is passed to the OMPD library as a parameter to
*ompd_initialize* (see Section 5.5.1.1). Each field points to a function that the OMPD library
must use either to interact with the OpenMP program or for memory operations.

The *alloc_memory* and *free_memory* fields are pointers to functions the OMPD library uses to
allocate and to release dynamic memory.

The *print_string* field points to a function that prints a string.

The architecture on which the OMPD library and third-party tool execute may be different from the
architecture on which the OpenMP program that is being examined executes. The *sizeof_type* field
points to a function that allows the OMPD library to determine the sizes of the basic integer and
pointer types that the OpenMP program uses. Because of the potential differences in the targeted
architectures, the conventions for representing data in the OMPD library and the OpenMP program
may be different. The *device_to_host* field points to a function that translates data from the
conventions that the OpenMP program uses to those that the third-party tool and OMPD library
use. The reverse operation is performed by the function to which the *host_to_device* field points.

The *symbol_addr_lookup* field points to a callback that the OMPD library can use to find the
address of a global or thread local storage symbol. The *read_memory*, *read_string* and
*write_memory* fields are pointers to functions for reading from and writing to global memory or
thread local storage in the OpenMP program.

The *get_thread_context_for_thread_id* field is a pointer to a function that the OMPD library can
use to obtain a thread context that corresponds to a native thread identifier.
Cross References

- `ompd_callback_memory_alloc_fn_t` type, see Section 5.4.1.1.
- `ompd_callback_memory_free_fn_t` type, see Section 5.4.1.2.
- `ompd_callback_get_thread_context_for_thread_id_fn_t` type, see Section 5.4.2.1.
- `ompd_callback_sizeof_fn_t` type, see Section 5.4.2.2.
- `ompd_callback_symbol_addr_fn_t` type, see Section 5.4.3.1.
- `ompd_callback_memory_read_fn_t` type, see Section 5.4.3.2.
- `ompd_callback_memory_write_fn_t` type, see Section 5.4.3.3.
- `ompd_callback_device_host_fn_t` type, see Section 5.4.4.
- `ompd_callback_print_string_fn_t` type, see Section 5.4.5

5.5 OMPD Tool Interface Routines

This section defines the interface provided by the OMPD library to be used by the third-party tool.

Description of Return Codes

All of the OMPD Tool Interface Routines must return function specific return codes or any of the following return codes:

- `ompd_rc_stale_handle` if a provided handle is stale;
- `ompd_rc_bad_input` if NULL is provided for any input argument unless otherwise specified;
- `ompd_rc_callback` if a callback returned an unexpected error, which leads to a failure of the query;
- `ompd_rc_needs_state_tracking` if the information cannot be provided while the `debug-var` is disabled;
- `ompd_rc_ok` on success; or
- `ompd_rc_error` for any other error.

5.5.1 Per OMPD Library Initialization and Finalization

The OMPD library must be initialized exactly once after it is loaded, and finalized exactly once before it is unloaded. Per OpenMP process or core file initialization and finalization are also required.
Once loaded, the tool can determine the version of the OMPD API that the library supports by calling `ompd_get_api_version` (see Section 5.5.1.2). If the tool supports the version that `ompd_get_api_version` returns, the tool starts the initialization by calling `ompd_initialize` (see Section 5.5.1.1) using the version of the OMPD API that the library supports. If the tool does not support the version that `ompd_get_api_version` returns, it may attempt to call `ompd_initialize` with a different version.

### 5.5.1.1 ompd_initialize

#### Summary

The `ompd_initialize` function initializes the OMPD library.

#### Format

```c
ompd_rc_t ompd_initialize(
    ompd_word_t api_version,
    const ompd_callbacks_t *callbacks
);
```

#### Description

A tool that uses OMPD calls `ompd_initialize` to initialize each OMPD library that it loads. More than one library may be present in a third-party tool, such as a debugger, because the tool may control multiple devices, which may use different runtime systems that require different OMPD libraries. This initialization must be performed exactly once before the tool can begin to operate on an OpenMP process or core file.

#### Description of Arguments

The `api_version` argument is the OMPD API version that the tool requests to use. The tool may call `ompd_get_api_version` to obtain the latest OMPD API version that the OMPD library supports.

The tool provides the OMPD library with a set of callback functions in the `callbacks` input argument which enables the OMPD library to allocate and to deallocate memory in the tool’s address space, to lookup the sizes of basic primitive types in the device, to lookup symbols in the device, and to read and to write memory in the device.

#### Description of Return Codes

This routine must return any of the general return codes listed at the beginning of Section 5.5 or any of the following return codes:

- `ompd_rc_bad_input` if invalid callbacks are provided; or
- `ompd_rc_unsupported` if the requested API version cannot be provided.
Cross References

- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_callbacks_t` type, see Section 5.4.6.
- `ompd_get_api_version` routine, see Section 5.5.1.2.

5.5.1.2 ompd_get_api_version

Summary

The `ompd_get_api_version` function returns the OMPD API version.

Format

```c
ompd_rc_t ompd_get_api_version(ompd_word_t *version);
```

Description

The tool may call the `ompd_get_api_version` function to obtain the latest OMPD API version number of the OMPD library. The OMPD API version number is equal to the value of the `_OPENMP` macro defined in the associated OpenMP implementation, if the C preprocessor is supported. If the associated OpenMP implementation compiles Fortran codes without the use of a C preprocessor, the OMPD API version number is equal to the value of the Fortran integer parameter `openmp_version`.

Description of Arguments

The latest version number is returned into the location to which the `version` argument points.

Description of Return Codes

This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References

- `ompd_rc_t` type, see Section 5.3.12.

5.5.1.3 ompd_get_version_string

Summary

The `ompd_get_version_string` function returns a descriptive string for the OMPD library version.

Format

```c
ompd_rc_t ompd_get_version_string(const char **string);
```
Description
The tool may call this function to obtain a pointer to a descriptive version string of the OMPD library vendor, implementation, internal version, date, or any other information that may be useful to a tool user or vendor. An implementation should provide a different string for every change to its source code or build that could be visible to the interface user.

Description of Arguments
A pointer to a descriptive version string is placed into the location to which the string output argument points. The OMPD library owns the string that the OMPD library returns; the tool must not modify or release this string. The string remains valid for as long as the library is loaded. The `ompd_get_version_string` function may be called before `ompd_initialize` (see Section 5.5.1.1). Accordingly, the OMPD library must not use heap or stack memory for the string.

The signatures of `ompd_get_api_version` (see Section 5.5.1.2) and `ompd_get_version_string` are guaranteed not to change in future versions of the API. In contrast, the type definitions and prototypes in the rest of the API do not carry the same guarantee. Therefore a tool that uses OMPD should check the version of the API of the loaded OMPD library before it calls any other function of the API.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
- `ompd_rc_t` type, see Section 5.3.12.

5.5.1.4 ompd_finalize

Summary
When the tool is finished with the OMPD library it should call `ompd_finalize` before it unloads the library.

Format

```
ompd_rc_t ompd_finalize(void);
```

Description
The call to `ompd_finalize` must be the last OMPD call that the tool makes before it unloads the library. This call allows the OMPD library to free any resources that it may be holding.

The OMPD library may implement a finalizer section, which executes as the library is unloaded and therefore after the call to `ompd_finalize`. During finalization, the OMPD library may use the callbacks that the tool provided earlier during the call to `ompd_initialize`. 
Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- ompd_rc_unsupported if the OMPD library is not initialized.

Cross References
- ompd_rc_t type, see Section 5.3.12.

5.5.2 Per OpenMP Process Initialization and Finalization

5.5.2.1 ompd_process_initialize

Summary
A tool calls ompd_process_initialize to obtain an address space handle when it initializes a session on a live process or core file.

Format

```c
ompd_rc_t ompd_process_initialize(
        ompd_address_space_context_t *context,
        ompd_address_space_handle_t **handle
    );
```

Description
A tool calls ompd_process_initialize to obtain an address space handle when it initializes a session on a live process or core file. On return from ompd_process_initialize, the tool owns the address space handle, which it must release with ompd_rel_address_space_handle. The initialization function must be called before any OMPD operations are performed on the OpenMP process or core file. This call allows the OMPD library to confirm that it can handle the OpenMP process or core file that context identifies.

Description of Arguments
The context argument is an opaque handle that the tool provides to address an address space. On return, the handle argument provides an opaque handle to the tool for this address space, which the tool must release when it is no longer needed.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- ompd_rc_incompatible if the OMPD library is incompatible with the runtime library loaded in the process.
Cross References
  • ompd_address_space_handle_t type, see Section 5.3.8.
  • ompd_address_space_context_t type, see Section 5.3.11.
  • ompd_rc_t type, see Section 5.3.12.
  • ompd_rel_address_space_handle routine, see Section 5.5.2.3.

5.5.2.2 ompd_device_initialize

Summary
A tool calls `ompd_device_initialize` to obtain an address space handle for a device that has at least one active target region.

Format

```c
ompd_rc_t ompd_device_initialize(
    ompd_address_space_handle_t *process_handle,
    ompd_address_space_context_t *device_context,
    ompd_device_t kind,
    ompd_size_t sizeof_id,
    void *id,
    ompd_address_space_handle_t **device_handle
);
```

Description
A tool calls `ompd_device_initialize` to obtain an address space handle for a device that has at least one active target region. On return from `ompd_device_initialize`, the tool owns the address space handle.

Description of Arguments
The `process_handle` argument is an opaque handle that the tool provides to reference the address space of the OpenMP process or core file. The `device_context` argument is an opaque handle that the tool provides to reference a device address space. The `kind`, `sizeof_id`, and `id` arguments represent a device identifier. On return the `device_handle` argument provides an opaque handle to the tool for this address space.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

• `ompd_rc_unsupported` if the OMPD library has no support for the specific device.
Cross References
- oompd_size_t type, see Section 5.3.1.
- oompd_device_t type, see Section 5.3.6.
- oompd_address_space_handle_t type, see Section 5.3.8.
- oompd_address_space_context_t type, see Section 5.3.11.
- oompd_rc_t type, see Section 5.3.12.

5.5.2.3 oompd_rel_address_space_handle

Summary
A tool calls oompd_rel_address_space_handle to release an address space handle.

Format

```c
ompd_rc_t oompd_rel_address_space_handle(
    oompd_address_space_handle_t *handle
);```

Description
When the tool is finished with the OpenMP process address space handle it should call oompd_rel_address_space_handle to release the handle, which allows the OMPD library to release any resources that it has related to the address space.

Description of Arguments
The handle argument is an opaque handle for the address space to be released.

Restrictions
Restrictions to the oompd_rel_address_space_handle routine are as follows:
- An address space context must not be used after the corresponding address space handle is released.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
- oompd_address_space_handle_t type, see Section 5.3.8.
- oompd_rc_t type, see Section 5.3.12.
5.5.3 Thread and Signal Safety

The OMPD library does not need to be reentrant. The tool must ensure that only one thread enters the OMPD library at a time. The OMPD library must not install signal handlers or otherwise interfere with the tool’s signal configuration.

5.5.4 Address Space Information

5.5.4.1 ompd_get_omp_version

Summary

The tool may call the `ompd_get_omp_version` function to obtain the version of the OpenMP API that is associated with an address space.

Format

```
ompd_rc_t ompd_get_omp_version(
    ompd_address_space_handle_t *address_space,
    ompd_word_t *omp_version
);  
```

Description

The tool may call the `ompd_get_omp_version` function to obtain the version of the OpenMP API that is associated with the address space.

Description of Arguments

The `address_space` argument is an opaque handle that the tool provides to reference the address space of the OpenMP process or device.

Upon return, the `omp_version` argument contains the version of the OpenMP runtime in the `_OPENMP` version macro format.

Description of Return Codes

This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References

- `ompd_address_space_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.
### 5.5.4.2 ompd_get_omp_version_string

#### Summary
The `ompd_get_omp_version_string` function returns a descriptive string for the OpenMP API version that is associated with an address space.

#### Format

```c
ompd_rc_t ompd_get_omp_version_string(
    ompd_address_space_handle_t *address_space,
    const char **string
);```

#### Description
After initialization, the tool may call the `ompd_get_omp_version_string` function to obtain the version of the OpenMP API that is associated with an address space.

#### Description of Arguments
The `address_space` argument is an opaque handle that the tool provides to reference the address space of the OpenMP process or device. A pointer to a descriptive version string is placed into the location to which the `string` output argument points. After returning from the call, the tool owns the string. The OMPD library must use the memory allocation callback that the tool provides to allocate the string storage. The tool is responsible for releasing the memory.

#### Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

#### Cross References
- `ompd_address_space_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.
5.5.5 Thread Handles

5.5.5.1 ompd_get_thread_in_parallel

Summary
The `ompd_get_thread_in_parallel` function enables a tool to obtain handles for OpenMP threads that are associated with a parallel region.

Format
```c
ompd_rc_t ompd_get_thread_in_parallel(
    ompd_parallel_handle_t *parallel_handle,
    int thread_num,
    ompd_thread_handle_t **thread_handle
);
```

Description
A successful invocation of `ompd_get_thread_in_parallel` returns a pointer to a thread handle in the location to which `thread_handle` points. This call yields meaningful results only if all OpenMP threads in the parallel region are stopped.

Description of Arguments
The `parallel_handle` argument is an opaque handle for a parallel region and selects the parallel region on which to operate. The `thread_num` argument selects the thread, the handle of which is to be returned. On return, the `thread_handle` argument is an opaque handle for the selected thread.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- `ompd_rc_bad_input` if the `thread_num` argument is greater than or equal to the `team-size-var` ICV or negative.

Restrictions
Restrictions on the `ompd_get_thread_in_parallel` function are as follows:

- The value of `thread_num` must be a non-negative integer smaller than the team size that was provided as the `team-size-var` ICV from `ompd_get_icv_from_scope`.

Cross References
- `ompd_parallel_handle_t` type, see Section 5.3.8.
- `ompd_thread_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_get_icv_from_scope` routine, see Section 5.5.9.2.
5.5.5.2 ompd_get_thread_handle

Summary
The **ompd_get_thread_handle** function maps a native thread to an OMPD thread handle.

Format

```c
ompd_rc_t ompd_get_thread_handle(
    ompd_address_space_handle_t *handle,
    ompd_thread_id_t kind,
    ompd_size_t sizeof_thread_id,
    const void *thread_id,
    ompd_thread_handle_t **thread_handle
);
```

Description
The **ompd_get_thread_handle** function determines if the native thread identifier to which `thread_id` points represents an OpenMP thread. If so, the function returns **ompd_rc_ok** and the location to which `thread_handle` points is set to the thread handle for the OpenMP thread.

Description of Arguments
The `handle` argument is an opaque handle that the tool provides to reference an address space. The `kind`, `sizeof_thread_id`, and `thread_id` arguments represent a native thread identifier. On return, the `thread_handle` argument provides an opaque handle to the thread within the provided address space.

The native thread identifier to which `thread_id` points is guaranteed to be valid for the duration of the call. If the OMPD library must retain the native thread identifier, it must copy it.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or any of the following return codes:

- **ompd_rc_bad_input** if a different value in `sizeof_thread_id` is expected for a thread kind of `kind`.
- **ompd_rc_unsupported** if the `kind` of thread is not supported.
- **ompd_rc_unavailable** if the thread is not an OpenMP thread.

Cross References
- **ompd_size_t** type, see Section 5.3.1.
- **ompd_thread_id_t** type, see Section 5.3.7.
- **ompd_address_space_handle_t** type, see Section 5.3.8.
- **ompd_thread_handle_t** type, see Section 5.3.8.
- **ompd_rc_t** type, see Section 5.3.12.
5.5.3 ompd_rel_thread_handle

Summary
The *ompd_rel_thread_handle* function releases a thread handle.

Format

```c
ompd_rc_t ompd_rel_thread_handle(
    ompd_thread_handle_t *thread_handle
);
```

Description
Thread handles are opaque to tools, which therefore cannot release them directly. Instead, when the tool is finished with a thread handle it must pass it to *ompd_rel_thread_handle* for disposal.

Description of Arguments
The *thread_handle* argument is an opaque handle for a thread to be released.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
- *ompd_thread_handle_t* type, see Section 5.3.8.
- *ompd_rc_t* type, see Section 5.3.12.

5.5.4 ompd_thread_handle_compare

Summary
The *ompd_thread_handle_compare* function allows tools to compare two thread handles.

Format

```c
ompd_rc_t ompd_thread_handle_compare(
    ompd_thread_handle_t *thread_handle_1,
    ompd_thread_handle_t *thread_handle_2,
    int *cmp_value
);
```

Description
The internal structure of thread handles is opaque to a tool. While the tool can easily compare pointers to thread handles, it cannot determine whether handles of two different addresses refer to the same underlying thread. The *ompd_thread_handle_compare* function compares thread handles.
On success, `ompd_thread_handle_compare` returns in the location to which `cmp_value` points a signed integer value that indicates how the underlying threads compare: a value less than, equal to, or greater than 0 indicates that the thread corresponding to `thread_handle_1` is, respectively, less than, equal to, or greater than that corresponding to `thread_handle_2`.

**Description of Arguments**
The `thread_handle_1` and `thread_handle_2` arguments are opaque handles for threads. On return the `cmp_value` argument is set to a signed integer value.

**Description of Return Codes**
This routine must return any of the general return codes listed at the beginning of Section 5.5.

**Cross References**
- `ompd_thread_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.

### 5.5.5.5 ompd_get_thread_id

**Summary**
The `ompd_get_thread_id` maps an OMPD thread handle to a native thread.

**Format**
```c
ompd_rc_t ompd_get_thread_id(
    ompd_thread_handle_t *thread_handle,
    ompd_thread_id_t kind,
    ompd_size_t sizeof_thread_id,
    void *thread_id
);
```

**Description**
The `ompd_get_thread_id` function maps an OMPD thread handle to a native thread identifier.

**Description of Arguments**
The `thread_handle` argument is an opaque thread handle. The `kind` argument represents the native thread identifier. The `sizeof_thread_id` argument represents the size of the native thread identifier. On return, the `thread_id` argument is a buffer that represents a native thread identifier.
Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or any
of the following return codes:

- `ompd_rc_bad_input` if a different value in `sizeof_thread_id` is expected for a thread kind of
  `kind`; or
- `ompd_rc_unsupported` if the `kind` of thread is not supported.

Cross References
- `ompd_size_t` type, see Section 5.3.1.
- `ompd_thread_id_t` type, see Section 5.3.7.
- `ompd_thread_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.

5.5.6 Parallel Region Handles
5.5.6.1 ompd_get_curr_parallel_handle

Summary
The `ompd_get_curr_parallel_handle` function obtains a pointer to the parallel handle for
an OpenMP thread’s current parallel region.

Format
```c
ompd_rc_t ompd_get_curr_parallel_handle(
    ompd_thread_handle_t *thread_handle,
    ompd_parallel_handle_t **parallel_handle
);
```

Description
The `ompd_get_curr_parallel_handle` function enables the tool to obtain a pointer to the
parallel handle for the current parallel region that is associated with an OpenMP thread. This call is
meaningful only if the associated thread is stopped. The parallel handle is owned by the tool and it
must be released by calling `ompd_rel_parallel_handle`.

Description of Arguments
The `thread_handle` argument is an opaque handle for a thread and selects the thread on which to
operate. On return, the `parallel_handle` argument is set to a handle for the parallel region that the
associated thread is currently executing, if any.
Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- `ompd_rc_unavailable` if the thread is not currently part of a team.

Cross References
- `ompd_thread_handle_t` type, see Section 5.3.8.
- `ompd_parallel_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_rel_parallel_handle` routine, see Section 5.5.6.4.

5.5.6.2 `ompd_get_enclosing_parallel_handle`

Summary
The `ompd_get_enclosing_parallel_handle` function obtains a pointer to the parallel handle for an enclosing parallel region.

Format
```
C
ompd_rc_t ompd_get_enclosing_parallel_handle(
    ompd_parallel_handle_t *parallel_handle,
    ompd_parallel_handle_t **enclosing_parallel_handle
);
```

Description
The `ompd_get_enclosing_parallel_handle` function enables a tool to obtain a pointer to the parallel handle for the parallel region that encloses the parallel region that `parallel_handle` specifies. This call is meaningful only if at least one thread in the parallel region is stopped. A pointer to the parallel handle for the enclosing region is returned in the location to which `enclosing_parallel_handle` points. After the call, the tool owns the handle; the tool must release the handle with `ompd_rel_parallel_handle` when it is no longer required.

Description of Arguments
The `parallel_handle` argument is an opaque handle for a parallel region that selects the parallel region on which to operate. On return, the `enclosing_parallel_handle` argument is set to a handle for the parallel region that encloses the selected parallel region.
Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- \texttt{ompd\_rc\_unavailable} if no enclosing parallel region exists.

Cross References
- \texttt{ompd\_parallel\_handle\_t} type, see Section 5.3.8.
- \texttt{ompd\_rc\_t} type, see Section 5.3.12.
- \texttt{ompd\_rel\_parallel\_handle} routine, see Section 5.5.6.4.

5.5.6.3 \texttt{ompd\_get\_task\_parallel\_handle}

Summary
The \texttt{ompd\_get\_task\_parallel\_handle} function obtains a pointer to the parallel handle for the parallel region that encloses a task region.

Format

```c
ompd\_rc\_t \texttt{ompd\_get\_task\_parallel\_handle}(
    \texttt{ompd\_task\_handle\_t} *\texttt{task\_handle},
    \texttt{ompd\_parallel\_handle\_t} **\texttt{task\_parallel\_handle}
);
```

Description
The \texttt{ompd\_get\_task\_parallel\_handle} function enables a tool to obtain a pointer to the parallel handle for the parallel region that encloses the task region that \texttt{task\_handle} specifies. This call is meaningful only if at least one thread in the parallel region is stopped. A pointer to the parallel regions handle is returned in the location to which \texttt{task\_parallel\_handle} points. The tool owns that parallel handle, which it must release with \texttt{ompd\_rel\_parallel\_handle}.

Description of Arguments
The \texttt{task\_handle} argument is an opaque handle that selects the task on which to operate. On return, the \texttt{parallel\_handle} argument is set to a handle for the parallel region that encloses the selected task.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
- \texttt{ompd\_task\_handle\_t} type, see Section 5.3.8.
- \texttt{ompd\_parallel\_handle\_t} type, see Section 5.3.8.
- \texttt{ompd\_rc\_t} type, see Section 5.3.12.
- \texttt{ompd\_rel\_parallel\_handle} routine, see Section 5.5.6.4.
5.5.6.4 ompd_rel_parallel_handle

Summary
The `ompd_rel_parallel_handle` function releases a parallel region handle.

Format
```
C
ompd_rc_t ompd_rel_parallel_handle(
    ompd_parallel_handle_t *parallel_handle
);
```

Description
Parallel region handles are opaque so tools cannot release them directly. Instead, a tool must pass a parallel region handle to the `ompd_rel_parallel_handle` function for disposal when finished with it.

Description of Arguments
The `parallel_handle` argument is an opaque handle to be released.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
- `ompd_parallel_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.

5.5.6.5 ompd_parallel_handle_compare

Summary
The `ompd_parallel_handle_compare` function compares two parallel region handles.

Format
```
C
ompd_rc_t ompd_parallel_handle_compare(
    ompd_parallel_handle_t *parallel_handle_1,
    ompd_parallel_handle_t *parallel_handle_2,
    int *cmp_value
);
```
Description
The internal structure of parallel region handles is opaque to tools. While tools can easily compare
pointers to parallel region handles, they cannot determine whether handles at two different
addresses refer to the same underlying parallel region and, instead must use the
ompd_parallel_handle_compare function.

On success, ompd_parallel_handle_compare returns a signed integer value in the location
to which cmp_value points that indicates how the underlying parallel regions compare. A value less
than, equal to, or greater than 0 indicates that the region corresponding to parallel_handle_1 is,
respectively, less than, equal to, or greater than that corresponding to parallel_handle_2. This
function is provided since the means by which parallel region handles are ordered is
implementation defined.

Description of Arguments
The parallel_handle_1 and parallel_handle_2 arguments are opaque handles that correspond to
parallel regions. On return the cmp_value argument points to a signed integer value that indicates
how the underlying parallel regions compare.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
• ompd_parallel_handle_t type, see Section 5.3.8.
• ompd_rc_t type, see Section 5.3.12.

5.5.7 Task Handles
5.5.7.1 ompd_get_curr_task_handle

Summary
The ompd_get_curr_task_handle function obtains a pointer to the task handle for the
current task region that is associated with an OpenMP thread.

Format

C

ompd_rc_t ompd_get_curr_task_handle(
    ompd_thread_handle_t *thread_handle,
    ompd_task_handle_t **task_handle
);

C

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Description
The `ompd_get_curr_task_handle` function obtains a pointer to the task handle for the current task region that is associated with an OpenMP thread. This call is meaningful only if the thread for which the handle is provided is stopped. The task handle must be released with `ompd_rel_task_handle`.

Description of Arguments
The `thread_handle` argument is an opaque handle that selects the thread on which to operate. On return, the `task_handle` argument points to a location that points to a handle for the task that the thread is currently executing.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- `ompd_rc_unavailable` if the thread is currently not executing a task.

Cross References
- `ompd_thread_handle_t` type, see Section 5.3.8.
- `ompd_task_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_rel_task_handle` routine, see Section 5.5.7.5.

5.5.7.2 ompd_get_generating_task_handle

Summary
The `ompd_get_generating_task_handle` function obtains a pointer to the task handle of the generating task region.

Format
```c
ompd_rc_t ompd_get_generating_task_handle(
    ompd_task_handle_t *task_handle,
    ompd_task_handle_t **generating_task_handle
);
```

Description
The `ompd_get_generating_task_handle` function obtains a pointer to the task handle for the task that encountered the OpenMP task construct that generated the task represented by `task_handle`. The generating task is the OpenMP task that was active when the task specified by `task_handle` was created. This call is meaningful only if the thread that is executing the task that `task_handle` specifies is stopped. The generating task handle must be released with `ompd_rel_task_handle`.
Description of Arguments
The task_handle argument is an opaque handle that selects the task on which to operate. On return, the generating_task_handle argument points to a location that points to a handle for the generating task.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

• ompd_rc_unavailable if no generating task region exists.

Cross References
• ompd_task_handle_t type, see Section 5.3.8.
• ompd_rc_t type, see Section 5.3.12.
• ompd_rel_task_handle routine, see Section 5.5.7.5.

5.5.7.3 ompd_get_scheduling_task_handle

Summary
The ompd_get_scheduling_task_handle function obtains a task handle for the task that was active at a task scheduling point.

Format

```c
ompd_rc_t ompd_get_scheduling_task_handle(
    ompd_task_handle_t *task_handle,
    ompd_task_handle_t **scheduling_task_handle
);
```

Description
The ompd_get_scheduling_task_handle function obtains a task handle for the task that was active when the task that task_handle represents was scheduled. This call is meaningful only if the thread that is executing the task that task_handle specifies is stopped. The scheduling task handle must be released with ompd_rel_task_handle.

Description of Arguments
The task_handle argument is an opaque handle for a task and selects the task on which to operate. On return, the scheduling_task_handle argument points to a location that points to a handle for the task that is still on the stack of execution on the same thread and was deferred in favor of executing the selected task.
**Description of Return Codes**

This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- `ompd_rc_unavailable` if no scheduling task region exists.

**Cross References**

- `ompd_task_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_rel_task_handle` routine, see Section 5.5.7.5.

### 5.5.7.4 `ompd_get_task_in_parallel`

**Summary**

The `ompd_get_task_in_parallel` function obtains handles for the implicit tasks that are associated with a parallel region.

**Format**

```c
ompd_rc_t ompd_get_task_in_parallel(
    ompd_parallel_handle_t *parallel_handle,
    int thread_num,
    ompd_task_handle_t **task_handle
);
```

**Description**

The `ompd_get_task_in_parallel` function obtains handles for the implicit tasks that are associated with a parallel region. A successful invocation of `ompd_get_task_in_parallel` returns a pointer to a task handle in the location to which `task_handle` points. This call yields meaningful results only if all OpenMP threads in the parallel region are stopped.

**Description of Arguments**

The `parallel_handle` argument is an opaque handle that selects the parallel region on which to operate. The `thread_num` argument selects the implicit task of the team to be returned. The `thread_num` argument is equal to the `thread-num-var` ICV value of the selected implicit task. On return, the `task_handle` argument points to a location that points to an opaque handle for the selected implicit task.
Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- `ompd_rc_bad_input` if the `thread_num` argument is greater than or equal to the `team-size-var` ICV or negative.

Restrictions
Restrictions on the `ompd_get_task_in_parallel` function are as follows:

- The value of `thread_num` must be a non-negative integer that is smaller than the size of the team size that is the value of the `team-size-var` ICV that `ompd_get_icv_from_scope` returns.

Cross References
- `ompd_parallel_handle_t` type, see Section 5.3.8.
- `ompd_task_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_get_icv_from_scope` routine, see Section 5.5.9.2.

5.5.7.5 `ompd_rel_task_handle`
Summary
This `ompd_rel_task_handle` function releases a task handle.

Format
```
C
ompd_rc_t ompd_rel_task_handle(
    ompd_task_handle_t *task_handle
);
```

Description
Task handles are opaque to tools; thus tools cannot release them directly. Instead, when a tool is finished with a task handle it must use the `ompd_rel_task_handle` function to release it.

Description of Arguments
The `task_handle` argument is an opaque task handle to be released.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
- `ompd_task_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.
5.5.7.6 ompd_task_handle_compare

Summary
The ompd_task_handle_compare function compares task handles.

Format
```c
ompd_rc_t ompd_task_handle_compare(
    ompd_task_handle_t *task_handle_1,
    ompd_task_handle_t *task_handle_2,
    int *cmp_value
);
```

Description
The internal structure of task handles is opaque; so tools cannot directly determine if handles at two
different addresses refer to the same underlying task. The ompd_task_handle_compare
function compares task handles. After a successful call to ompd_task_handle_compare, the
value of the location to which cmp_value points is a signed integer that indicates how the underlying
tasks compare: a value less than, equal to, or greater than 0 indicates that the task that corresponds
to task_handle_1 is, respectively, less than, equal to, or greater than the task that corresponds to
task_handle_2. The means by which task handles are ordered is implementation defined.

Description of Arguments
The task_handle_1 and task_handle_2 arguments are opaque handles that correspond to tasks. On
return, the cmp_value argument points to a location in which a signed integer value indicates how
the underlying tasks compare.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
• ompd_task_handle_t type, see Section 5.3.8.
• ompd_rc_t type, see Section 5.3.12.

5.5.7.7 ompd_get_task_function

Summary
This ompd_get_task_function function returns the entry point of the code that corresponds
to the body of a task.

Format
```c
ompd_rc_t ompd_get_task_function (    ompd_task_handle_t *task_handle,
ompd_address_t *entry_point
);
```
Description
The `ompd_get_task_function` function returns the entry point of the code that corresponds to the body of code that the task executes.

Description of Arguments
The `task_handle` argument is an opaque handle that selects the task on which to operate. On return, the `entry_point` argument is set to an address that describes the beginning of application code that executes the task region.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
- `ompd_address_t` type, see Section 5.3.4.
- `ompd_task_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.

5.5.7.8 `ompd_get_task_frame`

Summary
The `ompd_get_task_frame` function extracts the frame pointers of a task.

Format
```c
ompd_rc_t ompd_get_task_frame (  
    ompd_task_handle_t *task_handle,  
    ompd_frame_info_t *exit_frame,  
    ompd_frame_info_t *enter_frame  
);  
```

Description
An OpenMP implementation maintains an `ompt_frame_t` object for every implicit or explicit task. The `ompd_get_task_frame` function extracts the `enter_frame` and `exit_frame` fields of the `ompt_frame_t` object of the task that `task_handle` identifies.

Description of Arguments
The `task_handle` argument specifies an OpenMP task. On return, the `exit_frame` argument points to an `ompd_frame_info_t` object that has the frame information with the same semantics as the `exit_frame` field in the `ompt_frame_t` object that is associated with the specified task. On return, the `enter_frame` argument points to an `ompd_frame_info_t` object that has the frame information with the same semantics as the `enter_frame` field in the `ompt_frame_t` object that is associated with the specified task.
Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
• ompt_frame_t type, see Section 4.4.28.
• ompd_address_t type, see Section 5.3.4.
• ompd_frame_info_t type, see Section 5.3.5.
• ompd_task_handle_t type, see Section 5.3.8.
• ompd_rc_t type, see Section 5.3.12.

5.5.7.9 ompdEnumerateStates

Summary
The ompdEnumerateStates function enumerates thread states that an OpenMP implementation supports.

Format

```c
ompd_rc_t ompdEnumerateStates (  
    ompd_address_space_handle_t *address_space_handle,  
    ompd_word_t current_state,  
    ompd_word_t *next_state,  
    const char **next_state_name,  
    ompd_word_t *more Enums

);```

Description
An OpenMP implementation may support only a subset of the states that the ompt_state_t enumeration type defines. In addition, an OpenMP implementation may support implementation-specific states. The ompdEnumerateStates call enables a tool to enumerate the thread states that an OpenMP implementation supports.

When the current_state argument is a thread state that an OpenMP implementation supports, the call assigns the value and string name of the next thread state in the enumeration to the locations to which the next_state and next_state_name arguments point.

On return, the third-party tool owns the next_state_name string. The OMPD library allocates storage for the string with the memory allocation callback that the tool provides. The tool is responsible for releasing the memory.

On return, the location to which the more Enums argument points has the value 1 whenever one or more states are left in the enumeration. On return, the location to which the more Enums argument points has the value 0 when current_state is the last state in the enumeration.
Description of Arguments
The `address_space_handle` argument identifies the address space. The `current_state` argument must be a thread state that the OpenMP implementation supports. To begin enumerating the supported states, a tool should pass `ompt_state_undefined` as the value of `current_state`. Subsequent calls to `ompd_enumerate_states` by the tool should pass the value that the call returned in the `next_state` argument. On return, the `next_state` argument points to an integer with the value of the next state in the enumeration. On return, the `next_state_name` argument points to a character string that describes the next state. On return, the `more Enums` argument points to an integer with a value of 1 when more states are left to enumerate and a value of 0 when no more states are left.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- `ompd_rc_bad_input` if an unknown value is provided in `current_state`.

Cross References
- `ompt_state_t` type, see Section 4.4.27.
- `ompd_address_space_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.

5.5.7.10 ompd_get_state
Summary
The `ompd_get_state` function obtains the state of a thread.

Format

```c
ompd_rc_t ompd_get_state (    ompd_thread_handle_t *thread_handle,    ompd_word_t *state,    ompd_wait_id_t *wait_id    );
```

Description
The `ompd_get_state` function returns the state of an OpenMP thread.

Description of Arguments
The `thread_handle` argument identifies the thread. The `state` argument represents the state of that thread as represented by a value that `ompd_enumerate_states` returns. On return, if the `wait_id` argument is non-null then it points to a handle that corresponds to the `wait_id` wait identifier of the thread. If the thread state is not one of the specified wait states, the value to which `wait_id` points is undefined.
Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
• ompd_wait_id_t type, see Section 5.3.2.
• ompd_thread_handle_t type, see Section 5.3.8.
• ompd_rc_t type, see Section 5.3.12.
• ompdEnumerateStates routine, see Section 5.5.7.9.

5.5.8 Display Control Variables
5.5.8.1 ompd_get_display_control_vars

Summary
The ompd_get_display_control_vars function returns a list of name/value pairs for OpenMP control variables.

Format

```c
ompd_rc_t ompd_get_display_control_vars (  
    ompd_address_space_handle_t *address_space_handle,  
    const char * const **control_vars  
);  
```

Description
The ompd_get_display_control_vars function returns a NULL-terminated vector of NULL-terminated strings of name/value pairs of control variables that have user controllable settings and are important to the operation or performance of an OpenMP runtime system. The control variables that this interface exposes include all OpenMP environment variables, settings that may come from vendor or platform-specific environment variables, and other settings that affect the operation or functioning of an OpenMP runtime.

The format of the strings is "icv-name=icv-value".

On return, the third-party tool owns the vector and the strings. The OMPD library must satisfy the termination constraints; it may use static or dynamic memory for the vector and/or the strings and is unconstrained in how it arranges them in memory. If it uses dynamic memory then the OMPD library must use the allocate callback that the tool provides to ompd_initialize. The tool must use the ompdRelDisplayControlVars function to release the vector and the strings.

Description of Arguments
The address_space_handle argument identifies the address space. On return, the control_vars argument points to the vector of display control variables.
Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
- `ompd_address_space_handle_t` type, see Section 5.3.8.
- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_initialize` routine, see Section 5.5.1.1.
- `ompd_rel_display_control_vars` routine, see Section 5.5.8.2.

5.5.8.2 ompd_rel_display_control_vars

Summary
The `ompd_rel_display_control_vars` releases a list of name/value pairs of OpenMP control variables previously acquired with `ompd_get_display_control_vars`.

Format
```
ompd_rc_t ompd_rel_display_control_vars (const char * const **control_vars);
```

Description
The third-party tool owns the vector and strings that `ompd_get_display_control_vars` returns. The tool must call `ompd_rel_display_control_vars` to release the vector and the strings.

Description of Arguments
The `control_vars` argument is the vector of display control variables to be released.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5.

Cross References
- `ompd_rc_t` type, see Section 5.3.12.
- `ompd_get_display_control_vars` routine, see Section 5.5.8.1.
5.5.9 Accessing Scope-Specific Information

5.5.9.1 ompd_enumerate_icvs

Summary
The **ompd_enumerate_icvs** function enumerates ICVs.

Format

```c
ompd_rc_t ompd_enumerate_icvs (
    ompd_address_space_handle_t *handle,
    ompd_icv_id_t current,
    ompd_icv_id_t *next_id,
    const char **next_icv_name,
    ompd_scope_t *next_scope,
    int *more
);
```

Description
An OpenMP implementation must support all ICVs listed in Section 2.4.1. An OpenMP implementation may support additional implementation-specific variables. An implementation may store ICVs in a different scope than Table 2.3 indicates. The **ompd_enumerate_icvs** function enables a tool to enumerate the ICVs that an OpenMP implementation supports and their related scopes. The ICVs `num-procs-var`, `thread-num-var`, `final-task-var`, `implicit-task-var` and `team-size-var` must also be available with an **ompd-** prefix.

When the `current` argument is set to the identifier of a supported ICV, **ompd_enumerate_icvs** assigns the value, string name, and scope of the next ICV in the enumeration to the locations to which the `next_id`, `next_icv_name`, and `next_scope` arguments point. On return, the third-party tool owns the `next_icv_name` string. The OMPD library uses the memory allocation callback that the tool provides to allocate the string storage; the tool is responsible for releasing the memory.

On return, the location to which the `more` argument points has the value of 1 whenever one or more ICV are left in the enumeration. On return, that location has the value 0 when `current` is the last ICV in the enumeration.

Description of Arguments
The `address_space_handle` argument identifies the address space. The `current` argument must be an ICV that the OpenMP implementation supports. To begin enumerating the ICVs, a tool should pass `ompd_icv_undefined` as the value of `current`. Subsequent calls to **ompd_enumerate_icvs** should pass the value returned by the call in the `next_id` output argument. On return, the `next_id` argument points to an integer with the value of the ID of the next ICV in the enumeration. On return, the `next_icv_name` argument points to a character string with the name of the next ICV. On return, the `next_scope` argument points to the scope enum value of the scope of the next ICV. On return, the `more_enums` argument points to an integer with the value of 1 when more ICVs are left to enumerate and the value of 0 when no more ICVs are left.
Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- *ompd_rc_bad_input* if an unknown value is provided in *current*.

Cross References
- *ompd_address_space_handle_t* type, see Section 5.3.8.
- *ompd_scope_t* type, see Section 5.3.9.
- *ompd_icv_id_t* type, see Section 5.3.10.
- *ompd_rc_t* type, see Section 5.3.12.

5.5.9.2 ompd_get_icv_from_scope

Summary
The *ompd_get_icv_from_scope* function returns the value of an ICV.

Format

```c
ompd_rc_t ompd_get_icv_from_scope (  
    void *handle,  
    ompd_scope_t scope,  
    ompd_icv_id_t icv_id,  
    ompd_word_t *icv_value  
);  
```

Description
The *ompd_get_icv_from_scope* function provides access to the ICVs that *ompd Enumerate ICVs* identifies.

Description of Arguments
The *handle* argument provides an OpenMP scope handle. The *scope* argument specifies the kind of scope provided in *handle*. The *icv_id* argument specifies the ID of the requested ICV. On return, the *icv_value* argument points to a location with the value of the requested ICV.

Constraints on Arguments
The provided *handle* must match the *scope* as defined in Section 5.3.10.

The provided *scope* must match the scope for *icv_id* as requested by *ompd Enumerate ICVs*. 
**Description of Return Codes**

This routine must return any of the general return codes listed at the beginning of Section 5.5 or any of the following return codes:

- **ompd_rc_incompatible** if the ICV cannot be represented as an integer;
- **ompd_rc_incomplete** if only the first item of the ICV is returned in the integer (e.g., if `nthreads-var` is a list); or
- **ompd_rc_bad_input** if an unknown value is provided in `icv_id`.

**Cross References**

- **ompd_address_space_handle_t** type, see Section 5.3.8.
- **ompd_thread_handle_t** type, see Section 5.3.8.
- **ompd_parallel_handle_t** type, see Section 5.3.8.
- **ompd_task_handle_t** type, see Section 5.3.8.
- **ompd_scope_t** type, see Section 5.3.9.
- **ompd_icv_id_t** type, see Section 5.3.10.
- **ompd_rc_t** type, see Section 5.3.12.
- **ompdEnumerateICVs** routine, see Section 5.5.9.1.

**5.5.9.3 ompd_get_icv_string_from_scope**

**Summary**

The **ompd_get_icv_string_from_scope** function returns the value of an ICV.

**Format**

```c
ompd_rc_t ompd_get_icv_string_from_scope ( 
    void *handle, 
    ompd_scope_t scope, 
    ompd_icv_id_t icv_id, 
    const char **icv_string 
); 
```

**Description**

The **ompd_get_icv_string_from_scope** function provides access to the ICVs that **ompdEnumerateICVs** identifies.
**Description of Arguments**

The `handle` argument provides an OpenMP scope handle. The `scope` argument specifies the kind of scope provided in `handle`. The `icv_id` argument specifies the ID of the requested ICV. On return, the `icv_string` argument points to a string representation of the requested ICV.

On return, the third-party tool owns the `icv_string` string. The OMPD library allocates the string storage with the memory allocation callback that the tool provides. The tool is responsible for releasing the memory.

**Constraints on Arguments**

The provided `handle` must match the `scope` as defined in Section 5.3.10.

The provided `scope` must match the scope for `icv_id` as requested by `ompdEnumerateICVs`.

**Description of Return Codes**

This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- `ompdRCBadInput` if an unknown value is provided in `icv_id`.

**Cross References**

- `ompdAddressSpaceHandle_t` type, see Section 5.3.8.
- `ompdThreadHandle_t` type, see Section 5.3.8.
- `ompdParallelHandle_t` type, see Section 5.3.8.
- `ompdTaskHandle_t` type, see Section 5.3.8.
- `ompdScope_t` type, see Section 5.3.9.
- `ompdICVID_t` type, see Section 5.3.10.
- `ompdRC_t` type, see Section 5.3.12.
- `ompdEnumerateICVs` routine, see Section 5.5.9.1.

5.5.9.4 `ompdGetToolData`

**Summary**

The `ompdGetToolData` function provides access to the OMPT data variable stored for each OpenMP scope.

**Format**

```c
ompdRC_t ompdGetToolData(
   void* handle,
   ompdScope_t scope,
   ompdWord_t *value,
   ompdAddress_t *ptr
);
```
Description
The `ompd_get_tool_data` function provides access to the OMPT tool data stored for each scope. If the runtime library does not support OMPT then the function returns `ompd_rc_unsupported`.

Description of Arguments
The `handle` argument provides an OpenMP scope handle. The `scope` argument specifies the kind of scope provided in `handle`. On return, the `value` argument points to the `value` field of the `ompt_data_t` union stored for the selected scope. On return, the `ptr` argument points to the `ptr` field of the `ompt_data_t` union stored for the selected scope.

Description of Return Codes
This routine must return any of the general return codes listed at the beginning of Section 5.5 or the following return code:

- `ompd_rc_unsupported` if the runtime library does not support OMPT.

Cross References
- `ompt_data_t` type, see Section 4.4.4.4.
- `ompd_address_space_handle_t` type, see Section 5.3.8.
- `ompd_thread_handle_t` type, see Section 5.3.8.
- `ompd_parallel_handle_t` type, see Section 5.3.8.
- `ompd_task_handle_t` type, see Section 5.3.8.
- `ompd_scope_t` type, see Section 5.3.9.
- `ompd_rc_t` type, see Section 5.3.12.

5.6 Runtime Entry Points for OMPD

The OpenMP implementation must define several entry point symbols through which execution must pass when particular events occur and data collection for OMPD is enabled. A tool can enable notification of an event by setting a breakpoint at the address of the entry point symbol.

Entry point symbols have external C linkage and do not require demangling or other transformations to look up their names to obtain the address in the OpenMP program. While each entry point symbol conceptually has a function type signature, it may not be a function. It may be a labeled location.
5.6.1 Beginning Parallel Regions

**Summary**

Before starting the execution of an OpenMP parallel region, the implementation executes `ompd_bp_parallel_begin`.

**Format**

```c
void ompd_bp_parallel_begin(void);
```

**Description**

The OpenMP implementation must execute `ompd_bp_parallel_begin` at every `parallel-begin` event. At the point that the implementation reaches `ompd_bp_parallel_begin`, the binding for `ompd_get_curr_parallel_handle` is the parallel region that is beginning and the binding for `ompd_get_curr_task_handle` is the task that encountered the `parallel` construct.

**Cross References**

- `parallel` construct, see Section 2.6.
- `ompd_get_curr_parallel_handle` routine, see Section 5.5.6.1.
- `ompd_get_curr_task_handle` routine, see Section 5.5.7.1.

5.6.2 Ending Parallel Regions

**Summary**

After finishing the execution of an OpenMP parallel region, the implementation executes `ompd_bp_parallel_end`.

**Format**

```c
void ompd_bp_parallel_end(void);
```

**Description**

The OpenMP implementation must execute `ompd_bp_parallel_end` at every `parallel-end` event. At the point that the implementation reaches `ompd_bp_parallel_end`, the binding for `ompd_get_curr_parallel_handle` is the parallel region that is ending and the binding for `ompd_get_curr_task_handle` is the task that encountered the `parallel` construct.

After execution of `ompd_bp_parallel_end`, any `parallel_handle` that was acquired for the `parallel` region is invalid and should be released.
5.6.3 Beginning Task Regions

Summary
Before starting the execution of an OpenMP task region, the implementation executes `ompd_bp_task_begin`.

Format
```
C
void ompd_bp_task_begin(void);
```

Description
The OpenMP implementation must execute `ompd_bp_task_begin` immediately before starting execution of a `structured-block` that is associated with a non-merged task. At the point that the implementation reaches `ompd_bp_task_begin`, the binding for `ompd_get_curr_task_handle` is the task that is scheduled to execute.

Cross References
• `ompd_get_curr_task_handle` routine, see Section 5.5.7.1.

5.6.4 Ending Task Regions

Summary
After finishing the execution of an OpenMP task region, the implementation executes `ompd_bp_task_end`.

Format
```
C
void ompd_bp_task_end(void);
```
The OpenMP implementation must execute `ompd_bp_task_end` immediately after completion of a structured-block that is associated with a non-merged task. At the point that the implementation reaches `ompd_bp_task_end`, the binding for `ompd_get_curr_task_handle` is the task that finished execution. After execution of `ompd_bp_task_end`, any `task_handle` that was acquired for the task region is invalid and should be released.

Cross References

- `ompd_get_curr_task_handle` routine, see Section 5.5.7.1.
- `ompd_rel_task_handle` routine, see Section 5.5.7.5.

5.6.5 Beginning OpenMP Threads

Summary

When starting an OpenMP thread, the implementation executes `ompd_bp_thread_begin`.

Format

```c
void ompd_bp_thread_begin(void);
```

Description

The OpenMP implementation must execute `ompd_bp_thread_begin` at every `native-thread-begin` and `initial-thread-begin` event. This execution occurs before the thread starts the execution of any OpenMP region.

Cross References

- `parallel` construct, see Section 2.6.
- Initial task, see Section 2.12.5.

5.6.6 Ending OpenMP Threads

Summary

When terminating an OpenMP thread, the implementation executes `ompd_bp_thread_end`.

Format

```c
void ompd_bp_thread_end(void);
```
Description
The OpenMP implementation must execute `ompd_bp_thread_end` at every native-thread-end and initial-thread-end event. This execution occurs after the thread completes the execution of all OpenMP regions. After executing `ompd_bp_thread_end`, any thread_handle that was acquired for this thread is invalid and should be released.

Cross References
- `parallel` construct, see Section 2.6.
- Initial task, see Section 2.12.5.
- `ompd_rel_thread_handle` routine, see Section 5.5.5.3.

5.6.7 Initializing OpenMP Devices

Summary
The OpenMP implementation must execute `ompd_bp_device_begin` at every device-initialize event.

Format
```c
void ompd_bp_device_begin(void);
```

Description
When initializing a device for execution of a target region, the implementation must execute `ompd_bp_device_begin`. This execution occurs before the work associated with any OpenMP region executes on the device.

Cross References
- Device Initialization, see Section 2.14.1.

5.6.8 Finalizing OpenMP Devices

Summary
When terminating an OpenMP thread, the implementation executes `ompd_bp_device_end`.

Format
```c
void ompd_bp_device_end(void);
```
Description
The OpenMP implementation must execute `ompd_bp_device_end` at every `device-finalize` event. This execution occurs after the thread executes all OpenMP regions. After execution of `ompd_bp_device_end`, any `address_space_handle` that was acquired for this device is invalid and should be released.

Cross References
- Device Initialization, see Section 2.14.1.
- `ompd_rel_address_space_handle` routine, see Section 5.5.2.3.
6 Environment Variables

This chapter describes the OpenMP environment variables that specify the settings of the ICVs that affect the execution of OpenMP programs (see Section 2.4). The names of the environment variables must be upper case. Unless otherwise specified, the values assigned to the environment variables are case insensitive and may have leading and trailing white space. Modifications to the environment variables after the program has started, even if modified by the program itself, are ignored by the OpenMP implementation. However, the settings of some of the ICVs can be modified during the execution of the OpenMP program by the use of the appropriate directive clauses or OpenMP API routines.

The following examples demonstrate how the OpenMP environment variables can be set in different environments:

- csh-like shells:
  ```plaintext
  setenv OMP_SCHEDULE "dynamic"
  ```

- bash-like shells:
  ```plaintext
  export OMP_SCHEDULE="dynamic"
  ```

- Windows Command Line:
  ```plaintext
  set OMP_SCHEDULE=dynamic
  ```

As defined following Table 2.1 in Section 2.4.2, device-specific environment variables extend many of the environment variables defined in this chapter. If the corresponding environment variable for a specific device number, including the host device, is set, then the setting for that environment variable is used to set the value of the associated ICV of the device with the corresponding device number. If the corresponding environment variable that includes the __DEV suffix but no device number is set, then the setting of that environment variable is used to set the value of the associated ICV of any non-host device for which the device-number-specific corresponding environment variable is not set. In all cases the setting of an environment variable for which a device number is specified takes precedence.

Restrictions

Restrictions to device-specific environment variables are as follows:

- Device-specific environment variables must not correspond to environment variables that initialize ICVs with global scope.
6.1 **OMP_SCHEDULE**

The **OMP_SCHEDULE** environment variable controls the schedule kind and chunk size of all loop directives that have the schedule kind *runtime*, by setting the value of the *run-sched-var* ICV.

The value of this environment variable takes the form:

```
[modifier:]kind[, chunk]
```

where

- *modifier* is one of **monotonic** or **nonmonotonic**;
- *kind* is one of **static**, **dynamic**, **guided**, or **auto**;
- *chunk* is an optional positive integer that specifies the chunk size.

If the *modifier* is not present, the *modifier* is set to **monotonic** if *kind* is **static**; for any other *kind* it is set to **nonmonotonic**.

If *chunk* is present, white space may be on either side of the “,”. See Section 2.11.4 for a detailed description of the schedule kinds.

The behavior of the program is implementation defined if the value of **OMP_SCHEDULE** does not conform to the above format.

Examples:

```
setenv OMP_SCHEDULE "guided,4"
setenv OMP_SCHEDULE "dynamic"
setenv OMP_SCHEDULE "nonmonotonic:dynamic,4"
```

**Cross References**

- *run-sched-var* ICV, see Section 2.4.
- Worksharing-Loop construct, see Section 2.11.4.
- Parallel worksharing-loop construct, see Section 2.16.1.
- **omp_set_schedule** routine, see Section 3.2.11.
- **omp_get_schedule** routine, see Section 3.2.12.

6.2 **OMP_NUM_THREADS**

The **OMP_NUM_THREADS** environment variable sets the number of threads to use for parallel regions by setting the initial value of the *nthreads-var* ICV. See Section 2.4 for a comprehensive set of rules about the interaction between the **OMP_NUM_THREADS** environment variable, the *num_threads* clause, the **omp_set_num_threads** library routine and dynamic adjustment of
threads, and Section 2.6.1 for a complete algorithm that describes how the number of threads for a
\texttt{parallel} region is determined.

The value of this environment variable must be a list of positive integer values. The values of the
list set the number of threads to use for \texttt{parallel} regions at the corresponding nested levels.

The behavior of the program is implementation defined if any value of the list specified in the
\texttt{OMP_NUM_THREADS} environment variable leads to a number of threads that is greater than an
implementation can support, or if any value is not a positive integer.

The \texttt{OMP_NUM_THREADS} environment variable sets the \texttt{max-active-levels-var} ICV to the number
of active levels of parallelism that the implementation supports if the \texttt{OMP_NUM_THREADS}
environment variable is set to a comma-separated list of more than one value. The value of the
\texttt{max-active-level-var} ICV may be overridden by setting \texttt{OMP_MAX_ACTIVE_LEVELS} or
\texttt{OMP_NESTED}. See Section 6.8 and Section 6.9 for details.

Example:

\begin{verbatim}
| setenv OMP_NUM_THREADS 4,3,2
\end{verbatim}

\section*{Cross References}

\begin{itemize}
  \item \texttt{nthreads-var} ICV, see Section 2.4.
  \item \texttt{num_threads} clause, see Section 2.6.
  \item \texttt{omp_set_num_threads} routine, see Section 3.2.1.
  \item \texttt{omp_get_num_threads} routine, see Section 3.2.2.
  \item \texttt{omp_get_max_threads} routine, see Section 3.2.3.
  \item \texttt{omp_get_team_size} routine, see Section 3.2.19.
\end{itemize}

\section*{6.3 OMP_DYNAMIC}

The \texttt{OMP_DYNAMIC} environment variable controls dynamic adjustment of the number of threads
to use for executing \texttt{parallel} regions by setting the initial value of the \texttt{dyn-var} ICV.

The value of this environment variable must be one of the following:

\begin{verbatim}
| true | false
\end{verbatim}

If the environment variable is set to \texttt{true}, the OpenMP implementation may adjust the number of
threads to use for executing \texttt{parallel} regions in order to optimize the use of system resources. If
the environment variable is set to \texttt{false}, the dynamic adjustment of the number of threads is
disabled. The behavior of the program is implementation defined if the value of \texttt{OMP_DYNAMIC} is
neither \texttt{true} nor \texttt{false}.
Example:

```
setenv OMP_DYNAMIC true
```

Cross References

- *dyn-var ICV*, see Section 2.4.
- *omp_set_dynamic* routine, see Section 3.2.6.
- *omp_get_dynamic* routine, see Section 3.2.7.

### 6.4 OMP_PROC_BIND

The **OMP_PROC_BIND** environment variable sets the initial value of the *bind-var ICV*. The value of this environment variable is either `true`, `false`, or a comma separated list of `primary`, `master` (*master* has been deprecated), `close`, or `spread`. The values of the list set the thread affinity policy to be used for parallel regions at the corresponding nested level.

If the environment variable is set to `false`, the execution environment may move OpenMP threads between OpenMP places, thread affinity is disabled, and **proc_bind** clauses on **parallel** constructs are ignored.

Otherwise, the execution environment should not move OpenMP threads between OpenMP places, thread affinity is enabled, and the initial thread is bound to the first place in the *place-partition-var ICV* prior to the first active parallel region. An initial thread that is created by a **teams** construct is bound to the first place in its *place-partition-var ICV* before it begins execution of the associated structured block.

If the environment variable is set to `true`, the thread affinity policy is implementation defined but must conform to the previous paragraph. The behavior of the program is implementation defined if the value in the **OMP_PROC_BIND** environment variable is not `true`, `false`, or a comma separated list of `primary`, `master` (*master* has been deprecated), `close`, or `spread`. The behavior is also implementation defined if an initial thread cannot be bound to the first place in the *place-partition-var ICV*.

The **OMP_PROC_BIND** environment variable sets the *max-active-levels-var ICV* to the number of active levels of parallelism that the implementation supports if the **OMP_PROC_BIND** environment variable is set to a comma-separated list of more than one element. The value of the *max-active-level-var ICV* may be overridden by setting **OMP_MAX_ACTIVE_LEVELS** or **OMP_NESTED**. See Section 6.8 and Section 6.9 for details.

Examples:

```
setenv OMP_PROC_BIND false
setenv OMP_PROC_BIND "spread, spread, close"
```
6.5 OMP_PLACES

The OMP_PLACES environment variable sets the initial value of the place-partition-var ICV. A list of places can be specified in the OMP_PLACES environment variable. The value of OMP_PLACES can be one of two types of values: either an abstract name that describes a set of places or an explicit list of places described by non-negative numbers.

The OMP_PLACES environment variable can be defined using an explicit ordered list of comma-separated places. A place is defined by an unordered set of comma-separated non-negative numbers enclosed by braces, or a non-negative number. The meaning of the numbers and how the numbering is done are implementation defined. Generally, the numbers represent the smallest unit of execution exposed by the execution environment, typically a hardware thread.

Intervals may also be used to define places. Intervals can be specified using the \(<lower-bound> : <length> : <stride>\) notation to represent the following list of numbers: \("<lower-bound>, <lower-bound> + <stride>, ..., <lower-bound> + (<length> - 1)\times <stride>\)." When \(<stride>\) is omitted, a unit stride is assumed. Intervals can specify numbers within a place as well as sequences of places.

An exclusion operator \("!\) can also be used to exclude the number or place immediately following the operator.

Alternatively, the abstract names listed in Table 6.1 should be understood by the execution and runtime environment. The precise definitions of the abstract names are implementation defined. An implementation may also add abstract names as appropriate for the target platform.

The abstract name may be appended by a positive number in parentheses to denote the length of the place list to be created, that is \(abstract\_name(num\_places)\). When requesting fewer places than available on the system, the determination of which resources of type \(abstract\_name\) are to be included in the place list is implementation defined. When requesting more resources than available, the length of the place list is implementation defined.
### Table 6.1: Predefined Abstract Names for `OMP_PLACES`

<table>
<thead>
<tr>
<th>Abstract Name</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>threads</code></td>
<td>Each place corresponds to a single hardware thread on the device.</td>
</tr>
<tr>
<td><code>cores</code></td>
<td>Each place corresponds to a single core (having one or more hardware threads) on the device.</td>
</tr>
<tr>
<td><code>ll_caches</code></td>
<td>Each place corresponds to a set of cores that share the last level cache on the device.</td>
</tr>
</tbody>
</table>
| `numa_domains`| Each place corresponds to a set of cores for which their closest memory on the device is:  
  - the same memory; and  
  - at a similar distance from the cores. |
| `sockets`     | Each place corresponds to a single socket (consisting of one or more cores) on the device. |

The behavior of the program is implementation defined when the execution environment cannot map a numerical value (either explicitly defined or implicitly derived from an interval) within the `OMP_PLACES` list to a processor on the target platform, or if it maps to an unavailable processor. The behavior is also implementation defined when the `OMP_PLACES` environment variable is defined using an abstract name.

The following grammar describes the values accepted for the `OMP_PLACES` environment variable.

\[
\begin{align*}
\langle \text{list} \rangle & \mid= \langle \text{p-list} \rangle \mid \langle \text{aname} \rangle \\
\langle \text{p-list} \rangle & \mid= \langle \text{p-interval} \rangle \mid \langle \text{p-list} \rangle,\langle \text{p-interval} \rangle \\
\langle \text{p-interval} \rangle & \mid= \langle \text{place} \rangle:\langle \text{len} \rangle:\langle \text{stride} \rangle \mid \langle \text{place} \rangle:\langle \text{len} \rangle \mid \langle \text{place} \rangle \mid !\langle \text{place} \rangle \\
\langle \text{place} \rangle & \mid= \{ \langle \text{res-list} \rangle \} \mid \langle \text{res} \rangle \\
\langle \text{res-list} \rangle & \mid= \langle \text{res-interval} \rangle \mid \langle \text{res-list} \rangle,\langle \text{res-interval} \rangle \\
\langle \text{res-interval} \rangle & \mid= \langle \text{res} \rangle:\langle \text{num-places} \rangle:\langle \text{stride} \rangle \mid \langle \text{res} \rangle:\langle \text{num-places} \rangle \mid \langle \text{res} \rangle \mid !\langle \text{res} \rangle \\
\langle \text{aname} \rangle & \mid= \langle \text{word} \rangle(\langle \text{num-places} \rangle) \mid \langle \text{word} \rangle \\
\langle \text{word} \rangle & \mid= \text{sockets} \mid \text{cores} \mid \text{ll_caches} \mid \text{numa_domains} \mid \text{threads} \mid \text{<implementation-defined abstract name>} \\
\langle \text{res} \rangle & \mid= \text{non-negative integer} \\
\langle \text{num-places} \rangle & \mid= \text{positive integer} \\
\langle \text{stride} \rangle & \mid= \text{integer} \\
\langle \text{len} \rangle & \mid= \text{positive integer}
\end{align*}
\]
Examples:

```
setenv OMP_PLACES threads
setenv OMP_PLACES "threads(4)"
setenv OMP_PLACES 
    "{0,1,2,3},{4,5,6,7},{8,9,10,11},{12,13,14,15}"
setenv OMP_PLACES "{0:4},{4:4},{8:4},{12:4}"
setenv OMP_PLACES "{0:4}:4:4"
```

where each of the last three definitions corresponds to the same 4 places including the smallest
units of execution exposed by the execution environment numbered, in turn, 0 to 3, 4 to 7, 8 to 11,
and 12 to 15.

Cross References

- `place-partition-var`, see Section 2.4.
- Controlling OpenMP thread affinity, see Section 2.6.2.
- `omp_get_num_places` routine, see Section 3.3.2.
- `omp_get_place_num_procs` routine, see Section 3.3.3.
- `omp_get_place_proc_ids` routine, see Section 3.3.4.
- `omp_get_place_num` routine, see Section 3.3.5.
- `omp_get_partition_num_places` routine, see Section 3.3.6.
- `omp_get_partition_place_nums` routine, see Section 3.3.7.

6.6 OMP_STACKSIZE

The `OMP_STACKSIZE` environment variable controls the size of the stack for threads created by
the OpenMP implementation, by setting the value of the `stacksize-var` ICV. The environment
variable does not control the size of the stack for an initial thread.

The value of this environment variable takes the form:

```
size | sizeB | sizeK | sizeM | sizeG
```

where:

- `size` is a positive integer that specifies the size of the stack for threads that are created by the
  OpenMP implementation.
- `B, K, M, and G` are letters that specify whether the given size is in Bytes, Kilobytes (1024 Bytes),
  Megabytes (1024 Kilobytes), or Gigabytes (1024 Megabytes), respectively. If one of these letters
  is present, white space may occur between `size` and the letter.
If only size is specified and none of B, K, M, or G is specified, then size is assumed to be in Kilobytes.

The behavior of the program is implementation defined if OMP_STACKSIZE does not conform to the above format, or if the implementation cannot provide a stack with the requested size.

Examples:

<table>
<thead>
<tr>
<th>Command</th>
</tr>
</thead>
<tbody>
<tr>
<td>setenv OMP_STACKSIZE 2000500B</td>
</tr>
<tr>
<td>setenv OMP_STACKSIZE &quot;3000 k &quot;</td>
</tr>
<tr>
<td>setenv OMP_STACKSIZE 10M</td>
</tr>
<tr>
<td>setenv OMP_STACKSIZE &quot; 10 M &quot;</td>
</tr>
<tr>
<td>setenv OMP_STACKSIZE &quot;20 m &quot;</td>
</tr>
<tr>
<td>setenv OMP_STACKSIZE &quot; 1G&quot;</td>
</tr>
<tr>
<td>setenv OMP_STACKSIZE 20000</td>
</tr>
</tbody>
</table>

Cross References
- stacksize-var ICV, see Section 2.4.

6.7 OMP_WAIT_POLICY

The OMP_WAIT_POLICY environment variable provides a hint to an OpenMP implementation about the desired behavior of waiting threads by setting the wait-policy-var ICV. A compliant OpenMP implementation may or may not abide by the setting of the environment variable.

The value of this environment variable must be one of the following:

active | passive

The active value specifies that waiting threads should mostly be active, consuming processor cycles, while waiting. An OpenMP implementation may, for example, make waiting threads spin.

The passive value specifies that waiting threads should mostly be passive, not consuming processor cycles, while waiting. For example, an OpenMP implementation may make waiting threads yield the processor to other threads or go to sleep.

The details of the active and passive behaviors are implementation defined.

The behavior of the program is implementation defined if the value of OMP_WAIT_POLICY is neither active nor passive.

Examples:

<table>
<thead>
<tr>
<th>Command</th>
</tr>
</thead>
<tbody>
<tr>
<td>setenv OMP_WAIT_POLICY ACTIVE</td>
</tr>
<tr>
<td>setenv OMP_WAIT_POLICY active</td>
</tr>
<tr>
<td>setenv OMP_WAIT_POLICY PASSIVE</td>
</tr>
<tr>
<td>setenv OMP_WAIT_POLICY passive</td>
</tr>
</tbody>
</table>

Cross References
- wait-policy-var ICV, see Section 2.4.
6.8 OMP_MAX_ACTIVE_LEVELS

The OMP_MAX_ACTIVE_LEVELS environment variable controls the maximum number of nested active parallel regions by setting the initial value of the max-active-levels-var ICV.

The value of this environment variable must be a non-negative integer. The behavior of the program is implementation defined if the requested value of OMP_MAX_ACTIVE_LEVELS is greater than the maximum number of nested active parallel levels an implementation can support, or if the value is not a non-negative integer.

Cross References
- max-active-levels-var ICV, see Section 2.4.
- omp_set_max_active_levels routine, see Section 3.2.15.
- omp_get_max_active_levels routine, see Section 3.2.16.

6.9 OMP_NESTED (Deprecated)

The OMP_NESTED environment variable controls nested parallelism by setting the initial value of the max-active-levels-var ICV. If the environment variable is set to true, the initial value of max-active-levels-var is set to the number of active levels of parallelism supported by the implementation. If the environment variable is set to false, the initial value of max-active-levels-var is set to 1. The behavior of the program is implementation defined if the value of OMP_NESTED is neither true nor false.

If both the OMP_NESTED and OMP_MAX_ACTIVE_LEVELS environment variables are set, the value of OMP_NESTED is false, and the value of OMP_MAX_ACTIVE_LEVELS is greater than 1, then the behavior is implementation defined. Otherwise, if both environment variables are set then the OMP_NESTED environment variable has no effect.

The OMP_NESTED environment variable has been deprecated.

Example:

```
setenv OMP_NESTED false
```

Cross References
- max-active-levels-var ICV, see Section 2.4.
- omp_set_nested routine, see Section 3.2.9.
- omp_get_team_size routine, see Section 3.2.19.
- OMP_MAX_ACTIVE_LEVELS environment variable, see Section 6.8.
6.10 OMP_THREAD_LIMIT

The OMP_THREAD_LIMIT environment variable sets the maximum number of OpenMP threads to use in a contention group by setting the thread-limit-var ICV.

The value of this environment variable must be a positive integer. The behavior of the program is implementation defined if the requested value of OMP_THREAD_LIMIT is greater than the number of threads an implementation can support, or if the value is not a positive integer.

Cross References

• thread-limit-var ICV, see Section 2.4.
• omp_get_thread_limit routine, see Section 3.2.13.

6.11 OMP_CANCELLATION

The OMP_CANCELLATION environment variable sets the initial value of the cancel-var ICV.

The value of this environment variable must be one of the following:

true | false

If the environment variable is set to true, the effects of the cancel construct and of cancellation points are enabled and cancellation is activated. If the environment variable is set to false, cancellation is disabled and the cancel construct and cancellation points are effectively ignored. The behavior of the program is implementation defined if OMP_CANCELLATION is set to neither true nor false.

Cross References

• cancel-var, see Section 2.4.1.
• cancel construct, see Section 2.20.1.
• cancellation point construct, see Section 2.20.2.
• omp_get_cancellation routine, see Section 3.2.8.

6.12 OMP_DISPLAY_ENV

The OMP_DISPLAY_ENV environment variable instructs the runtime to display the information as described in the omp_display_env routine section (Section 3.15).

The value of the OMP_DISPLAY_ENV environment variable may be set to one of these values:

true | false | verbose
If the environment variable is set to `true`, the effect is as if the `omp_display_env` routine is called with the `verbose` argument set to `false` at the beginning of the program. If the environment variable is set to `verbose`, the effect is as if the `omp_display_env` routine is called with the `verbose` argument set to `true` at the beginning of the program. If the environment variable is undefined or set to `false`, the runtime does not display any information. For all values of the environment variable other than `true`, `false`, and `verbose`, the displayed information is unspecified.

Example:
```
% setenv OMP_DISPLAY_ENV true
```

For the output of the above example, see Section 3.15.

Cross References
- `omp_display_env` routine, see Section 3.15.

### 6.13 OMP_DISPLAY_AFFINITY

The `OMP_DISPLAY_AFFINITY` environment variable instructs the runtime to display formatted affinity information for all OpenMP threads in the parallel region upon entering the first parallel region and when any change occurs in the information accessible by the format specifiers listed in Table 6.2. If affinity of any thread in a parallel region changes then thread affinity information for all threads in that region is displayed. If the thread affinity for each respective parallel region at each nesting level has already been displayed and the thread affinity has not changed, then the information is not displayed again. Thread affinity information for threads in the same parallel region may be displayed in any order.

The value of the `OMP_DISPLAY_AFFINITY` environment variable may be set to one of these values:

```
true | false
```

The `true` value instructs the runtime to display the OpenMP thread affinity information, and uses the format setting defined in the `affinity-format-var ICV`.

The runtime does not display the OpenMP thread affinity information when the value of the `OMP_DISPLAY_AFFINITY` environment variable is `false` or undefined. For all values of the environment variable other than `true` or `false`, the display action is implementation defined.

Example:
```
setenv OMP_DISPLAY_AFFINITY TRUE
```

The above example causes an OpenMP implementation to display OpenMP thread affinity information during execution of the program, in a format given by the `affinity-format-var ICV`. The following is a sample output:
Cross References

- Controlling OpenMP thread affinity, see Section 2.6.2.
- `omp_set_affinity_format` routine, see Section 3.3.8.
- `omp_get_affinity_format` routine, see Section 3.3.9.
- `omp_display_affinity` routine, see Section 3.3.10.
- `omp_capture_affinity` routine, see Section 3.3.11.
- `OMP_AFFINITY_FORMAT` environment variable, see Section 6.14.

## 6.14 OMP_AFFINITY_FORMAT

The `OMP_AFFINITY_FORMAT` environment variable sets the initial value of the affinity-format-var ICV which defines the format when displaying OpenMP thread affinity information.

The value of this environment variable is case sensitive and leading and trailing whitespace is significant.

The value of this environment variable is a character string that may contain as substrings one or more field specifiers, in addition to other characters. The format of each field specifier is

```%
```

where an individual field specifier must contain the percent symbol (%) and a type. The type can be a single character short name or its corresponding long name delimited with curly braces, such as `%n` or `%{thread_num}`. A literal percent is specified as `%%`. Field specifiers can be provided in any order.

The 0 modifier indicates whether or not to add leading zeros to the output, following any indication of sign or base. The . modifier indicates the output should be right justified when `size` is specified. By default, output is left justified. The minimum field length is `size`, which is a decimal digit string with a non-zero first digit. If no `size` is specified, the actual length needed to print the field will be used. If the 0 modifier is used with type of `A`, `{thread_affinity}`, `H`, `{host}`, or a type that is not printed as a number, the result is unspecified. Any other characters in the format string that are not part of a field specifier will be included literally in the output.
<table>
<thead>
<tr>
<th>Short Name</th>
<th>Long Name</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>t</td>
<td>team_num</td>
<td>The value returned by <code>omp_get_team_num()</code></td>
</tr>
<tr>
<td>T</td>
<td>num_teams</td>
<td>The value returned by <code>omp_get_num_teams()</code></td>
</tr>
<tr>
<td>L</td>
<td>nesting_level</td>
<td>The value returned by <code>omp_get_level()</code></td>
</tr>
<tr>
<td>n</td>
<td>thread_num</td>
<td>The value returned by <code>omp_get_thread_num()</code></td>
</tr>
<tr>
<td>N</td>
<td>num_threads</td>
<td>The value returned by <code>omp_get_num_threads()</code></td>
</tr>
<tr>
<td>a</td>
<td>ancestor_tnum</td>
<td>The value returned by <code>omp_get_ancestor_thread_num(level)</code>, where <code>level</code> is <code>omp_get_level()</code> minus 1</td>
</tr>
<tr>
<td>H</td>
<td>host</td>
<td>The name for the host device on which the OpenMP program is running.</td>
</tr>
<tr>
<td>P</td>
<td>process_id</td>
<td>The process identifier used by the implementation.</td>
</tr>
<tr>
<td>i</td>
<td>native_thread_id</td>
<td>The native thread identifier used by the implementation.</td>
</tr>
<tr>
<td>A</td>
<td>thread_affinity</td>
<td>The list of numerical identifiers, in the format of a comma-separated list of integers or integer ranges, that represent processors on which a thread may execute, subject to OpenMP thread affinity control and/or other external affinity mechanisms.</td>
</tr>
</tbody>
</table>

Implementations may define additional field types. If an implementation does not have information for a field type, "undefined" is printed for this field when displaying the OpenMP thread affinity information.

Example:

```bash
setenv OMP_AFFINITY_FORMAT
"Thread Affinity: %0.3L %.8n %.15{thread_affinity} %.12H"
```

The above example causes an OpenMP implementation to display OpenMP thread affinity information in the following form:

```
Thread Affinity: 001 0 0-1,16-17 nid003
Thread Affinity: 001 1 2-3,18-19 nid003
```

Cross References
- Controlling OpenMP thread affinity, see Section 2.6.2.
- `omp_set_affinity_format` routine, see Section 3.3.8.
• `omp_get_affinity_format` routine, see Section 3.3.9.
• `omp_display_affinity` routine, see Section 3.3.10.
• `omp_capture_affinity` routine, see Section 3.3.11.
• `OMP_DISPLAY_AFFINITY` environment variable, see Section 6.13.

6.15 OMP_DEFAULT_DEVICE

The `OMP_DEFAULT_DEVICE` environment variable sets the device number to use in device constructs by setting the initial value of the `default-device-var` ICV.

The value of this environment variable must be a non-negative integer value.

Cross References
• `default-device-var` ICV, see Section 2.4.
• device directives, Section 2.14.

6.16 OMP_MAX_TASK_PRIORITY

The `OMP_MAX_TASK_PRIORITY` environment variable controls the use of task priorities by setting the initial value of the `max-task-priority-var` ICV. The value of this environment variable must be a non-negative integer.

Example:
```bash
% setenv OMP_MAX_TASK_PRIORITY 20
```

Cross References
• `max-task-priority-var` ICV, see Section 2.4.
• Tasking Constructs, see Section 2.12.
• `omp_get_max_task_priority` routine, see Section 3.5.1.

6.17 OMP_TARGET_OFFLOAD

The `OMP_TARGET_OFFLOAD` environment variable sets the initial value of the `target-offload-var` ICV. The value of the `OMP_TARGET_OFFLOAD` environment variable must be one of the following:

```bash
mandatory | disabled | default
```
The **mandatory** value specifies that program execution is terminated if a device construct or device memory routine is encountered and the device is not available or is not supported by the implementation. Support for the **disabled** value is implementation defined. If an implementation supports it, the behavior is as if the only device is the host device.

The **default** value specifies the default behavior as described in Section 1.3.

Example:

```bash
% setenv OMP_TARGET_OFFLOAD mandatory
```

**Cross References**

- `target-offload-var` ICV, see Section 2.4.
- Device Directives, see Section 2.14.
- Device Memory Routines, see Section 3.8.

### 6.18 OMP_TOOL

The **OMP_TOOL** environment variable sets the **tool-var** ICV, which controls whether an OpenMP runtime will try to register a first party tool.

The value of this environment variable must be one of the following:

- `enabled` | `disabled`

If **OMP_TOOL** is set to any value other than `enabled` or `disabled`, the behavior is unspecified. If **OMP_TOOL** is not defined, the default value for **tool-var** is `enabled`.

Example:

```bash
% setenv OMP_TOOL enabled
```

**Cross References**

- `tool-var` ICV, see Section 2.4.
- OMPT Interface, see Chapter 4.

### 6.19 OMP_TOOL_LIBRARIES

The **OMP_TOOL_LIBRARIES** environment variable sets the **tool-libraries-var** ICV to a list of tool libraries that are considered for use on a device on which an OpenMP implementation is being initialized. The value of this environment variable must be a list of names of dynamically-loadable libraries, separated by an implementation specific, platform typical separator. Whether the value of this environment variable is case sensitive is implementation defined.
If the `tool-var ICV` is not enabled, the value of `tool-libraries-var` is ignored. Otherwise, if `ompt_start_tool` is not visible in the address space on a device where OpenMP is being initialized or if `ompt_start_tool` returns `NULL`, an OpenMP implementation will consider libraries in the `tool-libraries-var` list in a left to right order. The OpenMP implementation will search the list for a library that meets two criteria: it can be dynamically loaded on the current device and it defines the symbol `ompt_start_tool`. If an OpenMP implementation finds a suitable library, no further libraries in the list will be considered.

Example:

```
% setenv OMP_TOOL_LIBRARIES libtoolXY64.so:/usr/local/lib/
   libtoolXY32.so
```

**Cross References**
- `tool-libraries-var ICV`, see Section 2.4.
- OMPT Interface, see Chapter 4.
- `ompt_start_tool` routine, see Section 4.2.1.

## 6.20 OMP_TOOL_VERBOSE_INIT

The `OMP_TOOL_VERBOSE_INIT` environment variable sets the `tool-verbose-init-var ICV`, which controls whether an OpenMP implementation will verbosely log the registration of a tool.

The value of this environment variable must be one of the following:

```
disabled | stdout | stderr | <filename>
```

If `OMP_TOOL_VERBOSE_INIT` is set to any value other than case insensitive `disabled`, `stdout` or `stderr`, the value is interpreted as a filename and the OpenMP runtime will try to log to a file with prefix `filename`. If the value is interpreted as a filename, whether it is case sensitive is implementation defined. If opening the logfile fails, the output will be redirected to stderr. If `OMP_TOOL_VERBOSE_INIT` is not defined, the default value for `tool-verbose-init-var` is `disabled`. Support for logging to `stdout` or `stderr` is implementation defined. Unless `tool-verbose-init-var` is `disabled`, the OpenMP runtime will log the steps of the tool activation process defined in Section 4.2.2 to a file with a name that is constructed using the provided filename prefix. The format and detail of the log is implementation defined. At a minimum, the log will contain the following:

- either that tool-var is disabled, or
- an indication that a tool was available in the address space at program launch, or
- the path name of each tool in `OMP_TOOL_LIBRARIES` that is considered for dynamic loading, whether dynamic loading was successful, and whether the `ompt_start_tool` function is found in the loaded library.
In addition, if an `ompt_start_tool` function is called the log will indicate whether or not the tool will use the OMPT interface.

Example:

```
% setenv OMP_TOOL_VERBOSE_INIT disabled
% setenv OMP_TOOL_VERBOSE_INIT STDERR
% setenv OMP_TOOL_VERBOSE_INIT ompt_load.log
```

Cross References
- `tool-verbose-init-var` ICV, see Section 2.4.
- OMPT Interface, see Chapter 4.

### 6.21 `OMP_DEBUG`

The `OMP_DEBUG` environment variable sets the `debug-var` ICV, which controls whether an OpenMP runtime collects information that an OMPD library may need to support a tool.

The value of this environment variable must be one of the following:

```
enabled | disabled
```

If `OMP_DEBUG` is set to any value other than `enabled` or `disabled` then the behavior is implementation defined.

Example:

```
% setenv OMP_DEBUG enabled
```

Cross References
- `debug-var` ICV, see Section 2.4.
- OMPD Interface, see Chapter 5.
- Enabling the Runtime for OMPD, see Section 5.2.1.

### 6.22 `OMP_ALLOCATOR`

The `OMP_ALLOCATOR` environment variable sets the initial value of the `def-allocator-var` ICV that specifies the default allocator for allocation calls, directives and clauses that do not specify an allocator.

The following grammar describes the values accepted for the `OMP_ALLOCATOR` environment variable.
\[\text{allocator} \mid (\text{predef-allocator}) \mid (\text{predef-mem-space}) \mid (\text{predef-mem-space}): (\text{traits})\]

\[\text{traits} \mid (\text{trait}) = (\text{value}) \mid (\text{trait}) = (\text{value}), (\text{traits})\]

\[\text{predef-allocator} \mid \text{one of the predefined allocators from Table 2.10}\]

\[\text{predef-mem-space} \mid \text{one of the predefined memory spaces from Table 2.8}\]

\[\text{trait} \mid \text{one of the allocator trait names from Table 2.9}\]

\[\text{value} \mid \text{one of the allowed values from Table 2.9} \mid \text{non-negative integer}\]

\[\text{allocator}\]

\[\langle\text{allocator}\rangle \mid \langle\text{predef-allocator}\rangle \mid \langle\text{predef-mem-space}\rangle \mid (\text{predef-mem-space}): (\text{traits})\]

\[\langle\text{traits}\rangle \mid (\text{trait}) = (\text{value}) \mid (\text{trait}) = (\text{value}), (\text{traits})\]

\[\langle\text{predef-allocator}\rangle \mid \langle\text{predef-mem-space}\rangle \mid (\text{traits})\]

\[\langle\text{predef-allocator}\rangle \mid \langle\text{predef-mem-space}\rangle \mid (\text{traits})\]

\[\langle\text{traits}\rangle \mid (\text{trait}) = (\text{value}) \mid (\text{trait}) = (\text{value}), (\text{traits})\]

\[\langle\text{predef-allocator}\rangle \mid \text{one of the predefined allocators from Table 2.10}\]

\[\langle\text{predef-mem-space}\rangle \mid \text{one of the predefined memory spaces from Table 2.8}\]

\[\langle\text{trait}\rangle \mid \text{one of the allocator trait names from Table 2.9}\]

\[\langle\text{value}\rangle \mid \text{one of the allowed values from Table 2.9} \mid \text{non-negative integer}\]

\[\langle\text{predef-allocator}\rangle \mid \langle\text{predef-mem-space}\rangle\]

\[\langle\text{value}\rangle \mid \text{can be an integer only if the \text{trait} accepts a numerical value, for the \text{fb_data} trait the \text{value} can only be \text{predef-allocator}. If the value of this environment variable is not a predefined allocator, then a new allocator with the given predefined memory space and optional traits is created and set as the \text{def-allocator-var} ICV. If the new allocator cannot be created, the \text{def-allocator-var} ICV will be set to \text{omp_default_mem_alloc}.\]

\[\text{Example:}\]

\[\text{setenv OMP_ALLOCATOR omp_high_bw_mem_alloc}\]

\[\text{setenv OMP_ALLOCATOR omp_large_cap_mem_space:alignment=16,}\]

\[\text{pinned=true}\]

\[\text{setenv OMP_ALLOCATOR omp_high_bw_mem_space:pool_size=1048576,}\]

\[\text{fallback=allocator_fb,fb_data=omp_low_lat_mem_alloc}\]

\[\text{Cross References}\]

\[\bullet \text{def-allocator-var ICV, see Section 2.4.}\]

\[\bullet \text{Memory allocators, see Section 2.13.2.}\]

\[\bullet \text{omp_set_default_allocator routine, see Section 3.13.4.}\]

\[\bullet \text{omp_get_default_allocator routine, see Section 3.13.5.}\]

\[\bullet \text{omp_alloc and omp_aligned_alloc routines, see Section 3.13.6}\]

\[\bullet \text{omp_calloc and omp_aligned_calloc routines, see Section 3.13.8}\]

\[\text{6.23 OMP_NUM_TEAMS}\]

\[\text{The OMP_NUM_TEAMS environment variable sets the maximum number of teams created by a}\]

\[\text{teams construct by setting the nteams-var ICV.}\]
The value of this environment variable must be a positive integer. The behavior of the program is implementation defined if the requested value of `OMP_NUM_TEAMS` is greater than the number of teams that an implementation can support, or if the value is not a positive integer.

**Cross References**
- `nteams-var` ICV, see Section 2.4.
- `omp_get_max_teams` routine, see Section 3.4.4.

### 6.24 `OMP_TEAMS_THREAD_LIMIT`

The `OMP_TEAMS_THREAD_LIMIT` environment variable sets the maximum number of OpenMP threads to use in each contention group created by a `teams` construct by setting the `teams-thread-limit-var` ICV.

The value of this environment variable must be a positive integer. The behavior of the program is implementation defined if the requested value of `OMP_TEAMS_THREAD_LIMIT` is greater than the number of threads that an implementation can support, or if the value is not a positive integer.

**Cross References**
- `teams-thread-limit-var` ICV, see Section 2.4.
- `omp_get_teams_thread_limit` routine, see Section 3.4.6.
A OpenMP Implementation-Defined Behaviors

This appendix summarizes the behaviors that are described as implementation defined in this API. Each behavior is cross-referenced back to its description in the main specification. An implementation is required to define and to document its behavior in these cases.

Chapter 1:
- Processor: A hardware unit that is implementation defined (see Section 1.2.1).
- Device: An implementation defined logical execution engine (see Section 1.2.1).
- Device pointer: an implementation defined handle that refers to a device address (see Section 1.2.6).
- Supported active levels of parallelism: The maximum number of active parallel regions that may enclose any region of code in the program is implementation defined (see Section 1.2.7).
- Memory model: The minimum size at which a memory update may also read and write back adjacent variables that are part of another variable (as array or structure elements) is implementation defined but is no larger than required by the base language. The manner in which a program can obtain the referenced device address from a device pointer, outside the mechanisms specified by OpenMP, is implementation defined (see Section 1.4.1).

Chapter 2:
- OpenMP context: Whether the dispatch construct is added to the construct set, the accepted isa-name values for the isa trait, the accepted arch-name values for the arch trait, and the accepted extension-name values for the extension trait are implementation defined (see Section 2.3.1).
- Metadirectives: The number of times that each expression of the context selector of a when clause is evaluated is implementation defined (see Section 2.3.4).
- Declare variant directive: If two replacement candidates have the same score, their order is implementation defined. The number of times each expression of the context selector of a match clause is evaluated is implementation defined. For calls to constexpr base functions that are evaluated in constant expressions, whether any variant replacement occurs is implementation defined. Any differences that the specific OpenMP context requires in the prototype of the variant from the base function prototype are implementation defined (see Section 2.3.5).
• **Internal control variables**: The initial values of `dyn-var`, `nthreads-var`, `run-sched-var`, `def-sched-var`, `bind-var`, `stacksize-var`, `wait-policy-var`, `thread-limit-var`, `max-active-levels-var`, `place-partition-var`, `affinity-format-var`, `default-device-var`, `num-procs-var` and `def-allocator-var` are implementation defined (see Section 2.4.2).

• **requires directive**: Support for any feature specified by a requirement clause on a `requires` directive is implementation defined (see Section 2.5.1).

• **Dynamic adjustment of threads**: Providing the ability to adjust the number of threads dynamically is implementation defined (see Section 2.6.1).

• **Thread affinity**: For the **close** thread affinity policy, if `T > P` and `P` does not divide `T` evenly, the exact number of threads in a particular place is implementation defined. For the **spread** thread affinity, if `T > P` and `P` does not divide `T` evenly, the exact number of threads in a particular subpartition is implementation defined. The determination of whether the affinity request can be fulfilled is implementation defined. If not, the mapping of threads in the team to places is implementation defined (see Section 2.6.2).

• **teams construct**: The number of teams that are created is implementation defined, it is greater than or equal to the lower bound and less than or equal to the upper bound values of the `num_teams` clause if specified or it is less than or equal to the value of the `nteams-var` ICV if its value is greater than zero. Otherwise it is greater than or equal to 1. The maximum number of threads that participate in the contention group that each team initiates is implementation defined if no `thread_limit` clause is specified on the construct. The assignment of the initial threads to places and the values of the `place-partition-var` and `default-device-var` ICVs for each initial thread are implementation defined (see Section 2.7).

• **sections construct**: The method of scheduling the structured blocks among threads in the team is implementation defined (see Section 2.10.1).

• **single construct**: The method of choosing a thread to execute the structured block each time the team encounters the construct is implementation defined (see Section 2.10.2).

• **Canonical loop nest form**: The particular integer type used to compute the iteration count for the collapsed loop is implementation defined (see Section 2.11.1).

• **Worksharing-loop directive**: The effect of the `schedule(runtime)` clause when the `run-sched-var` ICV is set to `auto` is implementation defined. The value of `simd_width` for the `simd` schedule modifier is implementation defined (see Section 2.11.4).

• **simd construct**: The number of iterations that are executed concurrently at any given time is implementation defined. If the `alignment` parameter is not specified in the `aligned` clause, the default alignments for the SIMD instructions are implementation defined (see Section 2.11.5.1).

• **declare simd directive**: If the parameter of the `simdlen` clause is not a constant positive integer expression, the number of concurrent arguments for the function is implementation defined. If the `alignment` parameter of the `aligned` clause is not specified, the default alignments for SIMD instructions are implementation defined (see Section 2.11.5.3).
• **distribute construct**: If no `dist_schedule` clause is specified then the schedule for the `distribute` construct is implementation defined (see Section 2.11.6.1).

• **unroll construct**: If the `partial` clause is specified without an argument, the unroll factor is a positive integer that is implementation defined. If neither the `partial` nor the `full` clause is specified, if and how the loop is unrolled is implementation defined (see Section 2.11.9.2).

• **taskloop construct**: The number of loop iterations assigned to a task created from a `taskloop` construct is implementation defined, unless the `grainsize` or `num_tasks` clause is specified (see Section 2.12.2).

```
C++

• **taskloop construct**: For `firstprivate` variables of class type, the number of invocations of copy constructors to perform the initialization is implementation defined (see Section 2.12.2).
```

• **Memory spaces**: The actual storage resources that each memory space defined in Table 2.8 represents are implementation defined (see Section 2.13.1).

• **Memory allocators**: The minimum partitioning size for partitioning of allocated memory over the storage resources is implementation defined. The default value for the `pool_size` allocator trait (see Table 2.9) is implementation defined. The associated memory space for each of the predefined `omp_cgroup_mem_alloc`, `omp_pteam_mem_alloc` and `omp_thread_mem_alloc` allocators (see Table 2.10) is implementation defined (see Section 2.13.2).

• **target construct**: The maximum number of threads that participate in the contention group that each team initiates is implementation defined if no `thread_limit` clause is specified on the construct (see Section 2.14.5).

• **is_device_ptr clause**: Support for pointers created outside of the OpenMP device data management routines is implementation defined (see Section 2.14.5).

• **interop directive**: The `foreign-runtime-id` that is used if the implementation does not support any of the items in `preference-list` is implementation defined (see Section 2.15.1).

• **interop Construct**: The `foreign-runtime-id` values for the `prefer_type` clause that the implementation supports, including non-standard names compatible with this clause, and the default choice when the implementation supports multiple values are implementation defined (see Section 2.15.1).

• The concrete types of the values of interop properties for implementation defined `foreign-runtime-ids` are implementation defined (see Section 2.15.1).

• **atomic construct**: A compliant implementation may enforce exclusive access between `atomic` regions that update different storage locations. The circumstances under which this occurs are implementation defined. If the storage location designated by `x` is not size-aligned (that is, if the byte alignment of `x` is not a multiple of the size of `x`), then the behavior of the `atomic` region is implementation defined (see Section 2.19.7).
• **Data-sharing attributes**: The data-sharing attributes of dummy arguments without the `VALUE` attribute are implementation defined if the associated actual argument is shared, except for the conditions specified (see Section 2.21.1.2).

• **`threadprivate` directive**: If the conditions for values of data in the `threadprivate` objects of threads (other than an initial thread) to persist between two consecutive active parallel regions do not all hold, the allocation status of an allocatable variable in the second region is implementation defined (see Section 2.21.2).

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**Chapter 3:**

• **Runtime library definitions**: The enum types for `omp_allocator_handle_t`, `omp_event_handle_t`, `omp_interop_type_t` and `omp_memspace_handle_t` are implementation defined. The integral or pointer type for `omp_interop_t` is implementation defined (see Section 3.1).

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• **Runtime library definitions**: Whether the include file `omp_lib.h` or the module `omp_lib` (or both) is provided is implementation defined. Whether the `omp_lib.h` file provides derived-type definitions or those routines that require an explicit interface is implementation defined. Whether any of the OpenMP runtime library routines that take an argument are extended with a generic interface so arguments of different `KIND` type can be accommodated is implementation defined (see Section 3.1).

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• **`omp_set_num_threads` routine**: If the argument is not a positive integer the behavior is implementation defined (see Section 3.2.1).

• **`omp_set_schedule` routine**: For implementation-specific schedule kinds, the values and associated meanings of the second argument are implementation defined (see Section 3.2.11).

• **`omp_get_schedule` routine**: The value returned by the second argument is implementation defined for any schedule kinds other than `static`, `dynamic` and `guided` (see Section 3.2.12).

• **`omp_get_supported_active_levels` routine**: The number of active levels of parallelism supported by the implementation is implementation defined, but must be greater than 0 (see Section 3.2.14).

• **`omp_set_max_active_levels` routine**: If the argument is not a non-negative integer then the behavior is implementation defined (see Section 3.2.15).
• **omp_get_place_proc_ids routine**: The meaning of the non-negative numerical identifiers returned by the `omp_get_place_proc_ids` routine is implementation defined. The order of the numerical identifiers returned in the array `ids` is implementation defined (see Section 3.3.4).

• **omp_set_affinity_format routine**: When called from within any `parallel` or `teams` region, the binding thread set (and binding region, if required) for the `omp_set_affinity_format` region and the effect of this routine are implementation defined (see Section 3.3.8).

• **omp_get_affinity_format routine**: When called from within any `parallel` or `teams` region, the binding thread set (and binding region, if required) for the `omp_get_affinity_format` region is implementation defined (see Section 3.3.9).

• **omp_display_affinity routine**: If the `format` argument does not conform to the specified format then the result is implementation defined (see Section 3.3.10).

• **omp_capture_affinity routine**: If the `format` argument does not conform to the specified format then the result is implementation defined (see Section 3.3.11).

• **omp_set_num_teams routine**: If the argument is not evaluated to a positive integer the behavior of this routine is implementation defined (see Section 3.4.3).

• **omp_set_teams_thread_limit routine**: If the argument is not a positive integer the behavior is implementation defined (see Section 3.4.5).

• **omp_target_memcpy_rect routine**: The maximum number of dimensions supported is implementation defined, but must be at least three (see Section 3.8.6).

• **Lock routines**: If a lock contains a synchronization hint, the effect of the hint is implementation defined (see Section 3.9 and Section 3.9.2).

**Chapter 4:**

• **ompt_callback_sync_region_wait, ompt_callback_mutex_released, ompt_callback_dependencies, ompt_callback_task_dependence, ompt_callback_work, ompt_callback_master** (deprecated), ompt_callback_masked, ompt_callback_target_map, ompt_callback_target_map_emi, ompt_callback_sync_region, ompt_callback_reduction, ompt_callback_lock_init, ompt_callback_lock_destroy, ompt_callback_mutex_acquire, ompt_callback_mutex_acquired, ompt_callback_nest_lock, ompt_callback_flush, ompt_callback_cancel and **ompt_callback_dispatch tool callbacks**: If a tool attempts to register a callback with the string name using the runtime entry point `ompt_set_callback` (see Table 4.3), whether the registered callback may never, sometimes or always invoke this callback for the associated events is implementation defined (see Section 4.2.4).
• **Device tracing**: Whether a target device supports tracing or not is implementation defined; if a target device does not support tracing, a `NULL` may be supplied for the `lookup` function to the device initializer of a tool (see Section 4.2.5).

• **ompt_set_trace_ompt** and **ompt_buffer_get_record_ompt runtime entry points**: Whether a device-specific tracing interface will define this runtime entry point, indicating that it can collect traces in OMPT format is implementation defined. The kinds of trace records available for a device is implementation defined (see Section 4.2.5).

• **Native record abstract type**: The meaning of a `hwid` value for a device is implementation defined (see Section 4.4.3.3).

• **ompt_record_abstract_t type**: The set of OMPT thread states supported is implementation defined (see Section 4.4.27).

• **ompt_callback_target_data_op_t callback type**: Whether in some operations `src_addr` or `dest_addr` might point to an intermediate buffer is implementation defined (see Section 4.5.2.25).

• **ompt_set_callback_t entry point type**: The subset of the associated event in which the callback is invoked is implementation defined (see Section 4.6.1.3).

• **ompt_get_place_proc_ids_t entry point type**: The meaning of the numerical identifiers returned is implementation defined. The order of `ids` returned in the array is implementation defined (see Section 4.6.1.8).

• **ompt_get_partition_place_nums_t entry point type**: The order of the identifiers returned in the array `place_nums` is implementation defined (see Section 4.6.1.10).

• **ompt_get_proc_id_t entry point type**: The meaning of the numerical identifier returned is implementation defined (see Section 4.6.1.11).

**Chapter 5:**

• **ompd_callback_print_string_fn_t callback function**: The value of `category` is implementation defined (see Section 5.4.5).

• **ompd_parallel_handle_compare operation**: The means by which parallel region handles are ordered is implementation defined (see Section 5.5.6.5).

• **ompd_task_handle_compare operation**: The means by which task handles are ordered is implementation defined (see Section 5.5.7.6).

**Chapter 6:**

• **OMP_SCHEDULE environment variable**: If the value does not conform to the specified format then the behavior of the program is implementation defined (see Section 6.1).

• **OMP_NUM_THREADS environment variable**: If any value of the list specified leads to a number of threads that is greater than the implementation can support, or if any value is not a positive integer, then the behavior of the program is implementation defined (see Section 6.2).
• **OMP_DYNAMIC environment variable**: If the value is neither `true` nor `false` the behavior of the program is implementation defined (see Section 6.3).

• **OMP_PROC_BIND environment variable**: If the value is not `true`, `false`, or a comma separated list of `primary` (master has been deprecated), `close`, or `spread`, the behavior is implementation defined. The behavior is also implementation defined if an initial thread cannot be bound to the first place in the OpenMP place list. The thread affinity policy is implementation defined if the value is `true` (see Section 6.4).

• **OMP_PLACES environment variable**: The meaning of the numbers specified in the environment variable and how the numbering is done are implementation defined. The precise definitions of the abstract names are implementation defined. An implementation may add implementation-defined abstract names as appropriate for the target platform. When creating a place list of n elements by appending the number n to an abstract name, the determination of which resources to include in the place list is implementation defined. When requesting more resources than available, the length of the place list is also implementation defined. The behavior of the program is implementation defined when the execution environment cannot map a numerical value (either explicitly defined or implicitly derived from an interval) within the **OMP_PLACES** list to a processor on the target platform, or if it maps to an unavailable processor. The behavior is also implementation defined when the **OMP_PLACES** environment variable is defined using an abstract name (see Section 6.5).

• **OMP_STACKSIZE environment variable**: If the value does not conform to the specified format or the implementation cannot provide a stack of the specified size then the behavior is implementation defined (see Section 6.6).

• **OMP_WAIT_POLICY environment variable**: The details of the active and passive behaviors are implementation defined (see Section 6.7).

• **OMP_MAX_ACTIVE_LEVELS environment variable**: If the value is not a non-negative integer or is greater than the maximum number of nested active parallel levels that an implementation can support then the behavior of the program is implementation defined (see Section 6.8).

• **OMP_NESTED environment variable (deprecated)**: If the value is neither `true` nor `false` the behavior of the program is implementation defined (see Section 6.9).

• **Conflicting OMP_NESTED (deprecated) and OMP_MAX_ACTIVE_LEVELS environment variables**: If both environment variables are set, the value of **OMP_NESTED** is `false`, and the value of **OMP_MAX_ACTIVE_LEVELS** is greater than 1, the behavior is implementation defined (see Section 6.9).

• **OMP_THREAD_LIMIT environment variable**: If the requested value is greater than the number of threads an implementation can support, or if the value is not a positive integer, the behavior of the program is implementation defined (see Section 6.10).

• **OMP_DISPLAY_AFFINITY environment variable**: For all values of the environment variables other than `true` or `false`, the display action is implementation defined (see Section 6.13).
• **OMP_AFFINITY_FORMAT environment variable**: If the value does not conform to the specified format then the result is implementation defined (see Section 6.14).

• **OMP_TARGET_OFFLOAD environment variable**: The support of disabled is implementation defined (see Section 6.17).

• **OMP_TOOL_LIBRARIES environment variable**: Whether the value of the environment variable is case sensitive or insensitive is implementation defined (see Section 6.19).

• **OMP_TOOL_VERBOSE_INIT environment variable**: Support for logging to stdout or stderr is implementation defined. Whether the value of the environment variable is case sensitive when it is treated as a filename is implementation defined. The format and detail of the log is implementation defined (see Section 6.20).

• **OMP_DEBUG environment variable**: If the value is neither disabled nor enabled the behavior is implementation defined (see Section 6.21).

• **OMP_NUM_TEAMS environment variable**: If the value is not a positive integer or is greater than the number of teams that an implementation can support, the behavior of the program is implementation defined (see Section 6.23).

• **OMP_TEAMS_THREAD_LIMIT environment variable**: If the value is not a positive integer or is greater than the number of threads that an implementation can support, the behavior of the program is implementation defined (see Section 6.24).
B Features History

This appendix summarizes the major changes between OpenMP API versions since version 2.5.

B.1 Deprecated Features

The following features have been deprecated in Version 5.1.

- Cray pointer support was deprecated.
- The use of clauses supplied to the requires directive as context traits was deprecated.
- The master affinity policy was deprecated.
- The master construct and all combined and composite constructs of which it is a constituent construct were deprecated.
- The constant omp_atv_sequential was deprecated.
- In Fortran, specifying list items that are not of type C_PTR in a use_device_ptr or is_device_ptr clause was deprecated.
- The ompt_sync_region_barrier and ompt_sync_region_barrier_implicit values of the ompt_sync_region_t enum were deprecated.
- The ompt_state_waitBarrier and ompt_state_waitBarrier_implicit values of the ompt_state_t enum were deprecated.

The following features have been deprecated in Version 5.0.

- The nest-var ICV, the OMP_NESTED environment variable, and the omp_set_nested and omp_get_nested routines were deprecated.
- Lock hints were renamed to synchronization hints. The following lock hint type and constants were deprecated:
  - the C/C++ type omp_lock_hint_t and the Fortran kind omp_lock_hint_kind;
  - the constants omp_lock_hint_none, omp_lock_hint_uncontended, omp_lock_hint_contended, omp_lock_hint_nonspeculative, and omp_lock_hint_speculative.
B.2 Version 5.0 to 5.1 Differences

- Full support of C11, C++11, C++14, C++17, C++20 and Fortran 2008 was completed (see Section 1.7).

- Various changes throughout the specification were made to provide initial support of Fortran 2018 (see Section 1.7).

- The OpenMP directive syntax was extended to include C++ attribute specifiers (see Section 2.1).

- The `omp_all_memory` reserved locator was added (see Section 2.1), and the `depend` clause was extended to allow its use (see Section 2.19.11).

- The `target_device` trait set was added to the OpenMP Context (see Section 2.3.1), and the `target_device` selector set was added to context selectors (see Section 2.3.2).

- For C/C++, the declare variant directive was extended to support elision of preprocessed code and to allow enclosed function definitions to be interpreted as variant functions (see Section 2.3.5).

- The `declare variant` directive was extended with new clauses (`adjust_args` and `append_args`) that support adjustment of the interface between the original function and its variants (see Section 2.3.5).

- The `dispatch` construct was added to allow users to control when variant substitution happens and to define additional information that can be passed as arguments to the function variants (see Section 2.3.6).

- To support device-specific ICV settings the environment variable syntax was extended to support device-specific variables (see Section 2.4.2 and Section 6).

- The assume directive was added to allow users to specify invariants (see Section 2.5.2).

- To support clarity in metadirectives, the `nothing` directive was added (see Section 2.5.3).

- To allow users to control the compilation process and runtime error actions, the `error` directive was added (see Section 2.5.4).

- The `masked` construct was added to support restricting execution to a specific thread (see Section 2.8).

- The `scope` directive was added to support reductions without requiring a `parallel` or worksharing region (see Section 2.9).

- Loop transformation constructs were added (see Section 2.11.9).

- The `grainsize` and `num_tasks` clauses for the `taskloop` construct were extended with a `strict` modifier to ensure a deterministic distribution of logical iterations to tasks (see Section 2.12.2).
• Support for the `align` clause on the `allocate` directive and `allocator` and `align` modifiers on the `allocate` clause was added (see Section 2.13).

• The `thread_limit` clause was added to the `target` construct to control the upper bound on the number of threads in the created contention group (see Section 2.14.5).

• The `has_device_addr` clause was added to the `target` construct to allow access to variables or array sections that already have a device address (see Section 2.14.5).

• Support was added so that iterators may be defined and used in a motion clause on a `target update` directive (see Section 2.14.6) or in a `map` clause (see Section 2.21.7.1).

• Support was added for indirect calls to the device version of a procedure or function in `target` regions. (see Section 2.14.7).

• The `interop` directive was added to enable portable interoperability with foreign execution contexts used to implement OpenMP (see Section 2.15.1). Runtime routines that facilitate use of `omp_interop_t` objects were also added (see Section 3.12).

• The `nowait` clause was added to the `taskwait` directive to support insertion of non-blocking join operations in a task dependence graph (see Section 2.19.5).

• Support was added for compare-and-swap and (for C and C++) minimum and maximum atomic operations through the `compare` clause. Support was also added for the specification of the memory order to apply to a failed comparing atomic operation with the `fail` clause (see Section 2.19.7).

• Specification of the `seq_cst` clause on a `flush` construct was allowed, with the same meaning as a `flush` construct without a list and without a clause (see Section 2.19.8).

• To support inout sets, the `inoutset` argument was added to the `depend` clause (see Section 2.19.11).

• Support for `private` and `firstprivate` as an argument to the `default` clause in C and C++ was added (see Section 2.21.4.1).

• The `present` argument was added to the `defaultmap` clause (see Section 2.21.7.3).

• The `omp_set_num_teams` and `omp_set_teams_thread_limit` runtime routines were added to control the number of teams and the size of those teams on the `teams` construct (see Section 3.4.3 and Section 3.4.5). Additionally, the `omp_get_max_teams` and `omp_get_teams_thread_limit` runtime routines were added to retrieve the values that will be used in the next `teams` construct (see Section 3.4.4 and Section 3.4.6).

• The `omp_target_is_accessible` runtime routine was added to test whether host memory is accessible from a given device (see Section 3.8.4).

• To support asynchronous device memory management, `omp_target_memcpy_async` and `omp_target_memcpy_rect_async` runtime routines were added (see Section 3.8.7 and Section 3.8.8).
• The `omp_get_mapped_ptr` runtime routine was added to support obtaining the device pointer that is associated with a host pointer for a given device (see Section 3.8.11).

• The `omp_calloc`, `omp_realloc`, `omp_aligned_alloc` and `omp_aligned_calloc` API routines were added (see Section 3.13).

• For the `omp_alloctrait_key_t` enum, the `omp_atv_serialized` value was added and the `omp_atv_default` value was changed (see Section 3.13.1).

• The `omp_display_env` runtime routine was added to provide information about ICVs and settings of environment variables (see Section 3.15).

• The `ompt_scope_beginend` value was added to the `ompt_scope_endpoint_t` enum to indicate the coincident beginning and end of a scope (see Section 4.4.4.11).

• The `ompt_sync_region_barrier_implicit_workshare`, `ompt_sync_region_barrier_implicit_parallel` and `ompt_sync_region_barrier_teams` values were added to the `ompt_sync_region_t` enum (see Section 4.4.4.13).

• Values for asynchronous data transfers were added to the `ompt_target_data_op_t` enum (see Section 4.4.4.14).

• The `ompt_state_wait_barrier_implementation` and `ompt_state_wait_barrier_teams` values were added to the `ompt_state_t` enum (see Section 4.4.4.27).

• The `ompt_callback_target_data_op_emi_t`, `ompt_callback_target_emi_t`, `ompt_callback_target_map_emi_t` and `ompt_callback_target_submit_emi_t` callbacks were added to support external monitoring interfaces (see Section 4.5.2.25, Section 4.5.2.26, Section 4.5.2.27 and Section 4.5.2.28).

• The `ompt_callback_error_t` type was added (see Section 4.5.2.30).

• The `OMP_PLACES` syntax was extended (see Section 6.5).

• The `OMP_NUM_TEAMS` and `OMP_TEAMS_THREAD_LIMIT` environment variables were added to control the number and size of teams on the `teams` construct (see Section 6.23 and Section 6.24).

**B.3 Version 4.5 to 5.0 Differences**

• The memory model was extended to distinguish different types of flush operations according to specified flush properties (see Section 1.4.4) and to define a happens before order based on synchronizing flush operations (see Section 1.4.5).
• Various changes throughout the specification were made to provide initial support of C11, C++11, C++14, C++17 and Fortran 2008 (see Section 1.7).

• Full support of Fortran 2003 was completed (see Section 1.7).

• Support for array shaping (see Section 2.1.4) and for array sections with non-unit strides in C and C++ (see Section 2.1.5) was added to facilitate specification of discontiguous storage and the target update construct (see Section 2.14.6) and the depend clause (see Section 2.19.11) were extended to allow the use of shape-operators (see Section 2.1.4).

• Iterators (see Section 2.1.6) were added to support expressions in a list that expand to multiple expressions.

• The metadirective directive (see Section 2.3.4) and declare variant directive (see Section 2.3.5) were added to support selection of directive variants and declared function variants at a call site, respectively, based on compile-time traits of the enclosing context.

• The target-offload-var internal control variable (see Section 2.4) and the OMP_TARGET_OFFLOAD environment variable (see Section 6.17) were added to support runtime control of the execution of device constructs.

• Control over whether nested parallelism is enabled or disabled was integrated into the max-active-levels-var internal control variable (see Section 2.4.2), the default value of which is now implementation defined, unless determined according to the values of the OMP_NUM_THREADS (see Section 6.2) or OMP_PROC_BIND (see Section 6.4) environment variables.

• The requires directive (see Section 2.5.1) was added to support applications that require implementation-specific features.

• The teams construct (see Section 2.7) was extended to support execution on the host device without an enclosing target construct (see Section 2.14.5).

• The canonical loop form was defined for Fortran and, for all base languages, extended to permit non-rectangular loop nests (see Section 2.11.1).

• The relational-op in the canonical loop form for C/C++ was extended to include != (see Section 2.11.1).

• The default loop schedule modifier for worksharing-loop constructs without the static schedule and the ordered clause was changed to nonmonotonic (see Section 2.11.4).

• The collapse of associated loops that are imperfectly nested loops was defined for the worksharing-loop (see Section 2.11.4), simd (see Section 2.11.5.1), taskloop (see Section 2.12.2) and distribute (see Section 2.11.6.2) constructs.

• The simd construct (see Section 2.11.5.1) was extended to accept the if, nontemporal and order(concurrent) clauses and to allow the use of atomic constructs within it.
• The `loop` construct and the `order(concurrent)` clause were added to support compiler optimization and parallelization of loops for which iterations may execute in any order, including concurrently (see Section 2.11.7).

• The `scan` directive (see Section 2.11.8) and the `inscan` modifier for the `reduction` clause (see Section 2.21.5.4) were added to support inclusive and exclusive scan computations.

• To support task reductions, the `task` (see Section 2.12.1) and `target` (see Section 2.14.5) constructs were extended to accept the `in_reduction` clause (see Section 2.21.5.6), the `taskgroup` construct (see Section 2.19.6) was extended to accept the `task_reduction` clause Section 2.21.5.5), and the `task` modifier was added to the `reduction` clause (see Section 2.21.5.4).

• The `affinity` clause was added to the `task` construct (see Section 2.12.1) to support hints that indicate data affinity of explicit tasks.

• The `detach` clause for the `task` construct (see Section 2.12.1) and the `omp_fulfill_event` runtime routine (see Section 3.11.1) were added to support execution of detachable tasks.

• To support taskloop reductions, the `taskloop` (see Section 2.12.2) and `taskloop simd` (see Section 2.12.3) constructs were extended to accept the `reduction` (see Section 2.21.5.4) and `in_reduction` (see Section 2.21.5.6) clauses.

• The `taskloop` construct (see Section 2.12.2) was added to the list of constructs that can be canceled by the `cancel` construct (see Section 2.20.1)).

• To support mutually exclusive inout sets, a `mutexinoutset` dependence-type was added to the `depend` clause (see Section 2.12.6 and Section 2.19.11).

• Predefined memory spaces (see Section 2.13.1), predefined memory allocators and allocator traits (see Section 2.13.2) and directives, clauses (see Section 2.13 and API routines (see Section 3.13) to use them were added to support different kinds of memories.

• The semantics of the `use_device_ptr` clause for pointer variables was clarified and the `use_device_addr` clause for using the device address of non-pointer variables inside the `target data` construct was added (see Section 2.14.2).

• To support reverse offload, the `ancestor` modifier was added to the `device` clause for `target` constructs (see Section 2.14.5).

• To reduce programmer effort implicit declare target directives for some functions (C, C++, Fortran) and subroutines (Fortran) were added (see Section 2.14.5 and Section 2.14.7).

• The `target update` construct (see Section 2.14.6) was modified to allow array sections that specify discontiguous storage.

• The `to` and `from` clauses on the `target update` construct (see Section 2.14.6), the `depend` clause on task generating constructs (see Section 2.19.11), and the `map` clause (see Section 2.21.7.1) were extended to allow any lvalue expression as a list item for C/C++.
- Support for nested `declare target` directives was added (see Section 2.14.7).
- New combined constructs `master taskloop` (see Section 2.16.7), `parallel master` (see Section 2.16.6), `parallel master taskloop` (see Section 2.16.9), `master taskloop simd` (see Section 2.16.8), `parallel master taskloop simd` (see Section 2.16.10) were added.
- The `depend` clause was added to the `taskwait` construct (see Section 2.19.5).
- To support acquire and release semantics with weak memory ordering, the `acq_rel`, `acquire`, and `release` clauses were added to the `atomic` construct (see Section 2.19.7) and `flush` construct (see Section 2.19.8), and the memory ordering semantics of implicit flushes on various constructs and runtime routines were clarified (see Section 2.19.8.1).
- The `atomic` construct was extended with the `hint` clause (see Section 2.19.7).
- The `depend` clause (see Section 2.19.11) was extended to support iterators and to support depend objects that can be created with the new `depobj` construct.
- Lock hints were renamed to synchronization hints, and the old names were deprecated (see Section 2.19.12).
- To support conditional assignment to lastprivate variables, the `conditional` modifier was added to the `lastprivate` clause (see Section 2.21.4.5).
- The description of the `map` clause was modified to clarify the mapping order when multiple `map-types` are specified for a variable or structure members of a variable on the same construct. The `close map-type-modifier` was added as a hint for the runtime to allocate memory close to the target device (see Section 2.21.7.1).
- The capability to map C/C++ pointer variables and to assign the address of device memory that is mapped by an array section to them was added. Support for mapping of Fortran pointer and allocatable variables, including pointer and allocatable components of variables, was added (see Section 2.21.7.1).
- The `defaultmap` clause (see Section 2.21.7.3) was extended to allow selecting the data-mapping or data-sharing attributes for any of the scalar, aggregate, pointer or allocatable classes on a per-region basis. Additionally it accepts the `none` parameter to support the requirement that all variables referenced in the construct must be explicitly mapped or privatized.
- The `declare mapper` directive was added to support mapping of data types with direct and indirect members (see Section 2.21.7.4).
- The `omp_set_nested` (see Section 3.2.9) and `omp_get_nested` (see Section 3.2.10) routines and the `OMP_NESTED` environment variable (see Section 6.9) were deprecated.
- The `omp_get_supported_active_levels` routine was added to query the number of active levels of parallelism supported by the implementation (see Section 3.2.14).
• Runtime routines `omp_set_affinity_format` (see Section 3.3.8),
  `omp_get_affinity_format` (see Section 3.3.9), `omp_set_affinity` (see Section 3.3.10), and `omp_capture_affinity` (see Section 3.3.11) and environment variables `OMP_DISPLAY_AFFINITY` (see Section 6.13) and `OMP_AFFINITY_FORMAT` (see Section 6.14) were added to provide OpenMP runtime thread affinity information.

• The `omp_pause_resource` and `omp_pause_resource_all` runtime routines were added to allow the runtime to relinquish resources used by OpenMP (see Section 3.6.1 and Section 3.6.2).

• The `omp_get_device_num` runtime routine (see Section 3.7.5) was added to support determination of the device on which a thread is executing.

• Support for a first-party tool interface (see Section 4) was added.

• Support for a third-party tool interface (see Section 5) was added.

• Support for controlling offloading behavior with the `OMP_TARGET_OFFLOAD` environment variable was added (see Section 6.17).

• Stubs for Runtime Library Routines (previously Appendix A) were moved to a separate document.

• Interface Declarations (previously Appendix B) were moved to a separate document.

**B.4 Version 4.0 to 4.5 Differences**

• Support for several features of Fortran 2003 was added (see Section 1.7).

• A parameter was added to the `ordered` clause of the worksharing-loop construct (see Section 2.11.4) and clauses were added to the `ordered` construct (see Section 2.19.9) to support doacross loop nests and use of the `simd` construct on loops with loop-carried backward dependences.

• The `linear` clause was added to the worksharing-loop construct (see Section 2.11.4).

• The `simdlen` clause was added to the `simd` construct (see Section 2.11.5.1) to support specification of the exact number of iterations desired per SIMD chunk.

• The `priority` clause was added to the `task` construct (see Section 2.12.1) to support hints that specify the relative execution priority of explicit tasks. The `omp_get_max_task_priority` routine was added to return the maximum supported priority value (see Section 3.5.1) and the `OMP_MAX_TASK_PRIORITY` environment variable was added to control the maximum priority value allowed (see Section 6.16).

• Taskloop constructs (see Section 2.12.2 and Section 2.12.3) were added to support nestable parallel loops that create OpenMP tasks.
• To support interaction with native device implementations, the `use_device_ptr` clause was added to the `target data` construct (see Section 2.14.2) and the `is_device_ptr` clause was added to the `target` construct (see Section 2.14.5).

• To support unstructured data mapping for devices, the `target enter data` (see Section 2.14.3) and `target exit data` (see Section 2.14.4) constructs were added and the `map` clause (see Section 2.21.7.1) was updated.

• The `nowait` and `depend` clauses were added to the `target` construct (see Section 2.14.5) to improve support for asynchronous execution of `target` regions.

• The `private, firstprivate` and `defaultmap` clauses were added to the `target` construct (see Section 2.14.5).

• The `declare target` directive was extended to allow mapping of global variables to be deferred to specific device executions and to allow an `extended-list` to be specified in C/C++ (see Section 2.14.7).

• To support a more complete set of device construct shortcuts, the `target parallel` (see Section 2.16.16), target parallel worksharing-loop (see Section 2.16.17), target parallel worksharing-loop SIMD (see Section 2.16.18), and `target simd` (see Section 2.16.20), combined constructs were added.

• The `if` clause was extended to take a `directive-name-modifier` that allows it to apply to combined constructs (see Section 2.18).

• The `hint` clause was added to the `critical` construct (see Section 2.19.1).

• The `source` and `sink` dependence types were added to the `depend` clause (see Section 2.19.11) to support doacross loop nests.

• The implicit data-sharing attribute for scalar variables in `target` regions was changed to `firstprivate` (see Section 2.21.1.1).

• Use of some C++ reference types was allowed in some data sharing attribute clauses (see Section 2.21.4).

• The `ref, val, and uval` modifiers were added to the `linear` clause (see Section 2.21.4.6).

• Semantics for reductions on C/C++ array sections were added and restrictions on the use of arrays and pointers in reductions were removed (see Section 2.21.5.4).

• Support was added to the map clauses to handle structure elements (see Section 2.21.7.1).

• Query functions for OpenMP thread affinity were added (see Section 3.3.2 to Section 3.3.7).

• Device memory routines were added to allow explicit allocation, deallocation, memory transfers and memory associations (see Section 3.8).

• The lock API was extended with lock routines that support storing a hint with a lock to select a desired lock implementation for a lock’s intended usage by the application code (see...
Section 3.9.2).

- C/C++ Grammar (previously Appendix B) was moved to a separate document.

### B.5 Version 3.1 to 4.0 Differences

- Various changes throughout the specification were made to provide initial support of Fortran 2003 (see Section 1.7).

- C/C++ array syntax was extended to support array sections (see Section 2.1.5).

- The `proc_bind` clause (see Section 2.6.2), the `OMP_PLACES` environment variable (see Section 6.5), and the `omp_get_proc_bind` runtime routine (see Section 3.3.1) were added to support thread affinity policies.

- SIMD directives were added to support SIMD parallelism (see Section 2.11.5).

- Implementation defined task scheduling points for untied tasks were removed (see Section 2.12.6).

- Device directives (see Section 2.14), the `OMP_DEFAULT_DEVICE` environment variable (see Section 6.15), and the `omp_set_default_device`, `omp_get_default_device`, `omp_get_num_devices`, `omp_get_num_teams`, `omp_get_team_num`, and `omp_is_initial_device` routines were added to support execution on devices.

- The `taskgroup` construct (see Section 2.19.6) was added to support more flexible deep task synchronization.

- The `atomic` construct (see Section 2.19.7) was extended to support atomic swap with the `capture` clause, to allow new atomic update and capture forms, and to support sequentially consistent atomic operations with a new `seq_cst` clause.

- The `depend` clause (see Section 2.19.11) was added to support task dependences.

- The `cancel` construct (see Section 2.20.1), the `cancellation point` construct (see Section 2.20.2), the `omp_get_cancellation` runtime routine (see Section 3.2.8) and the `OMP_CANCELLATION` environment variable (see Section 6.11) were added to support the concept of cancellation.

- The `reduction` clause (see Section 2.21.5.4) was extended and the `declare reduction` construct (see Section 2.21.5.7) was added to support user defined reductions.

- The `OMP_DISPLAY_ENV` environment variable (see Section 6.12) was added to display the value of ICVs associated with the OpenMP environment variables.

- Examples (previously Appendix A) were moved to a separate document.
B.6 Version 3.0 to 3.1 Differences

- The `bind-var` ICV (see Section 2.4.1) and the `OMP_PROC_BIND` environment variable (see Section 6.4) were added to support control of whether threads are bound to processors.

- The `nthreads-var` ICV was modified to be a list of the number of threads to use at each nested parallel region level and the algorithm for determining the number of threads used in a parallel region was modified to handle a list (see Section 2.6.1).

- The `final` and `mergeable` clauses (see Section 2.12.1) were added to the `task` construct to support optimization of task data environments.

- The `taskyield` construct (see Section 2.12.4) was added to allow user-defined task scheduling points.

- The `atomic` construct (see Section 2.19.7) was extended to include `read`, `write`, and `capture` forms, and an `update` clause was added to apply the already existing form of the `atomic` construct.

- Data environment restrictions were changed to allow `intent(in)` and `const`-qualified types for the `firstprivate` clause (see Section 2.21.4.4).

- Data environment restrictions were changed to allow Fortran pointers in `firstprivate` (see Section 2.21.4.4) and `lastprivate` (see Section 2.21.4.5).

- New reduction operators `min` and `max` were added for C and C++ (see Section 2.21.5).

- The nesting restrictions in Section 2.22 were clarified to disallow closely-nested OpenMP regions within an `atomic` region so that an `atomic` region can be consistently defined with other OpenMP regions to include all code in the `atomic` construct.

- The `omp_in_final` runtime library routine (see Section 3.5.2) was added to support specialization of final task regions.

- Descriptions of examples (previously Appendix A) were expanded and clarified.

- Incorrect use of `omp_integer_kind` in Fortran interfaces was replaced with `selected_int_kind(8)`.

B.7 Version 2.5 to 3.0 Differences

- The definition of active `parallel` region was changed so that a `parallel` region is active if it is executed by a team that consists of more than one thread (see Section 1.2.2).

- The concept of tasks was added to the OpenMP execution model (see Section 1.2.5 and Section 1.3).

- The OpenMP memory model was extended to cover atomicity of memory accesses (see Section 1.4.1). The description of the behavior of `volatile` in terms of `flush` was removed.
• The definition of the nest-var, dyn-var, nthreads-var and run-sched-var internal control variables (ICVs) were modified to provide one copy of these ICVs per task instead of one copy for the whole program (see Section 2.4). The omp_set_num_threads, omp_set_nested and omp_set_dynamic runtime library routines were specified to support their use from inside a parallel region (see Section 3.2.1, Section 3.2.6 and Section 3.2.9).

• The thread-limit-var ICV, the omp_get_thread_limit runtime library routine and the OMP_THREAD_LIMIT environment variable were added to support control of the maximum number of threads that participate in the OpenMP program (see Section 2.4.1, Section 3.2.13 and Section 6.10).

• The max-active-levels-var ICV, the omp_set_max_active_levels and omp_get_max_active_levels runtime library routine and the OMP_MAX_ACTIVE_LEVELS environment variable and were added to support control of the number of nested active parallel regions (see Section 2.4.1, Section 3.2.15, Section 3.2.16 and Section 6.8).

• The stacksize-var ICV and the OMP_STACKSIZE environment variable were added to support control of the stack size for threads that the OpenMP implementation creates (see Section 2.4.1 and Section 6.6).

• The wait-policy-var ICV and the OMP_WAIT_POLICY environment variable were added to control the desired behavior of waiting threads (see Section 2.4.1 and Section 6.7).

• The rules for determining the number of threads used in a parallel region were modified (see Section 2.6.1).

• The assignment of iterations to threads in a loop construct with a static schedule kind was made deterministic (see Section 2.11.4).

• The worksharing-loop construct was extended to support association with more than one perfectly nested loop through the collapse clause (see Section 2.11.4).

• Iteration variables for worksharing-loops were allowed to be random access iterators or of unsigned integer type (see Section 2.11.4).

• The schedule kind auto was added to allow the implementation to choose any possible mapping of iterations in a loop construct to threads in the team (see Section 2.11.4).

• The task construct (see Section 2.12) was added to support explicit tasks.

• The taskwait construct (see Section 2.19.5) was added to support task synchronization.

• Predetermined data-sharing attributes were defined for Fortran assumed-size arrays (see Section 2.21.1.1).

• Static class members variables were allowed to appear in a threadprivate directive (see Section 2.21.2).
Invocations of constructors and destructors for private and threadprivate class type variables was clarified (see Section 2.21.2, Section 2.21.4.3, Section 2.21.4.4, Section 2.21.6.1 and Section 2.21.6.2).

- The use of Fortran allocatable arrays was allowed in `private`, `firstprivate`, `lastprivate`, `reduction`, `copyin` and `copyprivate` clauses (see Section 2.21.2, Section 2.21.4.3, Section 2.21.4.4, Section 2.21.4.5, Section 2.21.5.4, Section 2.21.6.1 and Section 2.21.6.2).

- The `firstprivate` argument was added for the `default` clause in Fortran (see Section 2.21.4.1).

- Implementations were precluded from using the storage of the original list item to hold the new list item on the primary thread for list items in the `private` clause and the value was made well defined on exit from the `parallel` region if no attempt is made to reference the original list item inside the `parallel` region (see Section 2.21.4.3).

- The runtime library routines `omp_set_schedule` and `omp_get_schedule` were added to set and to retrieve the value of the `run-sched-var` ICV (see Section 3.2.11 and Section 3.2.12).

- The `omp_get_level` runtime library routine was added to return the number of nested `parallel` regions that enclose the task that contains the call (see Section 3.2.17).

- The `omp_get_ancestor_thread_num` runtime library routine was added to return the thread number of the ancestor for a given nested level of the current thread, (see Section 3.2.18).

- The `omp_get_team_size` runtime library routine was added to return the size of the thread team to which the ancestor belongs for a given nested level of the current thread, (see Section 3.2.19).

- The `omp_get_active_level` runtime library routine was added to return the number of nested active `parallel` regions that enclose the task that contains the call (see Section 3.2.20).

- Lock ownership was defined in terms of tasks instead of threads (see Section 3.9).
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